Induction proofs: examples and nonexamples

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Exercise. Prove $1^2 + 2^2 + \dots + n^2 = n(n+1)(2n+1)/6$ in a similar way.

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Let S_n be the statement that "*n* is a product of prime numbers". Let's prove S_n by Strong Induction for n = 1, 2, 3, ...

Base Case(s). (n = 1, 2, 3, 4) S_1 : 1 is a product of zero primes. \checkmark S_2 : 2 is a product of one prime. \checkmark S_3 : 3 is a product of one prime. \checkmark S_4 : $4 = 2 \cdot 2$ is a product of two primes. \checkmark

Inductive Step. Assume that $n \ge 4$ and that 1, 2, ..., n are all products of primes. We must argue that n + 1 is a product of primes.

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$$N_k = \binom{k}{4} + \binom{k-1}{2} + \binom{k}{1} = \frac{k^4 - 6k^3 + 23k^2 - 18k + 24}{24}.$$

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We did not prove the Base Case. (In fact, the statement is false.)

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Corollary.

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Corollary. All horses are white.



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We only proved that all finite decimal approximations π are rational, but not that the limit π is rational.