

Moving the quantifiers to the front! (Prenex Form.)

We have discussed how to decide the truth of a statement with the quantifiers in the front. But what if they are not in the front?

$$((\exists y)(\forall x)(x = y)) \rightarrow ((\forall x)(\exists y)(x = y))$$

We have rules to move quantifiers to the front, without altering the meaning.

- (1) $\neg(\forall x)P \equiv (\exists x)(\neg P)$.
- (2) $\neg(\exists x)P \equiv (\forall x)(\neg P)$.
- (3) $P \vee ((\exists x)Q) \equiv (\exists x)(P \vee Q)$ if P does not depend on x .
- (4) $P \vee ((\forall x)Q) \equiv (\forall x)(P \vee Q)$ if P does not depend on x .
- (5) $P \wedge ((\exists x)Q) \equiv (\exists x)(P \wedge Q)$ if P does not depend on x .
- (6) $P \wedge ((\forall x)Q) \equiv (\forall x)(P \wedge Q)$ if P does not depend on x .

For example,

$$((\exists y)(\forall x)(x = y)) \rightarrow ((\forall x)(\exists y)(x = y)) \equiv (\forall s)(\exists t)(\forall x)(\exists y)((s = t) \rightarrow (x = y))$$

Practice!

Write the following statements in a logically equivalent form with quantifiers at the front.

- (1) If every lumberjack is hungry, then some lumberjack is hungry.

(2) $(\exists x)P(x) \leftrightarrow (\exists x)Q(x)$

(3) $(\forall x)(\forall y)((x < y) \rightarrow (\exists z)(x < z < y))$

Restricted quantifiers!

Often one sees something like

$$(\forall \varepsilon > 0)(\exists \delta > 0)(|x - a| < \varepsilon \rightarrow |f(x) - f(a)| < \delta).$$

Question: What does it mean to write $(\forall \varepsilon > 0)$? More generally, if $C(x)$ is a condition on x and $P(x)$ is a statement about x , what does $((\forall x)C(x))P(x)$ mean?

Answer: $((\forall x)C(x))P(x)$ is an abbreviation for

$$(\forall x)(C(x) \rightarrow P(x)),$$

and $((\exists x)C(x))P(x)$ is an abbreviation for

$$(\exists x)(C(x) \wedge P(x)).$$

We call $((\forall x)C(x))$ and $((\exists x)C(x))$ *restricted quantifiers*. They behave just like ordinary quantifiers in the sense that

- (1) Rules for logical equivalence are the same:
 - (a) $\neg((\forall x)C(x))P(x) \equiv ((\exists x)C(x))(\neg P(x))$.
 - (b) $P \vee ((\exists x)C(x))Q(x) \equiv ((\exists x)C(x))(P \vee Q(x))$ if P does not depend on x .
 - (c) ETC
- (2) Quantifier games are played the same way. For example, to determine the truth of

$$(\forall x > 0)(\exists y < 1)(x < y)$$

we play a game where \forall first chooses x *satisfying the condition* $x > 0$, then \exists chooses y *satisfying the condition* $y < 1$.

Practice!

- (1) Is $(\forall x > 0)(\exists y < 1)(x < y)$ true in \mathbb{R} ? In \mathbb{N} ? In each case give a strategy for the appropriate quantifier.

- (2) Move the restricted quantifiers to the front: $(\exists x > 0)P(x) \leftrightarrow (\exists x > 0)Q(x)$