

The structure of idempotent involutive residuated lattices and weakening relation algebras

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Part I, with O. Tuyt and D. Valota

- Commutative idempotent involutive residuated lattices
- Gluing construction
- Ungluing decomposition

Part II, with N. Galatos

- FL^2 -algebras and their congruences
- Weakening relation algebras
- Double-division conuclei

Involutive residuated lattices

Definition

A **pointed residuated lattice** $\mathbf{A} = \langle A, \wedge, \vee, \cdot, \backslash, /, 1, 0 \rangle$ is

- a lattice $\langle A, \wedge, \vee \rangle$ and a monoid $\langle A, \cdot, 1 \rangle$ such that

$$x \cdot y \leq z \iff x \leq z/y \iff y \leq x \backslash z \quad \text{for all } x, y, z \in A.$$

\mathbf{A} is **involutive** if $\sim -x = x = -\sim x$, where $\sim x = x \backslash 0$ and $-x = 0/x$.

$\backslash, /$ can be term-defined via $x \backslash y = \sim(-y \cdot x)$ and $x/y = -(y \cdot \sim x)$.

- \mathbf{A} is **commutative** if $x \cdot y = y \cdot x$ (hence $-x = \sim x$)
- \mathbf{A} is **idempotent** if $x \cdot x = x$ for all $x \in A$

CIdInRL denotes the variety of **commutative idempotent involutive residuated lattices**.

Examples of CIdInRLs

Let $\mathbf{A} \in \text{CIdInRL}$.

- $\langle A, \cdot, 1 \rangle$ is a meet-semilattice with top element 1 and order \sqsubseteq (**monoidal order**) defined as

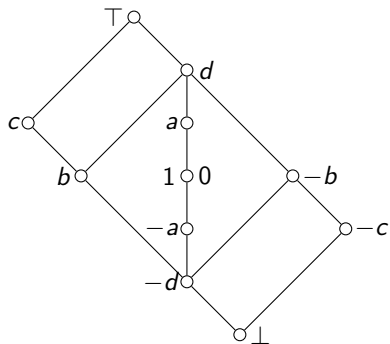
$$a \sqsubseteq b \iff a \cdot b = a.$$

Hence, the orders \leq and \sqsubseteq together with the involution $-$ completely determine \mathbf{A} , allowing us to work in the signature $\langle A, \vee, \cdot, -, 0, 1 \rangle$

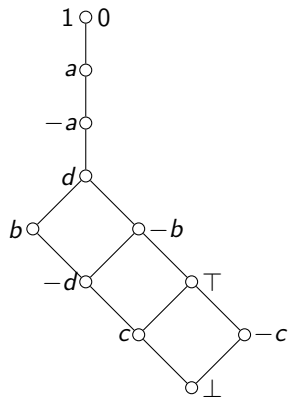
- **Boolean algebras** (where $\leq = \sqsubseteq$)
- **Sugihara monoids** defined as **distributive** CIdInRLs (= algebraic semantics for relevance logic RM^t)

Dunn [1970] proved that the subdirectly irreducible Sugihara monoids are linearly ordered. Up to isomorphism, there is one such algebra \mathbf{S}_n for each chain with n elements.

