

Last time:

- proving equivalences of multiple statements
- existence proofs

eg For any $a, b \in \mathbb{Z}_{>0}$, $\exists k, l \in \mathbb{Z}$ s.t. $\gcd(a, b) = ak + bl$.

We proved: the smallest positive integer, d , in the set

$$D = \{ax + by : x, y \in \mathbb{Z}\}, \text{ must divide } a \text{ and divide } b.$$

Today:

- finishing the proof, and the Euclidean algorithm.

- Constructive vs. nonconstructive proofs

1. Finishing the example.

Need: $\forall a, b \in \mathbb{Z}_{>0}, \exists k, l$ s.t. $\gcd(a, b) = ak + bl$.

Shown: $d := \min \{ ax + by : x, y \in \mathbb{Z}, ax + by > 0 \}$ divides both a and b .

To finish the proof, it suffices to show that $\gcd(a, b) = d$.

(1) " $d \in \gcd$ ": Since $d | a$ and $d | b$, we have $d \rightarrow$ a common divisor of a and b ,
(c.d.)
therefore $d \in \gcd(a, b)$.

(2) " $\gcd \leq d$ ": We have $\gcd(a, b) | a$ and $\gcd(a, b) | b$, so
 $\gcd(a, b) | ax + by \quad \forall x, y \in \mathbb{Z}$.

therefore $\gcd(a, b) | d$ since d is of the form $ax + by$ for some $x, y \in \mathbb{Z}$.

Thus, we have $\gcd(a, b) \leq d$.

By (1) and (2), we have $\gcd(a, b) = d$. so we are done.

Take-away: The gcd of two pos. int. equals the smallest positive integral lin. comb. of those two integers. ($\text{gcd}(a,b) = ak + bl$)

Aside*: The Euclidean algorithm (for getting "k" and "l".)

Examples: $a = \underline{43}$, $b = \underline{13}$.

$$43 = \underline{13} \cdot 3 + \underline{4}$$

$$13 = \underline{4} \cdot 3 + \underline{1} \rightarrow \text{the gcd}$$

$$4 = \underline{4} \cdot 1 + \underline{0} \rightarrow \text{halt when 0 appears}$$

$a = \underline{210}$, $b = \underline{45}$.

$$210 = \underline{45} \cdot 4 + \underline{30}$$

$$45 = \underline{30} \cdot 1 + \underline{15} \rightarrow \text{the gcd.}$$

$$30 = \underline{15} \cdot 2 + \underline{0} \rightarrow \text{halt at 0}$$

General algorithm: WLOG, assume $a > b$. Starting with $(x_0, y_0) = (a, b)$, keep dividing x_i by y_i and setting $(x_{i+1}, y_{i+1}) = (y_i, \text{the remainder } r \text{ in } x_i = y_i \cdot q_i + r)$ until $r = \underline{0}$.

Fact: The last nonzero "r" is the gcd.

The algorithm can also help us find "k" and "l" :

$$a = \underline{\underline{210}}, \quad b = \underline{\underline{45}}.$$

$$210 = \underline{\underline{45}} \cdot 4 + \underline{\underline{30}}$$

$$45 = \underline{\underline{30}} \cdot 1 + \underline{\underline{15}}$$

$$30 = \underline{\underline{15}} \cdot 2 + \underline{\underline{0}}$$

$$\gcd(a, b) = \underline{\underline{15}} = \underline{\underline{45}} - \underline{\underline{30}} \cdot 1$$

$$= 45 - (\underline{\underline{210}} - \underline{\underline{45}} \cdot 4)$$

$$= -\underline{\underline{210}} + \underline{\underline{45}} \cdot 5$$

$$= (-1) \cdot 210 + 5 \cdot 45. \quad \checkmark$$

EX. Try the Euclidean algorithm and find k, l for $a = 270$ and $b = 192$.

2. Constructive vs. nonconstructive proofs

$$\text{Recall: } ((x^a)^b)^c = x^{abc}$$

We'll prove the following prop. in two ways.

Prop. There exist irrational numbers x, y s.t. x^y is rational.

Pf 1 (constructive): Take $x = \sqrt{2}$ and $y = \log_2 9$. We know x is irrational.

We claim that y is also irrational and x^y is rational, so we'd be done.

$$(b) \checkmark: x^y = \sqrt{2}^{\log_2 9} \stackrel{(a)}{=} \sqrt{2}^{\log_2(3^2)} = (\sqrt{2})^{\log_2(3^2)} \stackrel{(b)}{=} \left(2^{\frac{1}{2}}\right)^{\log_2(3^2)} = \left(2^{\frac{1}{2} \cdot 2}\right)^{\log_2 3} = 2^{\log_2 3} = 3 \in \mathbb{Q}.$$

(a) \checkmark : We prove y is irr. by contradiction: Suppose $y = \log_2 9$ is rational. Then $\log_2 9 = \frac{m}{n}$ for some $m, n \in \mathbb{Z}_{>0}$. Thus, we have $(2)^{m/n} = 9$, so

$2^m = (2^{m/n})^n = 9^n$. This is impossible since 2^m is even while 9^n is odd. \square

Pf 2 (non-constructive): Consider $x = \sqrt{2}$, $y = \sqrt{2}$, so that $x^y = \sqrt{2}^{\sqrt{2}}$ and $x \cdot y \notin \mathbb{Q}$.

If $\sqrt{2}^{\sqrt{2}} \in \mathbb{Q}$, we are done.

If $\sqrt{2}^{\sqrt{2}} \notin \mathbb{Q}$, then we may take $x' = \sqrt{2}^{\sqrt{2}} \notin \mathbb{Q}$ and $y' = \sqrt{2} \notin \mathbb{Q}$,

and we have $(x')^{y'} = (\sqrt{2}^{\sqrt{2}})^{\sqrt{2}} = (\sqrt{2})^{\sqrt{2} \cdot \sqrt{2}} = \sqrt{2}^2 = 2 \in \mathbb{Q}$.

and we are again done. \square

Next time: more proof examples