

The Shape of Congruence Lattices

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tame congruence theory?

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ABSTRACT. We develop the theories of the strong commutator, the rectangular commutator, the strong rectangular commutator, as well as a solvability theory for the nonmodular TC commutator. These theories are used to show that each of the following sets of statements are equivalent for a variety \mathcal{V} of algebras.

are there analogs of Theorems 9.14 and 9.15 from Hobby-McKenzie that could be proved?

- (I)
 - (a) \mathcal{V} satisfies a nontrivial congruence identity.
 - (b) \mathcal{V} satisfies an idempotent Maltsev condition that fails in the variety of semilattices.
 - (c) The rectangular commutator is trivial throughout \mathcal{V} .
- (II)
 - (a) \mathcal{V} satisfies a nontrivial meet continuous congruence identity.
 - (b) \mathcal{V} satisfies an idempotent Maltsev condition that fails in the variety of sets.
 - (c) The strong commutator is trivial throughout \mathcal{V} .
 - (d) The strong rectangular commutator is trivial throughout \mathcal{V} .
- (III)
 - (a) \mathcal{V} is congruence semidistributive.
 - (b) \mathcal{V} satisfies an idempotent Maltsev condition that fails in the variety of semilattices and in any nontrivial variety of modules.
 - (c) The rectangular and TC commutators are both trivial throughout \mathcal{V} .

We prove that a residually small variety that satisfies a congruence identity is congruence modular.

*Dedicated to Bjarni Jónsson,
For his work on congruence varieties,*

*To Ralph McKenzie,
For his work on commutator theories,*

*And to Walter Taylor,
For his work on Maltsev conditions.*

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CHAPTER 1

Introduction

This monograph is concerned with the relationships between Mal'tsev conditions, commutator theories and the shapes of congruence lattices in varieties of algebras.

1.1. Shapes of Congruence Lattices

Carl F. Gauss, in [23], introduced the notation

$$(1.1) \quad a \equiv b \pmod{m},$$

which is read as “ a is congruent to b modulo m ”, to mean that the integers a and b have the same remainder upon division by the integer modulus m , equivalently that $a - b \in m\mathbb{Z}$. As the notation suggests, congruence modulo m is an equivalence relation on \mathbb{Z} . It develops that congruence modulo m is compatible with the ring operations of \mathbb{Z} , and that the only equivalence relations on \mathbb{Z} that are compatible with the ring operations are congruences modulo m for $m \in \mathbb{Z}$.

Richard Dedekind conceived of a more general notion of “integer”, which nowadays we call an ideal in a number ring. Dedekind extended the notation (1.1) to

$$(1.2) \quad a \equiv b \pmod{\mu}$$

where $a, b \in \mathbb{C}$ and $\mu \subseteq \mathbb{C}$; (1.2) is defined to hold if $a - b \in \mu$. Dedekind called a subset $\mu \subseteq \mathbb{C}$ a **module** if it could serve as the modulus of a congruence, i.e., if this relation of congruence modulo μ is an equivalence relation on \mathbb{C} . This happens precisely when μ is closed under subtraction. For Dedekind, therefore, a “module” was an additive subgroup of \mathbb{C} .

The set of Dedekind’s modules is closed under the operations of intersection and sum. These two operations make the set of modules into a lattice. Dedekind proposed and investigated the problem of determining the identities of this lattice (the “laws of congruence arithmetic”). In 1900, in [13], he published the discovery that if $\alpha, \beta, \gamma \subseteq \mathbb{C}$ are modules, then

$$(1.3) \quad \alpha \cap (\beta + (\alpha \cap \gamma)) = (\alpha \cap \beta) + (\alpha \cap \gamma).$$

quasi-identity and quasi-affine are hyphenated, but quasiorder and quasivariety aren't.

1. INTRODUCTION

This 3-variable law of the lattice of modules is now called the **modular law**. Dedekind went on to prove that any equational law of congruence arithmetic that can be expressed with at most 3-variables is a consequence of the modular law and the laws valid in all lattices.

Dedekind did not write the law in the form (1.3), which is an identity, but rather as a quasi-identity: for all modules $\alpha, \beta, \gamma \subseteq \mathbb{C}$

$$(1.4) \quad \alpha \supseteq \gamma \longrightarrow \alpha \cap (\beta + \gamma) = (\alpha \cap \beta) + \gamma.^1$$

Dedekind also discovered a useful “omitting sublattices” version of the modular law. It is the assertion that there do not exist modules $\alpha, \beta, \gamma \subseteq \mathbb{C}$ which generate a sublattice isomorphic to \mathbf{N}_5 (Theorem I.7.12 of [1]).

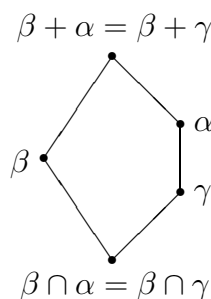


FIGURE 1.1. The lattice \mathbf{N}_5

More generally, a **congruence** on an arbitrary algebra \mathbf{A} is an equivalence relation on the universe of \mathbf{A} that is compatible with the operations of \mathbf{A} . Equivalently, it is the kernel of a homomorphism with domain \mathbf{A} . The set of all congruences is a sublattice of the lattice of equivalence relations on A . A **congruence identity** of \mathbf{A} is an identity that holds in the lattice $\mathbf{Con}(\mathbf{A})$ of all congruences of \mathbf{A} . Dedekind’s observation is that the modular law is a congruence identity satisfied by \mathbb{C} as an abelian group. In fact, it is a congruence identity of any group, ring, vector space, Boolean algebra or lattice.

Dedekind’s result initiated many lines of research in the 20th century, but to avoid losing focus we mention only a few. Garrett Birkhoff observed that the congruences of any group **permute**, meaning that $\alpha \circ \beta = \beta \circ \alpha$, and that any lattice of permuting equivalence relations is modular. This “explains” Dedekind’s result to some degree, but in [22] N. Funayama and T. Nakayama showed that the lattice of congruences of a lattice satisfies the **distributive law**: for all α, β, γ

$$(1.5) \quad \alpha \cap (\beta + \gamma) = (\alpha \cap \beta) + (\alpha \cap \gamma),$$

¹In fact, Dedekind used the symbols $+$ and $-$ instead of $+$ and \cap .

which is stronger than the modular law, yet congruences of lattices need not permute. Thus permutability implies modularity, but not vice versa. Bjarni Jónsson then refined the results of Birkhoff and Dedekind by showing in [37] that any lattice of permuting equivalence relations satisfies the arguesian law, which is stronger than the modular law. Conversely, he showed that any complemented lattice satisfying the arguesian law is embeddable into the lattice of congruences of an abelian group. Here the **arguesian law** is the 6-variable law asserting that for all $\alpha_i, \beta_i, i = 0, 1, 2$,

$$(1.6) \quad (\alpha_0 + \beta_0) \cap (\alpha_1 + \beta_1) \cap (\alpha_2 + \beta_2) \leq \alpha_0 \cap (\gamma + \alpha_1) + \beta_0 \cap (\gamma + \beta_1)$$

where $\gamma = (\alpha_0 + \alpha_1) \cap (\beta_0 + \beta_1) \cap (\delta_0 + \delta_1)$, $\delta_0 = (\alpha_0 + \alpha_2) \cap (\beta_0 + \beta_2)$ and $\delta_1 = (\alpha_1 + \alpha_2) \cap (\beta_1 + \beta_2)$. Thus Jónsson solved the special instance of Dedekind's problem (of determining the laws of modules) which concerns only complemented lattices of modules. The full problem is still open, although in the same paper Jónsson gave an example of a finite noncomplemented arguesian lattice that is not embeddable in the congruence lattice of a group, and in [30] Mark Haiman proved that lattices of permuting equivalence relations satisfy laws stronger than the arguesian law.

1.2. Maltsev Conditions

Let \mathcal{R} be the variety of rings and \mathcal{A} be the variety of abelian groups. The fact that rings have a term-definable underlying abelian group structure may be denoted $\mathcal{A} \leq \mathcal{R}$, i.e., rings have the structure of abelian groups and more. A. I. Maltsev proved in [62] that a variety \mathcal{V} of algebras has the property that the congruence lattices of its members consist of permuting equivalence relations if and only if $\mathcal{P} \leq \mathcal{V}$ where \mathcal{P} is the variety with one ternary basic operation symbol p that is axiomatized by the **Maltsev identities**:

$$(1.7) \quad p(x, y, y) \approx x \quad \text{and} \quad p(y, y, x) \approx x.$$

In other words, \mathcal{V} consists of congruence permutable algebras if and only if there is a ternary \mathcal{V} -term p , called a **Maltsev term**, such that the identities (1.7) hold in \mathcal{V} . Thus, one may refine Birkhoff's earlier explanation of Dedekind's modularity result to: groups are congruence modular because groups have a Maltsev term $p(x, y, z) = xy^{-1}z$, so groups have permuting congruences, and lattices of permuting equivalence relations are modular.

Let \mathcal{U} be a finitely presented² variety. The condition on a variety \mathcal{V} that $\mathcal{U} \leq \mathcal{V}$ may be expressed as “there is a finite set of \mathcal{V} -terms corresponding to the basic operation symbols of \mathcal{U} such that the finite set of identities corresponding to the axioms of \mathcal{U} hold in \mathcal{V} .” A condition of this type is called a **strong Maltsev condition**, and given a finitely presented variety \mathcal{U} the class $\{\mathcal{V} \mid \mathcal{U} \leq \mathcal{V}\}$ is the **strong Maltsev class** that is **defined** by this condition. Given a descending sequence $\cdots \leq \mathcal{U}_2 \leq \mathcal{U}_1 \leq \mathcal{U}_0$ of finitely presented varieties, the class $\{\mathcal{V} \mid \exists n(\mathcal{U}_n \leq \mathcal{V})\}$ is the **Maltsev class** that is defined by this sequence, and the **Maltsev condition** associated to this sequence is the assertion that for some n there is a finite set of \mathcal{V} -terms corresponding to the basic operation symbols of \mathcal{U}_n such that a finite set of identities corresponding to the axioms of \mathcal{U}_n hold in \mathcal{V} .

"the" instead of "a"

After Maltsev published the condition from (1.7) defining the class of varieties with permuting congruences, Alden Pixley found in [72] a strong Maltsev condition defining the class of varieties with distributive and permuting congruences, B. Jónsson found in [38] a Maltsev condition defining the class of varieties with distributive congruence lattices, and Alan Day found in [8] a Maltsev condition defining the class of varieties with modular congruence lattices. These results (and many others like them) culminated in the theorem, obtained independently by Pixley [73] and Rudolf Wille [78], that if ε is any lattice identity, then the class of varieties whose congruence lattices satisfy ε is the intersection of countably many Maltsev classes. This proves, in particular, that if two varieties satisfy the same Maltsev conditions, then they satisfy the same congruence identities. This result of Pixley and Wille includes an algorithm for generating Maltsev conditions associated with congruence identities, paving the way for a deeper study of congruence identities.

J. B. Nation discovered in [69] that there exist lattice identities that are strictly weaker than the modular law when considered as lattice identities, but equivalent to the modular law when considered as congruence identities. That is, there exists a variety \mathcal{W} of lattices that strictly contains all modular lattices, yet any variety \mathcal{V} of algebras with the property that the congruence lattice of every algebra from \mathcal{V} is found in \mathcal{W} is actually congruence modular. Nation’s Theorem brought attention to the possibility that there might be “few” different **congruence varieties**, which are varieties of lattices of the form

$$(1.8) \quad \text{CON}(\mathcal{V}) = \text{HSP}(\{\text{Con}(\mathbf{A}) \mid \mathbf{A} \in \mathcal{V}\}).$$

²I.e., \mathcal{U} is a finitely axiomatized variety with finitely many basic operation symbols.

For example, it suddenly became plausible that the only proper nondistributive congruence varieties might be those of the form $\text{Con}(\mathcal{M})$ where \mathcal{M} is a variety of modules. Nation's result led Ralph McKenzie to conjecture in [65] that if ε is any nontrivial lattice identity, then ε implies modularity when considered as a congruence identity. Studies of identities of small complexity yielded positive evidence for McKenzie's Conjecture in [9, 10, 21, 39, 68]. The most general of these results is the theorem of Ralph Freese and J. B. Nation that any lattice identity ε that can be written as an inclusion of the form

$$\bigwedge \bigvee \bigwedge \text{ variables} \leq \bigvee \bigwedge \bigvee \text{ variables}$$

implies modularity as a congruence identity. Similar results were proved at the same time about congruence identities implying congruence distributivity, and in [17] one finds the result that any congruence identity implying congruence modularity must also imply the stronger arguesian identity as a congruence identity.

McKenzie's Conjecture was refuted by S. V. Polin in his famous paper [74]. Polin constructed a locally finite variety \mathcal{P} that is not congruence modular, but satisfies a nontrivial congruence identity. To describe a congruence identity that holds in Polin's variety, let x, y and z be lattice variables and let $y_1 = y, z_1 = z, y_{n+1} = y + (x \cap z_n)$, and $z_{n+1} = z + (x \cap y_n)$. Let δ_n be the weakened distributive law:

$$x \cap (y + z) \approx (x \cap y_n) + (x \cap z_n).$$

The identity δ_1 is the usual distributive law. Although δ_2 is strictly weaker than distributivity as a lattice identity, it implies distributivity as a congruence identity.³ Polin's variety satisfies δ_3 as a congruence identity.⁴ Since it does not satisfy δ_2 , and since any lattice variety that satisfies δ_3 and not δ_2 is nonmodular, $\text{Con}(\mathcal{P})$ is an example of a proper, nondistributive congruence variety that differs from $\text{Con}(\mathcal{M})$ for any variety of modules \mathcal{M} .⁵ Polin's variety was thoroughly investigated by Day and Freese in [11] with the result that we now have an efficient algorithm for determining if an identity implies modularity as a congruence identity. In an unpublished manuscript, Day showed that

³We know this because Ralph Freese showed us that δ_2 fails in $\text{Con}(\mathbf{F}_{\mathcal{P}}(1))$. Using this observation and the main result of [11] it is easy to see that δ_2 implies distributivity as a congruence identity.

⁴Polin showed that \mathcal{P} satisfies a different congruence identity. Alan Day was the one to emphasize the importance of the identities δ_n .

⁵We now know that there are continuum many different nonmodular congruence varieties (see [11]), and even some congruence varieties of groups that differ from any congruence variety of modules (see [71]).

the Polin construction can be iterated⁶ to produce varieties satisfying weaker and weaker nontrivial congruence identities.

Not all of the preceding results were originally proved via Maltsev conditions, but they could have been, and the analysis of Maltsev conditions is a powerful method for obtaining further results like these. Moreover, there are results on congruence identities that have been proved by a careful analysis of the Maltsev conditions defining a given class of varieties that seem unreachable by any other method, such as Paolo Lipparini's results on congruence identities satisfied by congruence n -permutable varieties⁷ in [55, 57, 59, 60, 61].

1.3. Commutator Theories

Birkhoff's HSP Theorem asserts that if \mathcal{K} is a class of similar algebras, then any model of the identities true in \mathcal{K} may be constructed as a homomorphic image of a subalgebra of a product of algebras in \mathcal{K} , that is as \mathbf{B}/θ where $\mathbf{B} \leq \prod \mathbf{A}_i$, $\mathbf{A}_i \in \mathcal{K}$. Replacing each \mathbf{A}_i by the projection of \mathbf{B} onto the i -th factor, we find that a typical model has the form \mathbf{B}/θ where \mathbf{B} is a subdirect product of subalgebras of members of \mathcal{K} . It is therefore important to know how to construct congruences on subdirect products. The case of two factors is already difficult. If $\mathbf{B} \leq_{\text{sd}} \mathbf{A}_1 \times \mathbf{A}_2$ and $\alpha_i \in \text{Con}(\mathbf{A}_i)$, then the restriction of $\alpha_1 \times \alpha_2$ to \mathbf{B} is a **product congruence** on \mathbf{B} . All other congruences are **skew**. For

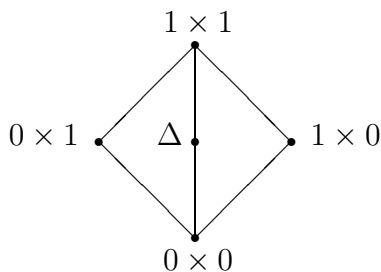


FIGURE 1.2. $\text{Con}(\mathbb{Z}_2 \times \mathbb{Z}_2)$

example, the congruence lattice of the group $\mathbf{B} = \mathbb{Z}_2 \times \mathbb{Z}_2$ is pictured in Figure 1.2. There is one skew congruence, Δ , which is the congruence that has the diagonal subgroup of $\mathbb{Z}_2 \times \mathbb{Z}_2$ as a class.

⁶This and other modifications of Polin's construction were rediscovered and appear in [53] and [70].

⁷Congruences α and β **n -permute** if the alternating composition $\alpha \circ_n \beta = \alpha \circ \beta \circ \alpha \cdots$, with $n - 1$ occurrences of \circ , equals $\beta \circ_n \alpha$. A variety is **congruence n -permutable** if the congruences on all members n -permute.

To understand skew congruences like this one, let us attempt to describe a typical “diagonal congruence” on a typical “diagonal subalgebra” \mathbf{B} of \mathbf{A}^2 for some algebra \mathbf{A} in a congruence permutable variety \mathcal{V} . Here a “diagonal subalgebra” is one containing the diagonal of \mathbf{A}^2 , and a “diagonal skew congruence” is a congruence generated by a set of pairs of elements from the diagonal. More formally, if $\delta: \mathbf{A} \rightarrow \mathbf{A}^2: a \mapsto (a, a)$ is the canonical diagonal embedding, then $\mathbf{B} \leq \mathbf{A}^2$ is a diagonal subalgebra if δ factors through the inclusion of \mathbf{B} into \mathbf{A}^2 . A diagonal congruence on \mathbf{B} is the extension to \mathbf{B} of a congruence on $\delta(\mathbf{A})$.

Since we have assumed that \mathbf{A} has a Maltsev term, the only diagonal subalgebras \mathbf{B} of \mathbf{A}^2 are those whose universe B is a congruence on \mathbf{A} , say $B = \beta$. Since the diagonal of \mathbf{A}^2 supports a subalgebra isomorphic to \mathbf{A} , the only diagonal congruences on \mathbf{B} are the extensions to \mathbf{B} of congruences $\delta(\alpha)$ for $\alpha \in \text{Con}(\mathbf{A})$. Thus, the construction of a typical diagonal congruence on a typical diagonal subalgebra $\mathbf{B} \leq \mathbf{A}^2$ involves a pair of congruences $\alpha, \beta \in \text{Con}(\mathbf{A})$. Specifically, we are considering \mathbf{B} to be the subalgebra supported by β , and a diagonal congruence $\Delta = \Delta_{\alpha, \beta}$ on \mathbf{B} that is generated by $\{ \langle (u, u), (v, v) \rangle \mid (u, v) \in \alpha \}$. We next try to understand how Δ is related to the product congruences of \mathbf{B} .

Since we have assumed that \mathbf{A} has a Maltsev term, and \mathbf{B} is in the variety generated by \mathbf{A} , $\text{Con}(\mathbf{B})$ is modular. It is a consequence of Dedekind’s analysis of 3-variable consequences of the modular law that the 3-generated free modular lattice has only 28 elements, hence any 3-generated modular lattice can be easily drawn. In particular, the

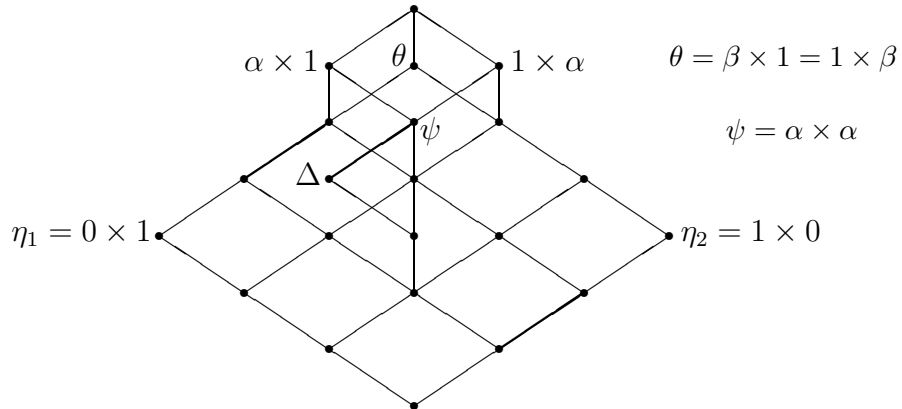


FIGURE 1.3. $\langle \eta_1, \eta_2, \Delta \mid \eta_1 \cap \eta_2 = 0 \rangle$

modular lattice with the presentation $\langle \eta_1, \eta_2, \Delta \mid \eta_1 \cap \eta_2 = 0 \rangle$ is the one

in Figure 1.3. Hence the sublattice of $\mathbf{Con}(\mathbf{B})$ that is generated by the projection kernels $\eta_1 = 0 \times 1$ and $\eta_2 = 1 \times 0$ together with $\Delta = \Delta_{\alpha, \beta}$ is a homomorphic image of this lattice. The smallest product congruence containing Δ , the “product cover” of Δ , is

$$(\Delta + (0 \times 1)) \cap (\Delta + (1 \times 0)) = \alpha \times \alpha = \psi.$$

We take the interval $I[\Delta, \psi]$, from Δ to its product cover, to be a measure of the skewness of Δ . Meeting with η_2 and joining with η_1 projects this interval in two steps to the isomorphic interval

$$(1.9) \quad I[\eta_1 + (\Delta \cap \eta_2), \eta_1 + (\psi \cap \eta_2)] = I[\eta_1 + (\Delta \cap \eta_2), (\alpha \cap \beta) \times 1],$$

so the interval in (1.9) can also be taken to be a measure of the skewness of Δ . But the interval in (1.9) lies entirely in the interval above the first coordinate projection kernel η_1 , which is naturally isomorphic to $\mathbf{Con}(\mathbf{A})$, so we can measure the skewness of Δ by considering the corresponding congruence interval of \mathbf{A} . The left lattice in Figure 1.4

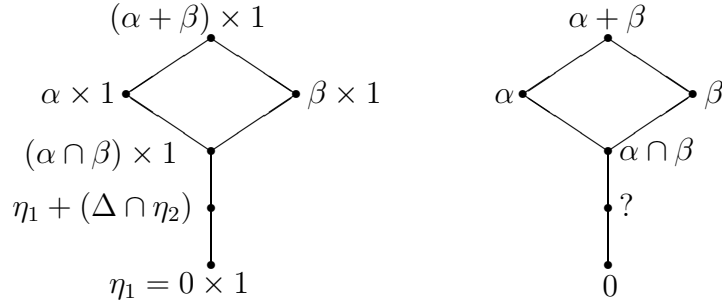


FIGURE 1.4. $I[0 \times 1, 1 \times 1]$ vs. $\mathbf{Con}(\mathbf{A})$

indicates the interval in $\mathbf{Con}(\mathbf{B})$ above η_1 , while the right lattice shows the corresponding congruences on \mathbf{A} . Note especially the question mark by the right lattice, which labels the congruence on \mathbf{A} induced by Δ . This induced congruence is a function of α and β , and will be written $[\alpha, \beta]$ and called the **commutator** of α and β . For groups it is the ordinary group commutator.⁸ The interval from the commutator $[\alpha, \beta]$ to the intersection $\alpha \cap \beta$ is the interval in $\mathbf{Con}(\mathbf{A})$ that measures the skewness of Δ .

In analogy with group theory, we call a congruence interval $I[\sigma, \tau]$ **abelian** if $[\tau, \tau] \leq \sigma$, and call \mathbf{A} abelian if $\mathbf{Con}(\mathbf{A})$ is an abelian congruence interval. If τ is defined to equal $\alpha \cap \beta$ in the last paragraph, then $\tau \leq \alpha$ and $\tau \leq \beta$ so $[\tau, \tau] \leq [\alpha, \beta]$. This shows that the interval $I[[\alpha, \beta], \alpha \cap \beta]$, which measures the skewness of Δ , is abelian.

⁸This means that if N_α and N_β are the normal subgroups corresponding to the subscripted congruences, then $[N_\alpha, N_\beta] = N_{[\alpha, \beta]}$.

Conversely, every abelian interval $I[\sigma, \tau]$ is contained in an interval that measures the skewness of some diagonal congruence, namely $\Delta_{\tau, \tau}$. Hence diagonal congruences fail to be skew in \mathcal{V} if and only if \mathcal{V} omits abelian congruence intervals if and only if the commutator trivializes throughout \mathcal{V} in the sense that

$$(1.10) \quad \forall \alpha, \beta ([\alpha, \beta] = \alpha \cap \beta)$$

holds throughout \mathcal{V} . It can be shown that for varieties with a Maltsev term the commutator identity (1.10) is equivalent to congruence distributivity, and from congruence distributivity one can derive that there are no skew congruences of any type in finite subdirect products of algebras in \mathcal{V} . Thus, although it seems that we examined only a very special instance of a diagonal skew congruence, which then resulted in the definition of the commutator, this commutator controls all skew congruences in finite subdirect products in varieties with a Maltsev term.

In fact, the assumption that \mathbf{A} has a Maltsev term was used only to deduce that every diagonal subalgebra of \mathbf{A}^2 has a congruence as its universe. This fact turns out to be inessential; the theory can be developed under the weaker hypothesis that \mathbf{A} lies in a congruence modular variety. The theory of this commutator was developed by Jonathan D. H. Smith (when \mathcal{V} has a Maltsev term) in [75], Joachim Hagemann and Christian Herrmann in [28], Heinz-Peter Gumm in [27], and Freese and McKenzie in [19]. The strength of the theory lies in the representation theorem for abelian algebras, proved by Herrmann in [31], which states that an abelian algebra in a congruence modular variety is **affine**. This theorem associates to an abelian algebra or congruence interval a module in the modern sense, i.e., a module over a ring. The fact that the commutator links every skew congruence with a module is the key to dealing with these congruences.

To emphasize the link with what we have said earlier, the commutator in congruence modular varieties encodes the existence of skew congruences. When the commutator trivializes, skew congruences are omitted, and this restricts the shape of the congruence lattices in some way. In the case of the modular commutator, this restriction on shapes is expressible by a congruence identity stronger than modularity, namely the distributive law.

When \mathcal{V} is not congruence modular the types of skew congruences multiply. Different commutator theories have been invented which deal with this. The **TC-commutator** is a commutator based on the **term condition** (Section 2.5). It was invented by Ralph McKenzie to generalize the modular commutator. As such it encodes the existence of

provide a reference
for this, if possible.

diagonal skew congruences. A useful and completely general representation theorem for abelian algebras and congruences is not likely to exist, but Keith Kearnes and Ágnes Szendrei extended Herrmann's representation theorem for congruence modular varieties to any variety satisfying some nontrivial idempotent Maltsev condition in [52].

McKenzie next introduced the **strong term condition**, [66, 34], a concept similar in spirit to the ordinary term condition. He proved its usefulness in many ways, but he did not develop a commutator theory for it. Later McKenzie defined the concept of **rectangulation**, [54], and established its usefulness. His definition was slightly short of a term condition for rectangulation, and he did not develop a commutator theory for it.

In the early 1990's we worked on the problem of developing a commutator theory for McKenzie's strong term condition. We found that certain "cross-diagonal" skew congruences occur naturally in algebras. These are congruences on symmetric diagonal subalgebras $\mathbf{B} \leq \mathbf{A}^2$ generated by pairs of the form $\langle (u, v), (v, u) \rangle$. The behavior of this type of skew congruence is encoded in a commutator we call the **strong rectangular commutator**. It is so named because it has an associated term condition related to McKenzie's description for rectangulation. We developed the theory of this commutator in [49] and proved a representation for its abelian algebras and congruences. We found that McKenzie's strong term condition was exactly the conjunction of his original term condition (TC) and our term condition for strong rectangulation, therefore a commutator theory for the strong term condition follows from the theories of the TC-commutator and the strong rectangular commutator without further effort.

In this monograph we introduce a term condition for rectangulation, develop the theory of the rectangular commutator, and prove a representation theorem for its abelian algebras and congruences (Chapter 5).

There are actually many more commutators than the four described above, but the others are approximations to these. Each of these commutators has an ideal model of an abelian algebra. For the TC commutator, the ideal model is an abelian group expanded by unary endomorphisms, i.e., a module over a ring. For the strong commutator it is a set expanded by unary endomorphisms, i.e., a unary algebra. For the rectangular commutator it is a semilattice expanded by unary endomorphisms, which is a natural type of semimodule. For the strong rectangular commutator it is a reduct to an antichain of a semilattice

expanded by unary endomorphisms. The strangeness of this fourth notion of ‘abelianness’ is a reminder that these commutator theories do not begin by postulating the structure of an ideal abelian algebra, but rather by identifying a natural type of skew congruence.

1.4. The Results of This Monograph

The theorem which prompted us to write this monograph was our 1999 discovery that a variety satisfies a nontrivial congruence identity if and only if it satisfies an idempotent Maltsev condition that fails in the variety of semilattices (Theorem 7.15). David Hobby and Ralph McKenzie proved in Chapter 9 of [34] that a certain Maltsev condition, that we will call in this introduction “HM”, defines the class of all varieties satisfying an idempotent Maltsev condition that fails in the variety of semilattices. HM is a disjunction $\bigvee \text{HM}_n$ of strong Maltsev conditions. The proof of the 1999 theorem is based on a careful analysis of various Maltsev conditions equivalent to HM. A novel feature of the proof is that it deals only with the local effects of HM on nonsolvable intervals in congruence lattices of algebras in a variety \mathcal{V} satisfying some HM_n . Namely, we prove that every nonmodular interval in a congruence lattice contains what we call a “solvability obstruction”. We use HM_n to introduce a rank function on solvability obstructions that appear in congruence lattices of algebras in \mathcal{V} . Then we prove that

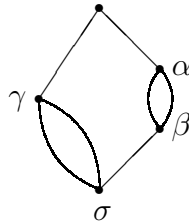
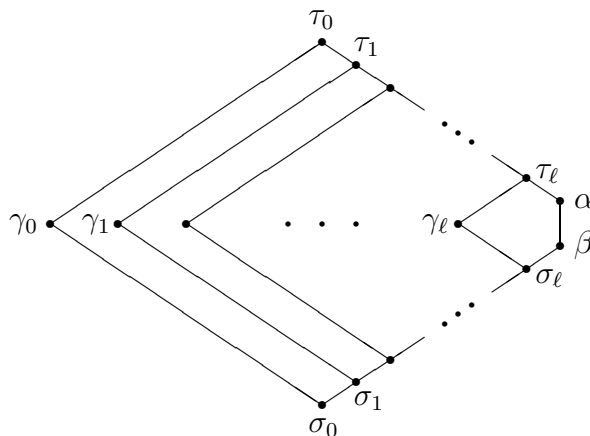


FIGURE 1.5.

for every solvability obstruction that appears in the interval $I[\beta, \alpha]$ of some copy of \mathbf{N}_5 in some congruence lattice there is associated a solvability obstruction of strictly smaller rank in the interval $I[\sigma, \gamma]$ (cf. Figure 1.5). By iterating this observation, one obtains that deeply nested copies of \mathbf{N}_5 , as depicted in Figure 1.6, can appear as a sublattice of a congruence lattice of an algebra in \mathcal{V} only if solvability obstructions of large rank exist. We complete the proof by proving that the complexity of HM_n induces a uniform (finite) bound on the rank of solvability obstructions throughout \mathcal{V} , hence for some ℓ it is not possible to embed $\mathbf{N}_{\ell+5}$ into the congruence lattice of any member of

FIGURE 1.6. The lattice $\mathbf{N}_{\ell+5}$

\mathcal{V} . This fact is converted into a nontrivial congruence identity satisfied by \mathcal{V} .⁹

This theorem gives a simple answer to the question “Which varieties satisfy nontrivial congruence identities?” But more importantly, the effort to prove the theorem answers the questions “Which varieties have trivial rectangular commutator?” and “What kinds of restrictions on congruence lattice shapes are linked to having trivial rectangular commutator?”, since the combination of Theorems 5.25 and 7.15 prove that a variety \mathcal{V} satisfies a nontrivial congruence identity if and only if the rectangular commutator is trivial throughout \mathcal{V} .

This monograph contains many characterizations of the class of varieties whose rectangular commutator is trivial: by the shapes of congruence lattices in these varieties and by the Maltsev conditions they satisfy (see, for example, Sections 5.3, 7.2 and 8.2). But it also contains characterizations of other classes of varieties with parallel descriptions. Namely, we consider the classes of varieties (i) whose strong commutator is trivial (Sections 3.1, 4.2 and 4.3), (ii) whose strong rectangular commutator is trivial (Sections 3.2, 4.2 and 4.3), (iii) whose TC commutator is trivial (Section 8.1), and (iv) whose rectangular and TC commutators are both trivial (Section 8.3). These characterizations of classes of varieties are also in terms of the shapes of congruence lattices

⁹It is interesting that Dedekind’s investigation into the arithmetic of congruences of $\langle \mathbb{C}; +, -, 0 \rangle$ led to the law postulating the omission of sublattices isomorphic to \mathbf{N}_5 , while the natural extension of that investigation to arbitrary algebras leads to “generalized modular laws” postulating the omission of sublattices isomorphic to $\mathbf{N}_{\ell+5}$ for some ℓ .

in these varieties and the Maltsev conditions they satisfy. One byproduct of this investigation is our first real understanding of congruence join semidistributivity: a variety \mathcal{V} is congruence join semidistributive if and only if it is congruence meet semidistributive and satisfies a nontrivial congruence identity (equivalently, if all commutators trivialize throughout \mathcal{V}).

We end the monograph with a chapter on residually small varieties satisfying congruence identities (Chapter 9). The main result of this chapter is that a residually small variety satisfies a congruence identity if and only if it is congruence modular, hence McKenzie's Conjecture is true for residually small varieties. We apply the main result of the chapter to show that there is no "almost congruence distributive" variety that satisfies a nontrivial idempotent Maltsev condition.

Ten open problems are posed in this monograph (six in Chapter 4, one in Chapter 6, and three in Chapter 9).

CHAPTER 2

Preliminary Notions

Here we fix notation and introduce definitions and results with the purpose of providing a bridge from standard material to the material in this monograph.