Another example

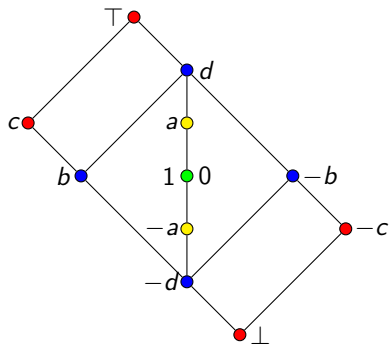


$\langle A, \leq \rangle$

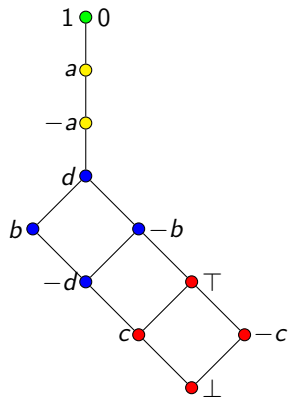


$\langle A, \subseteq \rangle$

Another example

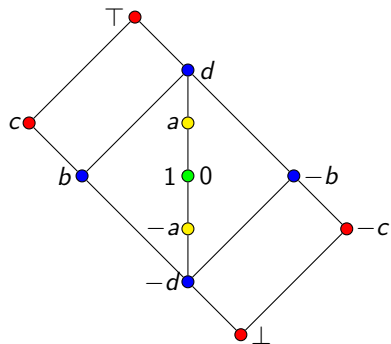


$\langle A, \leq \rangle$

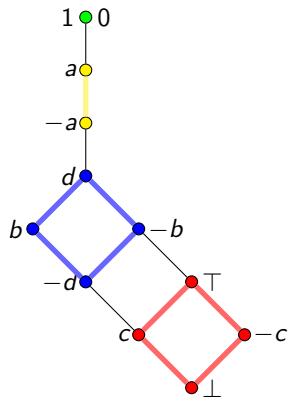


$\langle A, \subseteq \rangle$

Another example



$\langle A, \leq \rangle$



$\langle A, \sqsubseteq \rangle$

Some properties

For each $x \in A$, let

$$0_x := x \wedge -x = x \cdot -x$$

$$1_x := x \vee -x = -(x \cdot -x) = x/x$$

$$\mathbb{B}_x := \{y \in A \mid 0_x \sqsubseteq y \sqsubseteq 1_x\}$$

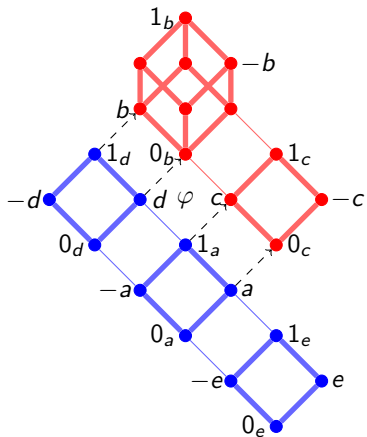
$$\downarrow 0 := \{y \in A \mid y \leq 0\} = \{0_x \mid x \in A\}$$

Lemma

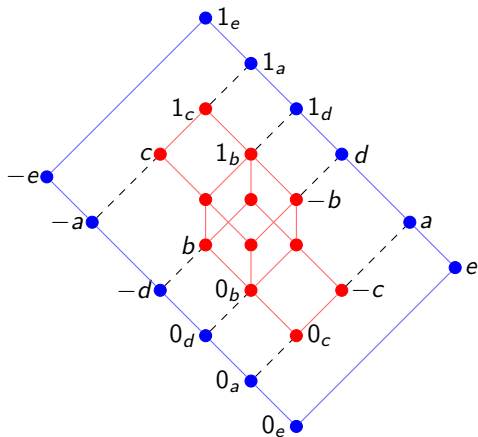
- For each $x \in A$, $\langle \mathbb{B}_x, \wedge, \vee, -, 0_x, 1_x \rangle$ is a **Boolean algebra**
- For each $x \in A$, the monoidal order and the lattice order **agree** on \mathbb{B}_x
- The monoidal intervals \mathbb{B}_x **partition** A
- $\langle \downarrow 0, \cdot, \vee \rangle$ is a **distributive lattice** with top element 0

Hence, the monoidal semilattice is a disjoint union of Boolean algebras over the 'skeleton' of a distributive lattice.

Construction: example of $\mathbf{C} = \mathbf{A} \oplus_{\phi} \mathbf{B}$



$\langle \mathbf{C}, \sqsubseteq \rangle$



$\langle \mathbf{C}, \leq \rangle$

Construction: formally

Let $\uparrow a = \{x \in A \mid a \sqsubseteq x\}$ and $\downarrow b = \{x \in B \mid x \sqsubseteq b\}$.

$\mathbf{A} = \langle A, \vee^A, \cdot^A, -^A, 0^A, 1^A \rangle$ (the bottom algebra) and

$\mathbf{B} = \langle B, \vee^B, \cdot^B, -^B, 0^B, 1^B \rangle$ (the top algebra) are **φ -compatible** if

- φ is a **bijection** $\uparrow a \rightarrow \downarrow b$ for some $a \leq 1^A$ and $0^B \leq b \leq 1^B$ such that
- φ **preserves join**, i.e. $\varphi(x \vee^A y) = \varphi(x) \vee^B \varphi(y)$
- φ **preserves fusion**, i.e. $\varphi(x \cdot^A y) = \varphi(x) \cdot^B \varphi(y)$ and
- $0^B = \varphi(a \vee^A 0^A)$.

For φ -compatible algebras we define a **glueing construction** \oplus_{φ}

Glueing construction

$$\mathbf{A} \oplus_{\varphi} \mathbf{B} := \langle A \uplus B, \vee, \cdot, -, 1^{\mathbf{B}}, 0^{\mathbf{B}} \rangle$$

$$x \vee y = \begin{cases} x \vee^A y & \text{if } x, y \in A \\ x \vee^B y & \text{if } x, y \in B \\ \varphi(x \vee^A a) \vee^B y & \text{if } x \in A, y \in B, x \leq^A -^A a \\ x \vee^A \varphi^{-1}(y \cdot^B b) & \text{if } x \in A, y \in B, x \not\leq^A -^A a \end{cases}$$

$$x \cdot y = \begin{cases} x \cdot^A y & \text{if } x, y \in A \\ x \cdot^B y & \text{if } x, y \in B \\ x \cdot^A \varphi^{-1}(y \cdot^B b) & \text{if } x \in A, y \in B \end{cases}$$

$$-x = \begin{cases} -^A x & \text{if } x \in A \\ -^B x & \text{if } x \in B \end{cases}$$

Theorem

For φ -**compatible** $\mathbf{A}, \mathbf{B} \in \text{CIdInRL}$ the algebra $\mathbf{A} \oplus_{\varphi} \mathbf{B}$ is in CIdInRL .

The proof is by case analysis and direct computation.

Unglueing decomposition

For finite $\mathbf{C} \in \text{CIdInRL}$, consider a co-atom c in the underlying distributive lattice with universe $\downarrow 0 = \{0_x \mid x \in C\}$.

By distributivity, there exists c^* such that $\langle c, c^* \rangle$ is a splitting pair of $\downarrow 0$.

