

Let  $t(n)$  be the least number  $N$  such that there exists some infinite structure with exactly  $N$   $n$ -types over the empty set. We will show there exists a single infinite structure with exactly  $t(n)$   $n$ -types for each  $n$  and classify these structures.

*Pf.* For an equivalence relation  $E$  on  $\{1, \dots, n\}$ , define  $\varepsilon_E(x_1, \dots, x_n)$  to be the conjunction of all atomic formulas  $(x_i = x_j)$ ,  $(i, j) \in E$ , and all negated atomic formulas  $\neg(x_i = x_j)$ ,  $(i, j) \notin E$ . Any  $n$ -tuple of any infinite structure satisfies exactly one formula of the form  $\varepsilon_E(x_1, \dots, x_n)$ , namely the one which expresses which of its coordinates are equal. For example, if  $a \neq b$ , then the triple  $(a, a, b)$  whose first two coordinates are equal and distinct from the third is associated to the equivalence relation  $E$  on coordinate set  $\{1, 2, 3\}$  that partitions this set into  $E$ -classes  $\{1, 2\}$  and  $\{3\}$ , the corresponding formula is

$$\varepsilon_E(x_1, x_2, x_3) : (x_1 = x_2) \wedge \neg(x_1 = x_3) \wedge \neg(x_2 = x_3).$$

Each complete type contains some formula of the form  $\varepsilon_E$ , no complete type contains more than one  $\varepsilon_E$ , but it is possible that different complete types contain the same formula  $\varepsilon_E$ . (That is, it is possible for  $\bar{a}$  and  $\bar{b}$  to have different complete types with respect to the language  $L$  even if they have the same type with respect to the language of equality.)

Altogether, this shows that  $t(n)$  is at least as large as the number of equivalence relations on an  $n$ -element set. (This number is called the  $n$ -th *Bell number*, and is sometimes written  $B_n$ .) On the other hand, an infinite pure  $M$  set has no more than  $B_n$  complete  $n$ -types. [To verify this we must argue that no two tuples with different complete types contain the same  $\varepsilon_E$ . Suppose to the contrary that  $\bar{a}, \bar{b} \in M^n$  satisfy the same  $\varepsilon_E$  but  $\text{tp}^M(\bar{a}) \neq \text{tp}^M(\bar{b})$ . The partial mapping  $a_i \mapsto b_i$ ,  $1 \leq i \leq n$ , can be extended to a permutation of the pure set  $M$ . A permutation of a pure set  $M$  is an automorphism of the structure  $\langle M; \emptyset \rangle$ , which must preserve types, hence  $\text{tp}^M(\bar{a}) = \text{tp}^M(\bar{b})$ . This contradicts the assumption that  $\text{tp}^M(\bar{a}) \neq \text{tp}^M(\bar{b})$ .] All in all, this shows that  $|S_n^M(\emptyset)| \geq B_n$  for any infinite structure and that equality holds when  $M$  is an infinite pure set. Hence the function  $t$  has been determined, it is  $t(n) = B_n$ . An infinite pure set  $M$  has exactly  $t(n)$  complete  $n$ -types for each  $n$ , and the map defined by  $p \mapsto E$  iff  $\varepsilon_E \in p$  is a bijection from  $S_n^M(\emptyset)$  to the set of equivalence relations on  $\{1, \dots, n\}$ . (Write  $p_E$  for  $p$  if  $\varepsilon_E \in p$ .)

Now we classify the structures  $M = \langle M; \mathcal{R}, \mathcal{F}, \mathcal{C} \rangle$  with  $|S_n^M(\emptyset)| = B_n$ . It follows from the previous paragraph that each complete  $n$ -type is isolated by some complete formula  $\varepsilon_E$ . We want to understand the meaning of all formulas. That is, we want to say that, because a typical clopen set  $O_{\varphi(\bar{x})}$  is determined by the points  $p_E$  that it contains, there should be some sense for which the meaning of the corresponding formula  $\varphi(\bar{x})$  is determined by the corresponding complete formulas  $\varepsilon_E$  for  $p_E \in O_{\varphi(\bar{x})}$ . We establish such a fact now.

