

1

11	12	6	10	9	8	5	6	2	4	1	8	3	7	2	3			
([([S ₁	→	(¬¬S ₂	→	S ₂)]	→	¬	¬	S ₀]	→	S ₃]	→	S ₄)	→	[(S ₀	→	S ₃)	→	S ₄]
0		1		1	0	1	0	0	0	1	0	0	0	1	0	0	0	

The implication on the far left is inconsistently assigned, so the formula is a tautology.

2 A formula built up from \neg and sentential variables must have the form $\neg \dots \neg S_i$ for some i . Such a formula is either equivalent to S_i (for an even number of \neg 's), or to $\neg S_i$ (for an odd number). Suppose that $S_0 \rightarrow S_0$ is equivalent to $\neg \dots \neg S_i$.

Case 1. An even number of \neg 's before S_i . Assigning 0 to all S_j 's gives the value 1 to $S_0 \rightarrow S_0$ and 0 to $\neg \dots \neg S_i$, contradiction.

Case 2. An odd number of \neg 's before S_i . Assigning 1 to all S_j 's gives the value 1 to $S_0 \rightarrow S_0$ and 0 to $\neg \dots \neg S_i$, contradiction.

3 We need several lemmas.

Lemma 1. For any formula φ , $\vdash \neg \neg \varphi \rightarrow \varphi$.

Proof.

$\{\neg \neg \varphi\} \vdash \neg \neg \neg \neg \varphi \rightarrow \neg \neg \varphi$	using axiom (1)
$\{\neg \neg \varphi\} \vdash \neg \varphi \rightarrow \neg \neg \neg \varphi$	using axiom (3)
$\{\neg \neg \varphi\} \vdash \neg \neg \varphi \rightarrow \varphi$	using axiom (3)
$\{\neg \neg \varphi\} \vdash \varphi$	
$\vdash \neg \neg \varphi \rightarrow \varphi$	□

Lemma 2. For any formula φ , $\vdash \varphi \rightarrow \neg \neg \varphi$.

Proof.

$\{\varphi, \neg \neg \varphi\} \vdash \neg \varphi$	using Lemma 1
$\{\varphi\} \vdash \neg \neg \neg \varphi \rightarrow \neg \varphi$	
$\{\varphi\} \vdash \varphi \rightarrow \neg \neg \varphi$	using axiom (3)
$\{\varphi\} \vdash \neg \neg \varphi$	
$\vdash \varphi \rightarrow \neg \neg \varphi$	

Lemma 3. For any formulas φ, ψ , $\vdash (\varphi \rightarrow \neg \psi) \rightarrow (\psi \rightarrow \neg \varphi)$.

$\{\varphi \rightarrow \neg \psi, \psi\} \vdash \neg \neg \varphi \rightarrow \varphi$	using Lemma 1
$\{\varphi \rightarrow \neg \psi, \psi\} \vdash \neg \neg \varphi \rightarrow \neg \psi$	using 1.14
$\{\varphi \rightarrow \neg \psi, \psi\} \vdash \psi \rightarrow \neg \varphi$	using axiom (3)
$\{\varphi \rightarrow \neg \psi, \psi\} \vdash \neg \varphi$	
$\{\varphi \rightarrow \neg \psi\} \vdash \psi \rightarrow \neg \varphi$	
$\vdash (\varphi \rightarrow \neg \psi) \rightarrow (\psi \rightarrow \neg \varphi)$	

Lemma 4. For any formulas φ, ψ , $\vdash \varphi \rightarrow (\psi \rightarrow (\varphi \wedge \psi))$.

Proof.

$$\begin{array}{ll}
\{\varphi, \psi, \varphi \rightarrow \neg\psi\} \vdash \psi & \\
\{\varphi, \psi, \varphi \rightarrow \neg\psi\} \vdash \neg\psi & \\
\{\varphi, \psi, \varphi \rightarrow \neg\psi\} \vdash \neg(\varphi \rightarrow \varphi) & \text{using 1.13} \\
\{\varphi, \psi\} \vdash (\varphi \rightarrow \neg\psi) \rightarrow \neg(\varphi \rightarrow \varphi) & \\
\{\varphi, \psi\} \vdash (\varphi \rightarrow \varphi) \rightarrow \neg(\varphi \rightarrow \neg\psi) & \text{using Lemma 3} \\
\{\varphi, \psi\} \vdash \neg(\varphi \rightarrow \neg\psi) & \text{using 1.11} \\
\vdash \varphi \rightarrow (\psi \rightarrow (\varphi \wedge \psi)) &
\end{array}$$

Now the exercise follows from Lemmas 1, 2, and 4.