Note: $c = 0_c$, hence $-c = 1_c$.

Lemma

The pair $\langle 1_c, c^ \rangle$ is a splitting pair of (C, \sqsubseteq) .*

Moreover, $\uparrow c^$ is a subuniverse of \mathbf{C} , and $\downarrow 1_c$ is closed under $\vee, \cdot, -$*

Unglueing decomposition

Let $\mathbf{A} = \langle \downarrow 1_c, \vee, \cdot, -, 1_c, 0_c \rangle$.

Let \mathbf{B} be the subalgebra of \mathbf{C} with subuniverse $\uparrow c^*$.

Choose $a = 1_c \cdot c^*$ and $b = (1_c \vee -a) \vee c^*$, and define

$\varphi(x) = (x \wedge -a) \vee c^*$ for $a \sqsubseteq x \sqsubseteq 1_c$.

Lemma

- $a \leq 1_c$ and $0 \leq b \leq 1$
- φ is a bijection to $\{y \mid c^* \sqsubseteq y \sqsubseteq b\}$ with $\varphi^{-1}(y) = y \cdot 1_c$
- $\varphi(c \vee a) = 0_b$

Theorem

The algebra $\mathbf{C} \in \text{CIdInRL}$ is isomorphic to $\mathbf{A} \oplus_{\varphi} \mathbf{B}$.

Structural characterization

The discovery of the previous theorem and the results below were guided by Prover9/Mace4 computations of all CIdInRLs with ≤ 16 elements.

Theorem

Any finite member \mathbf{A} of CIdInRL can be constructed using the gluing construction, starting from finite Boolean algebras.

Corollary

Any finite $\mathbf{A} \in \text{CIdInRL}$ is determined by its fusion semilattice and also by its lattice reduct.

To do: Implement an algorithm for constructing all finite CIdInRLs.

Fusion-distributivity

As an application, call an $\mathbf{A} \in \text{CldInRL}$ **fusion-distributive** if the meet-semilattice $\langle A, \cdot \rangle$ is **distributive**, i.e. if for all $x, y, z \in A$,

$$x \cdot y \sqsubseteq z \implies \exists x', y' \in A \text{ such that } x \sqsubseteq x', y \sqsubseteq y', \text{ and } z = x' \cdot y'.$$

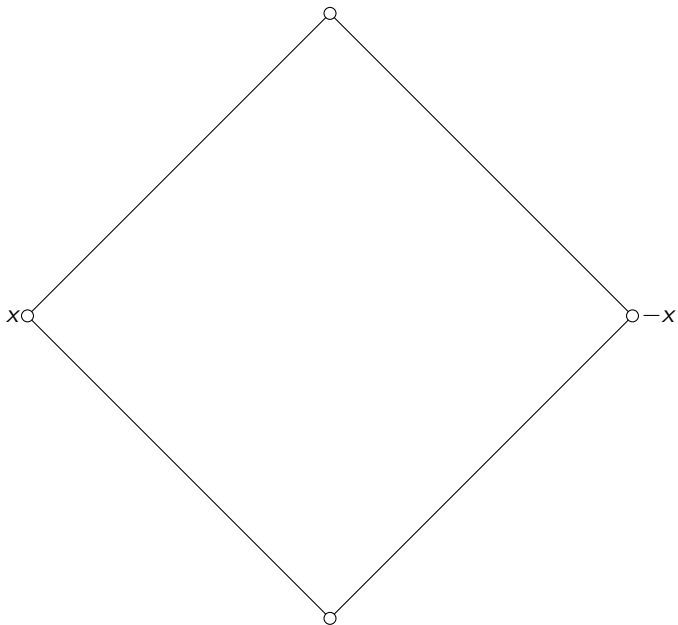
Lemma

For compatible fusion-distributive $\mathbf{A}, \mathbf{B} \in \text{CldInRL}$, their gluing \mathbf{C} is fusion-distributive.

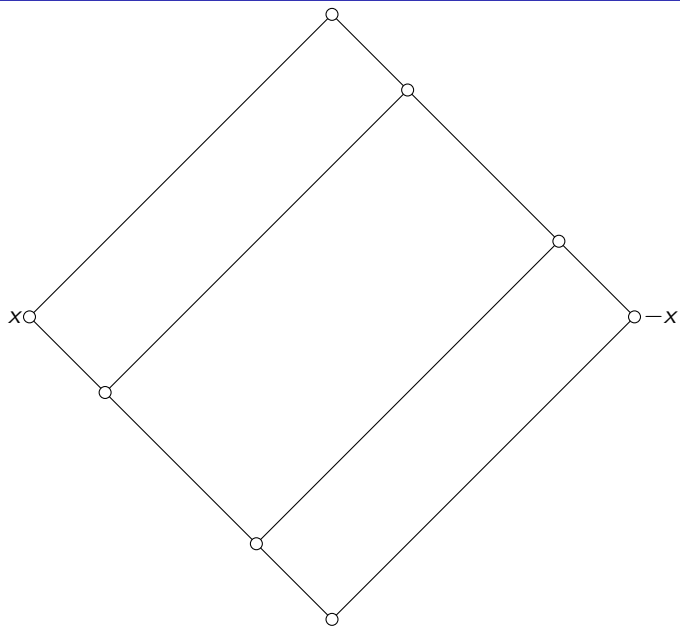
Corollary

- *Any finite $\mathbf{A} \in \text{CldInRL}$ is fusion-distributive.*
- *Every finite distributive lattice can occur as skeleton.*

