

# Complexes and Exact Sequences in Abelian Categories

Homological Algebra

Feb 10, 2010

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*Proof.*  $\mu = [\ker(\text{coker}(\varphi))]$  has this property. There is a unique  $\varepsilon$  such that

$$A \xrightarrow{\varepsilon} K \xrightarrow{\mu} B \xrightarrow{\gamma} C$$

and  $\mu\varepsilon = \varphi$ , so  $\varphi$  factors through  $\mu$ . HW2.5(a) shows that  $\varepsilon$  is epi.

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Now suppose that  $\varphi$  factors through some other monomorphism

$$A \xrightarrow{\varepsilon'} K' \xrightarrow{\mu'} B.$$

By factoring  $\varepsilon$  if necessary, we may assume that it is epi. Now the HW2.5(b) proves that  $\mu \leq \mu'$ . □

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**Thm.** (HW2.5(a))

$A \xrightarrow{\varphi} B$  factors as  $A \xrightarrow{\varepsilon} I \xrightarrow{\mu} B$  where  $\mu$  represents the image and  $\varepsilon$  represents the coimage. Either of the maps uniquely determines the other.

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**Cor.** (HW2.6(b))

If  $\varphi$  is a monomorphism, then  $\text{im}(\varphi) = [\varphi]$ .

# Complexes and exact sequences

A sequence

$$\dots \xrightarrow{d_{n+1}} C_n \xrightarrow{d_n} C_{n-1} \xrightarrow{d_{n-1}} \dots$$

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Exact sequences are complexes: recall that  $\operatorname{im}(d_{n+1})$  is the smallest subobject of  $C_n$  such that  $d_{n+1}$  factors as  $\cdots \xrightarrow{d_{n+2}} C_{n+1} \xrightarrow{\varepsilon} I = K \xrightarrow{\mu} C_n \xrightarrow{d_n} \cdots$  with  $\mu$  representing  $\operatorname{im}(d_{n+1})$ .

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- $0 \rightarrow A \rightarrow B \rightarrow C \rightarrow 0$  is exact iff  $A \rightarrow B$  is a monomorphism and  $B \rightarrow C$  is its cokernel.

# Pullbacks

A diagram

$$\begin{array}{ccc} P & \xrightarrow{\pi_B} & B \\ \pi_A \downarrow & & \beta \downarrow \\ A & \xrightarrow{\alpha} & C \end{array}$$

is a *pullback square* if for every pair of maps  $\mu: X \rightarrow A$  and  $\nu: X \rightarrow B$  such that

$$\begin{array}{ccc} X & \xrightarrow{\nu} & B \\ \mu \downarrow & & \beta \downarrow \\ A & \xrightarrow{\alpha} & C \end{array}$$

commutes there is a unique  $\lambda: X \rightarrow P$  such that  $\mu = \pi_A \lambda$  and  $\nu = \pi_B \lambda$ .  
 $(P, \pi_A, \pi_B)$  is the *pullback* of  $\alpha$  and  $\beta$ . *Pushouts* are defined dually.

# Converting squares to sequences

From the (possibly noncommutative) diagram

$$\begin{array}{ccc} P & \xrightarrow{\delta} & B \\ \gamma \downarrow & & \downarrow \beta \\ A & \xrightarrow{\alpha} & C \end{array}$$

we can construct the sequence  $P \xrightarrow{i_A\gamma + i_B\delta} A \oplus B \xrightarrow{\alpha\pi_A - \beta\pi_B} C$ . The composition of the maps in the sequence is  $\alpha\gamma - \beta\delta$ , so the diagram commutes iff the sequence is a complex. Conversely, given  $P \xrightarrow{\sigma} A \oplus B \xrightarrow{\tau} C$ , we can construct

$$\begin{array}{ccc} P & \xrightarrow{\pi_B\sigma} & B \\ \pi_A\sigma \downarrow & & \downarrow -\tau i_B \\ A & \xrightarrow{\tau i_A} & C \end{array}$$

where the difference of the two paths is  $\tau\sigma$ .

# Characterization of pullbacks

$$\begin{array}{ccc} P & \xrightarrow{\delta} & B \\ \gamma \downarrow & & \beta \downarrow \\ A & \xrightarrow{\alpha} & C \end{array} \text{ is a pullback square iff } 0 \longrightarrow P \longrightarrow A \oplus B \longrightarrow C \text{ is exact.}$$

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*Sketch of Proof.* ( $\Leftarrow$ )

$$\begin{array}{ccc}
 X & \xrightarrow{\nu} & B \\
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( $\Rightarrow$ ) Compare complex  $P \longrightarrow A \oplus B \longrightarrow C$  with exact sequence  $0 \longrightarrow K \xrightarrow{\kappa} A \oplus B \longrightarrow C$ . Use that  $P$  is a pullback to show that  $P \rightarrow A \oplus B$  is equivalent to  $K \rightarrow A \oplus B$ .



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*Idea of Proof.* Given  $\alpha: A \rightarrow C$  and  $\beta: B \rightarrow C$ , construct  $P$  as the kernel of the induced map  $A \oplus B \rightarrow C$ .

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Proofs given in class.