In this chapter and the sequel we use the following conventions concerning notation. The set of natural numbers is ω . The first natural number is 0. Sequences or tuples of elements from a set X are written in boldface, as in $\mathbf{x} \in X$ (or $\mathbf{x} \in X^n$ if the length n is to be specified). The i -th coordinate of the tuple \mathbf{x} is written in italic as x_i . If two sequences $\mathbf{x}, \mathbf{y} \in X$ have the same length and R is a binary relation on X , then we may write $\mathbf{x} \equiv \mathbf{y} \pmod{R}$ or $\mathbf{x} \equiv_R \mathbf{y}$ or $\mathbf{x} R \mathbf{y}$ to mean that $(x_i, y_i) \in R$ for all i . If R is a binary relation on X , then the set $X^n[R]$ consists of the tuples $\mathbf{x} \in X^n$ such that $x_i \equiv_R x_j$ for all $1 \leq i \leq j \leq n$. The projection of a subset $S \subseteq \prod_{i \in I} X_i$ onto a subset of coordinates $J \subseteq I$ is denoted $\pi_J: S \rightarrow \prod_{j \in J} X_j$, and its kernel is denoted η_J . If θ is an equivalence relation on $\prod_{j \in J} X_j$, then θ_J denotes $\pi_J^{-1}(\theta)$ (so $0_J = \eta_J$). When $J = \{j\}$, then we write π_j, η_j and θ_j instead of $\pi_{\{j\}}, \eta_{\{j\}}$ and $\theta_{\{j\}}$. Expressions $A := B$ or $B =: A$ mean “ A is defined by B ”.

non-empty set

2.1. Algebras, Varieties, and Clones

An algebra is a model of a 1-sorted first-order algebraic language. To fix conventions, an **algebraic signature** is a pair $\sigma := (F, \alpha)$ where F is a set (of operation symbols), and $\alpha: F \rightarrow \omega$ is a function (assigning arity). An **algebra** of signature σ is a pair $\mathbf{A} := \langle A; F \rangle$ where A is a set, called the **universe** of \mathbf{A} , and for each $f \in F$ with $\alpha(f) = k$ there is an assigned k -ary operation $f^{\mathbf{A}}: A^k \rightarrow A$. The operations $f^{\mathbf{A}}$ are called the **basic operations** of \mathbf{A} .

If Y is a set that is disjoint from F , then $T(Y)$ is the smallest set of words in the alphabet $F \cup Y$ satisfying

- (i) $Y \subseteq T(Y)$, and
- (ii) if $f \in F$, $\alpha(f) = k$ and $g_1, \dots, g_k \in T(Y)$, then $f g_1 \cdots g_k \in T(Y)$.

When $Y = X := \{x_i \mid 1 \leq i < \omega\}$, then $T(X)$ is the set of **terms of signature** σ , and its members are called **terms**. When $Y = X_n := \{x_1, \dots, x_n\}$, the members of $T(X_n)$ are **n -ary terms**.

$T(Y)$ has a canonical structure of an algebra of signature σ . If $f \in F$ has arity k , then in the term algebra $\mathbf{T} = \mathbf{T}(Y)$ the operation $f^{\mathbf{T}}$ is defined so that if $g_1, \dots, g_k \in T$, then $f^{\mathbf{T}}(g_1, \dots, g_k) = fg_1 \cdots g_k \in T$. If \mathbf{A} is an algebra of signature σ and $Y = X$ or X_n , then an **assignment in \mathbf{A}** of the variables Y is a function $v: Y \rightarrow A$. It follows from the unique readability of terms that any assignment v extends uniquely to an algebra homomorphism $\widehat{v}: \mathbf{T} \rightarrow \mathbf{A}$; i.e., \mathbf{T} is free over Y in the class of all algebras of signature σ . If $t \in T$, then both $\widehat{v}(t)$ and $t(v)$ denote the image of t under \widehat{v} .

If \mathbf{A} is an algebra of signature σ and $t \in T(X_n)$, then t determines an n -ary **term operation** on A , $t^{\mathbf{A}}: A^n \rightarrow A$, defined by $(a_1, \dots, a_n) \mapsto t(v)$ where $v: X_n \rightarrow A$ is the assignment $x_i \mapsto a_i$. These remarks apply in the case where $\mathbf{A} = \mathbf{T} = \mathbf{T}(X_n)$, and show that, if $t \in T(X_k)$ and $g_1, \dots, g_k \in T(X_n)$, then there is a term $t(g_1, \dots, g_k) := t^{\mathbf{T}}(g_1, \dots, g_k) \in T(X_n)$. It can be proved by induction that, in this notation, the term $t(x_1, \dots, x_k)$ is simply t itself.

The **language** L associated with the signature σ is the set of first-order formulas in this signature. Parentheses may be added to formulas to improve readability. An **identity** in the language L is an atomic L -formula, which is a formula of the form $p \approx q$ where $p, q \in T(X_n)$ for some n . \mathbf{A} **satisfies** $p \approx q$, written $\mathbf{A} \models p \approx q$, if $p(v) = q(v)$ for every assignment in \mathbf{A} (equivalently if $p^{\mathbf{A}} = q^{\mathbf{A}}$). A class of algebras in the language L is a **variety** if it is definable as the class of all L -algebras satisfying some set of identities. The smallest variety containing an algebra \mathbf{A} is denoted $\mathcal{V}(\mathbf{A})$.

A **quasi-identity** is an open first-order formula of the form

$$\bigwedge_{i=1}^n (p_i \approx q_i) \rightarrow (p \approx q)$$

where $p \approx q$ and each $p_i \approx q_i$ is an identity.

A **clone** is a multisorted structure $\mathcal{C} = \langle C_0, C_1, \dots; F \rangle$, with sorts indexed by ω , where

- (i) Each of the sorts C_i , $i \in \omega$, is a set.
- (ii) $F = \{\pi_i^n \mid n \in \omega, 1 \leq i \leq n\} \cup \{\text{comp}_n^m \mid m, n \in \omega\}$ is a set of operations between the C_i .
- (iii) Each π_i^n is a 0-ary operation (a constant) in C_n .
- (iv) $\text{comp}_n^m: C_m \times (C_n)^m \rightarrow C_n$ is an $(m+1)$ -ary operation.
- (v) $\text{comp}_n^n(f, \pi_1^n, \dots, \pi_n^n) = f$ and $\text{comp}_n^m(\pi_i^m, f_1, \dots, f_m) = f_i$.

$$(vi) \text{ comp}_m^p(f, \text{comp}_m^n(g_1, h_1, \dots, h_n), \dots, \text{comp}_m^n(g_p, h_1, \dots, h_n)) = \\ \text{comp}_m^n(\text{comp}_n^p(f, g_1, \dots, g_p), h_1, \dots, h_n).$$

For an example of a clone, take C_n to be the set of n -ary L -terms, take π_i^n to be x_i , and take comp_n^m to be the composition of an m -ary term with m n -ary terms. This is the **clone of L -terms**. The **clone of \mathbf{A}** when \mathbf{A} is an algebra is defined similarly using the term operations of \mathbf{A} instead of the L -terms.

It is evident that clones are multisorted algebras defined by identities, hence the class of clones is a multisorted variety. This implies that the usual algebraic notions apply to clones. In particular, a sequence

$$\mathbf{h} = (h_0, h_1, \dots): \mathcal{C} \rightarrow \mathcal{D}$$

is a **homomorphism** between clones if each $h_i: C_i \rightarrow D_i$ is a function, and the sequence \mathbf{h} preserves the clone operations of composition and projection. The notions of kernel and quotient are defined in the obvious way.

Let \mathcal{V} be a variety of 1-sorted algebras in the language L . The **clone of \mathcal{V}** , denoted $\text{Clo}(\mathcal{V})$, is that quotient of the clone of L -terms that is obtained by identifying terms $p, q \in C_n$ if $p \approx q$ is satisfied by all algebras in \mathcal{V} . The correspondence $\mathcal{V} \mapsto \text{Clo}(\mathcal{V})$ is essentially bijective. Namely, each variety \mathcal{V} is assigned the clone $\text{Clo}(\mathcal{V})$, and each clone \mathcal{C} is assigned a variety $\text{Var}(\mathcal{C})$ which we define now. The operation symbols of $\text{Var}(\mathcal{C})$ are $\bigcup_{i \in \omega} C_i$ where the arity of $f \in C_n$ is defined to be n . The identities defining $\text{Var}(\mathcal{C})$ are of two types:

(i) projection identities: $\pi_i^n(x_1, \dots, x_n) \approx x_i$ for each $1 \leq i \leq n$, and

(ii) composition identities:

$$f(x_1, \dots, x_n) \approx g(h_1(x_1, \dots, x_n), \dots, h_m(x_1, \dots, x_n))$$

whenever the equality $f = \text{comp}_n^m(g, h_1, \dots, h_m)$ holds in \mathcal{C} .

It can be shown that $\mathcal{C} = \text{Clo}(\text{Var}(\mathcal{C}))$ holds for any clone \mathcal{C} , while \mathcal{V} is definitionally equivalent to the variety $\text{Var}(\text{Clo}(\mathcal{V}))$ for any variety \mathcal{V} .

We write $\mathcal{U} \leq \mathcal{V}$ if there is a homomorphism $\mathbf{h}: \text{Clo}(\mathcal{U}) \rightarrow \text{Clo}(\mathcal{V})$. This notation is used to express the fact that the algebras in \mathcal{V} have an underlying \mathcal{U} -structure.

2.2. Lattice Theory

A **partial lattice** is a structure $\mathbf{P} = \langle P; \vee, \wedge, \leq \rangle$ where $\langle P; \leq \rangle$ is a partially ordered set and \vee and \wedge are partial binary operations on P (called **join** and **meet** respectively) such that, if $a \vee b$ is defined, then

it is the least upper bound of a and b in $\langle P; \leq \rangle$, and if $a \wedge b$ is defined, then it is the greatest lower bound of a and b in $\langle P; \leq \rangle$. A **lattice** is a total algebra $\mathbf{L} = \langle L; \vee, \wedge \rangle$ on which there is an order making $\langle L; \vee, \wedge, \leq \rangle$ a partial lattice.

If $x \leq y$ in some lattice \mathbf{L} , then the **interval** between them is $I[x, y] := \{z \mid x \leq z \leq y\}$. A **lattice ideal** of \mathbf{L} is a subset $I \subseteq L$ closed under join ($x, y \in I \implies x \vee y \in I$) and close downward ($x \in I$ & $z \leq x \implies z \in I$). The dual concept is a **lattice filter**. The **principal ideal** determined by $x \in L$ is $(x] := \{z \mid z \leq x\}$. The collection $\mathcal{I}(\mathbf{L})$ of all ideals of \mathbf{L} , ordered by inclusion, is the **ideal lattice** of \mathbf{L} . If $I, J \in \mathcal{I}(\mathbf{L})$, then $I \wedge J = I \cap J$ and

$$I \vee J = \{z \mid \exists x \in I, y \in J (z \leq x \vee y)\}.$$

Lattice terms may be called **words**. If p and q are n -ary lattice words, then the inclusion $p \leq q$ is **satisfied** by a lattice \mathbf{L} if and only if $p^{\mathbf{L}} \leq q^{\mathbf{L}}$ in the pointwise order. Therefore \mathbf{L} satisfies $p \approx q$ if and only if it satisfies both $p \leq q$ and $q \leq p$. Note also that $p \leq q$ is satisfied by \mathbf{L} if and only if the identity $p \approx p \wedge q$ is satisfied by \mathbf{L} . It follows that a class of lattices is definable by identities if and only if it is definable by inclusions. The variety of all lattices is denoted \mathcal{L} .

If Q is the lattice quasi-identity $\bigwedge (p_i \approx q_i) \rightarrow (p \approx q)$, then the **Q -configuration** is a pair $(\mathbf{P}(Q), p \approx q)$ where $\mathbf{P}(Q)$ is the natural partial lattice of subterms of terms of Q . That is, $\mathbf{P}(Q)$ is the partial lattice presented by $\langle G \mid R \rangle$ where G is the set of subterms of terms appearing in Q and R consists of relations of the following types: if s, t and $s \vee t$ are subterms, then R contains a relation expressing that $s \vee t$ is equal to the join of s and t , R contains similar relations for meet, and for each premise $p_i \approx q_i$ of Q the set R contains the relation $p_i = q_i$. It is not assumed that R contains $p = q$ where $p \approx q$ is the conclusion of Q . This definition exists to make the following statement true: an assignment of the variables of Q in a lattice \mathbf{L} which satisfies the premises of Q determines and is determined by a homomorphism of partial lattices $\varphi: \mathbf{P}(Q) \rightarrow \mathbf{L}$, and the assignment will satisfy Q precisely when $\varphi(p) = \varphi(q)$.

DEFINITION 2.1. The **meet semidistributive law** is the quasi-identity

$$(2.1) \quad ((p \wedge q) \approx s) \ \& \ ((p \wedge r) \approx s) \rightarrow ((p \wedge (q \vee r)) \approx s),$$

and the **join semidistributive law** is the dual quasi-identity.

Quasi-identity (2.1) is equivalent to

$$(2.2) \quad ((p \wedge q) \approx s) \ \& \ ((p \wedge r) \approx s) \ \& \ ((q \vee r) \geq p) \rightarrow (p \approx s).$$

The reason this is so is that any assignment of variables which fails to satisfy (2.2) will also fail to satisfy (2.1), while conversely if $p \mapsto a, q \mapsto b, r \mapsto c, s \mapsto d$ is an assignment of variables in some lattice that fails to satisfy (2.1), then $p \mapsto a \wedge (b \vee c), q \mapsto b, r \mapsto c, s \mapsto d$ is an assignment that fails to satisfy (2.2).

The SD_{\wedge} -configuration is the Q -configuration where Q is quasi-identity (2.2). More explicitly, let $\mathbf{P}(SD_{\wedge})$ be the partial lattice generated by $\{p, q, r, m, j\}$ where $m = p \wedge q = p \wedge r$ and $j = q \vee r \geq p$. Then the SD_{\wedge} -configuration is $(\mathbf{P}(SD_{\wedge}), p \approx m)$. An SD_{\wedge} -**failure** in a lattice \mathbf{L} is an interval of the form $I = I[\varphi(m), \varphi(p)]$ where $\varphi: \mathbf{P}(SD_{\wedge}) \rightarrow \mathbf{L}$ is a homomorphism of partial lattices. This SD_{\wedge} -failure is **trivial** if I has one element and is **nontrivial** otherwise. **The SD_{\vee} -configuration** and SD_{\vee} -**failures** are defined dually. Thus a lattice is meet semidistributive if and only if it has no nontrivial SD_{\wedge} -failures.

For finite lattices, or for varieties of lattices, the meet and join semidistributive laws are characterized in the following theorem.

THEOREM 2.2.

- (1) (Cf. [7]) *A finite lattice is meet semidistributive if and only if it has no sublattice isomorphic to $\mathbf{M}_3, \mathbf{D}_1, \mathbf{E}_1, \mathbf{E}_2$ and \mathbf{G} . (These are five of the following six lattices.)*

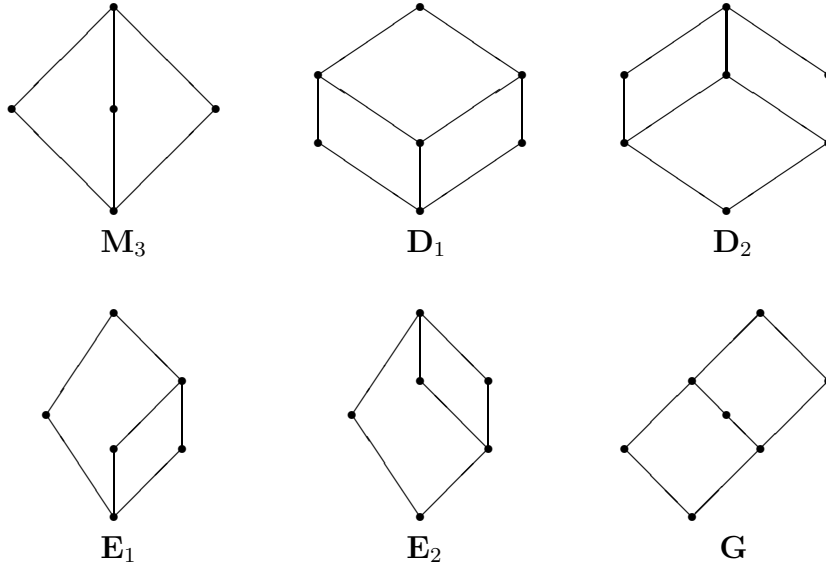


FIGURE 2.1.

- (2) (Cf. [41]) *A variety of lattices consists of meet semidistributive lattices if and only if it does not contain $\mathbf{M}_3, \mathbf{D}_1, \mathbf{E}_1, \mathbf{E}_2$ or \mathbf{G} .*

The dual results hold for join semidistributivity in place of meet semidistributivity.

DEFINITION 2.3. **Whitman's condition**, denoted by **(W)**, is the formula

$$(x \wedge y \leq u \vee v) \rightarrow ((x \leq u \vee v) \text{ or } (y \leq u \vee v) \text{ or } (x \wedge y \leq u) \text{ or } (x \wedge y \leq v)).$$

This condition may be written a little more compactly as:

$$x \wedge y \leq u \vee v \quad \text{implies} \quad \{x, y, u, v\} \cap I[x \wedge y, u \vee v] \neq \emptyset.$$

All lattices depicted in Theorem 2.2 satisfy **(W)**.

DEFINITION 2.4. A quasi-identity $\bigwedge_{i=1}^n (p_i \approx q_i) \rightarrow (p \approx q)$ in k variables **satisfies (W)** if the lattice with the presentation $\langle G \mid R \rangle$ satisfies **(W)**, where $G = \{g_1, \dots, g_k\}$, $\mathbf{g} = (g_1, \dots, g_k)$ and $R = \{p_1(\mathbf{g}) = q_1(\mathbf{g}), \dots, p_n(\mathbf{g}) = q_n(\mathbf{g})\}$.

Note that whether or not a quasi-identity satisfies **(W)** depends only on its premises. It is proved in [6] that a quasi-identity Q satisfies **(W)** iff the associated partial lattice $\mathbf{P}(Q)$ satisfies **(W)**. Since $\mathbf{P}(Q)$ is finite, this yields an algorithm for testing a quasi-identity for **(W)**. It is noted in [6] that the meet semidistributive law satisfies **(W)**, and that more generally any quasi-identity whose premises are free of joins (as in form (2.1) of the meet semidistributive law) satisfies **(W)**. These remarks also apply to the join semidistributive law.

DEFINITION 2.5. A surjective lattice homomorphism $h: \mathbf{K} \rightarrow \mathbf{L}$ is **upper bounded** if each kernel class has a largest element, is **lower bounded** if each kernel class $h^{-1}(a)$, $a \in L$, has a least element, and is **bounded** if it is both lower and upper bounded. A lattice \mathbf{L} is **(upper, lower) bounded** if there is a surjective (upper, lower) bounded homomorphism $h: \mathbf{F}_{\mathcal{L}}(x_1, \dots, x_n) \rightarrow \mathbf{L}$ from a finitely generated free lattice onto \mathbf{L} .

Any lower bounded lattice is join semidistributive (Theorem 2.20 of [16]). Chapter 2 of [16] describes an effective algorithm for testing if a finite lattice is lower or upper bounded. We shall not have cause to use this algorithm, but for later reference we point out that \mathbf{D}_1 is the only lower bounded lattice depicted in Theorem 2.2. The others are excluded because they are not join semidistributive, while the lower boundedness of \mathbf{D}_1 is proved one way in Example 2.74 of [16] and differently in Table 3 of [64].

DEFINITION 2.6. An algebra \mathbf{P} is **projective** relative to a variety \mathcal{V} if whenever $\sigma: \mathbf{A} \rightarrow \mathbf{B}$ is a surjective homomorphism between algebras

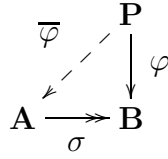


FIGURE 2.2.

of \mathcal{V} and $\varphi: \mathbf{P} \rightarrow \mathbf{B}$ is a homomorphism, then there is a homomorphism $\overline{\varphi}: \mathbf{P} \rightarrow \mathbf{A}$ such that $\sigma \circ \overline{\varphi} = \varphi$.

When \mathcal{V} is a variety of lattices, we may call a partial lattice \mathbf{P} **projective** relative to \mathcal{V} if it satisfies the homomorphism lifting property of this definition.

THEOREM 2.7. *If \mathcal{V} is a variety of algebras and \mathbf{P} is subdirectly irreducible and projective relative to \mathcal{V} , then there is an identity ε such that for all $\mathbf{A} \in \mathcal{V}$ it is the case that $\mathbf{A} \models \varepsilon$ if and only if \mathbf{A} has no subalgebra isomorphic to \mathbf{P} .*

PROOF. Let \mathcal{U} consist of all members of \mathcal{V} that have no subalgebra isomorphic to \mathbf{P} . \mathcal{U} is closed under the formation of homomorphic images because \mathbf{P} is projective relative to \mathcal{V} , \mathcal{U} is closed under subalgebras by definition, and \mathcal{U} is closed under products because \mathbf{P} is subdirectly irreducible. By Birkhoff's HSP Theorem, \mathcal{U} is a subvariety of \mathcal{V} . Since $\mathbf{P} \notin \mathcal{U}$, there is an identity ε that is satisfied in \mathcal{U} but not by \mathbf{P} . Necessarily ε fails in every member of $\mathcal{V} - \mathcal{U}$, so this identity has the property stated in the theorem. \square

We call the identity ε of Theorem 2.7 the **conjugate identity** for \mathbf{P} .

We shall have occasion to refer to the following closely-related result.

THEOREM 2.8. *If Q is a quasi-identity and $\mathbf{P}(Q)$ is projective relative to \mathcal{V} , then the class of lattices in \mathcal{V} that satisfy Q is a subvariety of \mathcal{V} .*

PROOF. As above, let \mathcal{U} consist of all members of \mathcal{V} that satisfy Q . \mathcal{U} is closed under the formation of homomorphic images because $\mathbf{P}(Q)$ is projective relative to \mathcal{V} , \mathcal{U} is closed under subalgebras and products because Q is a quasi-identity. \square

2.3. Meet Continuous Lattice Theory

A lattice is **meet continuous** if it is complete and the binary meet operation distributes over arbitrary up-directed suprema. That is, for every up-directed set \mathcal{D} ,

$$(2.3) \quad x \wedge \bigvee_{y \in \mathcal{D}} y = \bigvee_{y \in \mathcal{D}} x \wedge y.$$

Every algebraic lattice is meet continuous (Lemma VIII.5.2 of [1]).

The language of meet continuous lattices has operation symbols $\{\bigvee, \wedge\}$ where \bigvee is a class of κ -ary join operations for all $\kappa \geq 0$ and \wedge is binary meet. The 0-ary join is a constant that interprets as the least element and is denoted 0. The fact that \wedge is a meet operation is expressible by identities in the same way that it is expressed in lattice theory. As with binary join, it is easy to express with identities the fact that κ -ary \bigvee is κ -ary join with respect to the \wedge -order. To show that meet continuity is expressible by identities one must show that the distributive law in (2.3) can be rewritten with quantification over arbitrary sets rather than up-directed sets. This is done by replacing an arbitrary join with the up-directed supremum of its finite sub-joins:

$$x \wedge \bigvee_{y \in \mathcal{F}} y \approx \bigvee_{\substack{\mathcal{F}_0 \subseteq \mathcal{F} \\ \mathcal{F}_0 \text{ finite}}} \left(x \wedge \bigvee_{y \in \mathcal{F}_0} y \right).$$

The category \mathcal{L}_{MC} of meet continuous lattices with join complete homomorphisms has a forgetful functor $F: \mathcal{L}_{MC} \rightarrow \mathcal{L}$ to lattices that forgets all join operations except the binary join. Theorem 2.9, which is essentially due to Ralph Freese, asserts that this functor has a left adjoint given by the ideal lattice functor. Here if $h: \mathbf{K} \rightarrow \mathbf{L}$ is a homomorphism between lattices, then $\mathcal{I}(h): \mathcal{I}(\mathbf{K}) \rightarrow \mathcal{I}(\mathbf{L})$ is defined by

$$\mathcal{I}(h)(I) = \{z \in L \mid \exists x \in I (z \leq h(x))\}.$$

THEOREM 2.9. *Restriction to \mathbf{L} is a natural bijection*

$$|: \mathcal{L}_{MC}(\mathcal{I}(\mathbf{L}), \mathbf{K}) \rightarrow \mathcal{L}(\mathbf{L}, F(\mathbf{K})).$$

(We are identifying \mathbf{L} with the sublattice of $\mathcal{I}(\mathbf{L})$ consisting of principal ideals.)

PROOF. If $h: \mathcal{I}(\mathbf{L}) \rightarrow \mathbf{K}$ is a meet continuous lattice homomorphism, then the restriction $h|_{\mathbf{L}}: \mathbf{L} \rightarrow F(\mathbf{K})$ is evidently a lattice homomorphism (since lattices are reducts of meet continuous lattices).

remove "by"

has this notation been defined?

If $g: \mathbf{L} \rightarrow F(\mathbf{K})$ is a lattice homomorphism, then it is proved in Lemma 5.1 of [15] that the function

$$\widehat{g}: \mathcal{I}(\mathbf{L}) \rightarrow \mathbf{K}: I \mapsto \bigvee \{g(x) \mid x \in I\}$$

is a meet continuous lattice homomorphism whose restriction to \mathbf{L} is g . It is the unique extension of g to $\mathcal{I}(\mathbf{L})$, since $\mathcal{I}(\mathbf{L})$ is generated under \bigvee by L . This proves that restriction to \mathbf{L} is a bijection between hom-sets. The naturality is left as an exercise. (See Chapter IV of [63] for the method.) \square

remove "of"

The adjunction from \mathcal{L} to \mathcal{L}_{MC} may be composed with the adjunction from \mathcal{SET} to \mathcal{L} (given by the free lattice and forgetful functors) to produce a left adjoint $\mathcal{SET} \xrightarrow{\text{free}} \mathcal{L} \xrightarrow{\mathcal{I}} \mathcal{L}_{MC}$ to the composite forgetful functor $\mathcal{L}_{MC} \rightarrow \mathcal{L} \rightarrow \mathcal{SET}$. This adjoint establishes the existence and structure of free meet continuous lattices.

COROLLARY 2.10. *If $\mathbf{F} = \mathbf{F}_{\mathcal{L}}(X)$ is the free lattice generated by X , then $\mathcal{I}(\mathbf{F})$ is a free meet continuous lattice over the set $\{(x) \mid x \in X\}$.*

Since \mathcal{L}_{MC} is definable by identities and has free algebras of all ranks, Birkhoff's HSP theorem is valid for \mathcal{L}_{MC} : a subclass of \mathcal{L}_{MC} is definable by identities if and only if it is closed under the formation of homomorphic images, subalgebras and products. Hence the proof of Theorem 2.7 is valid for varieties of meet continuous lattices. Using known methods it can be shown that the projective subdirectly irreducible members of \mathcal{L}_{MC} are the finite, lower bounded, subdirectly irreducible lattices satisfying (W). This includes all lattices that are projective and subdirectly irreducible in \mathcal{L} , along with some others (such as \mathbf{D}_1).

2.4. Maltsev Conditions

A **strong Maltsev condition** is a primitive positive sentence in the language of clones. This means that it is a first-order sentence of the form " $\exists \bigwedge$ (atomic)" about clones. A **Maltsev condition** is a countably infinite disjunction $\bigvee_{i \in \omega} \sigma_i$ where the σ_i are strong Maltsev conditions that get weaker as i increases (i.e., $\sigma_i \vdash \sigma_{i+1}$ for all i). A variety **satisfies** a (strong) Maltsev condition if its clone does.

The concept of a Maltsev condition can be reformulated in terms of clone homomorphisms. To each strong Maltsev condition $\exists F \bigwedge \Sigma$ corresponds a finite presentation $\langle F \mid \Sigma \rangle$ of a clone. If \mathcal{U} is a variety whose clone has this presentation, then \mathcal{V} satisfies the strong Maltsev condition $\exists F \bigwedge \Sigma$ if and only if $\mathcal{U} \leq \mathcal{V}$. Similarly, to each Maltsev condition $\bigvee_{i \in \omega} \sigma_i$ corresponds a sequence $\cdots \leq \mathcal{U}_2 \leq \mathcal{U}_1 \leq \mathcal{U}_0$ of varieties

with finitely presented clones. A variety \mathcal{V} satisfies $\bigvee_{i \in \omega} \sigma_i$ if and only if $\mathcal{U}_i \leq \mathcal{V}$ for some i .

Our practice will be to express Maltsev conditions informally. For example, we will express the fact that variety \mathcal{V} satisfies the strong Maltsev condition

$$(2.4) \quad \exists p \left(\text{comp}_2^3(p, \pi_1^2, \pi_2^2, \pi_2^2) \approx \pi_1^2 \ \& \ \text{comp}_2^3(p, \pi_2^2, \pi_2^2, \pi_1^2) \approx \pi_1^2 \right),$$

which expresses the fact that \mathcal{V} has a Maltsev term, by saying that \mathcal{V} has a ternary term p such that the identities $p(x, y, y) \approx x$ and $p(y, y, x) \approx x$ hold in \mathcal{V} .

An n -ary element f of a clone \mathcal{C} is **idempotent** if

$$\text{comp}_1^n(f, \pi_1^1, \pi_1^1, \dots, \pi_1^1) = \pi_1^1.$$

If \mathcal{C} is the clone of a variety \mathcal{V} , this means that $f(x, x, \dots, x) \approx x$ is satisfied in \mathcal{V} . The idempotent elements of \mathcal{C} form a subclone, $\text{Id}(\mathcal{C})$. The **idempotent reduct** of \mathcal{V} is the variety $\text{Id}(\mathcal{V}) := \text{Var}(\text{Id}(\text{Clo}(\mathcal{V})))$ that is associated to the idempotent subclone of $\text{Clo}(\mathcal{V})$. A (strong) Maltsev condition is **idempotent** if for each term f in the condition the identity $f(x, x, \dots, x) \approx x$ is a consequence of the identities of the Maltsev condition.

Two (strong) Maltsev conditions are **equivalent** if they define the same class of varieties.

LEMMA 2.11. *Any idempotent strong Maltsev condition is equivalent to one of the form $\exists F \wedge \Sigma$ where $F = \{h, k\}$, h is n -ary and k is n^2 -ary, and Σ consists of the identities*

- (i) $h(x, x, \dots, x) \approx x$,
- (ii) $k(x_{11}, \dots, x_{nn}) \approx h(h(x_{11}, \dots, x_{1n}), \dots, h(x_{n1}, \dots, x_{nn}))$, plus
- (iii) *finitely many identities of the form $k(\text{variables}) \approx k(\text{variables})$.*

The proof of this lemma is part of the proof of Lemma 9.4 of [34].

A strong Maltsev condition $\exists F \wedge \Sigma$ is **linear** if each identity in Σ has the form $p(\text{variables}) \approx q(\text{variables})$ where $p, q \in F \cup \{\text{variables}\}$, while a Maltsev condition $\bigvee_{i \in \omega} \sigma_i$ is linear if each σ_i is.

LEMMA 2.12. *Any idempotent linear strong Maltsev condition is equivalent to one of the form $\exists F \wedge \Sigma$ where $F = \{f\}$, and Σ consists of the identities*

- (i) $f(x, x, \dots, x) \approx x$, plus
- (ii) *finitely many identities of the form $f(\text{variables}) \approx f(\text{variables})$.*

This is proved by slightly modifying the proof in [34] of the preceding lemma.

All of the Maltsev conditions considered in this monograph are idempotent and linear, and will usually be expressible in the form described in Lemma 2.12 using only two variables. We introduce a notation for such 2-variable Maltsev conditions now.

Let f be an n -ary operation symbol and let $N = \{1, \dots, n\}$. For each $U \subseteq N$, let $f_U(x, y)$ denote the term obtained from $f(x_1, \dots, x_n)$ by substituting x for x_i if $i \in U$ and y for x_j if $j \notin U$. Observe that any identity of the form

$$f(\text{variables}) \approx f(\text{variables})$$

that uses only the variable x and y may be rewritten in the form $f_U(x, y) \approx f_V(x, y)$ where U is the set of positions where x appears on the left of the identity and V is the set of positions where x appears on the right. This particular identity may be abbreviated by $U \equiv V$. In order to encode a family \mathcal{F} of 2-variable identities of the form $f(\text{variables}) \approx f(\text{variables})$, we define $\mathcal{B}(f)$ to be the Boolean algebra of subsets of N equipped with an equivalence relation E where $U \equiv_E V$ holds if and only if $f_U(x, y) \approx f_V(x, y)$ is a consequence of the identities in \mathcal{F} .

EXAMPLE 2.13. The statement that p is a Maltsev term for \mathcal{V} may be expressed by saying that \mathcal{V} satisfies identities of the form $p(\text{variables}) \approx p(\text{variables})$ if the idempotence of p is assumed, namely by $p(x, y, y) \approx p(x, x, x)$ and $p(y, y, x) \approx p(x, x, x)$. In this case, $N = \{1, 2, 3\}$ and $\mathcal{B}(p)$ is the Boolean algebra of Figure 2.3 equipped with

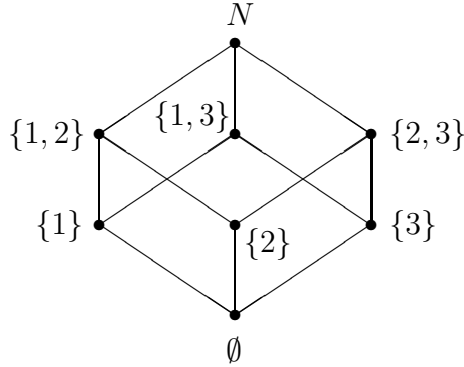


FIGURE 2.3. $\mathcal{B}(p)$

an equivalence relation E where $N \equiv_E \{1\} \equiv_E \{3\}$, since $p(x, x, x) \approx p(x, y, y) \approx p(y, y, x)$ are consequences of the starting identities, and $\emptyset \equiv_E \{1, 2\} \equiv_E \{2, 3\}$, since $p(y, y, y) \approx p(x, x, y) \approx p(y, x, x)$ are also consequences.

A subset of $\mathcal{B}(f)$ is **closed** if it is a union of E -classes. The next theorem is due to Walter Taylor.

THEOREM 2.14. *The following are equivalent for a variety \mathcal{V} .*

- (1) \mathcal{V} satisfies a nontrivial idempotent Maltsev condition (i.e., one that fails in some variety).
- (2) \mathcal{V} satisfies an idempotent Maltsev condition that fails in the variety of sets.
- (3) For some n , \mathcal{V} has an idempotent n -ary term f such that $\mathcal{B}(f)$ has no closed ultrafilter. 2.14

PROOF. See Corollary 5.3 of [77]. □

We will make frequent use of Theorem 2.15, so we introduce the following terminology.

DEFINITION 2.15. A term f is a **Taylor term** for a variety \mathcal{V} if \mathcal{V} satisfies $f(x, x, \dots, x) \approx x$ and enough other identities of the form $f_U(x, y) \approx f_V(x, y)$ so that $\mathcal{B}(f)$ has no closed ultrafilter.

Note that the word ‘linear’ does not appear in Conditions (1) and (2) of Theorem 2.14 although the Maltsev condition from (3) is linear. The fact that any variety satisfying a nontrivial idempotent Maltsev condition also satisfies one that is linear is the nontrivial part of the theorem.