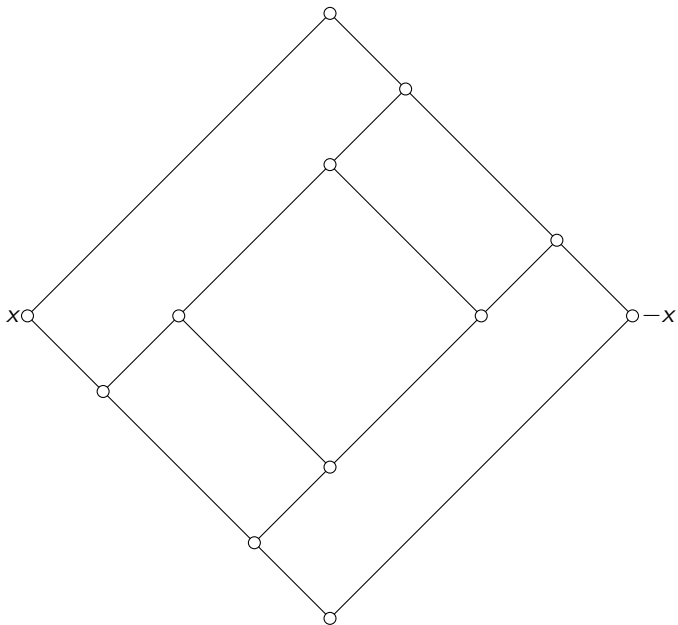
A one-generated infinite CIdInRL



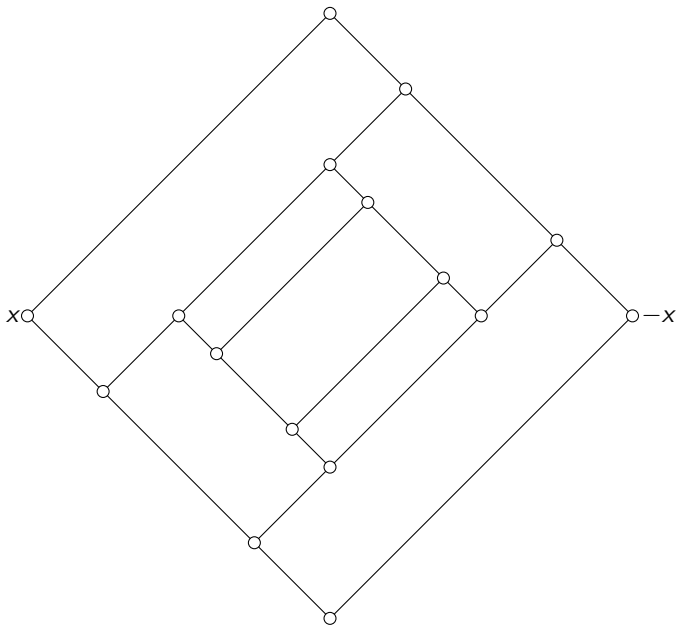
A one-generated infinite CIdInRL



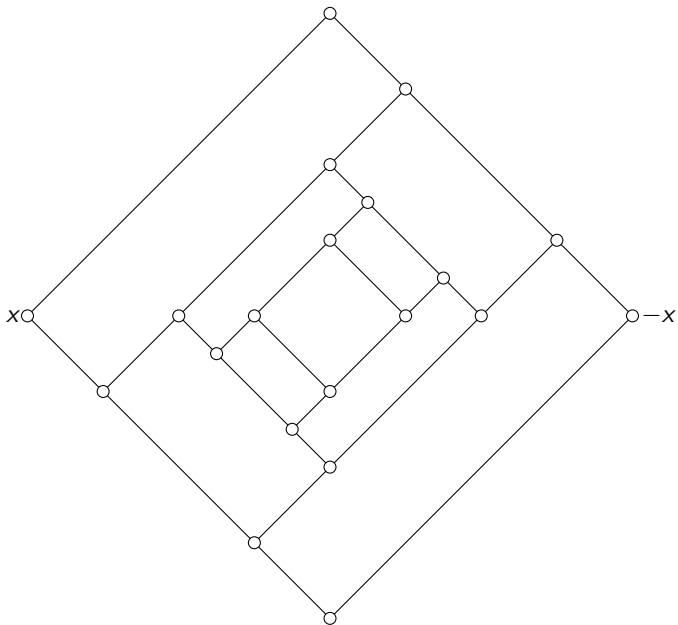
A one-generated infinite CIdInRL