4

$$\neg S_0 \vee \neg S_1 \vee (\neg S_2 \wedge \neg S_3) \vee (S_2 \wedge S_3)$$

5 Suppose that the given expression is a term. Then by Proposition 2.2(ii)(c) there are terms σ, τ such that

$$\langle +, v_0, +, v_0, v_1, v_2 \rangle = \langle + \rangle \frown \sigma \frown \tau$$

Hence

$$\langle v_0, +, v_0, v_1, v_2 \rangle = \sigma \frown \tau.$$

Now v_0 is a term, and it is a segment of σ . So by Proposition 2.2(iii) we have $v_0 = \sigma$. Hence

$$\langle +, v_0, v_1, v_2 \rangle = \tau$$

But then the term $\langle +, v_0, v_1 \rangle$ is a proper initial segment of τ , contradicting Proposition 2.2(iii).

6 Suppose that the given expression is a formula. Then by Proposition 2.6(ii)(e) there is a formula φ such that

$$\langle \forall, v_0, =, v_1, v_0, v_2, v_3 \rangle = \langle \forall, v_0 \rangle \frown \varphi.$$

It follows that

$$\langle =, v_1, v_0, v_2, v_3 \rangle = \varphi.$$

But $\langle =, v_1, v_0 \rangle$ is a formula which is a proper initial segment of φ , contradicting Proposition 2.6(iii).

7

$$\neg(v_0 = 0) \wedge \neg(v_0 = S0) \wedge \forall v_1 [\exists v_2 (v_1 \cdot v_2 = v_0) \rightarrow v_1 = S0 \vee v_1 = v_0]$$

8

$$\forall v_0 [0 < v_0 \rightarrow \exists v_1 (v_1 \cdot v_1 = v_0)]$$

9 Below each variable we write f for free occurrence and b for bound occurrence.

$$\forall \begin{array}{c} v_0 \\ b \end{array} \left(\begin{array}{c} v_0 \\ b \end{array} + \begin{array}{c} v_1 \\ f \end{array} = \begin{array}{c} v_2 \\ f \end{array} \vee \exists \begin{array}{c} v_1 \\ b \end{array} \left(\begin{array}{c} v_1 \\ b \end{array} + \begin{array}{c} v_0 \\ b \end{array} = \begin{array}{c} v_2 \\ f \end{array} \wedge \forall \begin{array}{c} v_2 \\ b \end{array} \left(\begin{array}{c} v_3 \\ f \end{array} = \begin{array}{c} v_2 \\ b \end{array} \right) \right)$$

10

$$\forall v_0 \exists v_2 \exists v_4 [v_0 > 0 \rightarrow (v_1 = v_3 \rightarrow v_0 = v_1)]$$

11 It suffices to find a model of the indicated set. Let $\overline{M} = (\{0, 1\}, f)$, where $f(0) = 1$ and $f(1) = 0$.

12 By 3.37 we have $\vdash \exists v_1 \forall v_2 \varphi \rightarrow \forall v_2 \exists v_1 \varphi$. Hence by generalization on v_0 and axiom (L2) we get $\vdash \forall v_0 \exists v_1 \forall v_2 \varphi \rightarrow \forall v_0 \forall v_2 \exists v_1 \varphi$.

13 In the language for $(\omega, S, 0, +, \cdot)$ let φ be the formula $\forall v_0 (\neg(Sv_0 = v_1))$. Then $\text{Subf}_0^{v_1} \varphi$ is $\forall v_0 (\neg(Sv_0 = 0))$. and $P \vdash \text{Subf}_0^{v_1} \varphi$. But $P \not\vdash \varphi$. In fact, let $a : \omega \rightarrow \omega$ be such that $a_1 = 1$. Then $\text{not}((\omega, S, 0, +, \cdot) \models \varphi[a])$, hence $\text{not}(P \models \varphi)$, hence $\text{not}(P \vdash \varphi)$.

14 We assume that x is v_0 and y is v_1 . Let \overline{M} be any \mathcal{L} -structure, and assume that $\overline{M} \models \exists v_0 Fv_0 \wedge \forall v_0 \forall v_1 [Fv_0 \wedge Fv_1 \rightarrow v_0 = v_1]$. This means that there is a $u \in M$ such that $F^{\overline{M}} = \{u\}$.