Condition (3) of Theorem 2.14 means that if \mathcal{U} is the ultrafilter of subsets of N containing the singleton $\{i\}$, then \mathcal{U} is not closed, i.e., there is some $U \in \mathcal{U}$ and some $V \notin \mathcal{U}$ such that \mathcal{V} satisfies $f_U(x, y) \approx f_V(x, y)$. Another way to express Condition (3) is that, for each $1 \leq i \leq n$, \mathcal{V} satisfies an identity of the type $f(\text{variables}) \approx f(\text{variables})$ where all variables are x or y , and x appears in the i -th position on the left while y appears in the i -th position on the right. An identity of this type will be called an **i -th Taylor identity** for f .

A result similar to Taylor’s was discovered by David Hobby and Ralph McKenzie.

THEOREM 2.16. *The following are equivalent for a variety \mathcal{V} .*

- (1) \mathcal{V} satisfies an idempotent Maltsev condition that fails in the variety of semilattices.
- (2) For some n , \mathcal{V} has an idempotent n -ary term f such that $\mathcal{B}(f)$ has no closed, proper, nonempty lattice filter.

PROOF. See Lemma 9.5 of [34]. □

DEFINITION 2.17. A term f is a **Hobby-McKenzie term** for a variety \mathcal{V} if \mathcal{V} satisfies $f(x, x, \dots, x) \approx x$ and enough other identities

since $\mathcal{B}(f)$ is a finite boolean algebra, a lattice filter is the same thing as a principal filter.

of the form $f_U(x, y) \approx f_V(x, y)$ so that $\mathcal{B}(f)$ has no closed, proper, nonempty lattice filter.

The term p from Example 2.13 is a Hobby–McKenzie term, and therefore also a Taylor term. To see this, suppose that \mathcal{F} is a closed, proper, nonempty lattice filter of $\mathcal{B}(p)$. Then since \mathcal{F} is a nonempty filter it contains the top element N . Since \mathcal{F} is closed it must contain $\{1\}$ ($\equiv_E N$) and also $\{3\}$ ($\equiv_E N$). Since \mathcal{F} is a lattice filter, it contains $\{1\} \cap \{3\} = \emptyset$. But any lattice filter containing \emptyset is improper.

2.5. The Term Condition

Let $\mathbf{A} = \langle A; F \rangle$ be an algebra. An n -ary relation $R \subseteq A^n$ is **compatible** if it is a subalgebra of \mathbf{A}^n . If \mathbf{B} is a subalgebra of \mathbf{A} , then the **restriction** of a relation $R \subseteq A^n$ to \mathbf{B} is $R|_{\mathbf{B}} := R \cap B^n$. If δ is a congruence on \mathbf{A} , then $R/\delta := \{(a_1/\delta, \dots, a_n/\delta) \mid (a_1, \dots, a_n) \in R\}$. Both $R|_{\mathbf{B}}$ and R/δ are compatible if R is. If δ is a congruence, then R is **δ -closed** if $R = \delta \circ R \circ \delta$. (I.e., if $\mathbf{a} \delta \mathbf{b}$ R $\mathbf{c} \delta \mathbf{d}$ implies $\mathbf{a} R \mathbf{d}$.)

A compatible, reflexive, symmetric binary relation is called a **tolerance**. We will usually denote tolerances by upper case italic letters: R, S, T, \dots . A compatible equivalence relation (i.e., a transitive tolerance) is a **congruence**, and congruences will usually be denoted by lower case Greek letters: $\alpha, \beta, \gamma, \dots$. The tolerance or congruence generated by set $X \subseteq A \times A$ is usually denoted by $\text{Tg}^{\mathbf{A}}(X)$ or $\text{Cg}^{\mathbf{A}}(X)$ respectively, although if X contains only a few pairs then we may write, for example, $\text{Cg}^{\mathbf{A}}(a, b)$ or $\text{Cg}^{\mathbf{A}}((a, b), (c, d))$ instead. If T is a tolerance on \mathbf{A} , then a maximal subset $B \subseteq A$ such that $B \times B \subseteq T$ is called a **block** of T . If T is in fact a congruence, then a block is the same thing as a congruence class. A tolerance or congruence is **trivial** if it is the equality relation and **nontrivial** otherwise.

The collection of congruences on \mathbf{A} , ordered by inclusion, is an algebraic lattice which is denoted $\mathbf{Con}(\mathbf{A})$. Its least and largest elements are denoted 0 and 1 . Meet and join are denoted \wedge and \vee and are computed by $\alpha \wedge \beta = \alpha \cap \beta$ and $\alpha \vee \beta = \text{tr.cl.}(\alpha \cup \beta)$ where tr.cl. represents transitive closure.

An m -ary **polynomial operation** of \mathbf{A} is an operation $f: A^m \rightarrow A$ such that $f(x_1, \dots, x_m) = t^{\mathbf{A}}(x_1, \dots, x_m, \mathbf{a})$ for some $(m+n)$ -ary term t and some tuple $\mathbf{a} \in A^n$.

If S and T are tolerances on \mathbf{A} , then an S, T -**matrix** is a 2×2 matrix of elements of \mathbf{A} of the form

$$\begin{bmatrix} p & q \\ r & s \end{bmatrix} = \begin{bmatrix} f(\mathbf{a}, \mathbf{u}) & f(\mathbf{a}, \mathbf{v}) \\ f(\mathbf{b}, \mathbf{u}) & f(\mathbf{b}, \mathbf{v}) \end{bmatrix}$$

where $f(\mathbf{x}, \mathbf{y})$ is an $(m + n)$ -ary polynomial of \mathbf{A} , $\mathbf{a} S \mathbf{b}$, and $\mathbf{u} T \mathbf{v}$. The set of all S, T -matrices is denoted $M(S, T)$.

Since tolerances are compatible with all polynomial operations, any two elements in the same row of an S, T -matrix are T -related and any two elements in the same column are S -related.

The fact that S and T are symmetric relations implies that $M(S, T)$ is closed under interchanging rows or columns:

$$\begin{bmatrix} p & q \\ r & s \end{bmatrix} \in M(S, T) \Leftrightarrow \begin{bmatrix} r & s \\ p & q \end{bmatrix} \in M(S, T) \Leftrightarrow \begin{bmatrix} s & r \\ q & p \end{bmatrix} \in M(S, T).$$

If $S = T$, then $M(S, T) = M(T, T)$ is also closed under transpose, as one sees by interchanging the roles of \mathbf{x} and \mathbf{y} in the polynomial $f(\mathbf{x}, \mathbf{y})$ that defines a given matrix.

DEFINITION 2.18. Let S and T be tolerances on an algebra \mathbf{A} , and let δ be a congruence on \mathbf{A} . If $p \equiv_{\delta} q$ implies that $r \equiv_{\delta} s$ whenever

$$(2.5) \quad \begin{bmatrix} p & q \\ r & s \end{bmatrix} \in M(S, T),$$

then we say that $\mathbf{C}(S, T; \delta)$ **holds**, or S **centralizes** T **modulo** δ .

By interchanging the rows of matrices one sees that $\mathbf{C}(S, T; \delta)$ holds if and only if

$$p \equiv_{\delta} q \iff r \equiv_{\delta} s$$

for every S, T -matrix in (2.5).

The S, T -**term condition** is the condition $\mathbf{C}(S, T; 0)$. There are other similar conditions called term conditions that we will meet later, but this is the original one.

When establishing that the implication defining $\mathbf{C}(S, T; \delta)$ holds, or when making use of the fact, we may use underlining to highlight places in equations or expressions where changes are to be made. For example, we may write the implication defining $\mathbf{C}(S, T; \delta)$ in the following form: If

$$f(\underline{\mathbf{a}}, \mathbf{u}) \equiv_{\delta} f(\underline{\mathbf{a}}, \mathbf{v}),$$

then

$$f(\underline{\mathbf{b}}, \mathbf{u}) \equiv_{\delta} f(\underline{\mathbf{b}}, \mathbf{v}).$$

The relation $\mathbf{C}(\ , \ ; \)$ is called the **centralizer relation**. The reason that this terminology is used is that when \mathbf{A} is a group and S, T and δ are congruences on \mathbf{A} , then $\mathbf{C}(S, T; \delta)$ holds if and only if $[S, T] \leq \delta$ (see Chapter 1 of [19]).

The basic properties of the centralizer relation are enumerated in the following theorem.

THEOREM 2.19. *Let \mathbf{A} be an algebra with tolerances S, S', T, T' and congruences $\alpha, \alpha_i, \beta, \delta, \delta', \delta_j$. The following are true.*

- (1) (Monotonicity in the first two variables) *If $\mathbf{C}(S, T; \delta)$ holds and $S' \subseteq S, T' \subseteq T$, then $\mathbf{C}(S', T'; \delta)$ holds.*
- (2) *$\mathbf{C}(S, T; \delta)$ holds if and only if $\mathbf{C}(\text{Cg}^{\mathbf{A}}(S), T; \delta)$ holds.*
- (3) *$\mathbf{C}(S, T; \delta)$ holds if and only if $\mathbf{C}(S, \delta \circ T \circ \delta; \delta)$ holds.*
- (4) *If $T \cap \delta = T \cap \delta'$, then $\mathbf{C}(S, T; \delta)$ holds if and only if $\mathbf{C}(S, T; \delta')$ holds.*
- (5) (Semidistributivity in the first variable) *If $\mathbf{C}(\alpha_i, T; \delta)$ holds for all $i \in I$, then $\mathbf{C}(\bigvee_{i \in I} \alpha_i, T; \delta)$ holds.*
- (6) *If $\mathbf{C}(S, T; \delta_j)$ holds for all $j \in J$, then $\mathbf{C}(S, T; \bigwedge_{j \in J} \delta_j)$ holds.*
- (7) *If $T \cap (S \circ (T \cap \delta) \circ S) \subseteq \delta$, then $\mathbf{C}(S, T; \delta)$ holds.*
- (8) *If $\beta \wedge (\alpha \vee (\beta \wedge \delta)) \leq \delta$, then $\mathbf{C}(\alpha, \beta; \delta)$ holds.*
- (9) *Let \mathbf{B} be a subalgebra of \mathbf{A} . If $\mathbf{C}(S, T; \delta)$ holds in \mathbf{A} , then $\mathbf{C}(S|_{\mathbf{B}}, T|_{\mathbf{B}}; \delta|_{\mathbf{B}})$ holds in \mathbf{B} .*
- (10) *If $\delta' \leq \delta$, then the relation $\mathbf{C}(S, T; \delta)$ holds in \mathbf{A} if and only if $\mathbf{C}(S/\delta', T/\delta'; \delta/\delta')$ holds in \mathbf{A}/δ' .*

PROOF. Item (1) follows from the fact that $M(S', T') \subseteq M(S, T)$.

For (2), $\mathbf{C}(\text{Cg}^{\mathbf{A}}(S), T; \delta) \implies \mathbf{C}(S, T; \delta)$ follows from (1), since $S \subseteq \text{Cg}^{\mathbf{A}}(S)$. For the reverse implication (and also for the proof of item (5)), we will argue that if S_i is a tolerance, $\mathbf{C}(S_i, T; \delta)$ holds for all $i \in I$, and $\alpha := \text{tr.cl.}(\bigcup_{i \in I} S_i)$, then $\mathbf{C}(\alpha, T; \delta)$. (To complete the proof of (2) we need this only when $|I| = 1$, while in (5) we need it only when the S_i are congruences.)

Choose any matrix in $M(\alpha, T)$. If it is

$$\begin{bmatrix} p & q \\ r & s \end{bmatrix} = \begin{bmatrix} f(\mathbf{a}, \mathbf{u}) & f(\mathbf{a}, \mathbf{v}) \\ f(\mathbf{b}, \mathbf{u}) & f(\mathbf{b}, \mathbf{v}) \end{bmatrix},$$

then \mathbf{a} is related to \mathbf{b} by $\text{tr.cl.}(\bigcup_{i \in I} S_i)$, so there exist tuples $\mathbf{a} = \mathbf{a}_0 S_{i_1} \mathbf{a}_1 S_{i_2} \cdots S_{i_n} \mathbf{a}_n = \mathbf{b}$. These tuples determine matrices

$$\begin{bmatrix} p_k & q_k \\ p_{k+1} & q_{k+1} \end{bmatrix} := \begin{bmatrix} f(\mathbf{a}_k, \mathbf{u}) & f(\mathbf{a}_k, \mathbf{v}) \\ f(\mathbf{a}_{k+1}, \mathbf{u}) & f(\mathbf{a}_{k+1}, \mathbf{v}) \end{bmatrix} \in M(S_{i_{k+1}}, T).$$

We must show that $p \equiv_{\delta} q$ implies $r \equiv_{\delta} s$, so assume that $p \equiv_{\delta} q$. This is the same as $p_0 \equiv_{\delta} q_0$, and so by induction (using that $\mathbf{C}(S_{i_k}, T; \delta)$ holds for all k) we get that $p_k \equiv_{\delta} q_k$ for all k . Therefore $r = p_n \equiv_{\delta} q_n = s$. This completes the proofs of (2) and (5).

For (3), the implication $\mathbf{C}(S, \delta \circ T \circ \delta; \delta) \implies \mathbf{C}(S, T; \delta)$ follows from (1), since $T \subseteq \delta \circ T \circ \delta$. For the reverse implication, assume that

the presentation of the proof of this theorem could be considerably shortened, since much of the proof is elementary and could be left to the reader to work out.

$\mathbf{C}(S, T; \delta)$ holds, that

$$\begin{bmatrix} p & q \\ r & s \end{bmatrix} = \begin{bmatrix} f(\mathbf{a}, \mathbf{u}) & f(\mathbf{a}, \mathbf{v}) \\ f(\mathbf{b}, \mathbf{u}) & f(\mathbf{b}, \mathbf{v}) \end{bmatrix} \in M(S, \delta \circ T \circ \delta),$$

and that $p \equiv_{\delta} q$. There exist tuples \mathbf{u}' and \mathbf{v}' such that $\mathbf{u} \delta \mathbf{u}' T \mathbf{v}' \delta \mathbf{v}$. The matrix

$$\begin{bmatrix} p' & q' \\ r' & s' \end{bmatrix} = \begin{bmatrix} f(\mathbf{a}, \mathbf{u}') & f(\mathbf{a}, \mathbf{v}') \\ f(\mathbf{b}, \mathbf{u}') & f(\mathbf{b}, \mathbf{v}') \end{bmatrix}$$

is an S, T -matrix. Moreover,

$$p' = f(\mathbf{a}, \mathbf{u}') \delta f(\mathbf{a}, \mathbf{u}) = p \delta q = f(\mathbf{a}, \mathbf{v}) \delta f(\mathbf{a}, \mathbf{v}') = q'.$$

Since $\mathbf{C}(S, T; \delta)$ holds, it follows that $r' \equiv_{\delta} s'$. Hence

$$r = f(\mathbf{b}, \mathbf{u}) \delta f(\mathbf{b}, \mathbf{u}') = r' \delta s' = f(\mathbf{b}, \mathbf{v}') \delta f(\mathbf{b}, \mathbf{v}) = s,$$

or $r \equiv_{\delta} s$. This establishes $\mathbf{C}(S, \delta \circ T \circ \delta; \delta)$.

For (4), recall that elements in the same row of an $M(S, T)$ are T -related. Therefore, if $\begin{bmatrix} p & q \\ r & s \end{bmatrix} \in M(S, T)$, then since $T \cap \delta = T \cap \delta'$ we get that

$$p \equiv_{\delta} q \iff p \equiv_{T \cap \delta} q \iff p \equiv_{T \cap \delta'} q \iff p \equiv_{\delta'} q,$$

and

$$r \equiv_{\delta} s \iff r \equiv_{T \cap \delta} s \iff r \equiv_{T \cap \delta'} s \iff r \equiv_{\delta'} s.$$

Therefore the implication $p \equiv_{\delta} q \implies r \equiv_{\delta} s$ is equivalent to $p \equiv_{\delta'} q \implies r \equiv_{\delta'} s$.

For (6), assume that $\begin{bmatrix} p & q \\ r & s \end{bmatrix} \in M(S, T)$. If $p \equiv q \pmod{\bigwedge \delta_j}$, then $p \equiv q \pmod{\delta_j}$ for all j . Since $\mathbf{C}(S, T; \delta_j)$ holds for all j we get that $r \equiv s \pmod{\delta_j}$ for all j , or equivalently that $r \equiv s \pmod{\bigwedge \delta_j}$. This shows that $\mathbf{C}(S, T; \bigwedge_{j \in J} \delta_j)$ holds.

For (7), choose an S, T -matrix $M = \begin{bmatrix} p & q \\ r & s \end{bmatrix}$. Assume that $p \equiv_{\delta} q$. Since the elements in the same row of M are T -related and the elements in the same column are S -related, we have $r S p T \cap \delta q S s$. Moreover, $r T s$ since these elements belong to the same row. Together this yields that $r T \cap (S \circ (T \cap \delta) \circ S) s$. By the assumption in (7), this implies that $r \equiv_{\delta} s$. This proves (7).

For item (8), if $\beta \wedge (\alpha \vee (\beta \wedge \delta)) \leq \delta$, then $\beta \cap (\alpha \circ (\beta \cap \delta) \circ \alpha) \leq \delta$, so $\mathbf{C}(\alpha, \beta; \delta)$ holds by (7).

Item (9) holds because any instance of the implication in Definition 2.18 defining $\mathbf{C}(S|_{\mathbf{B}}, T|_{\mathbf{B}}; \delta|_{\mathbf{B}})$ in \mathbf{B} is an instance of the implication defining $\mathbf{C}(S, T; \delta)$ in \mathbf{A} .

For item (10), it suffices to observe that, when $\delta' \leq \delta$,

$$\begin{bmatrix} p'/\delta' & q'/\delta' \\ r'/\delta' & s'/\delta' \end{bmatrix} \in M(S/\delta', T/\delta')$$

if and only if there exist $p \equiv_{\delta'} p'$, $q \equiv_{\delta'} q'$, $r \equiv_{\delta'} r'$, and $s \equiv_{\delta'} s'$ with

$$\begin{bmatrix} p & q \\ r & s \end{bmatrix} \in M(S, T),$$

and that $p \equiv_{\delta} q \Leftrightarrow p'/\delta' \equiv_{\delta/\delta'} q'/\delta'$ and $r \equiv_{\delta} s \Leftrightarrow r'/\delta' \equiv_{\delta/\delta'} s'/\delta'$. \square

DEFINITION 2.20. The **commutator** of S and T , denoted by $[S, T]$, is the least congruence δ such that $\mathbf{C}(S, T; \delta)$ holds. T is **abelian** if $[T, T] = 0$. An algebra \mathbf{A} is **abelian** if its largest congruence is.

By Theorem 2.19 (6), the class of all δ such that $\mathbf{C}(S, T; \delta)$ holds is closed under complete meet, so there is a least such δ . This implies that $[S, T]$ exists for any two tolerances S and T .

It is a well known fact, easily derivable from the definitions, that \mathbf{A} is abelian if and only if the diagonal of $A \times A$ is a class of a congruence of $\mathbf{A} \times \mathbf{A}$.

DEFINITION 2.21. The **centralizer of T modulo δ** , denoted by $(\delta : T)$, is the largest congruence α on \mathbf{A} such that $\mathbf{C}(\alpha, T; \delta)$ holds.

By Theorem 2.19 (5), the class of all α such that $\mathbf{C}(\alpha, T; \delta)$ holds is closed under complete join, so there is a largest such α . This implies that $(\delta : T)$ exists for every δ and T . By Theorem 2.19 (2), the centralizer $(\delta : T)$ contains every tolerance S such that $\mathbf{C}(S, T; \delta)$ holds.

2.6. Congruence Identities

If \mathcal{V} is a variety of algebras, then any lattice identity that holds in the class $\{\mathbf{Con}(\mathbf{A}) \mid \mathbf{A} \in \mathcal{V}\}$ of congruence lattices of algebras in \mathcal{V} is called a **congruence identity** of \mathcal{V} . The **congruence variety** of \mathcal{V} , denoted $\text{CON}(\mathcal{V})$, is the subvariety of \mathcal{L} generated by $\{\mathbf{Con}(\mathbf{A}) \mid \mathbf{A} \in \mathcal{V}\}$, or alternatively is the variety of lattices axiomatized by the congruence identities that hold in \mathcal{V} . Similarly, a lattice quasi-identity that holds in congruence lattices of members of \mathcal{V} is a **congruence quasi-identity** of \mathcal{V} .

The following theorem will be used in several places in this monograph.

THEOREM 2.22. (Cf. [6]) *Let Q be a quasi-identity satisfying (W). The class of varieties satisfying Q as a congruence quasi-identity is definable by a set of idempotent Maltsev conditions.*

"of" lattices ...

Since congruence lattices are algebraic, and therefore meet continuous, it is reasonable to discuss **meet continuous congruence identities**, which we define to be meet continuous lattice identities considered as congruence identities. A meet continuous congruence identity is **trivial** if it holds in the congruence lattice of every algebra and **nontrivial** otherwise. It follows from Corollary 2.10 that free meet continuous lattices are algebraic, so they are isomorphic to congruence lattices according to the Grätzer-Schmidt Theorem [26]. Therefore a meet continuous lattice identity is nontrivial as a congruence identity if and only if it is nontrivial as a meet continuous lattice identity. Both statements mean the identity fails in some algebraic lattice.

It is proved in [5] that the class of varieties that satisfy the meet semidistributive law as a congruence identity is definable by a set of idempotent Maltsev conditions. An examination of that argument reveals that what is really proved is that, modulo the axioms defining \mathcal{L}_{MC} , the meet semidistributive law is equivalent to a meet continuous lattice identity.¹ In the next theorem, we prove that many different quasi-identities, including SD_{\wedge} , are equivalent to meet continuous identities.

THEOREM 2.23. *Let $p, q, r, s, x_1, \dots, x_n$ be lattice variables. Define $t := p \wedge (q \vee r)$. Let $w_i(p, q, r)$, $1 \leq i \leq m$, be ternary lattice words such that $w_i(p, q, r) \leq q$ for some i , $w_j(p, q, r) \leq r$ for some j , and $w_k(p, q, r) \leq t$ for all k in the free lattice $\mathbf{F}_{\mathcal{L}}(\{p, q, r\})$. For each i let $x_i^* = (x_i \vee s) \wedge t$. For arbitrary n -ary lattice words u and v , the quasi-identity*

$$(2.6) \quad \bigwedge_{i=1}^m (w_i(p, q, r) \approx s) \rightarrow u(x_1^*, \dots, x_n^*) \approx v(x_1^*, \dots, x_n^*)$$

is equivalent to an identity in \mathcal{L}_{MC} .

PROOF. For each i let $f_i(x)$ be the polynomial $w_i(p, q \vee x, r \vee x)$ of the free meet continuous lattice $\mathbf{F}_0 = \mathbf{F}_{\mathcal{L}_{MC}}(\{p, q, r\})$. If $z \leq t = p \wedge (q \vee r)$, then the idempotence of lattice words implies that for any i between 1 and m

$$(2.7) \quad \begin{aligned} \underline{z} &= w_i(z, z, z) \leq w_i(p, q \vee z, r \vee z) = \underline{f_i(z)} \\ &\leq w_i(p, q \vee t, r \vee t) \leq w_i(p, q \vee r, q \vee r) \\ &\leq p \wedge ((q \vee r) \vee (q \vee r)) = \underline{t}. \end{aligned}$$

The underlined elements indicate why the principal ideal generated by $t = p \wedge (q \vee r)$ is closed under each f_i , and each f_i is an increasing

¹In fact, what is *really* shown is that the partial lattice $\mathbf{P}(SD_{\wedge})$ in the SD_{\wedge} -configuration is projective relative to \mathcal{L}_{MC} .

function on this ideal. It follows from this that the set of all elements of the form $f_{i_1} \circ \dots \circ f_{i_\ell}(0)$, $i_j \in \{1, \dots, m\}$, is an updirected set of elements below t . The supremum of this set, $\delta^{\mathbf{w}}(p, q, r) \in F_0$, $\mathbf{w} := (w_1, \dots, w_m)$, is therefore the least common fixed point in \mathbf{F}_0 of the polynomials f_i . Since $t = p \wedge (q \vee r)$ is a common fixed point of the f_i it follows that $\delta^{\mathbf{w}} \leq t$ in \mathbf{F}_0 . Now, if $\varphi: \mathbf{F}_0 \rightarrow \mathbf{L}$ is a surjective homomorphism of meet continuous lattices for which $a = \varphi(p)$, $b = \varphi(q)$ and $c = \varphi(r)$, then $\varphi(\delta^{\mathbf{w}}(p, q, r)) = \delta^{\mathbf{w}}(a, b, c)$ is the supremum of all $f_{i_1}^{\mathbf{L}} \circ \dots \circ f_{i_k}^{\mathbf{L}}(0)$, where $f_i^{\mathbf{L}}(x) := w_i^{\mathbf{L}}(a, b \vee x, c \vee x)$. By an argument like the one used in (2.7), each $f_i^{\mathbf{L}}$ is increasing on the principal ideal generated by $t^{\mathbf{L}}(a, b, c) = a \wedge (b \vee c)$ and maps this ideal into itself. Hence $\delta^{\mathbf{w}}(a, b, c)$ is the least common fixed point of the polynomials $f_i^{\mathbf{L}}$ in \mathbf{L} .

For each variable x_i , set $x_i^+ := (x_i \vee \delta^{\mathbf{w}}(p, q, r)) \wedge t$, which is an element of $\mathbf{F} := \mathbf{F}_{\mathcal{L}_{MC}}(\{p, q, r, s, x_1, \dots, x_n\})$. We are now prepared to argue that quasi-identity (2.6) is equivalent to the identity

$$(2.8) \quad u(x_1^+, \dots, x_n^+) \approx v(x_1^+, \dots, x_n^+)$$

in \mathcal{L}_{MC} .

Assume first that \mathbf{L} is a meet continuous lattice satisfying (2.8). To show that \mathbf{L} satisfies (2.6), choose any assignment of the variables $\{p, q, r, s, \mathbf{x}\} \rightarrow L$ that satisfies the premises of quasi-identity (2.6), and let $\psi: \mathbf{F} \rightarrow \mathbf{L}$ be the extension of this assignment to a homomorphism. Equivalently, choose a homomorphism ψ so that if

$$(\psi(p), \psi(q), \psi(r), \psi(s), \psi(\mathbf{x})) = (a, b, c, d, \mathbf{g}),$$

then $w_i^{\mathbf{L}}(a, b, c) = d$ for all i . Define $e := \psi(t) = a \wedge (b \vee c)$ and $g_i^* := \psi(x_i^+) = (g_i \vee d) \wedge e$. We must show that $u^{\mathbf{L}}(\mathbf{g}^*) = v^{\mathbf{L}}(\mathbf{g}^*)$.

Since $w_i(p, q, r) \leq q$ for some i , $w_j(p, q, r) \leq r$ for some j , and $w_k(p, q, r) \leq (p \wedge (q \vee r)) \leq p$ for all k in \mathbf{F}_0 , it follows that

$$d = w_i(a, b, c) = w_j(a, b, c) = w_k(a, b, c) \leq a \wedge b \wedge c$$

in \mathbf{L} . Since $d = w_1(a, b, c) \geq a \wedge b \wedge c$ we even have $d = a \wedge b \wedge c$. Therefore $f_i^{\mathbf{L}}(d) = t_i^{\mathbf{L}}(a, b \vee d, c \vee d) = t_i^{\mathbf{L}}(a, b, c) = d$ for all i , yielding that d is common fixed point of the f_i in \mathbf{L} . Since the meet continuous sublattice $\mathbf{L}_0 \leq \mathbf{L}$ that is generated by a, b and c contains d and has $\delta^{\mathbf{w}}(a, b, c)$ as its least common fixed point of the f_i 's, it follows that $\delta^{\mathbf{w}}(a, b, c) \leq d$. Since $d = a \wedge b \wedge c \leq f_1(0) \leq \delta^{\mathbf{w}}(a, b, c)$, we even have $d = \delta^{\mathbf{w}}(a, b, c)$. This implies that $\psi(x_i^+) = \psi(x_i^*) = g_i^*$. This gives us the desired result

$$u^{\mathbf{L}}(\mathbf{g}^*) = \psi(u(\mathbf{x}^+)) = \psi(v(\mathbf{x}^+)) = v^{\mathbf{L}}(\mathbf{g}^*),$$

where the middle equality is from (2.8).

w_i, not t_i

... of the f_i ...

Now suppose conversely that \mathbf{L} satisfies (2.6). To verify that \mathbf{L} satisfies (2.8), choose any variable assignment in \mathbf{L} and extend it to a homomorphism $\psi: \mathbf{F} \rightarrow \mathbf{L}$. We must show that $\psi(u(\mathbf{x}^+)) = \psi(v(\mathbf{x}^+))$. Since the variable s does not appear (2.8) we need only consider homomorphisms ψ for which $\psi(s) = \psi(\delta^{\mathbf{w}}(p, q, r))$, i.e., those of the form

$$(\psi(p), \psi(q), \psi(r), \psi(s), \psi(\mathbf{x})) = (a, b, c, \delta^{\mathbf{w}}(a, b, c), \mathbf{g}).$$

Under such homomorphisms

$$\psi(x_i^+) = \psi(x_i^*) = (g_i \vee \delta^{\mathbf{w}}(a, b, c)) \wedge (a \wedge (b \vee c)).$$

The element $\delta := \delta^{\mathbf{w}}(a, b, c)$ is the least common fixed point of the polynomials $f_i^{\mathbf{L}}(x) := w^{\mathbf{L}}(a, b \vee x, c \vee x)$ in the meet continuous sublattice of \mathbf{L} generated by a, b and c . Thus, if $b^\delta := b \vee \delta$ and $c^\delta := c \vee \delta$, we get that

$$\delta = f_i^{\mathbf{L}}(\delta) = w_i(a, b \vee \delta, c \vee \delta) = w_i(a, b^\delta, c^\delta)$$

for all i . Thus, each premise $w_i(a, b^\delta, c^\delta) = \delta$ of quasi-identity (2.6) is satisfied by the variable assignment $\psi^\delta: (p, q, r, s, \mathbf{x}) \mapsto (a, b^\delta, c^\delta, \delta, \mathbf{g})$. Since \mathbf{L} satisfies (2.6), we get that

$$(2.9) \quad u^{\mathbf{L}}(\psi^\delta(\mathbf{x}^*)) = v^{\mathbf{L}}(\psi^\delta(\mathbf{x}^*))$$

holds in \mathbf{L} , where

$$\psi^\delta(x_i^*) = (g_i \vee \delta) \wedge (a \wedge (b^\delta \vee c^\delta)).$$

But $\delta^{\mathbf{w}}(p, q, r) \leq p \wedge (q \vee r)$ in $\mathbf{F}_{\mathcal{L}_{MC}}(\{p, q, r\})$, so $\delta = \delta^{\mathbf{w}}(a, b, c) \leq a \wedge (b \vee c)$ in \mathbf{L} . This yields the last equality in

$$a \wedge (b \vee c) \leq a \wedge (b^\delta \vee c^\delta) = a \wedge (b \vee c \vee \delta) = a \wedge (b \vee c).$$

This shows that $a \wedge (b^\delta \vee c^\delta) = a \wedge (b \vee c)$, so $\psi^\delta(x_i^*) = (g_i \vee \delta) \wedge (a \wedge (b \vee c)) = \psi(x_i^+)$. From this and (2.9) we derive that

$$\psi(u(\mathbf{x}^+)) = u^{\mathbf{L}}(\psi^\delta(\mathbf{x}^*)) = v^{\mathbf{L}}(\psi^\delta(\mathbf{x}^*)) = \psi(v(\mathbf{x}^+)),$$

as desired. This proves that (2.8) holds in \mathbf{L} . \square

REMARK 2.24. If \mathbf{L} is meet continuous and satisfies quasi-identity (2.6), then for every homomorphism $\psi: \mathbf{F} \rightarrow \mathbf{L}$ for which $(p, q, r, s, \mathbf{x}) \mapsto (a, b, c, d, \mathbf{g})$ and $w_i(a, b, c) = d$ for all i , it must be that d is below $a \wedge (b \vee c) =: e$. Call $I := I[d, e]$ a “bad interval” in \mathbf{L} if it arises in this way from elements satisfying the premises of (2.6). For each i , $\psi(x_i^*) = (g_i \vee d) \wedge e \in I$, and conversely for any $h \in I$ we have $h = \psi(x_i^*)$ if ψ assigns h to x_i . Thus, $\psi(x_i^*)$ may be viewed as a typical element of I . Altogether this means that the premises of (2.6) identify which intervals are bad, and the quasi-identity itself asserts that all bad intervals satisfy the identity $u \approx v$.

you just use that $w_i(p, q, r) \leq e$ for this, not that the quasi-identity holds in \mathbf{L} .

If we choose $m = n = 2$, $w_1(p, q, r) = p \wedge q$, $w_2(p, q, r) = p \wedge r$, then the bad intervals are the SD_\wedge -failures. If $u(x_1, x_2) = x_1$ and $v(x_1, x_2) = x_2$, then (2.6) asserts precisely that all SD_\wedge -failures satisfy $x_1 \approx x_2$, i.e., they are trivial. These choices for w_i, u and v show that the meet semidistributive law is equivalent to a quasi-identity of type (2.6), hence to a meet continuous identity.

Consider choosing the premises of (2.6) in the same way as in the previous paragraph, but allowing $u \approx v$ to be some other lattice identity. We obtain from Theorem 2.23 that the class of meet continuous lattices whose SD_\wedge -failures satisfy $u \approx v$ is a subvariety of \mathcal{L}_{MC} .

Now consider modifying the choice of the premises of (2.6) from our earlier choice, but keeping $u(x_1, x_2) = x_1$ and $v(x_1, x_2) = x_2$. In this case, (2.6) still asserts that the bad intervals are trivial, but the bad intervals need not be the SD_\wedge -failures. Instead, the bad intervals are those of the form $I[d, e]$ where for some $a, b, c \in L$ we have $w_i(a, b, c) = d$ for all i and $e = a \wedge (b \vee c)$. A more direct way of asserting that $I[d, e]$ is trivial is to assert that $e = d$ (where $e = a \wedge (b \vee c)$). This is achieved by rewriting quasi-identity (2.6) in the form

$$(2.10) \quad \bigwedge_{i=1}^m (w_i(p, q, r) \approx s) \rightarrow ((p \wedge (q \vee r)) = s).$$

Quasi-identity (2.10) is equivalent modulo the identities defining lattice theory to quasi-identity (2.6) when $u(x_1, x_2) = x_1$ and $v(x_1, x_2) = x_2$. In particular, Theorem 2.23 proves that a quasi-identity obtained from the meet semidistributive law,

$$((p \wedge q) \approx s) \& ((p \wedge r) \approx s) \rightarrow ((p \wedge (q \vee r)) = s),$$

by adding additional premises of the form $w_i(p, q, r) \approx s$ is equivalent to a meet continuous identity.

