

Proof. the class of well ordered sets is not elementary

Lemma. *A linear ordered set is Well Ordered IFF it does not contain an infinite and descending sequence.*

Let the language of well ordered sets be $\mathcal{L} = \{<\}$.

Suppose for the sake of contradiction some theory \mathcal{T} axiomatizes the class of well ordered sets.

Expand the language \mathcal{L} to \mathcal{L}_2 by including the constants $\{a_1, a_2, \dots\}$. Consider the set of sentences:

$$S_1 := \{(a_m > a_n) \mid n > m\}.$$

Let $\mathcal{T}_2 = \mathcal{T} \cup S_1$. We will show that any finite subset of this theory is satisfiable. Let \mathcal{T}_f be an arbitrary finite subset of \mathcal{T}_2 . \mathcal{T}_f will contain a finite number of sentences in S_1 , $\{s_{i_1}, \dots, s_{i_k}\}$. Let W be the finite set of constants that appear in $\{s_{i_1}, \dots, s_{i_k}\}$. The Sentences in S_1 define an order on W as a finite descending chain. This set is a well ordered set so W satisfies \mathcal{T}_f . Therefore W satisfies \mathcal{T} and $\{s_{i_1}, \dots, s_{i_k}\}$.

Because every finite subset of \mathcal{T}_2 is satisfiable, by the compactness theorem \mathcal{T}_2 is satisfiable. However the model of \mathcal{T}_2 is a well ordered set with an infinite descending chain and a well ordered set cannot contain an infinite descending chain, thus contradicting the initial assumption that there exists a theory of well ordered sets.

□