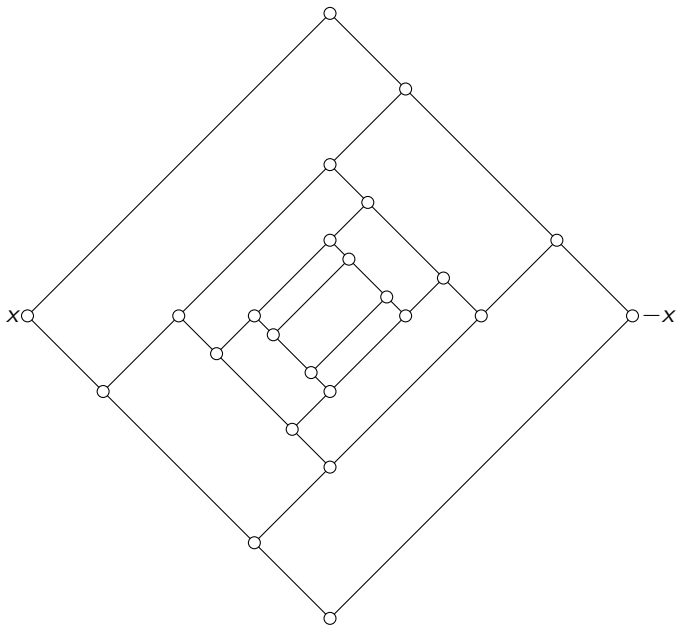
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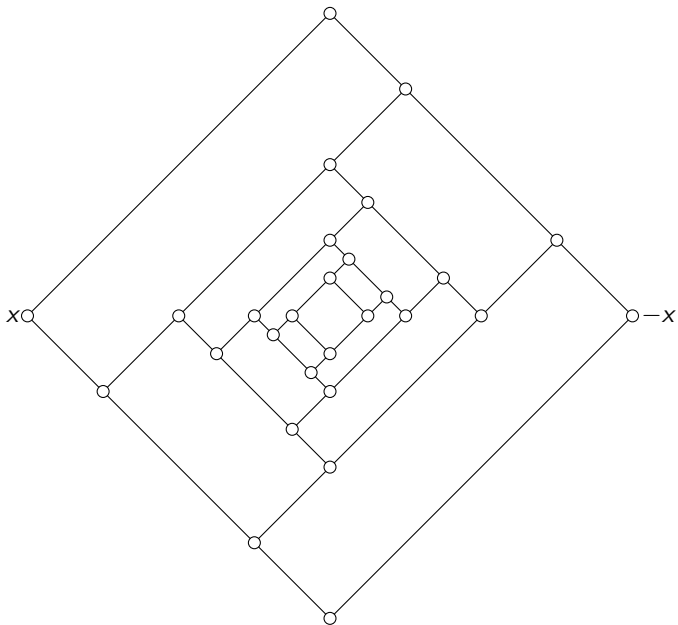
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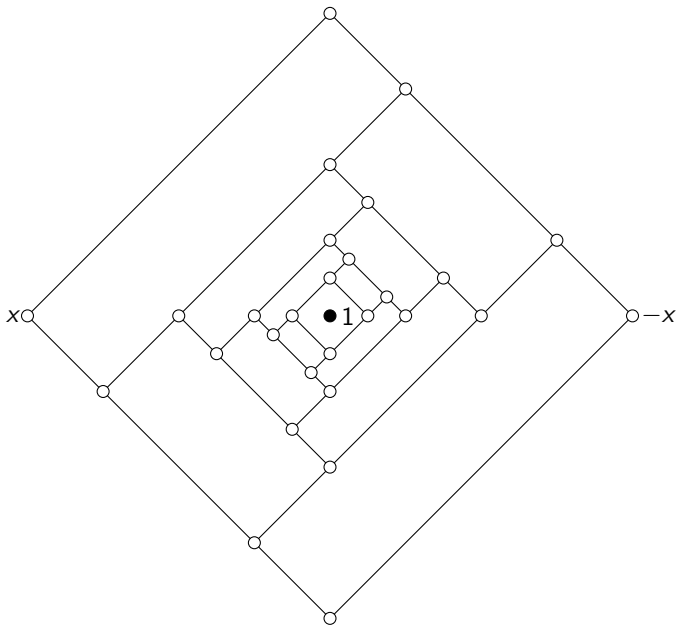
A one-generated infinite CIdInRL



