

Solutions to HW 5.

Addition

$$\begin{array}{ll} m + 0 & := m \quad (\text{Initial Condition for } +) \\ m + S(n) & := S(m + n) \quad (\text{Recurrence Relation (or Propagation Rule) for } +) \end{array}$$

Multiplication

$$\begin{array}{ll} m \cdot 0 & := 0 \quad (\text{Initial Condition for } \cdot) \\ m \cdot S(n) & := m \cdot n + m \quad (\text{Recurrence Relation for } \cdot) \end{array}$$

1. Prove that $m(n + k) = (mn) + (mk)$ holds for the natural numbers.

This is a proof by induction on k .

(Base Case: $k = 0$)

$$\begin{array}{ll} m(n + 0) & = mn \quad (\text{IC, } +) \\ & = mn + 0 \quad (\text{IC, } +) \\ & = mn + m0 \quad (\text{IC, } \cdot) \end{array}$$

(Inductive Step: Assume true for k , prove true for $S(k)$)

$$\begin{array}{ll} m(n + S(k)) & = mS(n + k) \quad (\text{RR, } +) \\ & = m(n + k) + m \quad (\text{RR, } \cdot) \\ & = (mn + mk) + m \quad (\text{IH}) \\ & = mn + (mk + m) \quad (\text{Associative Law, } +) \\ & = mn + mS(k) \quad (\text{RR, } \cdot) \end{array}$$

2. Prove that $m(nk) = (mn)k$ holds for the natural numbers.

This is a proof by induction on k .

(Base Case: $k = 0$)

$$\begin{array}{ll} m(n0) & = m0 \quad (\text{IC, } \cdot) \\ & = 0 \quad (\text{IC, } \cdot) \\ & = (mn)0 \quad (\text{IC, } \cdot) \end{array}$$

(Inductive Step: Assume true for k , prove true for $S(k)$)

$$\begin{array}{ll} m(n \cdot S(k)) & = m(nk + n) \quad (\text{RR, } \cdot) \\ & = m(nk) + mn \quad (\text{Distributive Law}) \\ & = (mn)k + mn \quad (\text{IH}) \\ & = (mn) \cdot S(k) \quad (\text{RR, } \cdot) \end{array}$$

3. Prove that $mn = nm$ holds for the natural numbers. (Some lemmas will be needed.)

Lemma 1. $0k = 0$.

Proof. This is a proof by induction on k .

(Base Case: $k = 0$)

$$00 = 0 \quad (\text{IC}, \cdot)$$

(Inductive Step: Assume true for k , prove true for $S(k)$)

$$\begin{aligned} 0 \cdot S(k) &= 0k + 0 & (\text{RR}, \cdot) \\ &= 0 + 0 & (\text{IH}) \\ &= 0 & (\text{IC}, +) \end{aligned}$$

Lemma 2. (Right distributivity) $(m + n)k = mk + nk$.

Proof. This is a proof by induction on k .

(Base Case: $k = 0$)

$$\begin{aligned} (m + n)0 &= 0 & (\text{IC}, \cdot) \\ &= 0 + 0 & (\text{IC}, +) \\ &= m0 + n0 & (\text{IC}, \cdot) \end{aligned}$$

(Inductive Step: Assume true for k , prove true for $S(k)$)

$$\begin{aligned} (m + n) \cdot S(k) &= (m + n)k + (m + n) & (\text{RR}, \cdot) \\ &= (mk + nk) + (m + n) & (\text{IH}) \\ &= mk + (nk + (m + n)) & (\text{Associative Law}, +) \\ &= mk + ((nk + m) + n) & (\text{Associative Law}, +) \\ &= mk + ((m + nk) + n) & (\text{Commutative Law}, +) \\ &= mk + (m + (nk + n)) & (\text{Associative Law}, +) \\ &= (mk + m) + (nk + n) & (\text{Associative Law}, +) \\ &= (m \cdot S(k)) + (n \cdot S(k)) & (\text{RR}, \cdot) \end{aligned}$$

Solution to the Problem.

Proof. We prove that $mn = nm$ by induction on n .

(Base Case: $n = 0$)

$$\begin{aligned} m0 &= 0 & (\text{IC}, \cdot) \\ &= 0m & (\text{Lemma 1}) \end{aligned}$$

(Inductive Step: Assume true for n , prove true for $S(n)$)

$$\begin{aligned} m \cdot S(n) &= mn + m & (\text{RR}, \cdot) \\ &= nm + m & (\text{IH}) \\ &= (n + 1)m & (\text{Lemma 2}) \\ &= S(n) \cdot m & (S(n) = n + 1) \end{aligned}$$