Now suppose that $a : \omega \rightarrow M$ is any assignment. First suppose that $\overline{M} \models \exists v_0 [Fv_0 \wedge Gv_0v_1][a]$. This means that $(u, a_1) \in G^{\overline{M}}$. Now take any $w \in M$. If $\overline{M} \models Fv_0[a_w^0]$, then $w = u$ and so $(w, a_1) \in R^{\overline{M}}$, and consequently $\overline{M} \models Gv_0v_1[a_w^0]$. Thus $\overline{M} \models (Fv_0 \rightarrow Gv_0v_1)[a_w^0]$. Since w is arbitrary, this proves that $\overline{M} \models \forall v_0 (Fv_0 \rightarrow Gv_0v_1)$.

Second suppose that $\overline{M} \models \forall v_0 (Fv_0 \rightarrow Gv_0v_1)[a]$. By definition of modeling of \forall it follows that $\overline{M} \models (Fv_0 \rightarrow Gv_0v_1)[a_u^0]$. Since $F^{\overline{M}} = \{u\}$, it follows that $\overline{M} \models Gv_0v_1[a_u^0]$. So $\overline{M} \models (Fv_0 \wedge Gv_0v_1)[a_u^0]$, so $\overline{M} \models \exists v_0 (Fv_0 \wedge Gv_0v_1)$.

15 We go through the proofs of 4.20–4.24 for this formula.

- (1) $\vdash v_0 + v_1 \cdot v_1 = 0 \leftrightarrow \exists v_2 [v_0 + v_1 \cdot v_1 = v_2 \wedge 0 = v_2]$ (1) in proof of 4.24
- (2) $\vdash 0 = v_2 \leftrightarrow \exists v_0 [v_0 = v_2 \wedge 0 = v_0]$ 4.20
- (3) $\vdash v_0 + v_1 \cdot v_1 = v_2 \leftrightarrow \exists v_3 \exists v_4 \exists v_5 [v_0 = v_3 \wedge v_1 \cdot v_1 = v_4 \wedge v_2 = v_5 \wedge v_3 + v_4 = v_5]$
 (*) in proof of 4.23 with $n = 3$
- (4) $\vdash v_1 \cdot v_1 = v_4 \leftrightarrow \exists v_5 \exists v_6 \exists v_7 [v_1 = v_5 \wedge v_1 = v_6 \wedge v_4 = v_7 \wedge v_5 \cdot v_6 = v_7]$
 (*) in proof of 4.23 with $n = 5$
- (5) $\vdash v_5 \cdot v_6 = v_7 \leftrightarrow \exists v_0 \exists v_1 \exists v_2 [v_0 = v_5 \wedge v_1 = v_6 \wedge v_2 = v_7 \wedge v_0 \cdot v_1 = v_2]$
 4.22
- (6) $\vdash v_3 + v_4 = v_5 \leftrightarrow \exists v_0 \exists v_1 \exists v_2 [v_0 = v_3 \wedge v_1 = v_4 \wedge v_2 = v_5 \wedge v_0 + v_1 = v_2]$
 4.22
- (7) $\vdash v_2 + v_3 = S0 \leftrightarrow \exists v_4 [v_2 + v_3 = v_4 \wedge S0 = v_4]$ (1) in proof of 4.24
- (8) $\vdash v_2 + v_3 = v_4 \leftrightarrow \exists v_5 \exists v_6 \exists v_7 [v_2 = v_5 \wedge v_3 = v_6 \wedge v_4 = v_7 \wedge v_5 + v_6 = v_7]$
 (*) in proof of 4.23 with $n = 5$
- (9) $\vdash v_5 + v_6 = v_7 \leftrightarrow \exists v_0 \exists v_1 \exists v_2 [v_5 = v_0 \wedge v_6 = v_1 \wedge v_7 = v_2 \wedge v_0 + v_1 = v_2]$
 4.22
- (10) $\vdash S0 = v_4 \leftrightarrow \exists v_5 \exists v_6 [0 = v_5 \wedge v_4 = v_6 \wedge Sv_5 = v_6]$
 (*) in proof of 4.23 with $n = 5$

$$(11) \quad \vdash 0 = v_5 \leftrightarrow \exists v_0[v_0 = v_5 \wedge 0 = v_0] \quad 4.20$$

$$(12) \quad \vdash Sv_5 = v_6 \leftrightarrow \exists v_0 \exists v_1[v_5 = v_0 \wedge v_6 = v_1 \wedge Sv_0 = v_1] \quad 4.22$$