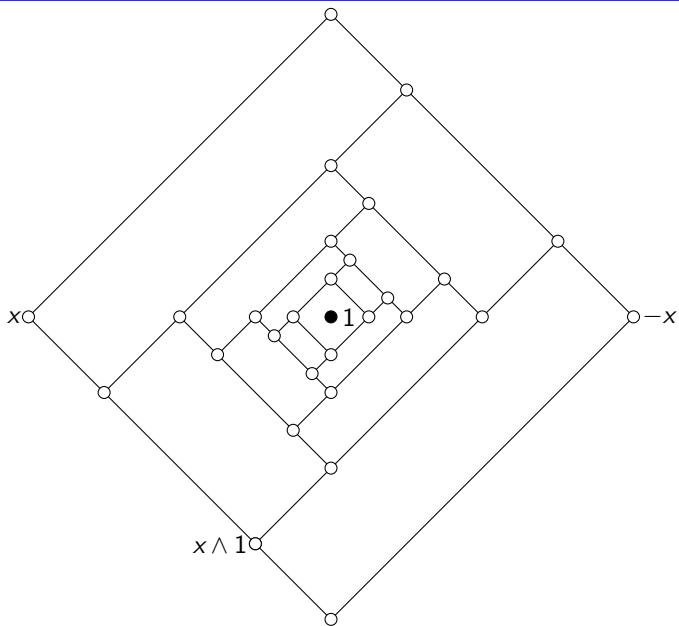
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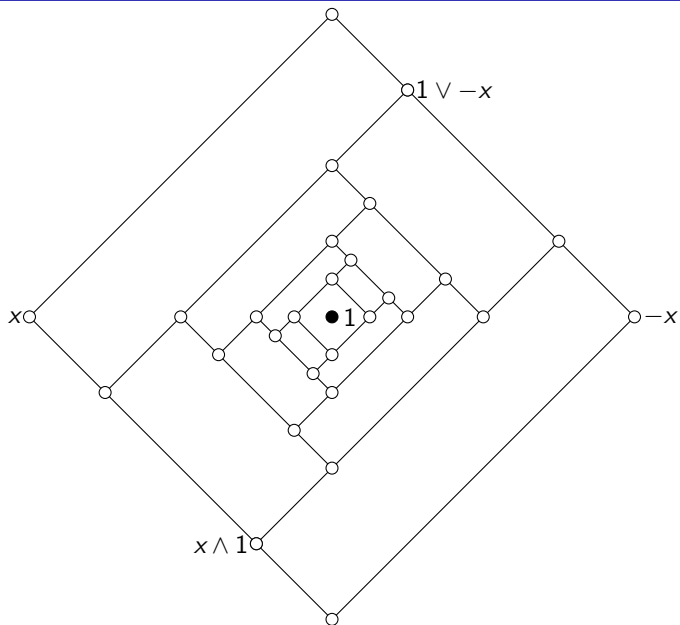
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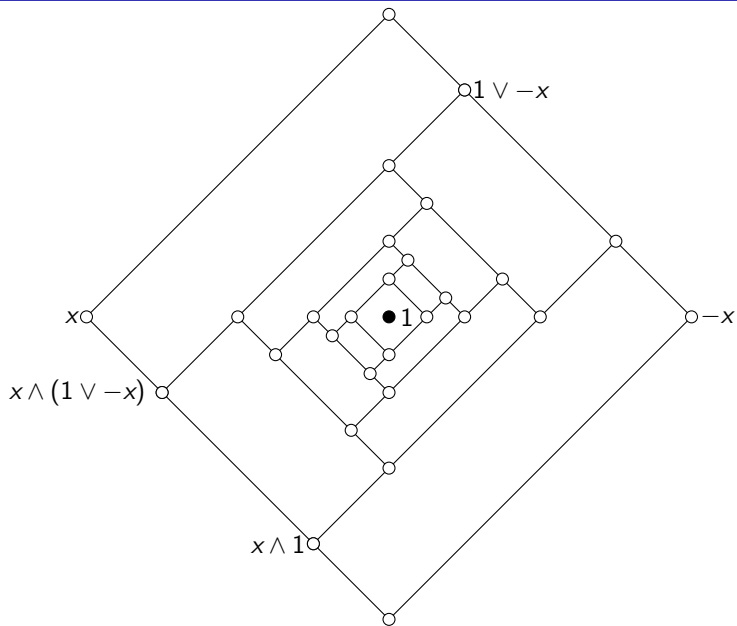
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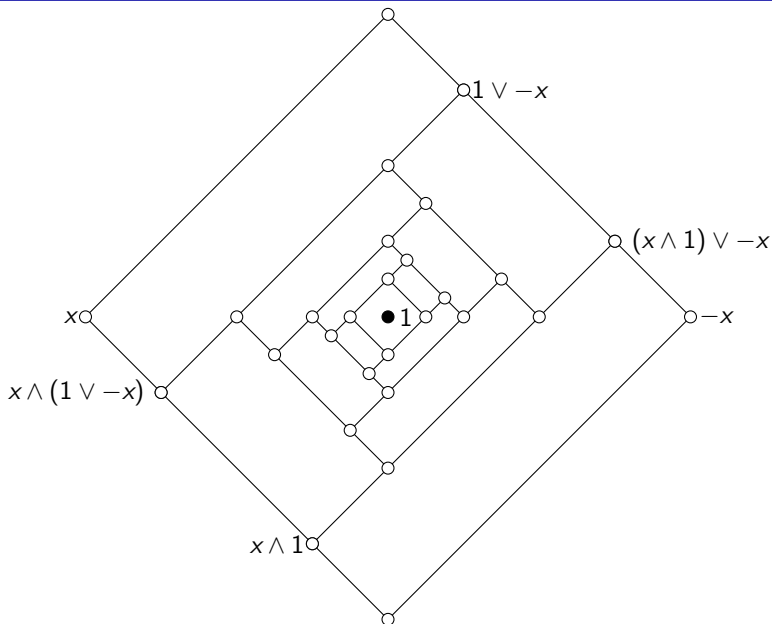
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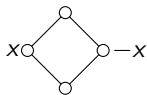
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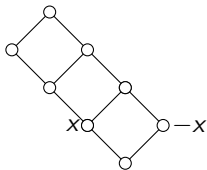
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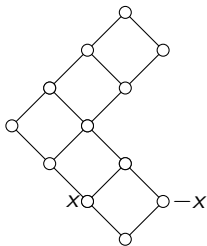
The fusion semilattice of a one-generated infinite CIdInRL



