RSA decoling.

(A) Recall the setup:

(1) Theorem RSA_1 : If $a,b \in \mathbb{N}$ are coprime (gcd(a,b)=1), then there exist $x,y \in \mathbb{N}$ with

ax-by=1.

Example: 570 and IIII are coprime (since 570 = 2.3.5.19 and IIII = 11.101); note that

570 · 115 - 1111 · 19 = 1.

(a) Theorem RSA2: Let m = pq where p, q are distinct primes; define p(m) = (p-1)(q-1). Then for any $a \in \mathbb{Z}$ with gcd(a, m) = 1, we have

 $a = 1 \pmod{m}$.

Example: Let a = 570 and m = 1111 = 11.101. We've seen that gcd(a,m) = 1. Since g(m) = (11-1)(101-1) = 1000, Theorem RSA2 says $570 = 1 \pmod{1111}$.

(B) Decodable encoding.

To assure a message n can be encoded in a decodable way:

- (1) Choose two (large) primes p and q; let m = pq.
 - (2) Make sure that n<m and gcd(n,m)=1.
 - (3) Choose an exponent k with gcd(k, p(m))=1.
 - (4) Compute n (mod m).

Remark: k and m (but not the factorization m=pql can be shared, so anyone can encode.

(C) Decoding.

You're given a message $b = n^k \pmod{m}$, coded as above. To decade it:

(1) Find x, y ∈ /N with

 $k \times - \varphi(m) y = 1.$ (*)

(possible by Theorem RSA1).

(a) Compute b (mod m) (by successive squarmal). The result is the original message n!

 $\begin{array}{c}
|Proof: \\
b^{\times} = (n^{k})^{\times} \equiv n^{k\times} \equiv n^{l+\varrho(m)\gamma} \\
\equiv n^{1}(n^{\varrho(m)})^{\gamma} \\
\equiv n \cdot 1^{\gamma} \equiv n \pmod{m}.
\end{array}$

by Thm. RSA2

(None of this works if you don't know p and q = without them, you don't know p(m), so you can't find x.)

(1) Example.

Decode the coded message b = 33, with k = 7 and m = 35 = 5.7.

Solution.

Since k=7 and p(m)=4.6=a4, we have gcd(k,p(m))=1. Note that

7.7-24.2=1, 1.7.7.24.2=1, 1.7.7.24.2=1, 1.7.7.24.2=1, 1.7.7.24.2=1, 1.7.7.24.2=1,

so we need to compute b (mod m), which is 33 (mod 35), by successive squaring:

7 = 4+2+1

 $33^{1} = 33 \pmod{35}$ $33^{2} = 1089 = 35 \cdot 31 + 4 = 4 \pmod{35}$ $33^{4} = (33^{2})^{2} = 4^{2} = 16 \pmod{35}$

 $33^{7} = 33^{4+2+1}$ $= 33^{4} \cdot 33^{2} \cdot 33^{1} = 16 \cdot 4 \cdot 33$ $= 64 \cdot 33 = 29 \cdot 33$ $= 957 = 35 \cdot 27 + 12 = 12 \pmod{35}$.

Compare with the last example from the class of 8/27.