We have  $M \models \forall \bar{x} (\varepsilon_E(\bar{x}) \rightarrow \varphi(\bar{x}))$  whenever  $\varphi(\bar{x}) \in p_E$ , so

$$M \models \forall \bar{x} \left( \left( \bigvee_{\varphi(\bar{x}) \in p_E} \varepsilon_E(\bar{x}) \right) \rightarrow \varphi(\bar{x}) \right).$$

We can replace  $\rightarrow$  with  $\leftrightarrow$  in this displayed line for the following reason. If

$$\mathcal{M} \not\models \forall \bar{x} \left( \varphi(\bar{x}) \rightarrow \left( \bigvee_{\varphi(\bar{x}) \in p_E} \varepsilon_E(\bar{x}) \right) \right),$$

then

$$\mathcal{M} \models \exists \bar{x} \left( \varphi(\bar{x}) \wedge \neg \left( \bigvee_{\varphi(\bar{x}) \in p_E} \varepsilon_E(\bar{x}) \right) \right),$$

which implies that  $\mathcal{M}$  has an  $n$ -tuple  $\bar{a}$  such that  $\text{tp}(\bar{a})$  contains  $\varphi(\bar{x})$ , but does not contain any  $\varepsilon_E(\bar{x})$  that strengthens  $\varphi(\bar{x})$ . This would force  $\text{tp}(\bar{a})$  to contain no  $\varepsilon_E(\bar{x})$  at all, which is impossible since every tuple satisfies some formula  $\varepsilon_E(\bar{x})$  that expresses the equalities among coordinates.

The previous paragraph shows that, when  $|S_n^{\mathcal{M}}(\emptyset)| = B_n$ , any formula  $\varphi(\bar{x})$  is equivalent modulo  $\text{Th}(\mathcal{M})$  to a formula  $\bigvee_{\varphi(\bar{x}) \in p_E} \varepsilon_E(\bar{x})$  in the language of equality.

Conversely, if  $\mathcal{M}$  is such that any formula  $\varphi(\bar{x})$  is equivalent modulo  $\text{Th}(\mathcal{M})$  to a formula expressed in the language of equality, then  $\mathcal{M}$  has the same types as the underlying pure set, hence  $|S_n^{\mathcal{M}}(\emptyset)| = B_n$ .

We have found the function  $t(n)$  and found that an infinite pure set has exactly  $t(n)$  complete  $n$ -types for each  $n$ . We have also fully characterized the infinite structures  $\mathcal{M} = \langle M; \mathcal{R}, \mathcal{F}, \mathcal{C} \rangle$  where  $|S_n^{\mathcal{M}}(\emptyset)| = B_n$ , namely they are the structures where every formula is equivalent modulo  $\text{Th}(\mathcal{M})$  to a formula expressed in the language of equality. Thus, our work is done, but let's briefly explore the implications of this conclusion.

First,  $\text{Aut}(\mathcal{M})$  must be the full symmetric group on  $M$ . (This is another characterization of those structures  $\mathcal{M}$  where  $|S_n^{\mathcal{M}}(\emptyset)| = B_n$ .) The fact that  $\text{Aut}(\mathcal{M})$  is the symmetric group on  $M$  implies, in particular, that  $\mathcal{M}$  has no constant symbols. Next, here is an example of a nontrivial ternary relation expressible in the language of equality:

$$(a, b, c) \in R^{\mathcal{M}} \quad \text{iff} \quad (a = b \neq c) \vee (a \neq b = c).$$

Here is an example of a nontrivial ternary function expressible in the language of equality:

$$F^{\mathcal{M}}(a, b, c) = \begin{cases} c & \text{if } a = b \\ a & \text{else.} \end{cases}$$

These examples suggest what a typical relation or function of  $\mathcal{M}$  might look like. (More specifically, if  $R$  is  $m$ -ary, then  $R^{\mathcal{M}}$  is a union of  $\text{Sym}(M)$ -orbits of  $M^m$ , while if  $F$  is  $n$ -ary, then  $F^{\mathcal{M}}$  must be a function for which the relation  $F^{\mathcal{M}}(x_1, \dots, x_n) = x_{n+1}$  is a union of  $\text{Sym}(M)$ -orbits of  $M^{n+1}$ .)