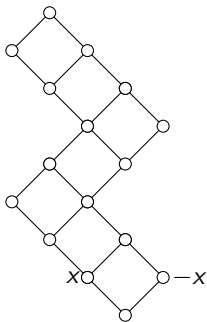
The fusion semilattice of a one-generated infinite CIdInRL



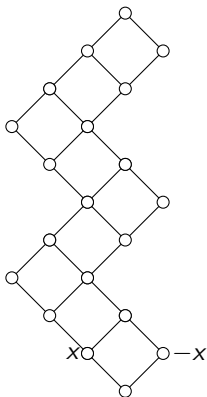
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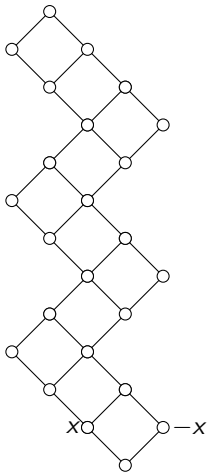
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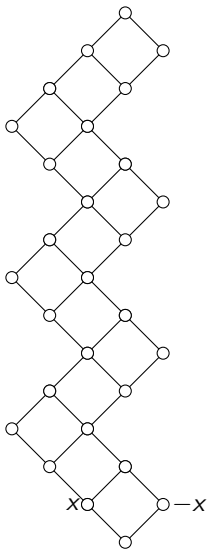
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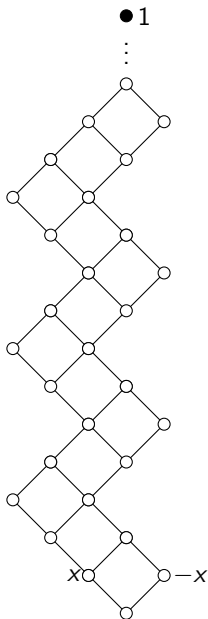
The fusion semilattice of a one-generated infinite CIdInRL



The fusion semilattice of a one-generated infinite CIdInRL



The fusion semilattice of a one-generated infinite CIdInRL



Part II: FL²-algebras and GBI-algebras

A **FL²-algebra** is of the form $\mathbf{A} = (A, \wedge, \vee, \diamond, \rightarrow, \leftarrow, t, f, \cdot, \backslash, /, 1, 0)$ s. t.

$$\mathbf{A}_t = (A, \wedge, \vee, \diamond, \rightarrow, \leftarrow, t, f) \quad \text{and} \quad \mathbf{A}_1 = (A, \wedge, \vee, \cdot, \backslash, /, 1, 0)$$

are pointed residuated lattices.

Relation algebras are examples of **classical** FL²-algebras: \mathbf{A}_t is a Boolean algebra with $x \wedge y = x \diamond y$.

A **bounded generalized bunched implication algebra** (**bGBI-algebra**) is a FL²-algebra that satisfies $x \wedge y = x \diamond y$, $t = \top$, $f = \perp$ and $0 = 1$.

A **bunched implication algebra**, or **BI-algebra**, is a commutative bGBI-algebra (i.e., $xy = yx$).

Congruences of residuated lattices

A **congruence filter** of a residuated lattice \mathbf{A} is a subset of the form $F = \uparrow([1]_\theta)$ where θ is a congruence.

Congruence filters satisfy the following **normality condition** for $a \in A$ (where quantifiers range over F):

$$\forall x \in F \exists x_1, x_2 \in F, x_1 a \leq ax \text{ and } ax_2 \leq xa. \quad (N_a)$$

A filter F satisfies (N) if (N_a) holds for all $a \in A$.

The set of **congruence filters** of \mathbf{A} is denoted by $\text{CF}(\mathbf{A})$.

Theorem (Blount-Tsinakis 2003)

For a residuated lattice \mathbf{A} , a subset F is a congruence-filter if and only if F is a lattice filter and a submonoid of \mathbf{A} that satisfies (N) .

Moreover, $\text{Con}(\mathbf{A})$ is isomorphic to the lattice $\text{CF}(\mathbf{A})$ of congruence-filters via the bijection $\theta \mapsto \uparrow([1]_\theta)$ and $F \mapsto \{(x, y) : x/y, y/x \in F\}$.

Congruences of FL^2 -algebras

For FL^2 the congruence 1-filters are determined by a stronger **t -normality** condition. For any $a \in A$

$$\forall x \in F, \exists x_1, x_2, x_3, x_4 \in F, \quad (N_a^t) \\ ax_1 \leq a \diamond xt, \quad x_2 a \leq xt \diamond a, \quad a \diamond x_3 t \leq xa, \quad x_4 t \diamond a \leq ax$$

A filter F satisfies (N^t) if (N_a^t) holds for all $a \in A$.

Theorem

For an FL^2 -algebra \mathbf{A} , a subset F is the 1-filter of some congruence θ of \mathbf{A} if and only if F is a lattice filter and $\cdot, 1$ -submonoid of \mathbf{A} that satisfies (N^t)

An analogous result holds for congruence t -filters $\uparrow([t]_\theta)$ of FL^2 -algebras.