By examining the proof of Theorem 2.23 one finds that the bad intervals in \mathbf{L} associated to a quasi-identity of type (2.6) or (2.10) are those of the form $I[d, e]$ where for some $a, b, c \in L$ it is the case that $d = \delta^{\mathbf{w}}(a, b, c)$ and $e = a \wedge (b \vee c)$. Since quasi-identity (2.10) asserts only that bad intervals are trivial, a simpler meet continuous identity equivalent to (2.10) is

$$(2.11) \quad \delta^{\mathbf{w}}(p, q, r) \approx (p \wedge (q \vee r)).$$

It is not hard to describe explicitly the meet continuous identity that is equivalent to the meet semidistributive law. Define lattice words in the variables p, q, r by $q_0 = q$, $r_0 = r$, $q_{n+1} = q \vee (p \wedge r_n)$, and $r_{n+1} = r \vee (p \wedge q_n)$. Let $q_\omega = \bigvee_{n < \omega} q_n$ and let $r_\omega = \bigvee_{n < \omega} r_n$.

do we need to assume that for these w_i that $w_i(p, q, r) \approx s$?

perhaps some mention of the origin of these lattices words should be made here. They appear in the work of Czedli, as well as in Hobby-McKenzie, and for all I know in other work.

COROLLARY 2.25. *A meet continuous lattice is meet semidistributive if and only if it satisfies the identity*

$$(2.12) \quad p \wedge (q \vee r) \approx p \wedge q_\omega .$$

PROOF. This is an instance of the comments in the final paragraph of Remark 2.24. For the choices $w_1(p, q, r) = p \wedge q$ and $w_2(p, q, r) = p \wedge r$ the polynomials f_i that occur in the proof of Theorem 2.23 are $f_1(x) = p \wedge (q \vee x) =: g(x)$ and $f_2(x) = p \wedge (r \vee x) =: h(x)$. Since these are idempotent polynomials of $\mathbf{F}_{\mathcal{L}_{MC}}(p, q, r)$, the words of the form $f_{i_1} \circ \dots \circ f_{i_\ell}(0)$ that occur in the preceding proof reduce to alternating compositions $(gh)^k(0)$, $(hg)^k(0)$, $g(hg)^k(0)$, or $h(gh)^k(0)$. The join of these words is $\delta^{\mathbf{w}}(p, q, r) = \bigvee_{k < \omega} (gh)^k(0) = \bigvee_{k < \omega} (hg)^k(0)$. But $g(0) = p \wedge q_0$, $h(0) = p \wedge r_0$, $g(p \wedge r_n) = p \wedge q_{n+1}$ and $h(p \wedge q_n) = p \wedge r_{n+1}$, so $\delta^{\mathbf{w}}(p, q, r) = \bigvee_{k \text{ odd}} p \wedge q_k = p \wedge q_\omega$ by meet continuity, and also $\delta^{\mathbf{w}}(p, q, r) = \bigvee_{k \text{ even}} p \wedge r_k = p \wedge r_\omega$. Hence the meet semidistributive law is equivalent to $p \wedge q_\omega \approx p \wedge (q \vee r)$ or to $p \wedge r_\omega \approx p \wedge (q \vee r)$ for meet continuous lattices. \square

This proof shows that $p \wedge q_\omega \approx p \wedge r_\omega$, so identity (2.12) can be written as a weakened distributive law:

$$p \wedge (q \vee r) \approx (p \wedge q_\omega) \vee (p \wedge r_\omega) .$$

CHAPTER 3

Strong Term Conditions

In this chapter we introduce the term conditions that define strongly abelian and strongly rectangular congruences. Our purpose is to show that a variety satisfies a nontrivial idempotent Maltsev condition if and only if it has no member with a nonzero strongly abelian congruence, or equivalently has no member with a nonzero strongly rectangular congruence.

3.1. Varieties Omitting Strongly Abelian Congruences

Ralph McKenzie introduced the notion of “strong abelianness” in [66], and, with David Hobby, showed in [34] that it is a natural and useful concept in the study of finite algebras. McKenzie also discovered a quite different abelianness concept, called “rectangularity” in [54], which made sense for finite subdirectly irreducible algebras with nonabelian monolith, and which has applications to the study of residually small varieties.

In [49], we required a commutator theory associated to the strong abelianness concept. In the process of developing one, we discovered a new natural and useful type of abelianness resembling McKenzie’s rectangularity concept that we will call in this monograph “strong rectangularity”. In [49] we developed a commutator theory for strong rectangularity, too. In this chapter we explore these two theories further. In Chapter 5 we will introduce a new definition for McKenzie’s original rectangularity concept that makes sense for any algebra, and initiate the development of a commutator theory for it.

It is important to draw attention to the fact that we used the term “rectangularity” in [49] for what will henceforth be called “strong rectangularity”.

DEFINITION 3.1. Let S and T be tolerances on an algebra \mathbf{A} , and let δ be a congruence on \mathbf{A} . If $q \equiv_{\delta} r$ implies that $r \equiv_{\delta} s$ whenever

$$(3.1) \quad \begin{bmatrix} p & q \\ r & s \end{bmatrix} \in M(S, T),$$

then we say that **SR**($S, T; \delta$) **holds**, or S **strongly rectangulates** T **modulo** δ . We say that **S**($S, T; \delta$) **holds**, or S **strongly centralizes** T **modulo** δ if **SR**($S, T; \delta$) and **C**($S, T; \delta$) both hold.

A tolerance T is **strongly rectangular** if **SR**($T, T; 0$) holds, and is **strongly abelian** if **S**($T, T; 0$) holds.

By interchanging the rows and columns of matrices one sees that **SR**($S, T; \delta$) holds if and only if

$$q \equiv_{\delta} r \iff p \equiv_{\delta} q \equiv_{\delta} r \equiv_{\delta} s \iff p \equiv_{\delta} s.$$

for every S, T -matrix in (3.1). It follows from this that **SR**($S, T; \delta$) is equivalent to **SR**($T, S; \delta$).

Now, and later, we must consider the issue of whether nonobvious relationships hold among the term conditions that we introduce. For example, is every abelian tolerance strongly abelian?¹ There is an all-purpose example for answering this type of question, which may be used to demonstrate the independence of any combination of the term conditions that we define in this monograph.

EXAMPLE 3.2. Let X be a set and let $\mathcal{M} \subseteq X^{2 \times 2}$ be a set of 2×2 matrices with entries in X . For $M \in \mathcal{M}$, let $Y_M = \{a_M, b_M, u_M, v_M\}$. Let $Y = \cup_{M \in \mathcal{M}} Y_M$. Let A be the disjoint union $X \cup Y \cup \{0\}$. For each $M = \begin{bmatrix} p & q \\ r & s \end{bmatrix} \in \mathcal{M}$ define a binary operation f_M on A by

$$f_M(x, y) := \begin{cases} p & \text{if } (x, y) = (a_M, u_M); \\ q & \text{if } (x, y) = (a_M, v_M); \\ r & \text{if } (x, y) = (b_M, u_M); \\ s & \text{if } (x, y) = (b_M, v_M); \text{ and} \\ 0 & \text{otherwise.} \end{cases}$$

should be A(M)

Let $\mathbf{A}(\mathcal{M}) = \langle A; \{f_M \mid M \in \mathcal{M}\} \rangle$, $R = \text{Tg}^{\mathbf{A}}(\{(a_M, b_M) \mid M \in \mathcal{M}\})$, $S = \text{Tg}^{\mathbf{A}}(\{(u_M, v_M) \mid M \in \mathcal{M}\})$, and $T = \text{Tg}^{\mathbf{A}}(R \cup S)$.

CLAIM 3.3. *The following hold.*

- (1) $M(R, S)$ consists only of the trivial R, S -matrices, namely those of the form

$$\begin{bmatrix} p & p \\ q & q \end{bmatrix} \quad \text{and} \quad \begin{bmatrix} r & s \\ r & s \end{bmatrix}$$

with $(p, q) \in R$ and $(r, s) \in S$, together with the closure of \mathcal{M} under interchanging rows and columns.

¹No.

- (2) $M(T, T)$ consists only of the trivial T, T -matrices together with the closure of \mathcal{M} under interchanging rows and columns and under transpose.

We leave the proof of this claim as an exercise, guided by the following comments.

- (i) The trivial matrices come from using a projection operation as the term in the matrix.
(ii) Any nontrivial composition of f_M 's is zero, so only $f_M(x, y)$ and $f_M(y, x)$ can be used to create a nontrivial matrix. However $f_M(y, x)$ does not produce any nontrivial R, S -matrices.

We also leave the verification of the following as an exercise: the claim remains true if you replace R, S and T by the congruences they generate. And, as a final exercise: any failure of any term condition introduced in this monograph must involve a nontrivial R, S or T, T -matrix.

What this means is that when considering relationships between term conditions we may ignore any algebraic considerations and think completely set-theoretically. If it is possible to write down matrices whose patterns of entries formally conflict with the satisfaction of one term condition but do not formally conflict with the satisfaction of another, then it is possible to build an algebra that has matrices with these properties (subject to the restrictions of (1) and (2) of Claim 3.3).

For example, to show that there exists an algebra that has an abelian tolerance that is not strongly abelian, it is enough to observe that any T, T -matrix of the form $M = \begin{bmatrix} p & q \\ r & s \end{bmatrix} = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}$ formally conflicts with the satisfaction of $\mathbf{SR}(T, T; 0)$ (since $q = r \neq s$), but does not formally conflict with $\mathbf{C}(T, T; 0)$ (even after interchanging the rows or columns, or taking the transpose). This suggests that we take $X = \{0, 1\}$ and $\mathcal{M} = \{M\}$. By Claim 3.3, the nontrivial T, T -matrices of $\mathbf{A}(\mathcal{M})$ are

$$\begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} \quad \text{and} \quad \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix}.$$

This implies that the tolerance T on $\mathbf{A}(\mathcal{M})$ is abelian but not strongly rectangular (hence not strongly abelian).

In Lemma 2.2 of [49] we state and prove a result for $\mathbf{SR}(\ , \ ; \)$ similar to Theorem 2.19 for $\mathbf{C}(\ , \ ; \)$. These two results can be combined to produce a similar result for $\mathbf{S}(\ , \ ; \)$. Here we describe only the results about the strong centrality relation and the strong rectangulation relation that will be used in this chapter.

to be consistent,
use "rectangularity"
instead of
"rectangulation"
here.

THEOREM 3.4. *Let \mathbf{A} be an algebra with tolerances S, S', T, T' and congruences $\delta, \delta', \delta_j$. Let $\mathbf{Q} = \mathbf{S}$ or \mathbf{SR} . The following are true.*

- (1) *If $\mathbf{Q}(S, T; \delta)$ holds and $S' \subseteq S, T' \subseteq T$, then $\mathbf{Q}(S', T'; \delta)$ holds.*
- (2) *If $\mathbf{Q}(S, T; \delta_j)$ holds for all $j \in J$, then $\mathbf{Q}(S, T; \bigwedge_{j \in J} \delta_j)$ holds.*
- (3) *Let \mathbf{B} be a subalgebra of \mathbf{A} . If $\mathbf{Q}(S, T; \delta)$ holds in \mathbf{A} , then $\mathbf{Q}(S|_{\mathbf{B}}, T|_{\mathbf{B}}; \delta|_{\mathbf{B}})$ holds in \mathbf{B} .*
- (4) *If $\delta' \leq \delta$, then the relation $\mathbf{Q}(S, T; \delta)$ holds in \mathbf{A} if and only if $\mathbf{Q}(S/\delta', T/\delta'; \delta/\delta')$ holds in \mathbf{A}/δ' .*
- (5) *If $T \cap (((S \circ T) \cap \delta) \circ S) \subseteq \delta$, then $\mathbf{SR}(S, T; \delta)$ holds. If $T \cap (S \circ (T \cap \delta) \circ S) \subseteq \delta$ also holds, then $\mathbf{S}(S, T; \delta)$ also holds.*

PROOF. Items (1)–(4) are proved in exactly the same way as items (1), (6), (9) and (10) of Theorem 2.19.

For item (5), choose any S, T -matrix $M = \begin{bmatrix} p & q \\ r & s \end{bmatrix}$. Assume that $r \equiv_{\delta} q$. Since the elements in the same row of M are T -related and the elements in the same column are S -related, we have $r S p T q$ and $r \delta q$, so $r (S \circ T) \cap \delta q$. Since $r (S \circ T) \cap \delta q S s$ and $r T s$ we have $(r, s) \in T \cap (((S \circ T) \cap \delta) \circ S) \subseteq \delta$. Since M was chosen arbitrarily, $\mathbf{SR}(S, T; \delta)$ holds. If now $T \cap (S \circ (T \cap \delta) \circ S) \subseteq \delta$ also holds, then $\mathbf{C}(S, T; \delta)$ also holds by Theorem 2.19 (7). Hence $\mathbf{S}(S, T; \delta)$ also holds. \square

Theorem 3.4 (2) implies that for $\mathbf{Q} = \mathbf{S}$ or \mathbf{SR} there is a least δ such that $\mathbf{Q}(S, T; \delta)$ holds, and we denote the least such δ by $[S, T]_{\mathbf{Q}}$.

LEMMA 3.5. *Let \mathcal{V} be a variety, let $\mathbf{F} = \mathbf{F}_{\mathcal{V}}(x, y)$ be the free \mathcal{V} -algebra over $X = \{x, y\}$, and let $\theta = \text{Cg}^{\mathbf{F}}(x, y)$. Let $\mathbf{Q} = \mathbf{S}$ or \mathbf{SR} . The following are equivalent.*

- (1) *There exists an algebra $\mathbf{A} \in \mathcal{V}$ that has a nonzero \mathbf{Q} -abelian congruence (i.e., a congruence α such that $[\alpha, \alpha]_{\mathbf{Q}} = 0$).*
- (2) *$(x, y) \notin [\theta, \theta]_{\mathbf{Q}}$.*

Now let T be the tolerance on \mathbf{F} generated by the pair (x, y) . The following are equivalent.

- (3) *There exists an algebra $\mathbf{A} \in \mathcal{V}$ that has a nonzero \mathbf{Q} -abelian tolerance (i.e., a tolerance S such that $[S, S]_{\mathbf{Q}} = 0$).*
- (4) *$(x, y) \notin [T, T]_{\mathbf{Q}}$.*

PROOF. Both halves of the lemma are proved in the same way, so we prove the first half only.

Assume that α is a nonzero \mathbf{Q} -abelian congruence on \mathbf{A} . Choose $(a, b) \in \alpha$ with $a \neq b$. Then from Theorem 3.4 (3) and (1) we get that if \mathbf{B} is the subalgebra of \mathbf{A} generated by $\{a, b\}$ and $\beta = \text{Cg}^{\mathbf{B}}(a, b)$,

then β is \mathbf{Q} -abelian. Let $h: \mathbf{F} \rightarrow \mathbf{B}$ be the homomorphism sending x to a and y to b . This homomorphism is surjective since \mathbf{B} is generated by a and b . Let $\kappa = \ker(h)$ and let $\lambda = h^{-1}(\beta)$. Since $(h(x), h(y)) = (a, b) \in \beta \setminus 0$ we have $(x, y) \in \lambda \setminus \kappa$, which means that $\theta \leq \lambda$ and $\theta \not\leq \kappa$. From Theorem 3.4 (4) we know that $\mathbf{Q}(\lambda, \lambda; \kappa)$ holds in \mathbf{F} if and only if $\mathbf{Q}(\lambda/\kappa, \lambda/\kappa; 0)$ holds in \mathbf{F}/κ . But it does, since λ/κ corresponds to β under the natural isomorphism of \mathbf{F}/κ with \mathbf{B} . Therefore, by Theorem 3.4 (1), we derive that $\mathbf{Q}(\theta, \theta; \kappa)$ holds. We have $\mathbf{Q}(\theta, \theta; \theta)$ by Theorem 3.4 (5), and therefore $\mathbf{Q}(\theta, \theta; \kappa \wedge \theta)$ by Theorem 3.4 (2). This proves that $[\theta, \theta]_{\mathbf{Q}} \leq \kappa \wedge \theta < \theta$, so $(x, y) \notin [\theta, \theta]_{\mathbf{Q}}$.

On the other hand, if $(x, y) \notin [\theta, \theta]_{\mathbf{Q}}$, then Theorem 3.4 (4) guarantees that $\theta/[\theta, \theta]_{\mathbf{Q}}$ is a \mathbf{Q} -abelian congruence on $\mathbf{F}/[\theta, \theta]_{\mathbf{Q}}$ that relates the distinct elements $x/[\theta, \theta]_{\mathbf{Q}}$ and $y/[\theta, \theta]_{\mathbf{Q}}$. This shows that \mathcal{V} contains an algebra with a nonzero \mathbf{Q} -abelian congruence. \square

Our goal in this section is to prove that the four conditions of Lemma 3.5 are equivalent to each other, and that when $\mathbf{Q} = \mathbf{S}$ the conditions hold for a variety if and only if the variety has no Taylor term (Definition 2.15). This is also true when $\mathbf{Q} = \mathbf{SR}$, and that will be proved in the next section. To prove the equivalence of (1) and (2) with (3) and (4) we start by analyzing the relationship between the congruence θ and the tolerance T from Lemma 3.5. This is done in a little more generality than is immediately necessary, in that the next lemma considers free algebras on more than two variables, but this will allow us to avoid the need to generalize the result later.

LEMMA 3.6. *Let \mathcal{V} be any variety and let $\mathbf{F} = \mathbf{F}_{\mathcal{V}}(X)$ be the free \mathcal{V} -algebra generated by X . Let E be an equivalence relation on X , and let $\tau: X \rightarrow X$ be a transversal for E . That is, τ is a function from X to X satisfying*

- (i) $x E \tau(x)$, and
- (ii) $x E y \implies \tau(x) = \tau(y)$.

Let $\theta = \text{Cg}^{\mathbf{F}}(E)$ and $T = \text{Tg}^{\mathbf{F}}(E)$. The following hold.

- (1) θ is the kernel of the endomorphism $h: \mathbf{F} \rightarrow \mathbf{F}$ induced by $\tau: X \rightarrow X$.
- (2) $(a, h(a)) \in T$ for any $a \in F$.
- (3) $T \circ T = \theta$.
- (4) If $B := \{u(\mathbf{x}) \in F \mid u \text{ is an idempotent term}\}$, $E = X \times X$, and $\tau(X) = \{x\}$, then $B = x/\theta$ and B is the unique T -block containing x .

PROOF. Observe that the fact that τ is a transversal for E implies that $\tau \circ \tau = \tau$ and $E = \ker(\tau)$.

We prove (2) first. If $a \in F$, then $a = t(x_1, \dots, x_n)$ for some term t . Since T is a compatible relation containing $(x_i, \tau(x_i))$ for every i , it also contains $t((x_1, \tau(x_1)), \dots, (x_n, \tau(x_n))) = (t(\mathbf{x}), h(t(\mathbf{x}))) = (a, h(a))$.

For (1), the fact that h is induced by an idempotent function on generators implies that h is an idempotent endomorphism, therefore h is a transversal for $\ker(h)$. Since $\ker(\tau) \subseteq \ker(h)$ we have $\theta \subseteq \ker(h)$. It follows that every θ -block contains at most one element of $\text{im}(h)$. On the other hand, every θ -block contains at least one element of $\text{im}(h)$, since if a is in a θ -block B , then $h(a)$ is certainly in $\text{im}(h)$, while $(a, h(a)) \in T \subseteq \theta$ by (2). Thus $h(a) \in \text{im}(h) \cap B$. This shows that h is also a transversal for θ . Since comparable equivalence relations with a common transversal are equal, $\theta = \ker(h)$.

For (3), $T \circ T \subseteq \text{Cg}^{\mathbf{F}}(T) = \theta$, so it is enough to show that $T \circ T \supseteq \theta$. If $(a, b) \in \theta$, then $h(a) = h(b)$. By (2), we have

$$a T h(a) = h(b) T b,$$

proving that $(a, b) \in T \circ T$, indeed.

For (4), it is clear that $(x, u(x_1, \dots, x_n)) \in \theta = \ker(h)$ if and only if

$$u(x, x, \dots, x) = h(u(x_1, \dots, x_n)) = h(x) = x$$

in \mathbf{F} , which holds if and only if u is idempotent. Thus $B = x/\theta$, as claimed. If $u(x_1, \dots, x_n)$ and $v(x_1, \dots, x_n)$ are arbitrarily chosen idempotent terms, then using the term

$$u(v(x_{11}, \dots, x_{1n}), \dots, v(x_{n1}, \dots, x_{nn})),$$

and the fact that $(x_i, x_j) \in E \subseteq T$ for all i and j , we derive that

$$\begin{aligned} u(x_1, \dots, x_n) &= u(v(x_1, \dots, x_1), \dots, v(x_n, \dots, x_n)) \\ &\equiv_T u(v(x_1, \dots, x_n), \dots, v(x_1, \dots, x_n)) \\ &= v(x_1, \dots, x_n). \end{aligned}$$

Thus B is contained in a T -block. That it is the unique T -block containing x follows from the fact that $B = x/\theta$ is a θ -block (so no element in $F - B$ is T -related to any element of B). \square

LEMMA 3.7. *Let \mathcal{V} be a variety and let $\mathbf{F} = \mathbf{F}_{\mathcal{V}}(x, y)$ be the free \mathcal{V} -algebra generated by $\{x, y\}$. Let $\theta = \text{Cg}^{\mathbf{F}}(x, y)$ and T be the tolerance of \mathbf{F} generated by (x, y) . Let $B = x/\theta$.*

(1) *Let ρ be a congruence on \mathbf{F} such that $\mathbf{SR}(T, T; \rho)$ holds. If*

$$\begin{bmatrix} p & q \\ r & s \end{bmatrix} \in M(\theta, \theta)$$

and $p, q, r, s \in B$, then $q \equiv_{\rho} r \implies r \equiv_{\rho} s$.

(2) Let ρ be a congruence on \mathbf{F} such that $\mathbf{S}(T, T; \rho)$ holds. If

$$\begin{bmatrix} p & q \\ r & s \end{bmatrix} \in M(\theta, \theta)$$

and $p, q, r, s \in B$, then $p \equiv_{\rho} q \implies r \equiv_{\rho} s$.

PROOF. Choose any θ, θ -matrix

$$\begin{bmatrix} p & q \\ r & s \end{bmatrix} = \begin{bmatrix} t(\mathbf{a}, \mathbf{c}) & t(\mathbf{a}, \mathbf{d}) \\ t(\mathbf{b}, \mathbf{c}) & t(\mathbf{b}, \mathbf{d}) \end{bmatrix}.$$

Let $h: \mathbf{F} \rightarrow \mathbf{F}$ be the endomorphism determined by $x \mapsto x, y \mapsto x$. The assumption $\mathbf{a} \equiv_{\theta} \mathbf{b}$ is equivalent to $h(\mathbf{a}) = h(\mathbf{b})$, and similarly $\mathbf{c} \equiv_{\theta} \mathbf{d}$ is equivalent to $h(\mathbf{c}) = h(\mathbf{d})$ (where by $h(\mathbf{a})$, etc., we mean that we apply the function h componentwise). Consider the following 4×4 matrix:

$$\begin{bmatrix} t(\mathbf{a}, \mathbf{c}) & t(\mathbf{a}, h(\mathbf{c})) & = & t(\mathbf{a}, h(\mathbf{d})) & t(\mathbf{a}, \mathbf{d}) \\ t(h(\mathbf{a}), \mathbf{c}) & t(h(\mathbf{a}), h(\mathbf{c})) & = & t(h(\mathbf{a}), h(\mathbf{d})) & t(h(\mathbf{a}), \mathbf{d}) \\ \parallel & \parallel & & \parallel & \parallel \\ t(h(\mathbf{b}), \mathbf{c}) & t(h(\mathbf{b}), h(\mathbf{c})) & = & t(h(\mathbf{b}), h(\mathbf{d})) & t(h(\mathbf{b}), \mathbf{d}) \\ t(\mathbf{b}, \mathbf{c}) & t(\mathbf{b}, h(\mathbf{c})) & = & t(\mathbf{b}, h(\mathbf{d})) & t(\mathbf{b}, \mathbf{d}) \end{bmatrix}.$$

We claim that the four elements in the middle are equal to x . Indeed, the assumption that $p, q, r, s \in B$ implies that $h(p) = h(q) = h(r) = h(s) = x$. Therefore, for example,

$$t(h(\mathbf{b}), h(\mathbf{d})) = h(t(\mathbf{b}, \mathbf{d})) = h(s) = x.$$

Note also that the four 2×2 matrices in the corners are T, T -matrices, since \mathbf{u} and $h(\mathbf{u})$ are T -related componentwise by Lemma 3.6 (2). All elements of this 4×4 matrix are contained in B , because all elements are θ -related, the corners are in B by assumption, and B is a block of θ .

Next, let $k: \mathbf{F} \rightarrow \mathbf{F}$ be the endomorphism of \mathbf{F} that sends x to y and y to y . Consider the analogous 4×4 matrix with respect to k . All the statements above remain true with k in place of h , with the exception that in this case the four elements in the middle are equal to y , not x . Now apply $q(x, y)$ to these two matrices, the first matrix in the first argument, the second matrix in the second argument. This produces a 4×4 matrix which, after deleting one of the doubled middle columns and one of the doubled middle rows, is a 3×3 matrix with certain known entries:

$$J = \begin{bmatrix} p & i & q \\ j & q & \ell \\ r & m & s \end{bmatrix}.$$

Indeed, the middle element is $q(x, y) = q$, the element in the upper left corner is $q(p, p) = p$, since q is idempotent, and a similar argument proves that in the other corners we get q , r and s . It will not matter in this argument what i, j, ℓ, m are. The four 2×2 matrices in the corners of the 3×3 matrix J are still T, T -matrices.

To prove (1), we have assumed that $q \equiv_{\rho} r$. We have also assumed that $\mathbf{SR}(T, T; \rho)$ holds, so from the T, T -matrix $\begin{bmatrix} i & q \\ q & \ell \end{bmatrix}$ in the upper right corner of J we derive that $q \equiv_{\rho} \ell$. From the matrix $\begin{bmatrix} j & q \\ r & m \end{bmatrix}$ in the lower left corner of J (and the assumption that $q \equiv_{\rho} r$) we get that $r \equiv_{\rho} m$. Thus $\ell \equiv_{\rho} q \equiv_{\rho} r \equiv_{\rho} m$, so since the cross diagonal entries of the T, T -matrix $\begin{bmatrix} q & \ell \\ m & s \end{bmatrix}$ in the lower right corner are ρ -related we deduce that $m \equiv_{\rho} s$. Thus $r \equiv_{\rho} m \equiv_{\rho} s$, and (1) is proved.

To prove (2), we have assumed that $p \equiv_{\rho} q$. We have also assumed that $\mathbf{S}(T, T; \rho)$ holds, and therefore $\mathbf{RS}(T, T; \rho)$ and $\mathbf{C}(T, T; \rho)$ both hold. From the T, T -matrix $\begin{bmatrix} i & q \\ q & \ell \end{bmatrix}$ in the upper right corner of J we get by strong rectangularity that $q \equiv_{\rho} \ell$. From the T, T -matrix $\begin{bmatrix} i & p \\ q & j \end{bmatrix}$, which is obtained from the T, T -matrix in the upper left corner of J by switching columns, we get (from $p \equiv_{\rho} q$) that $q \equiv_{\rho} j$. Now that we have $j \equiv_{\rho} q$ and $q \equiv_{\rho} \ell$, we apply $\mathbf{C}(T, T; \rho)$ to the matrices

$$\begin{bmatrix} j & q \\ r & m \end{bmatrix} \quad \text{and} \quad \begin{bmatrix} q & \ell \\ m & s \end{bmatrix}$$

from the bottom two corners of J to obtain $r \equiv_{\rho} m$ and $m \equiv_{\rho} s$. Therefore $r \equiv_{\rho} m \equiv_{\rho} s$, and (2) is proved. \square

Our next task is to describe how to construct the \mathbf{Q} -commutator $[S, T]_{\mathbf{Q}}$ of two tolerances where $\mathbf{Q} = \mathbf{S}$ or \mathbf{SR} . Fix the choice of $\mathbf{Q} \in \{\mathbf{S}, \mathbf{SR}\}$. Let \mathbf{A} be an algebra, and let S and T be tolerances of \mathbf{A} . Inductively define binary relations τ_n on \mathbf{A} as follows. Let τ_0 be the equality relation. If τ_n is defined, let τ_{n+1} be the symmetric and transitive closure of the set of all pairs that can be obtained in one of the following two ways. For each matrix

$$\begin{bmatrix} p & q \\ r & s \end{bmatrix} \in M(S, T),$$

- (1) if $\mathbf{Q} = \mathbf{S}$ or \mathbf{SR} and $q \equiv_{\tau_n} r$, then put the pair (r, s) into τ_{n+1} ;
- (2) if $\mathbf{Q} = \mathbf{S}$ and $p \equiv_{\tau_n} q$, then also put the pair (r, s) into τ_{n+1} .

"SR" not "RS"

Let $\tau = \bigcup_{n < \omega} \tau_n$.

LEMMA 3.8. $[S, T]_{\mathbf{Q}} = \tau$.

PROOF. The set of S, T -matrices is closed under the component-wise application of unary polynomials of \mathbf{A} , so inductively the same property holds for each of the τ_n 's. Since τ_n is a symmetric, transitive relation and is closed under the application of unary polynomials, it is a congruence. Since

$$\begin{bmatrix} p & q \\ p & q \end{bmatrix} \in M(S, T),$$

it follows that, if $(p, q) \in \tau_n$, then $(p, q) \in \tau_{n+1}$, by item (1) of the definition of τ_{n+1} above. Therefore $\tau_n \leq \tau_{n+1}$ for every n , which implies that the union $\tau = \bigcup_{n < \omega} \tau_n$ of this ascending chain of congruences on \mathbf{A} is also a congruence on \mathbf{A} .

Let $\delta = [S, T]_{\mathbf{Q}}$. We argue by induction that $\tau_n \leq \delta$ for all n . This is trivially true if $n = 0$, so suppose that $\tau_n \leq \delta$ for some fixed n and that (r, s) is one of the generators of τ_{n+1} . If $(r, s) \in \tau_{n+1}$ because of reason (1) preceding the statement of this lemma, then there is a matrix

$$\begin{bmatrix} p & q \\ r & s \end{bmatrix} \in M(S, T),$$

such that $q \equiv_{\tau_n} r$. We have $q \equiv_{\delta} r$ (since $\tau_n \leq \delta$) and also $\mathbf{Q}(S, T; \delta)$ (hence $\mathbf{SR}(S, T; \delta)$) since $\delta = [S, T]_{\mathbf{Q}}$. Therefore we have $(r, s) \in \delta$. If $(r, s) \in \tau_{n+1}$ because of reason (2), then we must have $\mathbf{Q} = \mathbf{S}$ and so $\mathbf{C}(S, T; \delta)$ holds. Now a similar argument shows that $(r, s) \in \delta$. This shows that δ contains all the generators of τ_{n+1} , so $\tau_{n+1} \leq \delta$. By induction $\tau_n \leq \delta$ for all n , so $\tau \leq \delta = [S, T]_{\mathbf{Q}}$.

To finish the proof we must verify that $[S, T]_{\mathbf{Q}} \leq \tau$, and for this it suffices to show that $\mathbf{Q}(S, T; \tau)$ holds. Choose any S, T -matrix

$$\begin{bmatrix} p & q \\ r & s \end{bmatrix},$$

and assume that $q \equiv_{\tau} r$. Then $q \equiv_{\tau_n} r$ for some n , so from item (1) preceding the lemma we have $(r, s) \in \tau_{n+1} \leq \tau$. Thus $\mathbf{SR}(S, T; \tau)$ holds, and $[S, T]_{\mathbf{Q}} = \tau$ when $\mathbf{Q} = \mathbf{SR}$. If $\mathbf{Q} = \mathbf{S}$, then a similar argument using (2) in place of (1) shows that $\mathbf{C}(S, T; \tau)$ holds, so $\mathbf{S}(S, T; \tau)$ holds (and so $[S, T]_{\mathbf{Q}} = \tau$ when $\mathbf{Q} = \mathbf{S}$). \square

Now we are in position to prove that the four conditions of Lemma 3.5 are equivalent.

LEMMA 3.9. *Let \mathcal{V} be any variety, and let $\mathbf{Q} = \mathbf{S}$ or \mathbf{SR} . The following are equivalent.*

It doesn't follow that if (p, q) is in τ_n then this matrix is in $M(S, T)$ and so one can't conclude that (p, q) is in τ_{n+1} from this. I think that you need to argue that the generators of τ_n all lie in τ_{n+1} instead.

- (1) *Some algebra in \mathcal{V} has a nonzero \mathbf{Q} -abelian congruence.*
- (2) *Some algebra in \mathcal{V} has a nonzero \mathbf{Q} -abelian tolerance.*

PROOF. Denote by \mathbf{F} the free \mathcal{V} -algebra generated by $\{x, y\}$, by T the tolerance of \mathbf{F} generated by (x, y) , and by θ the congruence of \mathbf{F} generated by (x, y) . Let $B = x/\theta$. According to Lemma 3.5, we must establish that $(x, y) \notin [\theta, \theta]_{\mathbf{Q}}$ if and only if $(x, y) \notin [T, T]_{\mathbf{Q}}$. Since $x, y \in B$, this will follow if we show that the restrictions of the relevant congruences to B are equal, i.e., if we show that,

$$[\theta, \theta]_{\mathbf{Q}}|_B = [T, T]_{\mathbf{Q}}|_B.$$

We proceed to do this.

Since $T \subseteq \theta$ we have $[T, T]_{\mathbf{Q}}|_B \subseteq [\theta, \theta]_{\mathbf{Q}}|_B$, so we only need to verify the reverse inclusion. Let $\rho = [T, T]_{\mathbf{Q}}$. Build $\tau = [\theta, \theta]_{\mathbf{Q}}$ according to the recipe described preceding Lemma 3.8. We argue by induction that $\tau_n|_B \subseteq \rho$. This is clear for $n = 0$. Suppose that $n \geq 0$ and that $\tau_n|_B \subseteq \rho$. Choose any $(u, v) \in \tau_{n+1}|_B$. By the definition of τ_{n+1} , there exists a chain $u = u_1, u_2, \dots, u_k = v$ such that for every $1 \leq i < k$ the pair (u_i, u_{i+1}) or the pair (u_{i+1}, u_i) is produced by rule (1) or (2) of the definition of τ_{n+1} . That is, there exists a θ, θ -matrix

$$\begin{bmatrix} p & q \\ r & s \end{bmatrix}$$

such that $\{u_i, u_{i+1}\} = \{r, s\}$, and either (q, r) or (p, q) is contained in τ_n . We want to show that $u_i \equiv_{\rho} u_{i+1}$.

As $u \in B$ and $\tau \leq \theta$, every u_i belongs to B . Hence the entries of the θ, θ -matrix above are in B , too. In the first case, when $(q, r) \in \tau_n$ we therefore have that $(q, r) \in \rho$ by the induction assumption. Then (1) of Lemma 3.7 shows that $(r, s) \in \rho$ as desired. In the second case, when $(p, q) \in \tau_n$ (which only arises when $\mathbf{Q} = \mathbf{S}$) we use (2) of Lemma 3.7 to reach the same conclusion. Thus the induction step, and hence the proof of the lemma is complete. \square

The next task we set for ourselves is to show that a variety has no member with a strongly abelian tolerance if and only if the variety has a Taylor term. To prove this we need a description of the T, T -matrices where T is the tolerance on $\mathbf{F}_{\mathcal{V}}(x, y)$ generated by (x, y) .

