

## Moving the quantifiers to the front! (Prenex Form.)

We have discussed how to decide the truth of a statement with the quantifiers in the front. But what if they are not in the front?

$$((\exists y)(\forall x)(x = y)) \rightarrow ((\forall x)(\exists y)(x = y))$$

We have rules to move quantifiers to the front, without altering the meaning.

- (1)  $\neg(\forall x)P \equiv (\exists x)(\neg P)$ .
- (2)  $\neg(\exists x)P \equiv (\forall x)(\neg P)$ .
- (3)  $P \vee ((\exists x)Q) \equiv (\exists x)(P \vee Q)$  if  $P$  does not depend on  $x$ .
- (4)  $P \vee ((\forall x)Q) \equiv (\forall x)(P \vee Q)$  if  $P$  does not depend on  $x$ .
- (5)  $P \wedge ((\exists x)Q) \equiv (\exists x)(P \wedge Q)$  if  $P$  does not depend on  $x$ .
- (6)  $P \wedge ((\forall x)Q) \equiv (\forall x)(P \wedge Q)$  if  $P$  does not depend on  $x$ .

For example,

$$((\exists y)(\forall x)(x = y)) \rightarrow ((\forall x)(\exists y)(x = y)) \equiv (\forall s)(\exists t)(\forall x)(\exists y)((s = t) \rightarrow (x = y))$$

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### Practice!

Write the following statements in a logically equivalent form with quantifiers at the front.

- (1) If every lumberjack is hungry, then some lumberjack is hungry.

- (2)  $(\exists x)P(x) \leftrightarrow (\exists x)Q(x)$

- (3)  $(\forall x)(\forall y)((x < y) \rightarrow (\exists z)(x < z < y))$

## Restricted quantifiers!

Often one sees something like

$$(\forall \varepsilon > 0)(\exists \delta > 0)(|x - a| < \varepsilon \rightarrow |f(x) - f(a)| < \delta).$$

Question: What does it mean to write  $(\forall \varepsilon > 0)$ ? More generally, if  $C(x)$  is a condition on  $x$  and  $P(x)$  is a statement about  $x$ , what does  $((\forall x)C(x))P(x)$  mean?

Answer:  $((\forall x)C(x))P(x)$  is an abbreviation for

$$(\forall x)(C(x) \rightarrow P(x)),$$

and  $((\exists x)C(x))P(x)$  is an abbreviation for

$$(\exists x)(C(x) \wedge P(x)).$$

We call  $((\forall x)C(x))$  and  $((\exists x)C(x))$  *restricted quantifiers*. They behave just like ordinary quantifiers in the sense that

- (1) Rules for logical equivalence are the same:
  - (a)  $\neg((\forall x)C(x))P(x) \equiv ((\exists x)C(x))(\neg P(x))$ .
  - (b)  $P \vee (((\exists x)C(x))Q(x)) \equiv ((\exists x)C(x))(P \vee Q(x))$  if  $P$  does not depend on  $x$ .
  - (c) ETC
- (2) Quantifier games are played the same way. For example, to determine the truth of

$$(\forall x > 0)(\exists y < 1)(x < y)$$

we play a game where  $\forall$  first chooses  $x$  *satisfying the condition*  $x > 0$ , then  $\exists$  chooses  $y$  *satisfying the condition*  $y < 1$ .

## Practice!

- (1) Is  $(\forall x > 0)(\exists y < 1)(x < y)$  true in  $\mathbb{R}$ ? In  $\mathbb{N}$ ? In each case give a strategy for the appropriate quantifier.

- (2) Move the restricted quantifiers to the front:  $(\exists x > 0)P(x) \leftrightarrow (\exists x > 0)Q(x)$