It remains only to put these together using 3.20.

$$(13) \quad \vdash S0 = v_4 \leftrightarrow \exists v_5 \exists v_6[0 = v_5 \wedge v_4 = v_6 \wedge \\ \exists v_0 \exists v_1[v_5 = v_0 \wedge v_6 = v_1 \wedge Sv_0 = v_1]] \quad (10), (12)$$

$$(14) \quad \vdash S0 = v_4 \leftrightarrow \exists v_5 \exists v_6[\exists v_0[v_0 = v_5 \wedge 0 = v_0] \wedge v_4 = v_6 \wedge \\ \exists v_0 \exists v_1[v_5 = v_0 \wedge v_6 = v_1 \wedge Sv_0 = v_1]] \quad (11), (13)$$

$$(15) \quad \vdash v_2 + v_3 = v_4 \leftrightarrow \exists v_5 \exists v_6 \exists v_7[v_2 = v_5 \wedge v_3 = v_6 \wedge v_4 = v_7 \wedge \\ \exists v_0 \exists v_1 \exists v_2[v_5 = v_0 \wedge v_6 = v_1 \wedge v_7 = v_2 \wedge v_0 + v_1 = v_2]] \quad (8), (9)$$

$$(16) \quad \vdash v_2 + v_3 = S0 \leftrightarrow \exists v_4[\exists v_5 \exists v_6 \exists v_7[v_2 = v_5 \wedge v_3 = v_6 \wedge v_4 = v_7 \wedge \\ \exists v_0 \exists v_1 \exists v_2[v_5 = v_0 \wedge v_6 = v_1 \wedge v_7 = v_2 \wedge v_0 + v_1 = v_2]] \\ \wedge \exists v_5 \exists v_6[\exists v_0[v_0 = v_5 \wedge 0 = v_0] \wedge v_4 = v_6 \wedge \\ \exists v_0 \exists v_1[v_5 = v_0 \wedge v_6 = v_1 \wedge Sv_0 = v_1]]] \quad (7), (14), (15)$$

$$(17) \quad \vdash v_0 + v_1 \cdot v_1 = 0 \leftrightarrow \exists v_2[\\ \exists v_3 \exists v_4 \exists v_5[v_0 = v_3 \wedge v_1 \cdot v_1 = v_4 \wedge v_2 = v_5 \wedge v_3 + v_4 = v_5] \wedge \\ \exists v_0[v_0 = v_2 \wedge 0 = v_0]] \quad (1), (2), (3)$$

$$(18) \quad \vdash v_0 + v_1 \cdot v_1 = 0 \leftrightarrow \exists v_2[\\ \exists v_3 \exists v_4 \exists v_5[v_0 = v_3 \wedge \exists v_5 \exists v_6 \exists v_7[v_1 = v_5 \wedge v_1 = v_6 \wedge v_4 = v_7 \wedge v_5 \cdot v_6 = v_7] \\ \wedge v_2 = v_5 \wedge \exists v_0 \exists v_1 \exists v_2[v_0 = v_3 \wedge v_1 = v_4 \wedge v_2 = v_5 \wedge v_0 + v_1 = v_2]]] \wedge \\ \exists v_0[v_0 = v_2 \wedge 0 = v_0]] \quad (17), (4), (6)]$$

$$(19) \quad \vdash v_0 + v_1 \cdot v_1 = 0 \leftrightarrow \exists v_2[\\ \exists v_3 \exists v_4 \exists v_5[v_0 = v_3 \wedge \exists v_5 \exists v_6 \exists v_7[v_1 = v_5 \wedge v_1 = v_6 \wedge v_4 = v_7 \wedge \\ \exists v_0 \exists v_1 \exists v_2[v_0 = v_5 \wedge v_1 = v_6 \wedge v_2 = v_7 \wedge v_0 \cdot v_1 = v_2] \\ \wedge v_2 = v_5 \wedge \exists v_0 \exists v_1 \exists v_2[v_0 = v_3 \wedge v_1 = v_4 \wedge v_2 = v_5 \wedge v_0 + v_1 = v_2]]] \wedge \\ \exists v_0[v_0 = v_2 \wedge 0 = v_0]] \quad (18), (5)]$$