LEMMA 3.10. *Let \mathcal{V} be a variety, $\mathbf{F} = \mathbf{F}_{\mathcal{V}}(x, y)$ be the free \mathcal{V} -algebra generated by $\{x, y\}$, and let T be the tolerance of \mathbf{F} generated by (x, y) . The T, T -matrices in \mathbf{F} are exactly the matrices of the form*

$$\begin{bmatrix} f(x, y, x, y, x, y) & f(x, x, x, y, y, y) \\ f(y, y, x, x, x, y) & f(y, x, x, x, y, y) \end{bmatrix},$$

where f is a sixary term.

PROOF. A typical T, T -matrix has the form

$$\begin{bmatrix} g(\mathbf{a}, \mathbf{c}) & g(\mathbf{a}, \mathbf{d}) \\ g(\mathbf{b}, \mathbf{c}) & g(\mathbf{b}, \mathbf{d}) \end{bmatrix}$$

where $(a_i, b_i) \in T$, $(c_j, d_j) \in T$, and (we may assume) g is a term. Since T is the subuniverse of F^2 generated by (x, y) , (y, x) , (x, x) and (y, y) , a typical pair in T has the form

$$r((x, y), (y, x), (x, x), (y, y)) = (r(x, y, x, y), r(y, x, x, y))$$

for some 4-ary term r . Therefore a typical T, T -matrix may be written as

$$\begin{bmatrix} g(\mathbf{r}(x, y, x, y), \mathbf{s}(x, y, x, y)) & g(\mathbf{r}(x, y, x, y), \mathbf{s}(y, x, x, y)) \\ g(\mathbf{r}(y, x, x, y), \mathbf{s}(x, y, x, y)) & g(\mathbf{r}(y, x, x, y), \mathbf{s}(y, x, x, y)) \end{bmatrix}.$$

If we define $f(x_1, \dots, x_6) := g(\mathbf{r}(x_1, x_4, x_3, x_6), \mathbf{s}(x_5, x_2, x_3, x_6))$, then this matrix is the one described in the lemma. \square

LEMMA 3.11. *If no algebra in the variety \mathcal{V} has a nonzero strongly abelian tolerance, then \mathcal{V} satisfies a nontrivial idempotent Maltsev condition.*

PROOF. Let \mathbf{F} be the free \mathcal{V} -algebra generated by $\{w, z\}$, let T be the tolerance of \mathbf{F} generated by (w, z) , and let B be the T -block $w/\text{Cg}^{\mathbf{F}}(w, z)$. For each element $u \in F$ choose and fix a binary term $u(x, y)$ such that $u(w, z) = u$. Build $\tau = \bigcup_{n < \omega} \tau_n = [T, T]_{\mathbf{F}}$ using the method described before Lemma 3.8. For each pair $(u, v) \in \tau$ we shall define a finite set of equations $\Sigma(u, v)$. If $(u, v) \in \tau_0$, let $\Sigma(u, v) = \{u(x, y) \approx v(x, y)\}$. To define $\Sigma(u, v)$ when $(u, v) \in \tau_{n+1} - \tau_n$, choose and fix a finite sequence $u = u_1, u_2, \dots, u_{k+1} = v$ of elements of F such that for each i either (u_i, u_{i+1}) or (u_{i+1}, u_i) is a generating pair for τ_{n+1} . Generating pairs for τ_{n+1} come from T, T -matrices, and we know from Lemma 3.10 that the T, T -matrices in \mathbf{F} are exactly the matrices of the form

$$\begin{bmatrix} p & q \\ r & s \end{bmatrix} = \begin{bmatrix} f(w, z, w, z, w, z) & f(w, w, w, z, z, z) \\ f(z, z, w, w, w, z) & f(z, w, w, w, z, z) \end{bmatrix}.$$

Therefore for each i between 1 and k there is a sixary term f_i such that

$$(3.2) \quad \{u_i, u_{i+1}\} = \{r_i, s_i\} = \{f_i(z, z, w, w, w, z), f_i(z, w, w, w, z, z)\}$$

where one of the following ‘‘side conditions’’ holds:

$$(3.3) \quad (f_i(w, w, w, z, z, z), f_i(z, z, w, w, w, z)) = (q_i, r_i) \in \tau_n, \quad \text{or}$$

$$(3.4) \quad (f_i(w, z, w, z, w, z), f_i(w, w, w, z, z, z)) = (p_i, q_i) \in \tau_n.$$

Let $(r_i^\circ, s_i^\circ) := (r_i, s_i)$ if $(r_i, s_i) = (u_i, u_{i+1})$ and let $(r_i^\circ, s_i^\circ) := (s_i, r_i)$ if $(r_i, s_i) = (u_{i+1}, u_i)$. In either case, $r_i^\circ = u_i$ and $s_i^\circ = u_{i+1}$ in \mathbf{F} .

Define $\Sigma(u, v)$ to be the set of the following identities.

- (i) $u(x, y) \approx r_1^\circ(x, y)$ and $s_k^\circ(x, y) \approx v(x, y)$.
- (ii) $s_i^\circ(x, y) \approx r_{i+1}^\circ(x, y)$ for each $1 \leq i \leq k$.
- (iii) $p_i(x, y) \approx f_i(x, y, x, y, x, y)$, $q_i(x, y) \approx f_i(x, x, x, y, y, y)$,
 $r_i(x, y) \approx f_i(y, y, x, x, x, y)$, and $s_i(x, y) \approx f_i(y, x, x, x, y, y)$.
- (iv) The identities in $\Sigma(q_i, r_i)$ for each i whose side condition is (3.3), along with $q_i(x, y) \approx f_i(x, x, x, y, y, y)$ and $r_i(x, y) \approx f_i(y, y, x, x, x, y)$.
- (v) The identities in $\Sigma(p_i, q_i)$ for each i whose side condition is (3.4), along with $p_i(x, y) \approx f_i(x, y, x, y, x, y)$ and $q_i(x, y) \approx f_i(x, x, x, y, y, y)$.

The identities in (i) and (ii) express that the pairs $(r_i^\circ(w, z), s_i^\circ(w, z))$ form a directed chain connecting u to v , the identities in (iii) express that the pair $(r_i^\circ(w, z), s_i^\circ(w, z))$ is the bottom row of a T, T -matrix, while the identities of type (iv) and (v) express that the relevant T, T -matrices satisfy side conditions sufficient to guarantee that $(r_i(w, z), s_i(w, z))$ is a generator of τ_{n+1} . These identities hold in \mathcal{V} , since they become equalities if we substitute the free generators w and z of \mathbf{F} for the variables x and y .

Induction on n may be used to prove that $\Sigma(u, v)$ is a finite set of identities for every $(u, v) \in \tau_{n+1} - \tau_n$, and that if $u \in B$ then all terms appearing in the identities in $\Sigma(u, v)$ are idempotent. Moreover, induction on n also shows that if $(u, v) \in \tau$, then for any pair (a, b) from any tolerance S on an algebra $\mathbf{A} \in \mathcal{V}$ the pair $(u(a, b), v(a, b))$ belongs to the strong commutator $[S, S]_{\mathbf{S}}$.

We have assumed that no algebra in \mathcal{V} has a nontrivial strongly abelian tolerance, so by Lemma 3.5 we have $(w, z) \in [T, T]_{\mathbf{S}} = \bigcup_{n < \omega} \tau_n$, or $(w, z) \in \tau_{n+1} - \tau_n$ for some n . Therefore \mathcal{V} satisfies the finite set of idempotent identities in $\Sigma(w, z)$. These identities constitute an idempotent Maltsev condition which, when satisfied by a variety \mathcal{W} , forces the following property: for any pair (c, d) from any tolerance R on an algebra $\mathbf{B} \in \mathcal{W}$ the pair $(w(c, d), z(c, d)) = (c, d)$ belongs to the strong commutator $[R, R]_{\mathbf{S}}$. Using the monotonicity of the strong commutator (Theorem 3.4 (1)), this is equivalent to the property that $(c, d) \in [\text{Tg}^{\mathbf{B}}(c, d), \text{Tg}^{\mathbf{B}}(c, d)]_{\mathbf{S}}$. Using Lemma 3.5 again we derive that if this Maltsev condition holds in \mathcal{W} , then no algebra in \mathcal{W} has a nontrivial strongly abelian tolerance. Since every tolerance on every member of the variety of sets is strongly abelian, this idempotent Maltsev condition fails in the variety of sets, hence is nontrivial. \square

THEOREM 3.12. *Let \mathcal{V} be a variety. The following are equivalent.*

- (1) \mathcal{V} has no member with a nonzero strongly abelian congruence.
- (2) \mathcal{V} has no member with a nonzero strongly abelian tolerance.
- (3) \mathcal{V} satisfies a nontrivial idempotent Maltsev condition.

PROOF. We proved in Lemma 3.9 that items (1) and (2) are equivalent, and in Lemma 3.11 that (2) \implies (3). It remains to prove that (3) \implies (1).

Assume that (3) holds, and let $f(x_1, \dots, x_n)$ be a Taylor term for \mathcal{V} . Arbitrarily choose and fix an algebra $\mathbf{A} \in \mathcal{V}$, a strongly abelian congruence $\alpha \in \text{Con}(\mathbf{A})$, and a pair $(a, b) \in \alpha$. Let $f_i(x, \mathbf{y})$ be an n -ary term obtained from f by reordering the variables so that x_i is first. (That is, $f(x_1, \dots, x_n) = f_i(x_i, \mathbf{y})$ where \mathbf{y} is an ordering of $\{x_1, \dots, x_n\} \setminus \{x_i\}$.) The i -th Taylor identity implies that $f_i(a, \mathbf{u}) = f_i(b, \mathbf{v})$ for some $(n-1)$ -tuples \mathbf{u} and \mathbf{v} consisting of a 's and b 's. Therefore

$$\begin{bmatrix} p & q \\ r & s \end{bmatrix} = \begin{bmatrix} f_i(a, \mathbf{v}) & f_i(b, \mathbf{v}) \\ f_i(a, \mathbf{u}) & f_i(b, \mathbf{u}) \end{bmatrix} \in M(\alpha, \alpha)$$

is an α, α -matrix with $q = r$. Since α is strongly abelian it is strongly rectangular, so $f_i(a, \mathbf{u}) = r = s = f_i(b, \mathbf{u})$. Now choose $\mathbf{w} \in \{a, b\}^{n-1}$ arbitrarily. The matrix

$$\begin{bmatrix} P & Q \\ R & S \end{bmatrix} = \begin{bmatrix} f_i(a, \mathbf{u}) & f_i(b, \mathbf{u}) \\ f_i(a, \mathbf{w}) & f_i(b, \mathbf{w}) \end{bmatrix} \in M(\alpha, \alpha)$$

is an α, α -matrix with $P = Q$. Since α is strongly abelian it is abelian, so $f_i(a, \mathbf{w}) = R = S = f_i(b, \mathbf{w})$.

What this says about the original Taylor term is that, for any $1 \leq i \leq n$, if $\mathbf{c}, \mathbf{d} \in \{a, b\}^n$ are tuples that agree in every coordinate except possibly the i -th, then $f(\mathbf{c}) = f(\mathbf{d})$. This statement is true for all i , so

$$\begin{aligned} a &= f(a, a, a, \dots, a) \\ &= f(b, a, a, \dots, a) \\ &= f(b, b, a, \dots, a) \\ &\quad \vdots \\ &= f(b, b, b, \dots, b) = b, \end{aligned}$$

forcing $a = b$. Since (a, b) was chosen arbitrarily, we conclude that no $\mathbf{A} \in \mathcal{V}$ has a nontrivial strongly abelian congruence. \square

3.2. Join Terms

In the previous section we showed that a variety omits strongly abelian congruences or tolerances if and only if it satisfies a nontrivial idempotent Maltsev condition (Theorem 3.12). In this section we will

"that"

prove the the varieties with these properties are exactly the varieties that omit strongly rectangular congruences or tolerances. The obvious approach to this result is to note that if a variety omits strongly rectangular congruences or tolerances, then it must omit strongly abelian congruences or tolerances, and therefore must satisfy a nontrivial idempotent Maltsev condition. Thus, the only thing to show is that it is possible to modify the proof of Theorem 3.12 (3) \implies (1) so that it only uses the assumption of strong rectangulation.

It must be possible to obtain the result in this way, but in this section we choose another path. We instead develop machinery for dealing with stongly rectangular congruences and tolerances that makes the proof of the desired result a triviality. Although the path we choose does not lead to the result more quickly or more simply than the path sketched in the previous paragraph, we feel that developing the machinery for strong rectangulation is more important than the target result.

DEFINITION 3.13. Let \mathbf{A} be an algebra, let T be a tolerance on \mathbf{A} and let δ be a congruence on \mathbf{A} . A (T, T) -**triple** is a triple $(b, c; d)$ such that there is a matrix

$$\begin{bmatrix} a & b \\ c & d \end{bmatrix} \in M(T, T).$$

A $(T, T; \delta)$ -**pair** is a pair $(c, d) \in A^2$ such that there is a (T, T) -triple $(b, c; d)$ with $b \equiv_{\delta} c$.

Observe that the set of all (T, T) -triples is a subuniverse of \mathbf{A}^3 , and the set of all $(T, T; \delta)$ -pairs is a subuniverse of \mathbf{A}^2 .

DEFINITION 3.14. Let T be a tolerance on an algebra \mathbf{A} . For any block B of T write $a \geq_B b$ (or $b \leq_B a$) to mean that there are $(T, T; 0)$ -pairs $(u_1, u_2), \dots, (u_{n-1}, u_n)$ such that $a = u_1$, $b = u_n$, and $u_i \in B$ for all i . Write $a \sim_B b$ to mean $a \geq_B b$ and $a \leq_B b$.

Write $a \geq_T b$ (or $b \leq_T a$) if there is some T -block B for which $a \geq_B b$. Similarly, write $a \sim_T b$ if $\exists B(a \sim_B b)$.

There is a potential ambiguity in the notation \geq_B , because a subset B may be a block of more than one tolerance, but this issue will not arise in this monograph.

LEMMA 3.15. *Let T be a tolerance on an algebra \mathbf{A} . The following hold.*

- (1) *If B is a block of T , then \geq_B is a quasiorder on B , and \sim_B is the induced equivalence relation on B . Both relations are compatible with the idempotent polynomials of \mathbf{A} .*

- (2) The relation \geq_T is a reflexive compatible relation of \mathbf{A} , and \sim_T is a tolerance on \mathbf{A} .
- (3) If $\mathbf{SR}(T, T; 0)$ holds, then \geq_T is the equality relation on \mathbf{A} .
- (4) If B is a block of T and is also a block of $\text{Cg}^{\mathbf{A}}(T)$, then \geq_B equals the restriction to B of \geq_T and also equals the restriction to B of the transitive closure of the set of $(T, T; 0)$ -pairs.
- (5) If T is a congruence, then \geq_T is a quasiorder of \mathbf{A} , which is the transitive closure of the set of $(T, T; 0)$ -pairs. The relation \sim_T is the equivalence relation induced by \geq_T , which is a congruence on \mathbf{A} contained in T .

PROOF. For this proof, let S denote the set of $(T, T; 0)$ -pairs of \mathbf{A} .

For item (1), the matrix $\begin{bmatrix} u & u \\ u & u \end{bmatrix}$ belongs to $M(T, T)$ for any $u \in A$, so the trivial pair (u, u) is a $(T, T; 0)$ -pair for any $u \in A$. This is enough to show that \geq_B is reflexive on B . It is clear from Definition 3.14 that \geq_B equals $\text{tr.cl.}(S|_B)$, hence \geq_B is transitive as well. This proves that \geq_B is a quasiorder on B ; it follows from Definition 3.14 that \sim_B is the induced equivalence relation. If $p(x_1, \dots, x_n)$ is an idempotent polynomial of \mathbf{A} , then $p(B, \dots, B)$ is a subset of A containing B and consisting of pairwise T -related elements. Since B is a block, $p(B, \dots, B) = B$. Since S is compatible with all polynomials of \mathbf{A} , the fact that $p(B, \dots, B) = B$ implies that $S|_B$ is compatible with p . Therefore $\geq_B = \text{tr.cl.}(S|_B)$ is also compatible with p , as is the intersection \sim_B of this relation with its converse.

Now we prove item (2). Since \geq_T equals the union over all blocks B of the relations \geq_B , it follows that \geq_T is reflexive. Similarly, \sim_T is reflexive and symmetric. It is only necessary to prove that \geq_T and \sim_T are compatible relations. We will prove the following slightly stronger fact.

CLAIM 3.16. Let $f(x_1, \dots, x_n)$ be an n -ary basic operation of \mathbf{A} . Assume that B_1, \dots, B_n, B are T -blocks such that $f(B_1, \dots, B_n) \subseteq B$. If $a_i \geq_{B_i} b_i$ holds for all i , then

$$f(a_1, \dots, a_n) \geq_B f(b_1, \dots, b_n).$$

The same statement is true with \sim in place of \geq .

The arguments are essentially the same for \geq and \sim , so we prove the claim only for \geq . By assumption $a_i \geq_{B_i} b_i$, so there are $(T, T; 0)$ -pairs $(u_i^1, u_i^2), \dots, (u_i^{n-1}, u_i^n)$ such that $a_i = u_i^1$, $b_i = u_i^n$, and all u_i^j lie in B_i . There is no a priori reason to expect that the number of such pairs is the same for all i , but this can be arranged to be so by adding

this should be something other than n , the arity of f .

m

trivial pairs of the form (u, u) to the end of any sequence of pairs that is shorter than the longest sequence. For each i , the pairs (u_i^j, u_i^{j+1}) define a chain of n elements in B_i connecting a_i to b_i . We write these chains as rows of an array, and apply f to the columns of the array.

n

$$\begin{array}{cccccc}
 a_1 & = & u_1^1 & & u_1^2 & & u_1^3 & & \cdots & & u_1^n & = & b_1 \\
 a_2 & = & u_2^1 & & u_2^2 & & u_2^3 & & \cdots & & u_2^n & = & b_2 \\
 & & & & & & \vdots & & & & & & \\
 a_m & = & u_m^1 & & u_m^2 & & u_m^3 & & \cdots & & u_m^n & = & b_m \\
 \hline
 f(\mathbf{a}) & = & f(\mathbf{u}^1) & & f(\mathbf{u}^2) & & f(\mathbf{u}^3) & & \cdots & & f(\mathbf{u}^n) & = & f(\mathbf{b})
 \end{array}
 \tag{3.5}$$

Each pair $(f(\mathbf{u}^j), f(\mathbf{u}^{j+1}))$ is a $(T, T; 0)$ -pair, since the set S of $(T, T; 0)$ -pairs is a subalgebra of \mathbf{A}^2 . The block B contains $f(\mathbf{u}^j)$ for all j , so the bottom row of (3.5) is a sequence that witnesses the fact that $f(\mathbf{a}) \geq_B f(\mathbf{b})$.

The compatibility of \geq_T follows immediately from this claim. For example, if $a_i \geq_T b_i$ for $i = 1, \dots, n$, then there exist T -blocks B_i such that $a_i \geq_{B_i} b_i$ for all i . If f is an n -ary basic operation of \mathbf{A} , then the set $f(B_1, \dots, B_n)$ consists of T -related elements, hence can be extended to a block B . The claim now applies to prove that $f(\mathbf{a}) \geq_B f(\mathbf{b})$, hence that $f(\mathbf{a}) \geq_T f(\mathbf{b})$.

Item (3) follows from the fact that $\mathbf{SR}(T, T; 0)$ holds iff all $(T, T; 0)$ -pairs are trivial.

For item (4), for any T -block B the transitive closure of $S|_B$ equals \geq_B , which is contained in \geq_T , which is contained in the transitive closure of S . If we restrict these relations to B we get

$$\text{tr.cl.}(S|_B) = (\geq_B) \subseteq (\geq_T)|_B \subseteq (\text{tr.cl.}(S))|_B.
 \tag{3.6}$$

But since B is also assumed to be a block of $\text{Cg}^{\mathbf{A}}(T)$ in item (4), we get that $\text{tr.cl.}(S|_B) = (\text{tr.cl.}(S))|_B$, so all of the relations in (3.6) are equal.

For item (5), the fact that T is an equivalence relation containing \geq_T , together with the fact proved in item (4) that \geq_T is transitive on any block of $T = \text{Cg}^{\mathbf{A}}(T)$, implies that \geq_T is transitive on \mathbf{A} . Since T contains the set S of all $(T, T; 0)$ -pairs and the relation \geq_T agrees with $\text{tr.cl.}(S)$ on each block of T , it follows that \geq_T agrees with $\text{tr.cl.}(S)$ on \mathbf{A} . Since each element of \mathbf{A} is contained in a unique block of the congruence T , it follows that \sim_T agrees with \sim_B on any block of T . Thus $a \sim_B b$ holds if and only if $a \geq_T b$ and $a \leq_T b$, which implies that \sim_T is the equivalence relation induced by \geq_T . It is a congruence by item (2), and is contained in T because all $(T, T; 0)$ -pairs are contained in T . \square

LEMMA 3.17. *Let T be a tolerance on an algebra \mathbf{A} . The following hold.*

- (1) *If $(b, c) \in T$, then $(b, c; b)$ and $(b, c; c)$ are (T, T) -triples.*
- (2) *Let p be a k -ary polynomial, and suppose that $(b_i, c_i; d_i)$ are (T, T) -triples for $1 \leq i \leq k$. If $p(b_1, \dots, b_n) = p(c_1, \dots, c_n)$, then $(p(c_1, \dots, c_n), p(d_1, \dots, d_n))$ is a $(T, T; 0)$ -pair.*
- (3) *Let f be a binary idempotent polynomial of \mathbf{A} , and assume*

$$\begin{bmatrix} a & b \\ c & d \end{bmatrix} \in M(T, T).$$

If B is a block of T containing a, b, c and d , then

$$f(b, c) \geq_B f(f(b, d), f(a, c)).$$

- (4) *Let q be an n -ary idempotent polynomial of \mathbf{A} . If B is a block of T containing elements $a_1, \dots, a_{n-1}, b_1, \dots, b_{n-1}$ and c , then*

$$q(\mathbf{a}, q(\mathbf{b}, c)) \geq_B q(\mathbf{a}, q(\mathbf{a}, c)).$$

- (5) *Let F be an n -ary idempotent polynomial of \mathbf{A} . Let B be a block of T containing elements u_{ij} ($1 \leq i, j \leq n$) and w , and let $\mathbf{u}^i = (u_{i1}, \dots, u_{in})$. If $u_{ii} \geq_B w$ for $1 \leq i \leq n$, then*

$$F(F(\mathbf{u}^1), F(\mathbf{u}^2), \dots, F(\mathbf{u}^n)) \geq_B w.$$

PROOF. Item (1) follows from the fact that

$$\begin{bmatrix} c & b \\ c & b \end{bmatrix} \quad \text{and} \quad \begin{bmatrix} b & b \\ c & c \end{bmatrix}$$

are (T, T) -matrices if $(b, c) \in T$.

To prove (2) suppose that $(b_i, c_i; d_i)$ are (T, T) -triples for $1 \leq i \leq k$. They come from some (T, T) -matrices

$$\begin{bmatrix} a_i & b_i \\ c_i & d_i \end{bmatrix}.$$

Apply p to these matrices componentwise. The resulting matrix is a T, T -matrix whose off-diagonal entries are $p(b_1, \dots, b_n)$ and $p(c_1, \dots, c_n)$. If these are equal, as is assumed in (2), then $(p(c_1, \dots, c_n), p(d_1, \dots, d_n))$ is a $(T, T; 0)$ -pair by definition.

To prove (3), we will apply item (2) to the polynomial

$$p(x_1, x_2, x_3, x_4) = f(f(x_1, x_2), f(x_3, x_4)).$$

The (T, T) -triples that we will use are $(b, b; b)$, $(b, c; d)$, $(c, b; a)$ and $(c, c; c)$. The first and last of these triples are (T, T) -triples by item (1),

the second is a (T, T) -triple since we have assumed that

$$\begin{bmatrix} a & b \\ c & d \end{bmatrix} \in M(T, T),$$

and the third is a (T, T) -triple since we can interchange the rows and columns of this matrix to obtain that

$$\begin{bmatrix} d & c \\ b & a \end{bmatrix} \in M(T, T).$$

Item (2) guarantees that since

$$p(b, b, c, c) = f(f(b, b), f(c, c)) = f(f(b, c), f(b, c)) = p(b, c, b, c)$$

we have that $(p(b, c, b, c), p(b, d, a, c))$ is a $(T, T; 0)$ -pair. But B is closed under f , since f is idempotent and B is a tolerance block, and $(f(b, c), f(f(b, d), f(a, c))) = (p(b, c, b, c), p(b, d, a, c))$, so

$$f(b, c) \geq_B f(f(b, d), f(a, c)),$$

as claimed.

For item (4), let $t = q(\mathbf{a}, q(\mathbf{b}, c))$ and write \mathbf{t} for the $(n-1)$ -tuple (t, t, \dots, t) . Since q is idempotent and all a_i, b_j and c belong to B , we get that both t and $q(\mathbf{b}, c)$ belong to B . Now apply item (2) using $p(\mathbf{x}, \mathbf{y}, z) = q(\mathbf{x}, q(\mathbf{y}, z))$ and the $2n-1$ (T, T) -triples $(a_1, t; a_1), \dots, (a_{n-1}, t; a_{n-1}), (b_1, a_1; a_1), \dots, (b_{n-1}, a_{n-1}; a_{n-1}), (c, q(\mathbf{b}, c); c)$ (which are (T, T) -triples, by item (1)). Item (2) guarantees that since

$$\begin{aligned} p(\mathbf{a}, \mathbf{b}, c) &= q(\mathbf{a}, q(\mathbf{b}, c)) \\ &= t \\ &= q(\mathbf{t}, t) \\ &= q(\mathbf{t}, q(\mathbf{a}, q(\mathbf{b}, c))) \\ &= p(\mathbf{t}, \mathbf{a}, q(\mathbf{b}, c)) \end{aligned}$$

extra ")" here.

the pair $(q(\mathbf{a}, q(\mathbf{b}, c)), q(\mathbf{a}, q(\mathbf{a}, c))) = (p(\mathbf{t}, \mathbf{a}, q(\mathbf{b}, c)), p(\mathbf{a}, \mathbf{a}, c))$ is a $(T, T; 0)$ -pair. Since these elements belong to B we have

$$q(\mathbf{a}, q(\mathbf{b}, c)) \geq_B q(\mathbf{a}, q(\mathbf{a}, c)).$$

This proves (4).

Finally we prove (5). We show that

$$F(F(\mathbf{u}^1), \dots, F(\mathbf{u}^i), w, w, \dots, w) \geq_B w$$

by induction on i . This statement follows from the idempotence of F when $i = 0$, and for $i = n$ yields the desired statement. All elements that we shall consider will be in B , because F is idempotent and B is a T -block.

As an inductive assumption suppose that

$$F(F(\mathbf{u}^1), \dots, F(\mathbf{u}^{i-1}), w, w, \dots, w) \geq_B w$$

holds. Now apply (4) in the i -th variable. By this we mean cyclically permute the last $n - i + 1$ variables of F to obtain an n -ary idempotent polynomial q whose last variable corresponds to the i -th variable of F :

$$q(x_1, \dots, x_n) := F(x_1, \dots, x_{i-1}, x_n, x_i, \dots, x_{n-1}),$$

then apply (4) to q . Choose $(n - 1)$ -tuples

$$\begin{aligned} \mathbf{a} &= (F(\mathbf{u}^1), \dots, F(\mathbf{u}^{i-1}), w, \dots, w), \\ \mathbf{b} &= (u_{i1}, \dots, u_{i(i-1)}, u_{i(i+1)}, \dots, u_{in}), \end{aligned}$$

and choose $c = u_{ii}$. With these choices $q(\mathbf{b}, c) = F(\mathbf{u}^i)$. The conclusion of (4) informs us that

$$q(\mathbf{a}, q(\mathbf{b}, c)) = F(F(\mathbf{u}^1), \dots, F(\mathbf{u}^{i-1}), F(\mathbf{u}^i), w, \dots, w)$$

is $\geq_B q(\mathbf{a}, q(\mathbf{a}, c))$, which equals

$$F(F(\mathbf{u}^1), \dots, F(\mathbf{u}^{i-1}), F(F(\mathbf{u}^1), \dots, F(\mathbf{u}^{i-1}), \underline{u_{ii}}, w, \dots, w), w, \dots, w).$$

Using the assumption that $u_{ii} \geq_B w$, we find that this element is \geq_B -related to the element

$$F(F(\mathbf{u}^1), \dots, F(\mathbf{u}^{i-1}), F(F(\mathbf{u}^1), \dots, F(\mathbf{u}^{i-1}), \underline{w}, w, \dots, w), w, \dots, w).$$

Using the inductive assumption twice we find that this element,

$$F(F(\mathbf{u}^1), \dots, F(\mathbf{u}^{i-1}), \underline{F(F(\mathbf{u}^1), \dots, F(\mathbf{u}^{i-1}), w, w, \dots, w)}, w, \dots, w),$$

is \geq_B -related to

$$F(F(\mathbf{u}^1), \dots, F(\mathbf{u}^{i-1}), \underline{w}, w, \dots, w) \geq_B w.$$

By the transitivity of \geq_B on B we get that

$$F(F(\mathbf{u}^1), \dots, F(\mathbf{u}^{i-1}), F(\mathbf{u}^i), w, \dots, w) \geq_B w,$$

which completes the proof of (5). \square

DEFINITION 3.18. A **join term** for a variety \mathcal{V} is an idempotent binary term $s(x, y)$ such that for every $\mathbf{A} \in \mathcal{V}$, tolerance T of \mathbf{A} , and pair $(u, v) \in T$ it is the case that $s(u, v) \geq_T u$ and $s(u, v) \geq_T v$.

In order to investigate this concept we need to understand the relation \geq_T in the generic setting.

LEMMA 3.19. *Let \mathcal{V} be a variety, let $\mathbf{F} = \mathbf{F}_{\mathcal{V}}(x, y)$ be the free \mathcal{V} -algebra generated by $\{x, y\}$, let T be the tolerance of \mathbf{F} generated by (x, y) , let $\theta = \text{Cg}^{\mathbf{F}}(x, y)$, and let $B = x/\theta$. If $u, v \in B$, then $u \geq_T v$ if and only if there exists an $m \geq 1$, and sixary terms f_1, \dots, f_m such that \mathcal{V} satisfies the following identities:*

- (i) $u(x, y) \approx f_1(y, y, x, x, x, y)$,
- (ii) $f_m(y, x, x, x, y, y) \approx v(x, y)$,
- (iii) $f_i(x, x, x, y, y, y) \approx f_i(y, y, x, x, x, y)$, $1 \leq i \leq m$
- (iv) $f_i(y, x, x, x, y, y) \approx f_{i+1}(y, y, x, x, x, y)$, $1 \leq i \leq m - 1$.

PROOF. By Theorem 3.6 (4), B is a block of the congruence θ . By Lemma 3.15 (4), the restriction of \geq_T to B equals the restriction of the transitive closure of the set of $(T, T; 0)$ -pairs to B . It follows from Lemma 3.10 that the $(T, T; 0)$ -pairs are exactly the pairs of the form $(f(y, y, x, x, x, y), f(y, x, x, x, y, y))$ for which \mathcal{V} satisfies the identity $f(x, x, x, y, y, y) \approx f(y, y, x, x, x, y)$. Therefore, the identities enumerated in the statement of this lemma express exactly that (u, v) is in the transitive closure of the set of $(T, T; 0)$ -pairs. \square

THEOREM 3.20. *The following are equivalent for a variety \mathcal{V} .*

- (1) \mathcal{V} satisfies a nontrivial idempotent Maltsev condition.
- (2) \mathcal{V} has a join term.
- (3) \mathcal{V} satisfies the following Maltsev condition: there exist $m \geq 1$, $n \geq 1$, idempotent sixary terms $f_1, \dots, f_m, g_1, \dots, g_n$, and a binary term s such that the following identities hold in \mathcal{V} :

n , not m .

- (i) $s(x, y) \approx f_1(y, y, x, x, x, y)$,
- (ii) $f_m(y, x, x, x, y, y) \approx x$,
- (iii) $f_i(x, x, x, y, y, y) \approx f_i(y, y, x, x, x, y)$, $1 \leq i \leq m$
- (iv) $f_i(y, x, x, x, y, y) \approx f_{i+1}(x, x, x, y, y, y)$, $1 \leq i \leq m - 1$.
- (v) $s(x, y) \approx g_1(y, y, x, x, x, y)$,
- (vi) $g_m(y, x, x, x, y, y) \approx y$,
- (vii) $g_i(x, x, x, y, y, y) \approx g_i(y, y, x, x, x, y)$, $1 \leq i \leq n$
- (viii) $g_i(y, x, x, x, y, y) \approx g_{i+1}(x, x, x, y, y, y)$, $1 \leq i \leq n - 1$.

these variable patterns don't match those in 3.19

PROOF. To prove (1) \implies (2) assume that a variety \mathcal{V} has a Taylor term f . The Taylor identities may be written in the form

$$f(x_{i1}, \dots, x_{in}) \approx f(y_{i1}, \dots, y_{in}),$$

where $1 \leq i \leq n$, $x_{ij}, y_{ij} \in \{x, y\}$, and $x_{ii} = x$ while $y_{ii} = y$. Consider the binary term

$$s(x, y) = f(f(x_{11}, \dots, x_{1n}), \dots, f(x_{n1}, \dots, x_{nn})).$$

The term s is idempotent, since f is. Applying Lemma 3.17 (5) to $F = f$, B a block of some tolerance T on some $\mathbf{A} \in \mathcal{V}$, and elements $u, v \in B$, shows that $s(u, v) \geq_B u$. On the other hand, the Taylor identities show that the identity

$$s(x, y) \approx f(f(y_{11}, \dots, y_{1n}), \dots, f(y_{n1}, \dots, y_{nn}))$$

holds in \mathcal{V} , and so Lemma 3.17 (5) shows that $s(u, v) \geq_B v$ holds as well. Thus, s is a join term in \mathcal{V} , and (2) is proved.

Now suppose that a variety \mathcal{V} has a join term s . Then we have $s(x, y) \geq_T x$ and $s(x, y) \geq_T y$ in the free algebra $\mathbf{F} = \mathbf{F}_{\mathcal{V}}(x, y)$ for the tolerance T generated by (x, y) . It follows from Lemma 3.19 that \mathcal{V} satisfies the idempotent Maltsev condition defined in (3). This proves that (2) \implies (3). This Maltsev condition expresses the fact that s is a join term in \mathcal{V} . The variety of sets does not have a join term, since the relation \geq_T is the equality relation in sets. Therefore this Maltsev condition is nontrivial, and (3) \implies (1). \square

According to Definitions 3.14 and 3.18, $s(x, y)$ is a join term for \mathcal{V} if it is idempotent and for every $\mathbf{A} \in \mathcal{V}$, tolerance T on \mathbf{A} , and pair $(u, v) \in T$, there exist blocks B_1 and B_2 such that $s(u, v) \geq_{B_1} u$ and $s(u, v) \geq_{B_2} v$. It is a consequence of the previous theorem that the blocks B_1 and B_2 may be chosen to be equal, as we now explain.