Congruences of GBI-algebras

The previous result specializes to generalized bunched implication algebras:

Corollary

The 1-filters of a GBI-algebra \mathbf{A} are the filter submonoids that are closed under the terms

$$u_a(x) = a \setminus (a \wedge x \top), \quad v_a(x) = (a \rightarrow xa) / \top \quad \text{and} \quad \rho_a(x) = ax / a,$$

A previously known characterization of the congruence classes of GBI-algebras used more complicated terms with two parameters.

Similar 1-parameter terms exist for congruence \top -filters of GBI-algebra.

Theorem

For an involutive GBI-algebra, a lattice filter F is a \top -filter if and only if for all $x \in F$ it follows that $\neg \sim x$, $\neg \neg x$, $\sim(\top(-x)\top) \in F$.

Weakening relation algebras

For a poset $\mathbf{P} = (P, \leq)$, let $\mathbf{Wk}(\mathbf{P}) = \{R \subseteq P^2 : \leq; R; \leq \subseteq R\}$.

Relations in $\mathbf{Wk}(\mathbf{P})$ are called **weakening closed relations** since

$$x \leq u \ R \ v \leq y \implies x \ R \ y$$

$\sim R := (R^c)^\smile = \{(y, x) \mid (x, y) \notin R\}$, the **complement-converse** of R .

Weakening relations are closed under **complement-converse**, **union**, **intersection**, Heyting **implication** \rightarrow (= residual of intersection), relation **composition** ; and **residuals** $\backslash, /$ of composition.

$1 := \leq$ is a weakening relation and is the identity of composition.

The **full weakening relation algebra** on a poset \mathbf{P} is

$$\mathbf{Wk}(\mathbf{P}) = (\mathbf{Wk}(\mathbf{P}), \cap, \cup, \rightarrow, P^2, \emptyset, ;, \sim, 1, 0), \text{ where } 0 = \sim 1.$$

Representable weakening relation algebras = $\mathbf{V}\{\mathbf{Wk}(\mathbf{P}) \mid \mathbf{P} \text{ is a poset}\}$.

Double division conuclei

An **interior operator** δ on a poset is an order-preserving map such that $\delta(\delta(x)) = \delta(x) \leq x$.

An interior operator δ is a **conucleus** if $\delta(x)\delta(y) \leq \delta(xy)$.

The conucleus **image** $\delta(\mathbf{A})$ of a residuated lattice is a residuated lattice $(\delta(A), \wedge_\delta, \vee, \cdot, \backslash_\delta, /_\delta)$ without 1, where $x *_\delta y = \delta(x * y)$ for $* \in \{\wedge, \vee, \cdot, \backslash, /\}$.

Let $p \in A$ be a **positive idempotent**, i.e., $p = p^2 \geq 1$.

Then $\delta_p(x) = p \backslash x / p$ is a conucleus called the **double division conucleus**.

Lemma

$\delta_p(\mathbf{A}) = \{pxp \mid x \in A\}$, and p is the identity element.

Double division conuclei of relation algebras

In a full relation algebra, a positive idempotent p is a **preorder** $\mathbf{P} = (P, \sqsubseteq)$ (i.e., $p = \sqsubseteq$ is reflexive and transitive).

If $p \wedge p^\smile = 1$ then \mathbf{P} is a poset and $\mathbf{Wk}(\mathbf{P}) = \delta_p(\text{Rel}(P))$.

Hence the variety RWkRA of representable weakening relation algebras contains all double division conucleus images of members of RRA.

For a class \mathcal{K} of algebras let $d\mathcal{K} = \{\delta_p(\mathbf{A}) : \mathbf{A} \in \mathcal{K}, 1 \leq p^2 = p \in A\}$.

Theorem

If \mathcal{V} is a variety of bounded GBI-algebras with $\top \setminus x / \top$ as unary discriminator on the subdirectly irreducible members then $S(d\mathcal{V})$ is a discriminator variety with the same unary discriminator term.

Applying this result to the variety RA produces the discriminator variety $S(d\text{RA})$ that contains both RA and RWkRA.

Some identities that hold in $S(\text{dRA})$

Recall that the variety RA of **relation algebras** is an abstract counterpart of the variety RRA of **representable relation algebras**.

The variety $S(\text{dRA})$ generated by double-division conucleus images of relation algebras is the abstract counterpart of RWkRA.

Open problem: Find a (finite?) axiomatization of $S(\text{dRA})$.

In a GBI-algebra let the **domain** $d(x) = x^\top \wedge 1$ and **range** $r(x) = \top x \wedge 1$.

Theorem

The identities

$$d(x)x = x, \quad xr(x) = x, \quad \top x \top x \top = \top x \top \quad \text{and} \quad \sim \neg(xy) \leq (\sim \neg y)(\sim \neg x)$$

hold in $S(\text{dRA})$.

- [1] K. Blount and C. Tsinakis: The structure of residuated lattices, *Internat. J. Algebra Comput.* **13**, no. 4, 437–461 (2003).
- [2] N. Galatos and P. Jipsen: The structure of generalized BI-algebras and weakening relation algebras, preprint, (2020).
- [3] N. Galatos, P. Jipsen, T. Kowalski, and H. Ono, *Residuated Lattices: An Algebraic Glimpse at Substructural Logics*, Studies in Logic and the Foundations of Mathematics, vol. 151, Elsevier B. V., 2007.
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Thank you!