

More NP-complete problems

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Cliques

- ▶ $K_n = (\{1, \dots, n\}, \neq)$... **complete (undirected) graph** on n -vertices
- ▶ A graph G has an n -**clique** if K_n embeds into G .
- ▶ $\text{CLIQUE} := \{(G, n) : G \text{ is a graph with } n\text{-clique}\}$

Theorem

CLIQUE is NP-complete.

Proof.

CLIQUE \in NP since a guessed n -clique can be verified in polynomial time.

Claim: 3-SAT \leq_m^p CLIQUE

Given a 3-SAT instance

$$\Phi = (a_1 \vee b_1 \vee c_1) \wedge \cdots \wedge (a_n \vee b_n \vee c_n)$$

with literals a_i, b_i, c_i .

For the reduction construct a graph G with

- ▶ $3n$ vertices labelled $a_1, b_1, c_1, \dots, a_n, b_n, c_n$
- ▶ edges between any 2 vertices except within any triple a_i, b_i, c_i representing a clause of Φ and between a variable and its negation.

Φ with $3n$ literals yields a graph with $3n$ vertices in polytime.

Example: $\Phi = (x_1 \vee x_1 \vee x_2) \wedge (x'_1 \vee x'_2 \vee x'_2) \wedge (x'_1 \vee x_2 \vee x_2)$

Claim: Φ is satisfiable iff G has an n -clique

\Rightarrow :

- ▶ Assume Φ has a satisfying assignment.
- ▶ In each triple a_i, b_i, c_i of G choose a vertex corresponding to a true literal in this satisfying assignment.
- ▶ These n vertices are each pairwise connected, hence a clique.

\Leftarrow :

- ▶ Assume G has an n -clique.
- ▶ Then any 2 vertices in that clique are in distinct clauses.
- ▶ Assign truth values to variables in Φ such that each literal labelling a vertex in the clique is true (possible since x_j, x'_j are not connected).
- ▶ Since each clause contains a vertex from the clique, this assignment satisfies Φ .



Graph coloring

- ▶ A graph G is n -**colorable** if its vertices can be colored in n colors such that any adjacent vertices have distinct colors.
- ▶ Equivalently, G is n -colorable iff there exists a homomorphism $G \rightarrow K_n$.
- ▶ n -Coloring $:= \{G : G \text{ is } n\text{-colorable}\}$.

Theorem

3-Coloring is NP-complete.

Proof.

3-Coloring \in NP since a guessed coloring can be verified in polynomial time.

Claim: 3-SAT \leq_m^p 3-Coloring

Given a 3-SAT instance

$$\Phi = (a_1 \vee b_1 \vee c_1) \wedge \cdots \wedge (a_n \vee b_n \vee c_n)$$

with literals a_i, b_i, c_i .

Construct G that is 3-colorable iff Φ is satisfiable as follows:

- ▶ Truth assignments of x_1, \dots, x_k correspond to colors 0, 1 of vertices.

- ▶ For each clause $a_i \vee b_i \vee c_i$ connect the vertices corresponding to a_i, b_i, c_i by a **gadget graph** implementing “or”.

Further NP-complete problems (Karp, 1972)

- ▶ SAT
- ▶ Circuit Satisfiability Problem
Given a Boolean circuit, is there an assignment of inputs that yields output 1?
- ▶ CLIQUE
- ▶ Graph k -Coloring for $k \geq 3$
- ▶ Graph Homomorphism Problem
Given graphs G, H , is there a homomorphism $G \rightarrow H$?
- ▶ Hamiltonian Cycle for digraphs
- ▶ Travelling Salesman Problem

- ▶ Exact Cover
Given subsets $A_1, \dots, A_n \subseteq X$, is X the disjoint union of some A_i ?
- ▶ Knapsack (Subset Sum)
Given integers a_1, \dots, a_n and $s \in \mathbb{Z}$, does a non-empty subset of the a_i sum to s ?
- ▶ MaxCut
Given a graph G and $k \in \mathbb{N}$, is there a cut of size at least k in G (a partition of vertices into 2 sets A, B with $\geq k$ edges between A and B)?
- ▶ Sudoku for n^2 numbers