COROLLARY 3.21. *Let \mathcal{V} be a variety with a join term $s(x, y)$. If $\mathbf{A} \in \mathcal{V}$, T is a tolerance on \mathbf{A} , and B is any T -block containing u and v , then $s(u, v) \geq_B u$ and $s(u, v) \geq_B v$.*

PROOF. Since B is a T -block, it is closed under all idempotent term operations of \mathbf{A} , in particular it is closed under s and all the f_i and g_i from Theorem 3.20 (3). If we substitute u for x and v for y in these term operations we obtain a sequence of elements of B which, according to the equations of that theorem, witness that $s(u, v) \geq_B u$ and $s(u, v) \geq_B v$. \square

THEOREM 3.22. *Let \mathcal{V} be a variety. The following are equivalent.*

- (1) \mathcal{V} has no member with a nonzero strongly rectangular congruence.
- (2) \mathcal{V} has no member with a nonzero strongly rectangular tolerance.
- (3) \mathcal{V} satisfies a nontrivial idempotent Maltsev condition.

PROOF. As noted at the beginning of this section, the only part of this theorem that remains to be proved is (3) \implies (1).

Assume that (3) holds, and let s be a join term for \mathcal{V} . Choose an algebra $\mathbf{A} \in \mathcal{V}$, a strongly rectangular congruence α on \mathbf{A} , and a pair $(a, b) \in \alpha$. Since s is a join term we have $s(a, b) \geq_{\alpha} a$ and $s(a, b) \geq_{\alpha} b$. By Lemma 3.15 (3) the relation \geq_{α} is the equality relation when α is strongly rectangular, so $a = b$. Hence $\alpha = 0$ and (1) holds. \square

3.3. Abelian Tolerances and Congruences

In the earlier part of this chapter we saw that a variety has a Taylor term if and only if its members have no strongly abelian tolerances or congruences. Taylor terms do not conflict with the existence of nontrivial (ordinary) abelian tolerances and congruences, but do improve their behavior. In this section we will see that if a variety has a Taylor term, then abelian tolerances generate abelian congruences, and also that a congruence interval perspective with an abelian congruence interval is itself abelian.

THEOREM 3.23. *Assume that \mathcal{V} satisfies a nontrivial idempotent Maltsev condition. If $\mathbf{A} \in \mathcal{V}$ has a tolerance T and a congruence δ such that $\mathbf{C}(T, T; \delta)$ holds, then $\mathbf{C}(\text{Cg}^{\mathbf{A}}(T), \text{Cg}^{\mathbf{A}}(T); \delta)$ holds.*

PROOF. It suffices to prove this theorem in the case when $\delta = 0$, for if $\mathbf{C}(T, T; \delta)$ holds, then $\mathbf{C}(T/\delta, T/\delta; 0)$ holds in \mathbf{A}/δ , according to Theorem 2.19 (10). If the theorem is true when $\delta = 0$, then we get that $\mathbf{C}(\alpha, \alpha; 0)$ holds for $\alpha := \text{Cg}^{\mathbf{A}/\delta}(T/\delta)$. But then $\alpha = \text{Cg}^{\mathbf{A}/\delta}(\text{Cg}^{\mathbf{A}}(T)/\delta)$, too, so again using Theorem 2.19 (10) we get that $\mathbf{C}(\text{Cg}^{\mathbf{A}}(T), \text{Cg}^{\mathbf{A}}(T); \delta)$ holds, which is the statement of the theorem. We therefore assume that T is an abelian tolerance, and make it our task to prove that T generates an abelian congruence.

Write \mathbf{T} for the subalgebra of \mathbf{A}^2 whose universe is $A^2[T]$, which is just the set T (cf. the first paragraph of Chapter 2), and write Δ for the congruence on \mathbf{T} generated by the set

$$\{\langle (a, a), (b, b) \rangle \mid (a, b) \in T\}.$$

Define tolerances on \mathbf{T} by $R_1 := \{\langle (a, b), (a, c) \rangle \in T^2 \mid (b, c) \in T\}$ and let $R_2 := \{\langle (a, c), (b, c) \rangle \in T^2 \mid (a, b) \in T\}$.

CLAIM 3.24. $\Delta \cap R_i$ is the equality relation for both $i = 1$ and 2 .

It suffices to prove this for $i = 1$ only, so choose a pair $\langle (a, b), (a, c) \rangle \in \Delta \cap R_1$. We must prove that $b = c$. Let $f(x_1, \dots, x_n)$ be a Taylor term for \mathcal{V} . The i -th Taylor identity has the form

$$f(x_1, \dots, x_n) \approx f(y_1, \dots, y_n)$$

where $(x_j, y_j) \in \{(x, x), (x, y), (y, x), (y, y)\}$ for all j , and (we may assume) $(x_i, y_i) = (x, y)$. Partition N into blocks $\{B_1, \dots, B_5\}$, some of which might be empty, as follows: B_1 consists of those $j \in N$ for which $(x_j, y_j) = (x, x)$, B_2 consists of those $j \neq i$ for which $(x_j, y_j) = (x, y)$, $B_3 = \{i\}$, B_4 consists of those $j \in N$ for which $(x_j, y_j) = (y, x)$, and B_5 consists of those j for which $(x_j, y_j) = (y, y)$. Now substitute the variable z_1 for the variable x_j of $f(x_1, \dots, x_n)$ if $j \in B_1$, z_2 if $j \in B_2$,

it is clear that N is supposed to be $\{1, 2, \dots, n\}$, but it would be helpful to mention this, rather than forcing the reader to look ahead to figure this out.

and so on. This produces a 5-ary term $f_i(z_1, z_2, z_3, z_4, z_5)$. The i -th Taylor identity is now $f_i(x, x, x, y, y) = f_i(x, y, y, x, y)$, where the third argument of f_i corresponds to the i -th argument of f . Define unary polynomials of \mathbf{T} by

$$p_i((x, y)) = f_i^{\mathbf{T}}((a, a), (a, b), (x, y), (b, a), (b, b)).$$

The i -th Taylor identity ensures that $p_i((a, b))$ lies on the diagonal of \mathbf{T} . The property that T is abelian is equivalent to the property that the diagonal of \mathbf{T} is a union of Δ -classes, so, since $p_i((a, b)) \Delta p_i((a, c))$, we get that $p_i((a, c))$ lies on the diagonal of \mathbf{T} . Since $p_i((a, b))$ and $p_i((a, c))$ have the same first coordinate, and both lie on the diagonal, we get $p_i((a, b)) = p_i((a, c))$. Equating second coordinates yields

$$f_i(\underline{a}, b, b, \underline{a}, b) = f_i(\underline{a}, b, c, \underline{a}, b).$$

Since $(a, b), (b, c) \in T$ and T is abelian, we still have equality after changing the underlined a 's to b 's:

$$f_i(\underline{b}, b, b, \underline{b}, b) = f_i(\underline{b}, b, c, \underline{b}, b).$$

For the original Taylor term this means that

$$b = f(b, b, \dots, b) = f(b, \dots, b, c, b, \dots, b)$$

where only one c occurs, and it occurs on the righthand side in the i -th position. We argue by induction that, for all i , $b = f(c, c, \dots, c, b, \dots, b)$, where c occurs in the first i positions and b occurs elsewhere. For $i = 1$ this follows from the last displayed line. For later values of i we apply the T, T -term condition to change b to c in the following underlined positions:

$$f(\underline{b}, \underline{b}, \dots, \underline{b}, b, b, \dots, b) = f(\underline{b}, \underline{b}, \dots, \underline{b}, c, b, \dots, b).$$

We obtain

$$f(\underline{c}, \underline{c}, \dots, \underline{c}, b, b, \dots, b) = f(\underline{c}, \underline{c}, \dots, \underline{c}, c, b, \dots, b),$$

which shows that the desired result with c in the first $i - 1$ positions implies the result for c in the first i positions.

Continuing to the end, we finally have $b = f(c, c, \dots, c) = c$ by the idempotence of f . This proves the claim.

CLAIM 3.25. Δ is abelian.

From Claim 3.24 and Theorem 2.19 (8) we get that $\mathbf{C}(R_i, \Delta; 0)$ holds for $i = 1$ and 2. This and Theorem 2.19 (2) and (5) imply that $\mathbf{C}(\text{Cg}^{\mathbf{T}}(R_1 \cup R_2), \Delta; 0)$ holds. If $(a, b) \in T$, then $(a, a) R_1 (a, b) R_2 (b, b)$, so $\langle (a, a), (b, b) \rangle \in \text{Cg}^{\mathbf{T}}(R_1 \cup R_2)$. This shows that each generator of Δ

belongs to $\text{Cg}^{\mathbf{T}}(R_1 \cup R_2)$, so $\text{Cg}^{\mathbf{T}}(R_1 \cup R_2) \geq \Delta$. From the monotonicity of the centralizer in its first two variables we get that $\mathbf{C}(\Delta, \Delta; 0)$ holds in \mathbf{T} . This proves the claim.

Let \mathbf{D} be the image of the diagonal embedding $i: \mathbf{A} \rightarrow \mathbf{T}: a \mapsto (a, a)$. Since Δ is an abelian congruence of \mathbf{T} , it follows from Theorem 2.19 (9) that $\Delta|_{\mathbf{D}}$ is an abelian congruence of \mathbf{D} . Since $i: \mathbf{A} \rightarrow \mathbf{D}$ is an isomorphism, the congruence $i^{-1}(\Delta|_{\mathbf{D}})$ is an abelian congruence of \mathbf{A} . Since $i(T)$ is the generating set for Δ , it follows that $T \subseteq i^{-1}(\Delta|_{\mathbf{D}})$, and therefore T is contained in the abelian congruence $i^{-1}(\Delta|_{\mathbf{D}})$. Hence $\text{Cg}^{\mathbf{A}}(T)$ is abelian. \square

THEOREM 3.26. *Assume that \mathcal{V} satisfies a nontrivial idempotent Maltsev condition and that $\mathbf{A} \in \mathcal{V}$. If $\alpha, \beta \in \text{Con}(\mathbf{A})$, then*

$$\mathbf{C}(\alpha, \alpha; \alpha \wedge \beta) \iff \mathbf{C}(\alpha \vee \beta, \alpha \vee \beta; \beta).$$

PROOF. If $\mathbf{C}(\alpha \vee \beta, \alpha \vee \beta; \beta)$ holds, then $\mathbf{C}(\alpha, \alpha; \beta)$ holds, by monotonicity. From this and $\mathbf{C}(\alpha, \alpha; \alpha)$ we get that $\mathbf{C}(\alpha, \alpha; \alpha \wedge \beta)$ holds, by Theorem 2.19 (4).

Now assume that $\mathbf{C}(\alpha, \alpha; \alpha \wedge \beta)$ holds. Using Theorem 2.19 (4) again we derive that $\mathbf{C}(\alpha, \alpha; \beta)$ holds. By Theorem 2.19 (8), $\mathbf{C}(\beta, \alpha; \beta)$ holds, so we get that $\mathbf{C}(\alpha \vee \beta, \alpha; \beta)$ and then $\mathbf{C}(\alpha \vee \beta, \beta \circ \alpha \circ \beta; \beta)$ hold by parts (5) and (3) of Theorem 2.19. This and the monotonicity of the centralizer yield that $\mathbf{C}(T, T; \beta)$ holds for $T = \beta \circ \alpha \circ \beta$. Since $\text{Cg}^{\mathbf{A}}(T) = \alpha \vee \beta$ we get $\mathbf{C}(\alpha \vee \beta, \alpha \vee \beta; \beta)$ holds from Theorem 3.23. \square

REMARK 3.27. It is not possible to strengthen the implication in Theorem 3.23 to

$$(3.7) \quad \mathbf{C}(S, T; \delta) \implies \mathbf{C}(\text{Cg}^{\mathbf{A}}(S), \text{Cg}^{\mathbf{A}}(T); \delta)$$

assuming only that \mathcal{V} has a Taylor term. If implication (3.7) holds for \mathcal{V} , then arguing as in the proof of Theorem 3.26 we would have

$$\mathbf{C}(\alpha, \beta; \delta) \implies \mathbf{C}(\alpha \vee \delta, \beta \vee \delta; \delta)$$

for any $\alpha, \beta, \delta \in \text{Con}(\mathbf{A})$, $\mathbf{A} \in \mathcal{V}$. In particular, since $\mathbf{C}(\alpha, \beta; \delta)$ holds whenever α, β and δ generate a sublattice isomorphic to \mathbf{N}_5 , labeled as in Figure 3.1, we would also have $\mathbf{C}(\alpha \vee \delta, \beta \vee \delta; \delta)$, and therefore $\mathbf{C}(\alpha, \alpha; \delta)$ by the monotonicity of the commutator. Altogether this shows that if \mathcal{V} has a Taylor term and satisfies implication (3.7), then the critical quotient of any \mathbf{N}_5 in a congruence lattice of a member of \mathcal{V} is abelian. Thus, for example, the variety of semilattices has the Taylor term $f(x, y) = x \wedge y$, but cannot satisfy implication (3.7) since some semilattices have sublattices isomorphic to \mathbf{N}_5 in their congruence lattices but no semilattice has an abelian congruence interval.

3. STRONG TERM CONDITIONS

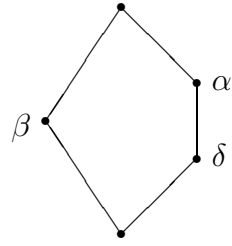


FIGURE 3.1.

you might want to ensure that this figure doesn't appear all by itself on a page in the final version.

Meet Continuous Congruence Identities

The class of all congruence lattices of algebras in a variety \mathcal{V} is usually too complicated to describe in any nontrivial way. There is more hope for a description of the class $\mathcal{L}(\mathcal{V})$ of lattices embeddable in congruence lattices of algebras in \mathcal{V} . It is evident that $\mathcal{L}(\mathcal{V})$ is closed under the formation of isomorphic lattices and sublattices. Moreover, the map

$$(4.1) \quad \prod_{i \in I} \mathbf{Con}(\mathbf{A}_i) \rightarrow \mathbf{Con}\left(\prod_{i \in I} \mathbf{A}_i\right) : (\theta_i)_{i \in I} \mapsto \Theta,$$

where $\mathbf{a} \Theta \mathbf{b}$ if $a_i \theta_i b_i$ for all i and $a_j = b_j$ for all but finitely many j , is an embedding, so $\mathcal{L}(\mathcal{V})$ is closed under the formation of products. This makes $\mathcal{L}(\mathcal{V})$ a prevariety, which we call the **congruence prevariety** of \mathcal{V} .

Any prevariety is axiomatizable by a class of formulas of the form

$$(4.2) \quad \bigwedge_{i \in I} \varepsilon_i(\mathbf{x}) \rightarrow \varepsilon(\mathbf{x})$$

where $\varepsilon(\mathbf{x})$ and each $\varepsilon_i(\mathbf{x})$ is an identity and \mathbf{x} and I may be infinite (cf. Theorem 9.2.2 of [35]). Therefore, in studying the shapes of congruence lattices in varieties, it is natural to start by considering the satisfaction of sentences of this form. The most obvious problems are still unsolved.

PROBLEM 4.1. Which congruence prevarieties are first-order axiomatizable?

In other words, for which varieties \mathcal{V} is it true that $\mathcal{L}(\mathcal{V})$ can be axiomatized by sentences like (4.2) where \mathbf{x} and I are finite? equivalently, when is $\mathcal{L}(\mathcal{V})$ a quasivariety? It is known that if \mathcal{V} is congruence n -permutable for some n , then $\mathcal{L}(\mathcal{V})$ is first-order axiomatizable, [5, 32, 33]. It is also known that if \mathcal{V} has a Taylor term but does not satisfy a nontrivial congruence identity, then $\mathcal{L}(\mathcal{V})$ is not first-order axiomatizable, [51]. The full answer to Problem 4.1 is not yet known. We do not even know the answer to the following special case.

PROBLEM 4.2. If the congruence prevariety of a congruence modular variety first-order axiomatizable?

"Is" not "If".

The congruence prevariety is first-order axiomatizable if and only if it is a quasivariety. So, it is a natural problem to consider the satisfaction of congruence quasi-identities.

PROBLEM 4.3. Is the class of varieties satisfying a given set of congruence quasi-identities definable by (idempotent, linear) Maltsev conditions?

Problem 4.3 is known to have a positive answer when posed for identities in place of quasi-identities. A. Pixley [73] and R. Wille [78] each described an algorithm to produce a family of idempotent linear Maltsev conditions defining the class of varieties satisfying a given congruence identity. By suitably modifying the Pixley–Wille algorithm, partial results have been obtained for special types of quasi-identities (cf. [3, 4, 5, 6, 40]). The strongest of these is the result we have recorded as Theorem 2.22, which states that Problem 4.3 has an affirmative answer for quasi-identities satisfying (W). The most that one could hope for is that if \mathcal{U} and \mathcal{V} satisfy the same idempotent Maltsev conditions then they have the same congruence prevariety. The most attractive special case of this is the case where $\mathcal{U} = \mathcal{SET}$, since it follows from Whitman’s Theorem (that every lattice is embeddable into a lattice of equivalence relations, cf. Section VII.8 of [1]) that $\mathcal{L}(\mathcal{SET})$ is the class \mathcal{L} of all lattices.

PROBLEM 4.4. Suppose that \mathcal{V} satisfies no nontrivial idempotent Maltsev condition. Is it true that $\mathcal{L}(\mathcal{V}) = \mathcal{L}$? Is it true at least that $\mathcal{L}(\mathcal{V})$ contains the lattice Π_4 of all equivalence relations on a 4-element set?

We expect a negative answer to both questions asked in Problem 4.4, but see Theorem 4.18 for a positive partial result.

The purpose of this chapter is to explain how the presence or absence of a Taylor term affects the shapes of congruence lattices in a variety. The Pixley–Wille algorithm shows that any variety that satisfies a nontrivial congruence identity in the language of lattices satisfies a nontrivial idempotent linear Maltsev condition and therefore has a Taylor term. The converse is not true; congruence identities are not flexible enough to say anything nontrivial about varieties satisfying very weak idempotent Maltsev conditions. In order to make nontrivial statements about the shapes of congruence lattices in arbitrary varieties with a Taylor term, we need to consider sentences other than lattice identities. In this chapter we will describe congruence lattice shapes in three different ways: by using quasi-identities in the language of

ordinary lattices, by using identities in the language of meet continuous lattices, and by identifying omitted sublattices. We will exhibit a nontrivial quasi-identity σ in the language of lattice theory such that \mathcal{V} has a Taylor term if and only if $\mathcal{L}(\mathcal{V}) \models \sigma$ (Theorem 4.12 (2)). We will prove that there is a nontrivial identity τ in the language of meet continuous lattice theory such that \mathcal{V} has a Taylor term if and only if $\mathcal{L}(\mathcal{V}) \models \tau$ (Theorem 4.12 (4)). Finally, we will exhibit a finite lattice \mathbf{L} such that \mathcal{V} has a Taylor term if and only if $\mathbf{L} \notin \mathcal{L}(\mathcal{V})$ (Theorem 4.23).

4.1. Maltsev Conditions From Congruence Identities

lattice \mathbf{L} for which ...

If \mathcal{V} is a variety, let $\mathcal{L}_{MC}(\mathcal{V})$ be the class of all meet continuous lattices for which there is a meet continuous embedding of \mathbf{L} into $\mathbf{Con}(\mathbf{A})$ for some $\mathbf{A} \in \mathcal{V}$. $\mathcal{L}_{MC}(\mathcal{V})$ is the meet continuous analogue of the congruence prevariety of \mathcal{V} . It is also a prevariety, since the mapping described in (4.1) preserves complete joins. It seems appropriate to us to identify the study of shapes of congruence lattices in varieties with the study of $\mathcal{L}_{MC}(\mathcal{V})$. In this short monograph we only begin to scratch the surface. In this chapter we will show that a variety has a Taylor term if and only if $\mathcal{L}_{MC}(\mathcal{V})$ satisfies a nontrivial meet continuous identity.

In this section we will show that the Pixley–Wille algorithm can be extended to meet continuous identities, which will show that the class of varieties satisfying a given meet continuous congruence identity is definable by idempotent linear Maltsev conditions. The proof is straightforward, but is included because there is a compactness condition associated to Maltsev conditions that we establish by proving the somewhat nonobvious fact that any meet continuous identity is equivalent to a set of identities involving only finitely many variables

third

piece. Using the same arguments as those used for ordinary lattice identities (cf. the ~~second~~ paragraph of Section 2.2), it can be shown that any family of meet continuous lattice identities is equivalent to a family of inclusions of the form $P \leq Q$ where P and Q are meet continuous lattice words. Modulo the identities defining meet continuous lattices, we may identify P and Q with elements from some free meet continuous lattice, hence with ideals from some free lattice. Now, an inclusion $P \leq Q$ between ideals is equivalent to a family of inclusions $(p] \leq Q, p \in P$, where the included ideal is principal, therefore any family of meet continuous lattice identities is equivalent to a family of inclusions of the form $p \leq Q$ where p and Q are meet continuous

lattice words with no infinitary join occurring in p . We refine this in the following lemma.

LEMMA 4.5. *Any family of meet continuous lattice identities is equivalent, modulo the identities defining the variety of meet continuous lattices, to a family of inclusions of the form $p \leq Q$ where both p and Q are meet continuous lattice words in finitely many variables, and no infinitary join occurs in p .*

continuous

PROOF. We work modulo the identities defining the class of meet continuous lattices, so we may assume that p and Q are elements in the free meet continuous lattice \mathbf{F} generated by some set $\{x_1, x_2, x_3, \dots\}$. A meet continuous lattice \mathbf{L} satisfies the inclusion $p \leq Q$ if for every meet continuous homomorphism $\varphi: \mathbf{F} \rightarrow \mathbf{L}$ it is the case that $\varphi(p) \leq \varphi(Q)$. Assume that this holds for \mathbf{L} . Since p has no infinitary joins, we may assume that the variables that appear in p are among $\{x_1, \dots, x_n\}$. Let ε be the meet continuous endomorphism of \mathbf{F} defined on the generators by $x_i \mapsto x_i$ if $1 \leq i \leq n$ and $x_i \mapsto 0$ otherwise. Then $\varphi \circ \varepsilon: \mathbf{F} \rightarrow \mathbf{L}$ is a meet continuous homomorphism, so

$$\varphi \circ \varepsilon(p) \leq \varphi \circ \varepsilon(Q)$$

by the assumption on \mathbf{L} . But ε is a retraction of \mathbf{F} onto its meet continuous sublattice generated by $\{x_1, \dots, x_n\}$, and this sublattice contains p , so $\varepsilon(p) = p$. Thus $\varphi(p) \leq \varphi(Q')$ where $Q' := \langle \varepsilon(Q) \rangle$. Since φ was arbitrary, this shows that if $p \leq Q$ holds in \mathbf{L} then $p \leq Q'$ holds in \mathbf{L} . The converse is obvious: ε is decreasing on generators, so $Q' = \varepsilon(Q) \leq Q$ in \mathbf{F} , and so $p \leq Q'$ implies $p \leq Q$. Therefore, a meet continuous lattice satisfies $p \leq Q$ if and only if it satisfies $p \leq Q'$. The inclusion $p \leq Q'$ involves elements of the (free) sublattice of \mathbf{F} generated by $\{x_1, \dots, x_n\}$, so our original inclusion $p \leq Q$ is equivalent to one in a finite set of variables. \square

Lemma 4.5 implies that there are at most $2^{2^{\aleph_0}}$ varieties of meet continuous lattices, since there are at most 2^{\aleph_0} inequivalent inclusions of the form $p \leq Q$ in finitely many variables. We do not know if this upper bound is attained, but it is not hard to see that there are at least 2^{\aleph_0} varieties of meet continuous lattices (Argument: If \mathcal{V} is a variety of lattices, then the subclass of meet continuous lattices in \mathcal{V} generates \mathcal{V} since it contains the ideal lattices of members of \mathcal{V} . This subclass is a variety in the language whose nonlogical symbols are $\{\vee, \wedge\}$, since it is closed under H, S and P. This proves that there are at least as many meet continuous lattice varieties as there are ordinary lattice varieties, and it is well known that there are 2^{\aleph_0} ordinary lattice varieties.)

PROBLEM 4.6. Determine the number of meet continuous lattice varieties.

THEOREM 4.7. *The class of varieties satisfying a given set of meet continuous congruence identities is definable by a family of idempotent linear Maltsev conditions.*

PROOF. It suffices to prove that the class of varieties whose congruence lattices satisfy a single inclusion $p \leq Q$, where p is an ordinary lattice word and Q is a meet continuous lattice word in the same set of variables, is definable by a family of idempotent linear Maltsev conditions.

If $p = p(x_1, \dots, x_n)$ and $\theta_1, \dots, \theta_n$ are congruences on an algebra, then

$$p(\theta_1, \dots, \theta_n) = \bigcup_{i \in \omega} p_i(\theta_1, \dots, \theta_n)$$

where p_i is the term in the signature $\{\circ, \wedge\}$ that is obtained from p by replacing each occurrence of binary join with the i -fold relational product and keeping each occurrence of \wedge . For example, if $p = x_1 \wedge (x_2 \vee x_3)$, then $p_1 = x_1 \wedge x_2$, $p_2 = x_1 \wedge (x_2 \circ x_3)$, $p_3 = x_1 \wedge (x_2 \circ x_3 \circ x_2)$, etc. Now, the inclusion $p(\theta_1, \dots, \theta_n) = \bigcup p_i(\theta_1, \dots, \theta_n) \leq Q(\theta_1, \dots, \theta_n)$ holds in some congruence lattice if and only if $p_i(\theta_1, \dots, \theta_n) \subseteq Q(\theta_1, \dots, \theta_n)$ holds for every i . We will show that for any single i the inclusion $p_i \subseteq Q$ can be characterized by a single Maltsev condition.

As noted in Corollary 2.10, there is a natural identification of Q with an ideal in the free lattice generated by $\{x_1, \dots, x_n\}$. This lattice is countable, so the ideal associated to Q has a cofinal sequence $q^0 \leq q^1 \leq q^2 \leq \dots$ of lattice words. Let q_j^k be the word in $\{\circ, \wedge\}$ obtained from q^k by replacing each join with j -fold alternating \circ -composition and keeping each occurrence of \wedge . For example, if $q^k = (x_1 \wedge x_2) \vee (x_1 \wedge x_3)$, then $q_3^k = (x_1 \wedge x_2) \circ (x_1 \wedge x_3) \circ (x_1 \wedge x_2)$. Let $r_j = q_j^0 \circ q_j^1 \circ \dots \circ q_j^j$. If $\theta_1, \dots, \theta_n$ are congruences on some algebra, \circ represents relational product, and \wedge represents intersection of relations, then $q_j^k(\theta_1, \dots, \theta_n) \subseteq r_j(\theta_1, \dots, \theta_n) \subseteq q^j(\theta_1, \dots, \theta_n)$ for all $k \leq j$, so

$$Q(\theta_1, \dots, \theta_n) = \bigcup_{j \in \omega} r_j(\theta_1, \dots, \theta_n).$$

We will argue that (a) any variety satisfying the congruence inclusion $p_i \subseteq Q = \bigcup r_j$ must satisfy a congruence inclusion of the form $p_i \subseteq r_j$ for some j , and that (b) the class of varieties satisfying any particular inclusion $p_i \subseteq r_j$ is definable by an idempotent linear strong Maltsev condition. If this strong Maltsev condition is denoted σ_j , then $\sigma_j \vdash$

σ_{j+1} (since $r_j(\theta_1, \dots, \theta_n) \subseteq r_{j+1}(\theta_1, \dots, \theta_n)$ for any choice of the θ 's). Therefore, (a) and (b) will show that the class of varieties satisfying $p_i \subseteq Q$ is definable by the Maltsev condition $\bigvee \sigma_j$.

Next we concern ourselves with the problem of constructing an idempotent linear strong Maltsev condition associated to a congruence inclusion $p \subseteq r$ where p and r are terms in the signature $\{\circ, \wedge\}$. The construction is a straightforward generalization of Maltsev's construction of a strong Maltsev condition defining congruence permutability, which many readers will already know. We will illustrate the construction by deriving the Maltsev condition associated with the congruence inclusion

$$(4.3) \quad \theta_1 \wedge (\theta_2 \circ \theta_3) \subseteq (\theta_1 \wedge \theta_2) \circ (\theta_1 \wedge \theta_3),$$

which is $p \subseteq r$ for $p(x_1, x_2, x_3) = x_1 \wedge (x_2 \circ x_3)$ and $r(x_1, x_2, x_3) = (x_1 \wedge x_2) \circ (x_1 \wedge x_3)$.

We begin by building a directed labeled graph associated to a $\{\circ, \wedge\}$ -term p . Start with a graph $\mathcal{G}_1(p)$ having two vertices y_1 and y_2 connected by an edge (y_1, y_2) directed from left to right and labeled with p , as depicted in Figure 4.1. From a partially completed graph $\mathcal{G}_i(p)$ the

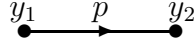


FIGURE 4.1.

construction continues by selecting an edge directed from left to right, labeled by a term w , and connecting vertices y_i and y_j . If $w = u \wedge v$,

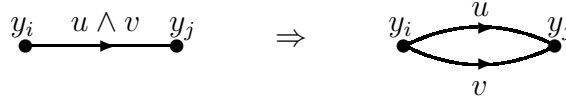


FIGURE 4.2.

then $\mathcal{G}_{i+1}(p)$ is the graph obtained from $\mathcal{G}_i(p)$ by replacing the edge labeled w with two edges directed from left to right, labeled u and v respectively, connecting the same vertices in parallel, as in Figure 4.2. If instead $w = u \circ v$, then construct $\mathcal{G}_{i+1}(p)$ by replacing the edge



FIGURE 4.3.

labeled w with two edges directed from left to right, labeled u and

v , connecting the same vertices in serial through a new vertex, as in Figure 4.3. If neither of the above steps can be performed, then all edges are labeled by variables and the construction is complete. The construction will eventually end since each step alters the complexity of exactly one edge label, and reduces the complexity of that label without introducing any more labels of the same or larger complexity.

In our illustrative example the term p equals $x_1 \wedge (x_2 \circ x_3)$ and the graphs that arise in the three steps of the construction are depicted in Figure 4.4.

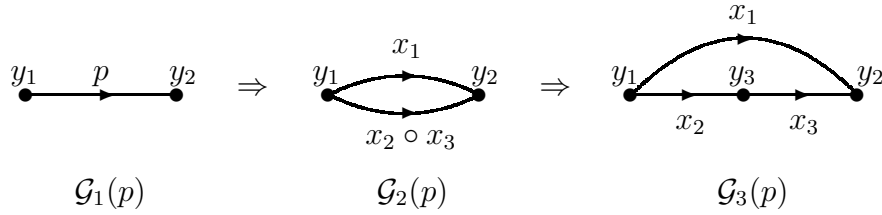


FIGURE 4.4. The construction of $\mathcal{G}(p)$ for $p = x_1 \wedge (x_2 \circ x_3)$

Call the final graph in the construction $\mathcal{G}(p)$. We may assume that the vertices are $Y = \{y_1, \dots, y_m\}$ and that the edge labels are the variables occurring in p , which are among the variables x_1, \dots, x_n . The key features of $\mathcal{G}(p)$ are stated in Claim 4.8.

CLAIM 4.8. *Let \mathbf{A} be an algebra and let $\alpha_i \in \text{Con}(\mathbf{A})$ for $1 \leq i \leq n$.*

- (1) *If $Y \rightarrow A: y_s \mapsto a_s$ is an assignment such that $(a_i, a_j) \in \alpha_k$ whenever (y_i, y_j) is an x_k -labeled edge of $\mathcal{G}(p)$, then $(a_1, a_2) \in p(\alpha_1, \dots, \alpha_n)$.*
- (2) *Conversely, given any $(a_1, a_2) \in p(\alpha_1, \dots, \alpha_n)$, there is an assignment $Y \rightarrow A: y_s \mapsto a_s$ extending $y_1 \mapsto a_1, y_2 \mapsto a_2$ such that $(a_i, a_j) \in \alpha_k$ whenever (y_i, y_j) is an x_k -labeled edge of $\mathcal{G}(p)$.*

Choose and fix assignments $Y \rightarrow A: y_s \mapsto a_s$ of vertex labels of $\mathcal{G}(p)$ to elements of A and $X \rightarrow \text{Con}(\mathbf{A}): x_s \mapsto \alpha_s$ of edge labels to congruences on \mathbf{A} . With respect to these choices, consider the following property of some $\mathcal{G}_\ell(p)$: “For each w -labeled edge (y_i, y_j) the pair (a_i, a_j) is in $w(\bar{\alpha}) := w(\alpha_1, \dots, \alpha_n)$.” Item (1) asserts that if this property holds for $\mathcal{G}(p)$, then it holds for $\mathcal{G}_1(p)$. Therefore, it is enough to show that if the property holds for $\mathcal{G}_{\ell+1}(p)$, then it holds for $\mathcal{G}_\ell(p)$. For a construction step of the ‘parallel type’ (Figure 4.2) this is a consequence of the fact that if $(a_i, a_j) \in u(\bar{\alpha})$ and $(a_i, a_j) \in v(\bar{\alpha})$, then

$(a_i, a_j) \in u \wedge v(\bar{\alpha})$. For a construction step of the ‘serial type’ (Figure 4.3) this is a consequence of the fact that if $(a_i, a_k) \in u(\bar{\alpha})$ and $(a_k, a_j) \in v(\bar{\alpha})$, then $(a_i, a_j) \in u \circ v(\bar{\alpha})$.

For (2), choose and fix an assignment $X \rightarrow \text{Con}(\mathbf{A}): x_s \mapsto \alpha_s$. Consider this property of an assignment $Y \rightarrow A: y_s \mapsto a_s$ of vertex labels of $\mathcal{G}_\ell(p)$ to elements of A : “For each w -labeled edge (y_i, y_j) the pair (a_i, a_j) is in $w(\bar{\alpha})$.” Item (2) asserts that any assignment $y_1 \mapsto a_1, y_2 \mapsto a_2$ of $\mathcal{G}_1(p)$ which satisfies this property can be extended to an assignment of the vertex labels of $\mathcal{G}(p)$ which satisfies this property. It is enough to prove that if an assignment of the vertex labels of $\mathcal{G}_\ell(p)$ has the property, then this assignment can be extended to an assignment of the vertex labels of $\mathcal{G}_{\ell+1}(p)$ that has the property. This is trivial for a construction step of the parallel type (Figure 4.2), since no new vertex labels are introduced and the assignment $y_s \mapsto a_s$ that witnesses the satisfaction of the property for $\mathcal{G}_\ell(p)$ also witnesses it for $\mathcal{G}_{\ell+1}(p)$. For a construction step of the serial type (Figure 4.3), a new vertex with label y_k is introduced to replace an edge (y_i, y_j) labeled $u \circ v$ with edges (y_i, y_k) and (y_k, y_j) labeled u and v respectively. In this case, extend the witnessing assignment $y_s \mapsto a_s$ for $\mathcal{G}_i(p)$ by defining $y_k \mapsto a_k$ where a_k is chosen so that $(a_i, a_k) \in u(\bar{\alpha})$ and $(a_k, a_j) \in v(\bar{\alpha})$. It is possible to make such a choice, since $u \circ v(\bar{\alpha}) = u(\bar{\alpha}) \circ v(\bar{\alpha})$, and such a choice witnesses the desired property.