Thus we get, finally, using (16) and (19)

$$\vdash \forall v_0 \forall v_1[v_0 + v_1 \cdot v_1 = 0 \wedge v_2 + v_3 = S0] \leftrightarrow \forall v_0 \forall v_1[\\ \exists v_4[\exists v_5 \exists v_6 \exists v_7[v_2 = v_5 \wedge v_3 = v_6 \wedge v_4 = v_7 \wedge \\ \exists v_0 \exists v_1 \exists v_2[v_5 = v_0 \wedge v_6 = v_1 \wedge v_7 = v_2 \wedge v_0 + v_1 = v_2]] \\ \wedge \exists v_5 \exists v_6[\exists v_0[v_0 = v_5 \wedge 0 = v_0] \wedge v_4 = v_6 \wedge \\ \exists v_0 \exists v_1[v_5 = v_0 \wedge v_6 = v_1 \wedge Sv_0 = v_1]]] \wedge \\ \exists v_2[\exists v_3 \exists v_4 \exists v_5[v_0 = v_3 \wedge \exists v_5 \exists v_6 \exists v_7[v_1 = v_5 \wedge v_1 = v_6 \wedge v_4 = v_7 \wedge \\ \exists v_0 \exists v_1 \exists v_2[v_0 = v_5 \wedge v_1 = v_6 \wedge v_2 = v_7 \wedge v_0 \cdot v_1 = v_2] \\ \wedge v_2 = v_5 \wedge \exists v_0 \exists v_1 \exists v_2[v_0 = v_3 \wedge v_1 = v_4 \wedge v_2 = v_5 \wedge v_0 + v_1 = v_2]]] \wedge \\ \exists v_0[v_0 = v_2 \wedge 0 = v_0]]$$

16 Let Γ consist of the following sentences:

$$\begin{aligned} & \exists v_0 \dots \exists v_{n-1} \left[\bigwedge_{i < j < n} (\neg(v_i = v_j)) \right] \quad \text{for each integer } n > 1; \\ & \forall v_0 (v_0 E v_0); \\ & \forall v_0 \forall v_1 [v_0 E v_1 \rightarrow v_1 E v_0]; \\ & \forall v_0 \forall v_1 \forall v_2 [v_0 E v_1 \wedge v_1 E v_2 \rightarrow v_0 E v_2]; \\ & \forall v_0 \exists v_1 [v_0 E v_1 \wedge (\neg(v_0 = v_1)) \wedge \forall v_2 (v_0 E v_2 \rightarrow v_2 = v_0 \vee v_2 = v_1)]. \end{aligned}$$

17 See Lemma 4 in the solution for exercise 3.

18 $\exists v_1 \exists v_2 [v_0 = v_1 \cdot v_1 + v_2 \cdot v_2]$.

19 We take x to be v_0 and y to be v_1 . Let \overline{M} be a structure and $a : \omega \rightarrow M$ be any assignment. First suppose that $\overline{M} \models \exists v_0 G v_0 v_1 [a]$. Choose $u \in M$ so that $\overline{M} \models G v_0 v_1 [a_u^0]$. So $(u, a_1) \in G^{\overline{M}}$. Hence $\overline{M} \models G v_1 v_0 [a_{a_1^1 u}^0]$ and so $\overline{M} \models \exists v_1 G v_1 v_0 [a_{a_1^1}^0]$; so $\overline{M} \models (v_0 = v_1 \wedge \exists v_1 G v_1 v_0) [a_{a_1^1}^0]$. So $\overline{M} \models \exists v_0 (v_0 = v_1 \wedge \exists v_1 G v_1 v_0) [a]$.

Second, suppose that $\overline{M} \models \exists v_0 (v_0 = v_1 \wedge \exists v_1 G v_1 v_0) [a]$. Choose $u \in M$ so that $\overline{M} \models (v_0 = v_1 \wedge \exists v_1 G v_1 v_0) [a_u^0]$. Thus $u = a_1$. Choose $w \in M$ so that $\overline{M} \models G v_1 v_0 [a_{a_1^1 w}^0]$. Hence $(w, a_1) \in G^{\overline{M}}$, so $\overline{M} \models G v_0 v_1 [a_w^0]$, hence $\overline{M} \models \exists v_0 G v_0 v_1 [a]$.

20 Expand the language by adding an individual constant c . Let Γ' consist of Γ together with all sentences

$$\exists v_0 \dots \exists v_{m-1} \left[\bigwedge_{i < j < m} \neg(v_i = v_j) \wedge \bigwedge_{i < m} c R v_i \right]$$

Clearly every finite subset of Γ' has a model, so Γ' itself has a model (M, R, d) . Clearly (M, R) is as desired.