Now construct a directed graph $\mathcal{G}(r)$ for r in exactly the same way, except use z ’s instead of y ’s for the vertices. See Figure 4.5 for $\mathcal{G}(r)$ when r equals $(x_1 \wedge x_2) \circ (x_1 \wedge x_3)$.

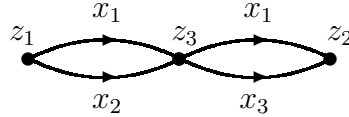


FIGURE 4.5. $\mathcal{G}(r)$ for $r = (x_1 \circ x_2) \wedge (x_1 \circ x_3)$

If the set of vertex labels for $\mathcal{G}(r)$ is $\{z_1, \dots, z_\ell\}$, then consider the following strong Maltsev condition associated to $p \subseteq r$: There exist m -ary terms F_1, \dots, F_ℓ for which the following identities are satisfied:

- (i) $F_1(y_1, \dots, y_m) \approx y_1$,
- (ii) $F_2(y_1, \dots, y_m) \approx y_2$, and
- (iii) an identity of the form $F_i(\text{variables}) \approx F_j(\text{variables})$ for each edge (z_i, z_j) in $\mathcal{G}(r)$. The pattern of the variables is determined as follows. Given a variable x_k , let E_k be the equivalence relation on the set Y of vertex labels of $\mathcal{G}(p)$ defined by $y_i E_k y_j$

if y_i and y_j belong to the same “ x_k -component”, by which we mean that there is an undirected path from y_i to y_j in $\mathcal{G}(p)$ consisting of x_k -labeled edges. Let $\tau_k: Y \rightarrow Y$ be a transversal for E_k . Now, if (z_i, z_j) is labeled by x_k , then the identity associated with (z_i, z_j) is

$$F_i(\tau_k(y_1), \dots, \tau_k(y_m)) \approx F_j(\tau_k(y_1), \dots, \tau_k(y_m)).$$

This is a linear strong Maltsev condition. The idempotence of F_1 , i.e., $F_1(x, \dots, x) \approx x$, is a consequence of identity (i) from above. There is a directed path in $\mathcal{G}(r)$ from the vertex labeled z_1 to any other vertex, and for any (z_i, z_j) in such a path the identity from (iii) above has the consequence $F_i(x, \dots, x) \approx F_j(x, \dots, x)$. Thus, the idempotence of each F_k follows from the identities (i) and (iii).

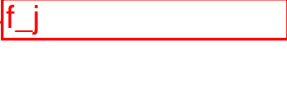
In our illustrative example, (4.3), we have $p = x_1 \wedge (x_2 \circ x_3)$ and $r = (x_1 \wedge x_2) \circ (x_1 \wedge x_3)$, and $\mathcal{G}(p)$ and $\mathcal{G}(r)$ are depicted in Figures 4.4 and 4.5. The associated Maltsev condition involves three ternary terms F_1, F_2 and F_3 , corresponding to the z 's in Figure 4.5. The equivalence relation E_k on the vertex labels of $\mathcal{G}(p)$ associated to the x_k -component can be read off Figure 4.4; we represent them by their partition of $\{y_1, y_2, y_3\}$: $E_1 \leftrightarrow y_1 y_2 / y_3$, $E_2 \leftrightarrow y_1 y_3 / y_2$, and $E_3 \leftrightarrow y_1 / y_2 y_3$. The transversals we choose for these equivalence relations are the functions τ_k which assign to each element of Y the E_k -related element whose subscript is smallest. With these choices, the Maltsev condition asserts that the following identities hold.

- $F_1(y_1, y_2, y_3) \approx y_1$,
- $F_2(y_1, y_2, y_3) \approx y_2$,
- and corresponding to τ_1, τ_2 , and τ_3 :
 - $F_1(y_1, y_1, y_3) \approx F_2(y_1, y_1, y_3) \approx F_3(y_1, y_1, y_3)$,
 - $F_1(y_1, y_2, y_1) \approx F_3(y_1, y_2, y_1)$, and
 - $F_2(y_1, y_2, y_2) \approx F_3(y_1, y_2, y_2)$.

Next we argue that the congruence lattices of algebras in a variety \mathcal{V} satisfy $p \subseteq r$ if and only if \mathcal{V} satisfies the strong Maltsev condition defined above in (i)—(iii). We begin by saying what we will mean by a “generic” pair $(u, v) \in p$. Let $\mathbf{F} = \mathbf{F}_{\mathcal{V}}(Y)$ be the free algebra in \mathcal{V} generated by $Y = \{y_1, \dots, y_m\}$. For each edge label x_k in $\mathcal{G}(p)$ let θ_k be the congruence on \mathbf{F} generated by all pairs (y_s, y_t) that are edges in $\mathcal{G}(p)$ with label x_k , i.e., $\theta_k = \text{Cg}^{\mathbf{F}}(E_k)$ where E_k is defined as above. This is the least congruence on \mathbf{F} that relates generators y_i and y_j if y_i is connected to y_j in $\mathcal{G}(p)$ by edges labeled x_k . By Lemma 3.6 (1), θ_k is the kernel of the retraction $\widehat{\tau}_k: \mathbf{F} \rightarrow \mathbf{F}$ induced by the function $\tau_k: Y \rightarrow Y$ defined in (iii) above. It is the pair $(y_1, y_2) =: (u, v)$ that

will be called a **generic** pair $(u, v) \in p$. The fact that a “generic pair in p ” actually is in $p(\theta_1, \dots, \theta_n)$ follows from Claim 4.8 (1).

If \mathcal{V} satisfies $p \subseteq r$, then the generic pair in p must belong to r , i.e., $(y_1, y_2) \in r(\theta_1, \dots, \theta_n)$ in \mathbf{F} . Using only the fact that $(y_1, y_2) \in r(\theta_1, \dots, \theta_n)$ we will prove that \mathcal{V} satisfies the Maltsev condition described above. By Claim 4.8 (2), the fact that $(y_1, y_2) \in r(\theta_1, \dots, \theta_n)$ implies that there exist $f_1 = y_1, f_2 = y_2, f_3, \dots, f_\ell \in F$ such that the assignment $z_s \mapsto f_s$ has the property that $(f_i, f_j) \in \theta_k$ whenever (z_i, z_j) is an x_k -labeled edge of $\mathcal{G}(r)$. Since \mathbf{F} is generated by $Y = \{y_1, \dots, y_m\}$, there exist terms F_i such that $f_i = F_i(y_1, \dots, y_m)$ for all i . Since $F_1(y_1, \dots, y_m) = y_1$ holds as a relation among the free generators of \mathbf{F} , it follows that $F_1(y_1, \dots, y_m) \approx y_1$ holds as an identity in \mathcal{V} . Similarly, $F_2(y_1, \dots, y_m) \approx y_2$ holds in \mathcal{V} . For each x_k -labeled edge (z_i, z_j) we have

$$\begin{aligned} F_i(\tau_k(y_1), \dots, \tau_k(y_m)) &= \widehat{\tau}_k(F_i(y_1, \dots, y_m)) \\ &= \widehat{\tau}_k(f_i) \\ &= \widehat{\tau}_k(f_k) \\ &= F_j(\tau_k(y_1), \dots, \tau_k(y_m)), \end{aligned}$$


since $(f_i, f_j) \in \theta_k = \ker(\widehat{\tau}_k)$. Since $F_i(\tau_k(\mathbf{y})) = F_j(\tau_k(\mathbf{y}))$ holds as a relation among free generators, $F_i(\tau_k(\mathbf{y})) \approx F_j(\tau_k(\mathbf{y}))$ holds as an identity in \mathcal{V} . Altogether this shows that \mathcal{V} satisfies the Maltsev condition described in (i)—(iii) above.

Conversely assume that \mathcal{V} satisfies the Maltsev condition we have described. Choose any $\mathbf{A} \in \mathcal{V}$ and congruences $\alpha_1, \dots, \alpha_n$. To show that $p(\alpha_1, \dots, \alpha_n) \subseteq r(\alpha_1, \dots, \alpha_n)$, choose $(a_1, a_2) \in p(\alpha_1, \dots, \alpha_n)$. According to Claim 4.8 (2), there is an assignment $v: Y \rightarrow A: y_s \mapsto a_s$ of the variables of $\mathcal{G}(p)$ in A such that $(a_i, a_j) \in \alpha_k$ if (y_i, y_j) is labeled by x_k . If $\widehat{v}: \mathbf{F}_{\mathcal{V}}(Y) \rightarrow \mathbf{A}$ is the extension of v to a homomorphism, and θ_k is the congruence generated by all pairs (y_i, y_j) whose edge label in $\mathcal{G}(p)$ is x_k , then $\widehat{v}(\theta_k) \subseteq \alpha_k$ since $\widehat{v}((y_i, y_j)) = (a_i, a_j) \in \alpha_k$ for any generating pair (y_i, y_j) of θ_k .

Now consider the assignment $Z \rightarrow F: z_s \mapsto F_s(y_1, \dots, y_m)$ where $Z = \{z_1, \dots, z_\ell\}$. The identities (i)—(iii) of our Maltsev condition guarantee respectively that

- (i) $z_1 \mapsto F_1(\mathbf{y}) = y_1$,
- (ii) $z_2 \mapsto F_2(\mathbf{y}) = y_2$, and
- (iii) if (z_i, z_j) is an x_k -labeled edge of $\mathcal{G}(r)$, then $(F_i(\mathbf{y}), F_j(\mathbf{y})) \in \theta_k$.

By Claim 4.8 (1), these conditions yield $(y_1, y_2) \in r(\theta_1, \dots, \theta_n)$, hence

$$(a_1, a_2) = \widehat{v}((y_1, y_2)) \in \widehat{v}(r(\theta_1, \dots, \theta_n)) \subseteq r(\alpha_1, \dots, \alpha_n).$$

Since the pair $(a_1, a_2) \in p(\alpha_1, \dots, \alpha_n)$ was arbitrary, we derive that $p \subseteq r$ holds throughout \mathcal{V} .

Recall that in the third paragraph of this proof we reduced the proof of this theorem to two claims: (a) any variety satisfying the congruence inclusion $p_i \subseteq Q = \bigcup r_j$ must satisfy a congruence inclusion of the form $p_i \subseteq r_j$ for some j , and that (b) the class of varieties satisfying any particular inclusion $p_i \subseteq r_j$ is definable by an idempotent linear strong Maltsev condition. We have just completed the proof of a strong form of (b): the class of varieties satisfying $p_i \subseteq r_j$ is definable by an idempotent linear strong Maltsev condition *which expresses the fact that $(u, v) \in r_j$ for a generic pair (u, v) of p_i* . Thus, we obtain (a) as well: a variety \mathcal{V} that satisfies $p_i \subseteq Q = \bigcup r_j$ must satisfy $(u, v) \in r_j$ for some j and some generic pair (u, v) of p_i , and for this i and j we have $p_i \subseteq r_j$. \square

For future reference we point out that this proof shows that the class of varieties satisfying a congruence inclusion of the form $p \subseteq r$, where p and r are $\{\circ, \wedge\}$ -terms, is definable by a strong Maltsev condition, and that the class of varieties satisfying a congruence inclusion of the form $p \subseteq \bigcup_{i \in \omega} r_i$, where p and r_i are $\{\circ, \wedge\}$ -terms is definable by a Maltsev condition, provided the r_i 's are increasing.

COROLLARY 4.9. *If a variety satisfies a nontrivial meet continuous congruence identity, then it satisfies a nontrivial idempotent Maltsev condition.*

4.2. Congruence Identities From Maltsev Conditions

We prove the converse to Corollary 4.9 in Theorem 4.12 of this section.

THEOREM 4.10. *Let \mathcal{V} be a variety that satisfies a nontrivial idempotent Maltsev condition. If $\mathbf{A} \in \mathcal{V}$ has congruences $\alpha, \beta, \gamma, \delta$ satisfying*

- (1) $\alpha \cap (\beta \circ \gamma) \cap (\gamma \circ \beta) \subseteq \delta \subseteq \alpha \subseteq \beta \vee \gamma$ and
- (2) $\mathbf{C}(\beta \vee \gamma, \alpha; \delta)$,

then $\delta = \alpha$.

PROOF. Let $f(x_1, \dots, x_n)$ be a Taylor term for \mathcal{V} , and let $N = \{1, \dots, n\}$ be the set of subscripts of variables in this term. The i -th Taylor identity has the form

$$f(x_1, \dots, x_n) \approx f(y_1, \dots, y_n)$$

where $(x_j, y_j) \in \{(x, x), (x, y), (y, x), (y, y)\}$ for all j , and $(x_i, y_i) = (x, y)$. As in the proof of Theorem 3.23, partition N into blocks

$\{B_1, \dots, B_5\}$ where B_1 consists of those $j \in N$ for which $(x_j, y_j) = (x, x)$, B_2 consists of those $j \neq i$ for which $(x_j, y_j) = (x, y)$, $B_3 = \{i\}$, B_4 consists of those $j \in N$ for which $(x_j, y_j) = (y, x)$, and B_5 consists of those j for which $(x_j, y_j) = (y, y)$. Let $f_i(z_1, z_2, z_3, z_4, z_5)$ be the 5-ary term obtained from f by substituting the variable z_1 for the variable x_j of $f(x_1, \dots, x_n)$ if $j \in B_1$, z_2 if $j \in B_2$, and so on. The i -th Taylor identity is now $f_i(x, x, x, y, y) \approx f_i(x, y, y, x, y)$.

Now suppose that $\mathbf{A} \in \mathcal{V}$ has congruences $\alpha, \beta, \gamma, \delta$ satisfying all the hypotheses of the theorem, and also satisfying $\alpha \neq \delta$. Then, since $\delta \leq \alpha \leq \beta \vee \gamma$, there is a least k such that $\alpha \cap (\beta \circ_k \gamma) \not\subseteq \delta$. This k is at least 2 since

$$(\alpha \cap \beta) \cup (\alpha \cap \gamma) \subseteq \alpha \cap (\beta \circ \gamma) \cap (\gamma \circ \beta) \subseteq \delta.$$

has this notation been defined?

By interchanging the roles of β and γ if necessary we may assume that $\alpha \cap (\beta \circ_{k-1} \gamma) \subseteq \delta$ and $\alpha \cap (\gamma \circ_{k-1} \beta) \subseteq \delta$.

Choose $(u, v) \in \alpha \cap (\beta \circ_k \gamma) - \delta$, and suppose that

$$u = u_0 \beta u_1 \gamma u_2 \cdots u_k = v$$

is a $\beta \circ_k \gamma$ -chain connecting u to v . If $p(x, y, z) = f_i(u_0, x, y, z, u_1)$, then for $k > 2$ the sequence

$$\begin{aligned} r = p(u_1, u_1, u_0) \gamma p(u_1, u_2, u_0) \beta p(u_0, u_3, u_1) \\ \gamma p(u_0, u_4, u_1) \beta \cdots p(u_0, u_k, u_1) = s \end{aligned}$$

is a $\gamma \circ_{k-1} \beta$ -chain. Moreover, the i -th Taylor identity implies that

$$p(u_0, u_0, u_1) = f_i(u_0, u_0, u_0, u_1, u_1) = f_i(u_0, u_1, u_1, u_0, u_1) = p(u_1, u_1, u_0),$$

so $r = p(u_0, u_0, u_1)$ is α -related to $s = p(u_0, u_k, u_1)$ (since $(u_0, u_k) \in \alpha$). Thus $(r, s) \in \alpha \cap (\gamma \circ_{k-1} \beta) \subseteq \delta$ when $k > 2$. If $k = 2$, then the same conclusion (that $(r, s) \in \delta$) may be reached as follows. Starting with the same $\beta \circ_k \gamma$ -chain $u = u_0 \beta u_1 \gamma u_2 = v$, and the same polynomial $p(x, y, z) = f_i(u_0, x, y, z, u_1)$, we obtain a chain

$$r = p(u_1, u_1, u_0) \gamma p(u_1, u_2, u_0) \beta p(u_0, u_2, u_1) = s.$$

This chain is not shorter (it has length $k = 2$ also), but it begins with γ instead of β . All other conclusions from the $k > 2$ case remain valid, in particular we have

$$r = p(u_1, u_1, u_0) = p(u_0, u_0, u_1) \alpha p(u_0, u_k, u_1) = s.$$

This shows that $(r, s) \in \alpha \cap (\gamma \circ \beta)$. But since (r, s) is an image of $(u_0, u_2) = (u, v)$ under the polynomial $p(u_0, x, u_1)$, and $(u, v) \in \beta \circ \gamma$, we even have $(r, s) \in \alpha \cap (\beta \circ \gamma) \cap (\gamma \circ \beta) \subseteq \delta$.

We now have

$$f_i(u_0, u_0, u_0, u_1, u_1) = p(u_0, u_0, u_1) = r$$

and

$$f_i(u_0, u_0, u_k, u_1, u_1) = p(u_0, u_k, u_1) = s,$$

so from $(r, s) \in \delta$ we get

$$f_i(u_0, u_0, u_0, \underline{u_1}, \underline{u_1}) \delta f_i(u_0, u_0, u_k, \underline{u_1}, \underline{u_1}).$$

Since $\mathbf{C}(\beta \vee \gamma, \alpha; \delta)$ holds, we may change underlined values from u_1 to u_0 to obtain

$$f_i(u_0, u_0, u_0, \underline{u_0}, \underline{u_0}) \delta f_i(u_0, u_0, u_k, \underline{u_0}, \underline{u_0}),$$

which we write as

$$f_i(u, u, u, u, u) \delta f_i(u, u, v, u, u).$$

This implies that for our original Taylor term f we have

$$f(u, \dots, u, v, u, \dots, u) \delta f(u, u, \dots, u) = u$$

where the v occurs in the i -th argument of f . This holds for every i . We now argue that the relation

$$f(v, \dots, v, u, \dots, u) \delta u,$$

where the last v occurs in the i -th argument of f , also holds for every i .

The facts that $(u, v) \in \alpha \leq \beta \vee \gamma$ and $\mathbf{C}(\beta \vee \gamma, \alpha; \delta)$ imply that we may change the underlined u to v in

$$f(v, \dots, v, \underline{u}, u, \dots, u) \delta u = f(u, \dots, u, \underline{u}, u, \dots, u)$$

to obtain

$$f(v, \dots, v, \underline{v}, u, \dots, u) \delta u = f(u, \dots, u, \underline{v}, u, \dots, u)$$

By our earlier conclusions, the value $f(u, \dots, u, v, u, \dots, u)$ on the right side is δ -related to u . By induction and idempotence we get $v = f(v, v, \dots, v) \delta u$, which is a contradiction. \square

The hypotheses of the previous theorem are consequences of purely order-theoretic conditions. For the statement of the next lemma, let (p, q, r) be a triple of lattice variables, and define lattice words $p_{[0]} := p$, $q_{[0]} := q$, $r_{[0]} := r$, and $p_{[n+1]} := p_{[n]} \wedge (q_{[n]} \vee r_{[n]})$, $q_{[n+1]} := q_{[n]} \wedge (p_{[n]} \vee r_{[n]})$, $r_{[n+1]} := r_{[n]} \wedge (p_{[n]} \vee q_{[n]})$. If (α, β, γ) is a triple of congruences on \mathbf{A} , then $\alpha_{[n]} := p_{[n]}(\alpha, \beta, \gamma)$, $\beta_{[n]} := q_{[n]}(\alpha, \beta, \gamma)$, $\gamma_{[n]} := r_{[n]}(\alpha, \beta, \gamma) \in \text{Con}(\mathbf{A})$.

LEMMA 4.11. *Let (α, β, γ) be a triple of congruences on \mathbf{A} , and let δ be a congruence on \mathbf{A} . The following two conditions imply the corresponding conditions of Theorem 4.10.*

- (1) $\alpha_{[n]} \leq \delta \leq \alpha \leq \beta \vee \gamma$ for some n .
- (2) $\alpha \wedge \beta = \alpha \wedge \gamma = \delta$.

should be lower case, or the start of a new sentence.

PROOF. Choose $(u, w) \in \alpha \cap (\beta \circ \gamma)$, There is a v such that $u \beta v \gamma w$. For $k = 0$, these three elements are related as in the following triangle: If u, v and w are related in this way for some $k \geq 0$, then

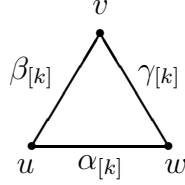


FIGURE 4.6.

$(u, w) \in \alpha_{[k]} \cap (\beta_{[k]} \vee \gamma_{[k]}) = \alpha_{[k+1]}$ and similarly $(u, v) \in \beta_{[k+1]}$ and $(v, w) \in \gamma_{[k+1]}$. This shows that $\alpha \cap (\beta \circ \gamma) \subseteq \alpha_{[k]}$ for all k . If $\alpha_{[n]} \leq \delta$ holds for some n , then we have

$$\alpha \cap (\beta \circ \gamma) \cap (\gamma \circ \beta) \subseteq \alpha \cap (\beta \circ \gamma) \subseteq \alpha_{[n]} \subseteq \delta,$$

so the first condition of this lemma implies the first condition of Theorem 4.10.

Now if $\alpha \wedge \beta = \alpha \wedge \gamma = \delta$, then we have $\mathbf{C}(\beta, \alpha; \delta)$ and $\mathbf{C}(\gamma, \alpha; \delta)$ according to item (8) of Theorem 2.19. By the semidistributivity of the centralizer in its first variable (Theorem 2.19 (5)) we derive that $\mathbf{C}(\beta \vee \gamma, \alpha; \delta)$ holds. This shows that the conditions of this lemma imply those of Theorem 4.10. \square

Now we are prepared to prove the main result of this section.

THEOREM 4.12. *The following conditions on a variety \mathcal{V} are equivalent.*

- (1) \mathcal{V} satisfies an idempotent Maltsev condition that fails in the variety of sets.
 - (2) The quasi-identity
- $$(4.4) \quad ((p \wedge q) \approx s) \ \& \ ((p \wedge r) \approx s) \ \& \ (p_{[2]} \approx s) \ \rightarrow \ ((p \wedge (q \vee r)) \approx s)$$
- holds in the congruence lattices of algebras in \mathcal{V} .*
- (3) Quasi-identity (4.4) holds in the meet continuous congruence variety of \mathcal{V} .
 - (4) \mathcal{V} satisfies a nontrivial meet continuous congruence identity.
 - (5) \mathcal{V} satisfies a congruence inclusion of the form

$$\alpha \cap (\beta \circ_4 \gamma) \subseteq w(\alpha, \beta, \gamma),$$

for some lattice word $w(p, q, r)$ such that $w(p, q, r) < p \wedge (q \vee r)$ in the free lattice $\mathbf{F}_{\mathcal{L}}(\{p, q, r\})$.

- (6) For some $k \geq 4$, \mathcal{V} satisfies a congruence inclusion of the form $\alpha \cap (\beta \circ_k \gamma) \subseteq w(\alpha, \beta, \gamma)$, where $w(p, q, r) < p \wedge (q \vee r)$ in $\mathbf{F}_{\mathcal{L}}(\{p, q, r\})$.

A seventh equivalent condition is described in the remarks following the proof.

PROOF. To see that (1) \implies (2), choose an assignment of variables $p \mapsto \alpha, q \mapsto \beta, r \mapsto \gamma, s \mapsto \delta$ in $\mathbf{Con}(\mathbf{A})$ for some $\mathbf{A} \in \mathcal{V}$ for which the premises of quasi-identity (4.4) are satisfied. Then the congruences $\alpha' := \alpha \wedge (\beta \vee \gamma), \beta, \gamma$ and δ satisfy the conditions of Lemma 4.11. (This uses the fact that $\alpha' \leq \alpha$, so $\alpha'_{[2]} \leq \alpha_{[2]} = \delta$.) Hence by Theorem 4.10 we have $\alpha \wedge (\beta \vee \gamma) = \alpha' = \delta$, which shows that the conclusion of quasi-identity (4.4) is satisfied by this assignment.

To show that (2) \iff (3) we argue that quasi-identity (4.4) is equivalent to a meet continuous lattice identity. This will follow from Theorem 2.23 if we show that the words $w_1(p, q, r) := p \wedge q, w_2(p, q, r) := p \wedge r$, and $w_3(p, q, r) := p_{[2]}$ satisfy the hypotheses of Theorem 2.23. For this we must check the comparabilities $w_1 \leq q, w_2 \leq r$ and $w_i \leq p \wedge (q \vee r)$ in the free lattice over $\{p, q, r\}$. The only nontrivial thing to verify is that $p_{[2]} \leq p \wedge (q \vee r)$. This follows from the definition $p_{[2]} := p_{[1]} \wedge (q_{[1]} \vee r_{[1]}) \leq p_{[1]} := p \wedge (q \vee r)$.

The implication (3) \implies (4) follows from the fact that quasi-identity (4.4) fails in some meet continuous lattice, e.g. the lattice \mathbf{D}_1 . The implication (4) \implies (1) is Corollary 4.9, so we have proved that the first four statements of this theorem are equivalent.

By Remark 2.24, quasi-identity (4.4) is equivalent to the meet continuous identity

$$(4.5) \quad \delta^{\mathbf{w}}(p, q, r) \approx p \wedge (q \vee r).$$

for some meet continuous lattice word $\delta^{\mathbf{w}}$, computed from the premises of (4.4), which satisfies $\delta^{\mathbf{w}}(p, q, r) \leq p \wedge (q \vee r)$ in $\mathbf{F}_{\mathcal{L}_{MC}}(\{p, q, r\})$. Since this identity is nontrivial, $\delta^{\mathbf{w}}(p, q, r)$ is in fact strictly smaller than $p \wedge (q \vee r)$ in $\mathbf{F}_{\mathcal{L}_{MC}}(\{p, q, r\})$. Because of the equivalence of the first four statements of this theorem, we may assume that the meet continuous congruence identity referred to in statement (4) is identity (4.5).

To prove that \mathcal{V} satisfies $\alpha \cap (\beta \circ_4 \gamma) \subseteq w(\alpha, \beta, \gamma)$ for some word w satisfying $w(p, q, r) < p \wedge (q \vee r)$ we begin with a generic occurrence of the situation “ $(u, v) \in \alpha \cap (\beta \circ_4 \gamma)$ ” in \mathcal{V} . That is, we let $\mathbf{F} = \mathbf{F}_{\mathcal{V}}(u, v, x, y, z)$ be the free \mathcal{V} -algebra on five generators, and let $\alpha = \mathbf{Cg}^{\mathbf{A}}(u, v), \beta = \mathbf{Cg}^{\mathbf{A}}((u, x), (y, z))$, and $\gamma = \mathbf{Cg}^{\mathbf{A}}((x, y), (z, v))$. The free generators are related as indicated in Figure 4.7. Since $\alpha \cap (\beta \circ_4 \gamma) \subseteq$

remind the reader where to find the definition of D_1.

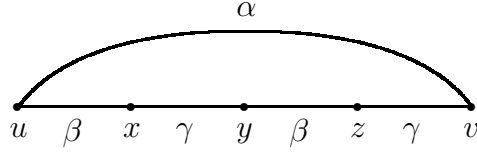


FIGURE 4.7.

$\alpha \cap (\beta \vee \gamma)$ and \mathcal{V} satisfies (4.5) as a congruence identity we get that

$$(u, v) \in \alpha \cap (\beta \circ_4 \gamma) \subseteq \alpha \cap (\beta \vee \gamma) = \delta^{\mathbf{w}}(\alpha, \beta, \gamma).$$

Under the identification of $\mathbf{F}_{\mathcal{L}_{MC}}(\{p, q, r\})$ with the ideal lattice of $\mathbf{F}_{\mathcal{L}}(\{p, q, r\})$ the word $\delta^{\mathbf{w}}(p, q, r)$ corresponds to an ideal D of ordinary lattice words $w(p, q, r)$, each of which is strictly less than $p \wedge (q \vee r)$ (since $w(p, q, r) \leq \delta^{\mathbf{w}}(p, q, r) < p \wedge (q \vee r)$). The assignment $p \mapsto \alpha, q \mapsto \beta, r \mapsto \gamma$ extends to a homomorphism $\varphi: \mathbf{F}_{\mathcal{L}_{MC}}(\{p, q, r\}) \rightarrow \mathbf{Con}(\mathbf{F})$ into the congruence lattice of the 5-generated free algebra of \mathcal{V} . This homomorphism assigns to each $w \in D$ the congruence $w(\alpha, \beta, \gamma)$, and assigns to $\delta^{\mathbf{w}}$ the union $\bigcup_{w \in D} w(\alpha, \beta, \gamma) = \delta^{\mathbf{w}}(\alpha, \beta, \gamma)$. Since $(u, v) \in \delta^{\mathbf{w}}(\alpha, \beta, \gamma)$, we must have $(u, v) \in w(\alpha, \beta, \gamma)$ for some $w \in D$. Since (u, v) is a generic element of $\alpha \cap (\beta \circ_4 \gamma)$ the inclusion $\alpha \cap (\beta \circ_4 \gamma) \subseteq w(\alpha, \beta, \gamma)$ holds throughout \mathcal{V} . This completes the proof that (4) \implies (5).

The implication (5) \implies (6) is trivial, so we now prove that (6) \implies (1). Item (6) asserts that \mathcal{V} satisfies a certain congruence inclusion in the signature $\{\vee, \wedge, \circ\}$. The class of varieties satisfying any such inclusion is definable by idempotent Maltsev conditions. To prove that at least one of the defining Maltsev conditions satisfied by \mathcal{V} is nontrivial it suffices to prove that the congruence inclusion in (6) fails in the variety of sets. To show this, let $\mathbf{A} = \langle \{a_0, a_1, \dots, a_k\}; \emptyset \rangle$ be a $(k+1)$ -element set, and let $\alpha = \text{Cg}^{\mathbf{A}}(a_0, a_k), \beta = \text{Cg}^{\mathbf{A}}(\{(a_i, a_{i+1}) \mid i \text{ even}\}), \gamma = \text{Cg}^{\mathbf{A}}(\{(a_i, a_{i+1}) \mid i \text{ odd}\})$. Then α, β and γ relate the elements of \mathbf{A} as indicated in Figure 4.8. These congruences generate a sublattice

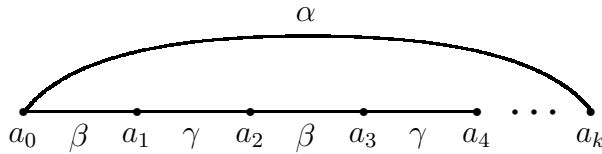


FIGURE 4.8.

of $\mathbf{Con}(\mathbf{A}) = \mathbf{Eq}(A)$ isomorphic to \mathbf{D}_1 , labeled as in Figure 4.9. Let $\varphi: \mathbf{F}_{\mathcal{L}}(\{p, q, r\}) \rightarrow \mathbf{Con}(\mathbf{A})$ be the lattice homomorphism determined by $p \mapsto \alpha, q \mapsto \beta, r \mapsto \gamma$. Since $\text{im}(\varphi) \cong \mathbf{D}_1$ is a finite lower bounded

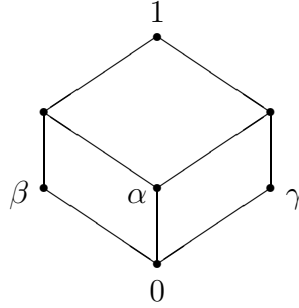


FIGURE 4.9.

lattice, for each $a \in \text{im}(\varphi)$ there is a least element $b \in \mathbf{F}_{\mathcal{L}}(\{p, q, r\})$ such that $\varphi(b) \geq a$. There is a procedure described in Lemma 2.7 of [16] for calculating this least b . Applying it to our situation shows that the least $b \in \mathbf{F}_{\mathcal{L}}(\{p, q, r\})$ such that $\varphi(b) \geq \alpha$ is $b = p \wedge (q \vee r)$. Since we assume in (4) that $w(p, q, r) < p \wedge (q \vee r)$, we get that

$$w(\alpha, \beta, \gamma) = \varphi(w(p, q, r)) \not\geq \alpha = \text{Cg}^{\mathbf{A}}(u, v).$$

Hence $(u, v) \in \alpha \cap (\beta \circ_4 \gamma) - w(\alpha, \beta, \gamma)$, which proves that the inclusion in (5) fails in the variety of sets. \square

REMARK 4.13. Quasi-identity (4.4) satisfies (W). (Recall from Section 2.2 that to check that a quasi-identity Q satisfies (W) it suffices to check that (W) holds in the partial lattice $\mathbf{P}(Q)$ of the Q -configuration. For (4.4) this partial lattice has only twelve distinct elements, so the verification is not difficult. In fact, if one first uses the axioms of lattice theory to simplify $p_{[2]} = p_{[1]} \wedge (q_{[1]} \vee r_{[1]})$ to

$$p \wedge (q_{[1]} \vee r_{[1]}) = p \wedge ((q \wedge (p \vee r)) \vee (r \wedge (p \vee q))),$$

then the associated partial lattice only has nine elements.) The theorem just proved, together with Theorem 2.22, shows that quasi-identity (4.4) is the weakest congruence quasi-identity in the collection of nontrivial congruence quasi-identities satisfying (W). That is, if Q is any nontrivial quasi-identity satisfying (W) and Q holds in all congruence lattices of algebras in some variety \mathcal{V} , then \mathcal{V} satisfies a nontrivial idempotent Maltsev condition by Theorem 2.22. Now, from (1) \implies (2) of Theorem 4.12 we get that quasi-identity (4.4) is also satisfied by all congruence lattices of algebras in \mathcal{V} . Thus any variety satisfying Q as a congruence quasi-identity also satisfies (4.4).

PROBLEM 4.14. Is there a weakest congruence quasi-identity (not necessarily satisfying (W))?

Problem 4.14 is related to Problems 4.3 and 4.4.

REMARK 4.15. It is natural to be curious about the effect of replacing $p_{[2]}$ with some other $p_{[n]}$ in item (2) of Theorem 4.12. The quasi-identities

$$Q_n : ((p \wedge q) \approx s) \ \& \ ((p \wedge r) \approx s) \ \& \ (p_{[n]} \approx s) \ \rightarrow \ ((p \wedge (q \vee r)) \approx s)$$

become formally stronger as n increases, since

$$((p \wedge q) \approx s) \ \& \ ((p \wedge r) \approx s) \ \& \ (p_{[n]} \approx s) \ \rightarrow \ (p_{[n+1]} \approx s)$$

holds in all lattices. But Q_1 is a tautology, and for $n \geq 2$ all Q_n are equivalent. For, if some assignment $p \mapsto a, q \mapsto b, r \mapsto c, s \mapsto d$ witnesses a failure of Q_n but not Q_{n-1} , then $p \mapsto a_{[n-2]}, q \mapsto b_{[n-2]}, r \mapsto c_{[n-2]}, s \mapsto d$ is an assignment witnessing a failure of Q_2 .

It is also natural to be curious about the effect of replacing the numbers 4 and k (≥ 4) in items (5) and (6) of Theorem 4.12 with smaller numbers. It happens that there is a slightly more complicated condition for $k = 3$ that is equivalent to the other conditions of Theorem 4.12, but there is no similar condition for $k = 2$.

To state the condition for $k = 3$, recursively define ternary lattice words as follows: $p^0 = p, r^0 = r, p^{n+1} = p \wedge (q \vee r^n)$, and $r^{n+1} = r \wedge (q \vee p^n)$. The equivalent condition for $k = 3$ is:

(7) \mathcal{V} satisfies a congruence inclusion of the form

$$\alpha \cap (\beta \circ_3 \gamma) \subseteq w(\alpha, \beta, \gamma),$$

where $w(p, q, r)$ is a lattice word such that $w(p, q, r) \not\geq p^n$ for any n in the free lattice $\mathbf{F}_{\mathcal{L}}(\{p, q, r\})$.

To see why it is necessary to phrase item (7) in this more complicated way, suppose that $(u, v) \in \alpha \cap (\beta \circ \gamma \circ \beta)$. Then, for $k = 0$, there are elements x and y related to u and v as in Figure 4.10. But it is

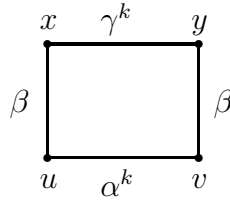


FIGURE 4.10.

easy to see that if u, v, x and y are related in this way for some k , then $(u, v) \in \alpha^k \cap (\beta \vee \gamma^k) = \alpha^{k+1}$ and $(x, y) \in \gamma^k \cap (\beta \vee \alpha^k) = \gamma^{k+1}$. By induction $\alpha \cap (\beta \circ_3 \gamma) \subseteq \alpha^k$ for all k . This shows that if a lattice word $w(p, q, r)$ is above p^n for some n , then $\alpha \cap (\beta \circ_3 \gamma) \subseteq \alpha^n \subseteq w(\alpha, \beta, \gamma)$

holds in any lattice of equivalence relations. Thus, for an inclusion $\alpha \cap (\beta \circ_3 \gamma) \subseteq w(\alpha, \beta, \gamma)$ to be nontrivial, we must have $w(p, q, r) \not\geq p^n$ for all n . This is different from inclusions of the form $\alpha \cap (\beta \circ_k \gamma) \subseteq w(\alpha, \beta, \gamma)$ for $k \geq 4$, since these inclusions are nontrivial as soon as $w(p, q, r) \not\geq p^1 = p \wedge (q \vee r)$.

Fortunately, that is the only modification of the condition on w that is needed. To prove (4) \implies (7), observe that the assignment of variables in \mathbf{E}_1 defined by $p \mapsto \alpha, q \mapsto \beta, r \mapsto \gamma, s \mapsto 0$ is one that satisfies the premises but not the conclusion of quasi-identity (4.4).

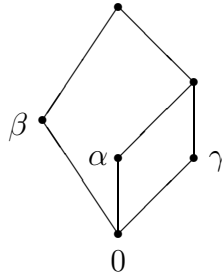


FIGURE 4.11. The lattice \mathbf{E}_1

An examination of the third paragraph of the proof of Theorem 2.23 reveals that the same assignment must therefore witness a failure of the equivalent identity (4.5). Thus, if $\varphi: \mathbf{F}_{\mathcal{L}_{MC}}(\{p, q, r\}) \rightarrow \mathbf{E}_1$ is the extension of the assignment $p \mapsto \alpha, q \mapsto \beta, r \mapsto \gamma$ to a homomorphism, then $\varphi(\delta^w(p, q, r)) < \varphi(p \wedge (q \vee r)) = \alpha$. On the other hand, an easy induction shows that $\varphi(p^n) = \alpha$ for all n . Hence $\delta^w(p, q, r) \not\geq p_n$ for all n in $\mathbf{F}_{\mathcal{L}_{MC}}(\{p, q, r\})$. Now, as in the proof of Theorem 4.12 (4) \implies (5), the word $w(p, q, r)$ lies below $\delta^w(p, q, r)$ in $\mathbf{F}_{\mathcal{L}_{MC}}(\{p, q, r\})$, hence satisfies $w(p, q, r) \not\geq p^n$ for all n in $\mathbf{F}_{\mathcal{L}}(\{p, q, r\})$. This proves (4) \implies (7).

One can prove that (7) \implies (1) in essentially the same way we proved that (6) \implies (1) using the number 3 in place of k . This time, define $\mathbf{A} = \langle \{a_0, a_1, a_2, a_3\}; \emptyset \rangle$, $\alpha = \text{Cg}^{\mathbf{A}}(a_0, a_3)$, $\beta = \text{Cg}^{\mathbf{A}}((a_0, a_1), (a_2, a_3))$ and $\gamma = \text{Cg}^{\mathbf{A}}(a_1, a_2)$. These congruences on \mathbf{A} generate a sublattice of $\mathbf{Con}(\mathbf{A})$ isomorphic to the lattice \mathbf{E}_1 , therefore \mathbf{E}_1 should be used in the argument in place of \mathbf{D}_1 . Unlike \mathbf{D}_1 , the lattice \mathbf{E}_1 is not lower bounded. However it can be shown (using Theorems 2.2 and 2.3 of [16]) that if $\varphi: \mathbf{F}_{\mathcal{L}}(\{p, q, r\}) \rightarrow \mathbf{Con}(\mathbf{A})$ is the lattice homomorphism determined by $p \mapsto \alpha, q \mapsto \beta, r \mapsto \gamma$, then the set of all $b \in \mathbf{F}_{\mathcal{L}}(\{p, q, r\})$ such that $\varphi(b) \geq \alpha$ is the filter $\{t \mid t(p, q, r) \geq p^n, \text{ for some } n\}$. If w is as in item (7), then w is not in this filter, so $\varphi(w) \not\geq \alpha$. As in the

proof we gave for (6) \implies (1), this implies that the congruence inclusion $\alpha \cap (\beta \circ_3 \gamma) \subseteq w(\alpha, \beta, \gamma)$ fails in \mathbf{A} , hence fails in the variety of sets.

There is no analogous condition for $k = 2$ that is equivalent to the other conditions of Theorem 4.12. The reason for this is that the condition that $\alpha \cap (\beta \circ \gamma) \subseteq w(\alpha, \beta, \gamma)$ is trivial if $w(p, q, r) \geq p_{[n]}$ for some n , as we showed in the first paragraph of the proof of Lemma 4.11. On the other hand, Theorems 2.2 and 2.3 of [16] can be used to show that if $w(p, q, r) \not\geq p_{[n]}$ for all n , then $\varphi(w) \not\geq \alpha$ for the homomorphism $h : \mathbf{F}_{\mathcal{L}}(\{p, q, r\}) \rightarrow \mathbf{M}_3 : p \mapsto \alpha, q \mapsto \beta, r \mapsto \gamma$, where \mathbf{M}_3 is labeled as in Figure 4.12. Therefore, if the congruence inclusion

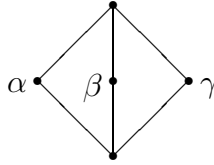


FIGURE 4.12. The lattice \mathbf{M}_3

$\alpha \cap (\beta \circ \gamma) \subseteq w(\alpha, \beta, \gamma)$ is nontrivial and holds throughout \mathcal{V} , then whenever α, β and γ are congruences on some $\mathbf{A} \in \mathcal{V}$ that generate a copy of \mathbf{M}_3 it must be that $\alpha \cap (\beta \circ \gamma) \subseteq \alpha \cap w(\alpha, \beta, \gamma) \not\subseteq \alpha$. In particular, the idempotent Maltsev condition associated with $\alpha \cap (\beta \circ \gamma) \subseteq w(\alpha, \beta, \gamma)$ precludes the possibility that α, β and γ generate a copy of \mathbf{M}_3 where $\beta \circ \gamma = \beta \vee \gamma$. Such an idempotent Maltsev condition must fail in any nontrivial variety of modules, since any nontrivial variety of modules contains members whose congruence lattices contain sublattices of permuting congruences isomorphic to \mathbf{M}_3 . But we will prove in Theorem 8.1 (10) \implies (1) that any idempotent Maltsev condition that fails in every nontrivial variety of modules implies congruence meet semidistributivity. This shows that a congruence inclusion of the type $\alpha \cap (\beta \circ \gamma) \subseteq w(\alpha, \beta, \gamma)$ for some lattice word w is trivial if $w(p, q, r) \geq p_{[n]}$ for some n , and implies congruence meet semidistributivity if $w(p, q, r) \not\geq p_{[n]}$ for all n . Therefore any such inclusion is either too weak or too strong to be equivalent to the existence of a Taylor term. We emphasize, however, that this conclusion was reached under the assumption that w is a lattice word. There do exist expressions $W(p, q, r)$ in the symbols \vee, \wedge and \circ such that the congruence inclusion $\alpha \cap (\beta \circ \gamma) \subseteq W(\alpha, \beta, \gamma)$ is nontrivial but does not imply congruence meet semidistributivity for a variety. (E.g., Theorem 8.13 (2) of this monograph, or Theorem 4.8 (2) and (3) of [52].)

4.3. Omitted Sublattices

In this section we show that if \mathcal{V} has a Taylor term, then certain lattices cannot be embedded in congruence lattices of algebras in \mathcal{V} .

THEOREM 4.16. *Let \mathcal{V} be a variety that satisfies a nontrivial idempotent Maltsev condition and let \mathbf{A} be a member of \mathcal{V} .*

- (1) *The lattices in Figure 4.13 do not occur in the meet continuous congruence variety of \mathcal{V} .*

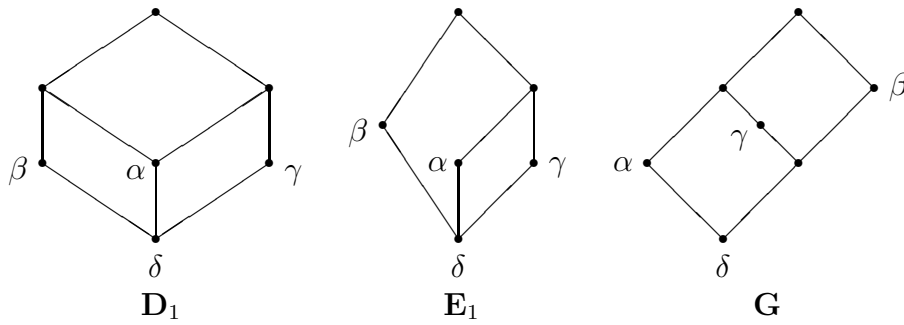


FIGURE 4.13.

- (2) *The lattices in Figure 4.14 with the given centralities do not occur in $\mathbf{Con}(\mathbf{A})$.*

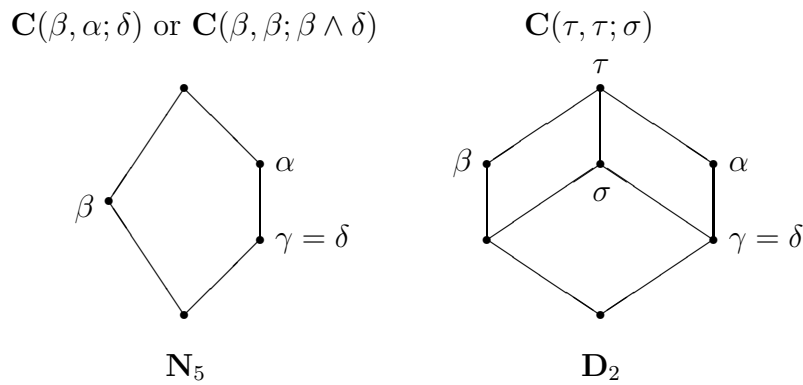


FIGURE 4.14.

- (3) *The lattices in Figure 4.15 do not occur as sublattices of $\mathbf{Con}(\mathbf{A})$.*

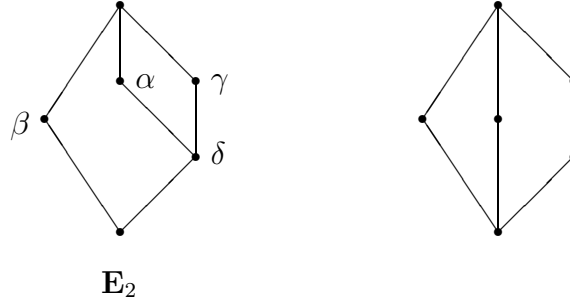


FIGURE 4.15.

PROOF. For the first claim of this theorem, the fact that \mathcal{V} satisfies a nontrivial idempotent Maltsev condition implies that the congruence lattices of algebras in \mathcal{V} satisfy quasi-identity (4.4) of Theorem 4.12 (2). But in the proof of Theorem 4.12 we explained why that quasi-identity is equivalent to a meet continuous identity, so in fact all members of the meet continuous congruence variety of \mathcal{V} satisfy quasi-identity (4.4). Item (1) of this theorem follows from the fact that \mathbf{D}_1 , \mathbf{E}_1 and \mathbf{G} all fail quasi-identity (4.4) of Theorem 4.12. (In each case the assignment $p \mapsto \alpha, q \mapsto \beta, r \mapsto \gamma, s \mapsto \delta$ witnesses a failure.)

For the second claim, observe that $\alpha_{[2]} \leq \delta \leq \alpha \leq \beta \vee \gamma$ holds for the indicated congruences in both \mathbf{N}_5 and \mathbf{D}_2 . By Lemma 4.11, this is enough to establish the first condition of Theorem 4.10. To show that the two configurations are forbidden it suffices to show that the second condition of Theorem 4.10 is satisfied, since $\delta \neq \alpha$. The second condition of Theorem 4.10, that $\mathbf{C}(\beta \vee \gamma, \alpha; \delta)$ holds, is equivalent to the conjunction of $\mathbf{C}(\beta, \alpha; \delta)$ and $\mathbf{C}(\gamma, \alpha; \delta)$ since the centralizer relation is monotone and meet semidistributive in its first variable. In the two examples $\mathbf{C}(\gamma, \alpha; \delta)$ holds simply because $\gamma = \delta$. In the first assertion concerning \mathbf{N}_5 we have $\mathbf{C}(\beta, \alpha; \delta)$ because we have assumed it. The second assertion concerning \mathbf{N}_5 can be reduced to the first, as follows. We start with our assumption that $\mathbf{C}(\beta, \beta; \beta \wedge \delta)$ holds, derive $\mathbf{C}(\beta \vee \delta, \beta \vee \delta; \delta)$ using Theorem 3.26, and then derive $\mathbf{C}(\beta, \alpha; \delta)$ from the monotonicity of the centralizer in its first two variables.

For \mathbf{D}_2 we have assumed only that $\mathbf{C}(\tau, \tau; \sigma)$ holds. From this and $\alpha, \beta \leq \tau$ we get that $\mathbf{C}(\beta, \alpha; \sigma)$ holds. We have $\mathbf{C}(\beta, \alpha; \alpha)$ trivially, so by item (6) of Theorem 2.19 and the fact that $\delta = \alpha \wedge \sigma$ we get that $\mathbf{C}(\beta, \alpha; \delta)$ holds. Since α, β and δ generate a sublattice isomorphic to \mathbf{N}_5 , we get a contradiction from the argument in the previous paragraph.

4.4. ADMITTED SUBLATTICES

For the part of the third claim concerning \mathbf{E}_2 , $\mathbf{C}(\alpha, \beta; \beta \wedge \delta)$ and $\mathbf{C}(\gamma, \beta; \beta \wedge \delta)$ hold by Theorem 2.19 (8). We can use Theorem 2.19 (5) and (1) to conclude first that $\mathbf{C}(\alpha \vee \gamma, \beta; \beta \wedge \delta)$ holds, and then that $\mathbf{C}(\beta, \beta; \beta \wedge \delta)$ holds. By applying item (2) of this theorem to the sublattice generated by α, β and δ we obtain a contradiction. One can handle the second lattice in the third claim by similar arguments. \square

consider labelling this lattice with alpha, beta, gamma, delta. my first attempt didn't work.

REMARK 4.17. The combination of Theorem 2.2 (1) and Theorem 4.16 shows that if \mathbf{L} is a finite sublattice of $\mathbf{Con}(\mathbf{A})$ where \mathbf{A} has a Taylor term, and \mathbf{M}_3 is not a sublattice of \mathbf{L} , then \mathbf{L} is meet semidistributive. This is the strongest conclusion about \mathbf{L} that can be derived from these hypotheses, since the variety \mathcal{S} of semilattices has a Taylor term, \mathbf{M}_3 is not embeddable in $\mathbf{Con}(\mathbf{A})$ for any $\mathbf{A} \in \mathcal{S}$, yet every finite meet semidistributive lattice is embeddable in the congruence lattice of some semilattice. (To show this, combine Theorem 4.4.12 of [25] with the main result of [14].)

4.4. Admitted Sublattices

Theorem 4.16 shows that if \mathcal{V} has a Taylor term, then certain lattices cannot be embedded in $\mathbf{Con}(\mathbf{A})$ for any $\mathbf{A} \in \mathcal{V}$. If, on the other hand, \mathcal{V} does not have a Taylor term, then many lattices must appear as sublattices of congruence lattices of algebras in \mathcal{V} . The strongest result of this sort that we know is the following.

THEOREM 4.18. *Let \mathcal{V} be a variety that satisfies no nontrivial idempotent Maltsev condition. If \mathbf{L} is a finitely presentable lattice satisfying Whitman's condition (W) (Definition 2.3), then $\mathbf{L} \in \mathcal{L}(\mathcal{V})$.*

PROOF. In order to prove that \mathbf{L} is a member of $\mathcal{L}(\mathcal{V})$ it suffices to prove that it is possible to separate the points of \mathbf{L} by homomorphisms into members of $\mathcal{L}(\mathcal{V})$, since $\mathcal{L}(\mathcal{V})$ is a prevariety. So, choose elements $p \neq q$ in \mathbf{L} which are to be separated.

Suppose that \mathbf{L} is presented by $\langle G \mid R \rangle$ where $G = \{g_1, \dots, g_n\}$, $\mathbf{g} = (g_1, \dots, g_n)$ and $R = \{p_1(\mathbf{g}) = q_1(\mathbf{g}), \dots, p_m(\mathbf{g}) = q_m(\mathbf{g})\}$. There exist n -ary terms $p(\mathbf{x})$ and $q(\mathbf{x})$ such that $p = p(\mathbf{g})$ and $q = q(\mathbf{g})$. The quasi-identity

$$(4.6) \quad \bigwedge_{i=1}^m (p_i(\mathbf{x}) \approx q_i(\mathbf{x})) \rightarrow (p(\mathbf{x}) = q(\mathbf{x})).$$

is nontrivial because it fails in \mathbf{L} when $\mathbf{x} = \mathbf{g}$. The quasi-identity (4.6) satisfies (W) (Definition 2.4), since \mathbf{L} does, so by Theorem 2.22 the class of varieties satisfying (4.6) is definable by idempotent Maltsev conditions.

in most, or maybe all other instances, the \approx symbol is used instead of $=$ in quasi-identities.

What this shows is that if \mathcal{V} satisfies no nontrivial idempotent Mal'tsev condition, then \mathcal{V} has an algebra \mathbf{A} such that $\mathbf{Con}(\mathbf{A})$ fails quasi-identity (4.6). If $(\theta_1, \dots, \theta_n)$ is a witness of the failure, then the function $g_i \mapsto \theta_i$ extends to a homomorphism $h: \mathbf{L} \rightarrow \mathbf{Con}(\mathbf{A})$ for which $h(p) \neq h(q)$. This proves the theorem. \square

In the case where \mathcal{V} is locally finite and \mathbf{L} is finite, Theorem 4.18 is part of Theorem 6.22 of [34].

For later reference, and to illustrate Theorem 4.18, we next describe a class of lattices satisfying (W).

DEFINITION 4.19. A **graph** is a pair $G = \langle V; E \rangle$ where V is a set, E is a set of 2-element subsets of V , and $V \cap E = \emptyset$. Elements of V are called **vertices** and elements of E are called **edges**. G is **discrete** if $E = \emptyset$.

If G is a graph, then the **lattice of G** (or **graph lattice**) is the partially ordered set $\mathbf{L}[G]$ obtained from $\langle V \cup E; \in \rangle$ by adding a new top element 1 and a new bottom element 0.

$\mathbf{L}[G]$ has height ≤ 3 and, when $V \neq \emptyset$, has V as its set of atoms and has each element of E as a coatom of height 2.

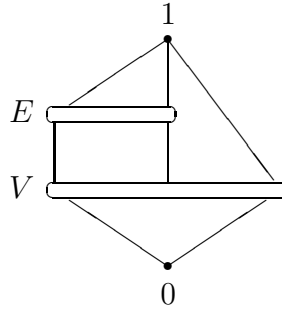


FIGURE 4.16. The shape of $\mathbf{L}[G]$

$\mathbf{L}[G]$ is indeed a lattice. For, $\langle V \cup E; \in \rangle$ is a poset of height 1 that does not have elements ordered as in Figure 4.17, since we have defined

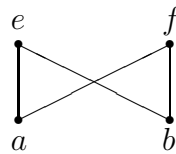


FIGURE 4.17.

graphs to exclude multiple edges e, f between distinct vertices a, b . It is easy to see that if one adds new top and bottom elements to any poset of height 1 with this property one always obtains a lattice.

LEMMA 4.20. *A lattice \mathbf{L} of height ≤ 3 satisfies (W) if and only if \mathbf{L} has no sublattice isomorphic to $\mathbf{M}_{3,3}$ (Figure 4.18).*

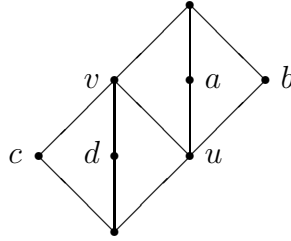


FIGURE 4.18. $\mathbf{M}_{3,3}$

In particular, any graph lattice or dual of a graph lattice satisfies (W).

PROOF. If \mathbf{L} fails (W), then there exist elements $a, b, c, d \in L$ such that $a \wedge b = u \leq v = c \vee d$ and $\{a, b, c, d\} \cap I[u, v] = \emptyset$. These conditions force $0 < u < x < 1$, $x \in \{a, b\}$, and $0 < y < v < 1$, $y \in \{c, d\}$. Since the height of \mathbf{L} is ≤ 3 , it must be that u, c and d are atoms, while v, a and b are coatoms (of height 2). In particular, this gives us that $u < v$. The interval $I[u, 1]$ has height 2, so it is isomorphic to \mathbf{M}_m for some m , and a, b and v are distinct coatoms in this interval. Similarly, $I[0, v]$ is isomorphic to \mathbf{M}_n for some $n \geq 3$, and c, d and u are distinct atoms in this interval. This, is enough to show that a, b, c and d generate a sublattice of \mathbf{L} isomorphic to the one in Figure 4.18.

Conversely, if \mathbf{L} satisfies (W), then any sublattice of \mathbf{L} must satisfy (W). Since $\mathbf{M}_{3,3}$ does not, it cannot be a sublattice of \mathbf{L} .

To see that any graph lattice omits $\mathbf{M}_{3,3}$ as a sublattice, it is enough to observe that the element v in Figure 4.18 has height 2 and dominates at least three distinct atoms. If this happened in a sublattice of a graph lattice, then v would represent an edge incident to at least three distinct vertices, which is impossible. A dual argument works for duals of graph lattices. \square

LEMMA 4.21. *If G is a nondiscrete graph with at least 4 vertices, then $\mathbf{L}[G]$ is a simple nonmodular lattice of height 3.*

PROOF. Since G has at least one edge, the height of G is exactly 3. If G is modular, then the join of 2 distinct elements of height 1 in $\mathbf{L}[G]$ (atoms, or vertices) is an element of height 2 (a coatom, or edge). This

I don't think that you get simple lattices in all cases. for example, when $G = \{a,b,c,d\}$ and a is adjacent to all other elements, and there are no other edges, then $\mathbf{L}[G]$ has $\mathbf{M}_{3,3}$ as a quotient (I think).

means that any two vertices of G are adjacent, hence G is a complete graph. Now suppose that a, b, c and d are 4 distinct vertices of G . Then $e := \{a, b\}$ and $f := \{c, d\}$ are disjoint edges of G , so they are coatoms that meet to zero in $\mathbf{L}[G]$; i.e., they are elements of coheight 1 whose meet has coheight > 2 . This is impossible in a modular lattice.

Let θ be a nontrivial congruence on $\mathbf{L}[G]$, and let $u < v$ be distinct, comparable, θ -related elements. The hypotheses guarantee that there is an atom $a \leq v$ for which $a \not\leq u$. (If v is an atom, choose $a = v$; if v is an edge, choose a to be any incident vertex not equal to u ; if v is the top, choose a to be one of the four atoms not below u .) Hence $(0, a) = (a \wedge u, a \wedge v) \in \theta$. If there is a vertex b not adjacent to a , then $(b, 1) = (0 \vee b, a \vee b) \in \theta$. Now for each vertex $c \neq b$ we have $(0, c) = (c \wedge b, c \wedge 1) \in \theta$. By hypothesis, G has at least 3 distinct vertices other than b , so $\mathbf{L}[G]$ has at least 3 distinct atoms θ -related to 0. But then the join of these atoms is θ -related to 0, and the join of any 3 atoms in a graph lattice is the top. Thus, θ is the total relation in the case that there exists a vertex b not adjacent to a .

In the remaining case G is a complete graph on at least 4 vertices. If ℓ is any edge not incident to a , then $(\ell, 1) = (0 \vee \ell, a \vee \ell) \in \theta$. From this it follows that the meet of all edges ℓ not incident to the vertex a is also θ -related to 1. These are the edges from the induced complete graph on the vertex set $V - \{a\}$. Since $V - \{a\}$ has at least 3 elements, and a complete graph on ≥ 3 vertices has no vertex incident to all edges, the meet of all edges not incident to a is zero. Again, θ is the total relation. \square

From Theorem 4.18 and Lemmas 4.20 and 4.21 we derive the following.

THEOREM 4.22. *If \mathcal{V} has no Taylor term, then any graph lattice is embeddable in the congruence lattice of some member of \mathcal{V} . Hence, there exist arbitrarily large finite, simple, nonmodular lattices that are embeddable in congruence lattices of members of \mathcal{V} .*

If we had defined graphs to allow loops on vertices, then we would have found that the corresponding graph lattices still satisfy (W), so Theorem 4.22 also applies to these lattices. Lattices of graphs with loops are still usually nonmodular, but they need not be simple. For example, this type of graph lattice is not simple if the underlying graph has at least one isolated vertex with a loop.

We have shown that certain lattices are not embeddable into congruence lattices of algebras in \mathcal{V} if \mathcal{V} has a Taylor term, while certain lattices must be embeddable if \mathcal{V} does not have a Taylor term. The

all we know at this point is that a is adjacent to all other vertices. this doesn't make G complete.

this conclusion is still valid, of course.

remove "of"

two classes that we have described overlap, allowing us to characterize the existence of a Taylor term according of the shapes of congruence lattices of algebras in \mathcal{V} .

THEOREM 4.23. *Let \mathcal{V} be a variety and let \mathbf{L} be a finitely presented lattice satisfying (W) but failing quasi-identity (4.4). (For example, \mathbf{L} could be \mathbf{D}_1 , \mathbf{E}_1 or \mathbf{G} .) The following conditions are equivalent.*

- (1) \mathcal{V} satisfies an idempotent Maltsev condition that fails in the variety of sets.
- (2) \mathcal{V} satisfies a nontrivial meet continuous congruence identity.
- (3) \mathbf{L} does not appear as a sublattice of $\mathbf{Con}(\mathbf{A})$ for any $\mathbf{A} \in \mathcal{V}$.
- (4) \mathbf{L} is not in the meet continuous congruence variety of \mathcal{V} .

PROOF. The equivalence of (1) and (2) follows from Theorems 2.14 and 4.12. That (1) \implies (4) follows from the argument we used in Theorem 4.16 (1), which only required that the lattice in question failed quasi-identity (4.4). The implication (4) \implies (3) is trivial. Finally, the contrapositive of (3) \implies (1) follows from Theorem 4.18. \square

REMARK 4.24. As we stated at the end of Section 2.3, \mathbf{D}_1 is subdirectly irreducible and projective in \mathcal{L}_{MC} . Therefore the equivalence of items (3) and (4) of Theorem 4.23 for \mathbf{D}_1 is to be expected. But \mathbf{E}_1 and \mathbf{G} are not projective in \mathcal{L}_{MC} , which makes the equivalence of (3) and (4) for these lattices somewhat surprising.

how does this definition differ from mckenzie's original definition?

Rectangulation

In this chapter we introduce a term condition for rectangularity and prove a representation theorem for rectangular tolerances. We explain how the structure of rectangular tolerances improves in the presence of a join term. We give three different Maltsev conditions that define the class of varieties that omit rectangular tolerances.

5.1. Rectangular Tolerances

DEFINITION 5.1. Let S and T be tolerances on an algebra \mathbf{A} , let \sqsupseteq be a compatible quasiorder on \mathbf{A} (i.e., \sqsupseteq is a reflexive, transitive, binary relation on A that is a subuniverse of \mathbf{A}^2). Then S **rectangulates** T **relative to** \sqsupseteq if for all $u \in A$ and all

$$(5.1) \quad \begin{bmatrix} p & q \\ r & s \end{bmatrix} \in M(S, T),$$

it is the case that

$$(5.2) \quad u \sqsupseteq q \ \& \ u \sqsupseteq r \implies u \sqsupseteq s.$$

We indicate this by saying that $\mathbf{R}(S, T; \sqsupseteq)$ **holds**.

By interchanging the rows and columns of matrices one sees that $\mathbf{R}(S, T; \sqsupseteq)$ holds if and only if

$$(5.3) \quad u \sqsupseteq q \ \& \ u \sqsupseteq r \iff u \sqsupseteq p \ \& \ u \sqsupseteq s$$

for every $u \in A$ and every S, T -matrix in (5.1).

THEOREM 5.2. Let \mathbf{A} be an algebra with tolerances S, S', T, T' , compatible quasiorders $\sqsupseteq, \sqsupseteq_j$ and a congruence δ . The following are true.

- (1) If $\mathbf{R}(S, T; \sqsupseteq)$ holds and $S' \subseteq S$, $T' \subseteq T$, then $\mathbf{R}(S', T'; \sqsupseteq)$ holds.
- (2) If $\mathbf{R}(S, T; \sqsupseteq_j)$ holds for all $j \in J$, then $\mathbf{R}(S, T; \bigcap_{j \in J} \sqsupseteq_j)$ holds.
- (3) Let \mathbf{B} be a subalgebra of \mathbf{A} . If $\mathbf{R}(S, T; \sqsupseteq)$ holds in \mathbf{A} , then $\mathbf{R}(S|_{\mathbf{B}}, T|_{\mathbf{B}}; \sqsupseteq|_{\mathbf{B}})$ holds in \mathbf{B} .
- (4) If \sqsupseteq is δ -closed (cf. Section 2.5), then $\mathbf{R}(S, T; \sqsupseteq)$ holds in \mathbf{A} if and only if $\mathbf{R}(S/\delta, T/\delta; \sqsupseteq/\delta)$ holds in \mathbf{A}/δ .
- (5) $\mathbf{R}(S, T; \sqsupseteq)$ holds if and only if $\mathbf{R}(T, S; \sqsupseteq)$ holds.

(6) $\mathbf{R}(S, T; \delta)$ holds if and only if $\mathbf{SR}(S, T; \delta)$ holds.

PROOF. Items (1)–(4) are proved in the same way as items (1), (6), (9) and (10) of Theorem 2.19.

Item (5) follows from the fact that the transpose of an S, T -matrix is a T, S -matrix, and the definition of rectangulation is invariant under transpose (as indicated in line (9) following Definition 5.1).

Item (6) follows from the fact that the definition of $\mathbf{SR}(S, T; \delta)$ is the restriction of the definition of $\mathbf{R}(S, T; \sqsubseteq)$ to the case where $\sqsubseteq = \delta$ is a congruence. \square

It follows from Theorem 5.2 (2) that for tolerances S and T on an algebra \mathbf{A} there is a least compatible quasiorder \sqsubseteq such that $\mathbf{R}(S, T; \sqsubseteq)$ holds, namely the intersection of all \sqsubseteq_j such that $\mathbf{R}(S, T; \sqsubseteq_j)$ holds. If some quasiorder \sqsubseteq_j is antisymmetric (equivalently, is a partial order), then the intersection of all \sqsubseteq_j will also be antisymmetric. It is this situation that interests us, so we make the following definition.

DEFINITION 5.3. A tolerance S **rectangulates** a tolerance T if $\mathbf{R}(S, T; \sqsubseteq)$ holds for some compatible partial order \sqsubseteq . If $S = T$ we say that T is **rectangular**.

If \sqsubseteq is the least compatible quasiorder such that $\mathbf{R}(S, T; \sqsubseteq)$ holds, then $[S, T]_{\mathbf{R}}$ denotes the intersection of \sqsubseteq with its converse \sqsupseteq ($:= \sqsubseteq^{\cup}$).

If S and T are tolerances on \mathbf{A} , \sqsubseteq is the least compatible quasiorder such that $\mathbf{R}(S, T; \sqsubseteq)$ holds, and $\delta = [S, T]_{\mathbf{R}}$ ($= \sqsubseteq \cap \sqsupseteq$), then δ is a congruence on \mathbf{A} and \sqsubseteq is δ -closed. Since $\mathbf{R}(S, T; \sqsubseteq)$ holds, we get from Theorem 5.2 (4) that $\mathbf{R}(S/\delta, T/\delta; \sqsubseteq/\delta)$ holds. Since \sqsubseteq/δ is a compatible partial order on \mathbf{A}/δ , this means that S/δ rectangulates T/δ in \mathbf{A}/δ . Moreover, it follows from Theorem 5.2 (4) that $\delta = [S, T]_{\mathbf{R}}$ is the least congruence on \mathbf{A} for which this statement is true. We record this as:

THEOREM 5.4. *If S and T are tolerances on \mathbf{A} , then $[S, T]_{\mathbf{R}}$ is the least congruence $\delta \in \text{Con}(\mathbf{A})$ such that S/δ rectangulates T/δ in \mathbf{A}/δ .*

If S and T are disjoint tolerances, then it follows from Definition 5.1 that $\mathbf{R}(S, T; 0)$ holds, and therefore $[S, T]_{\mathbf{R}} = 0$. This describes a trivial instance of rectangulation. If it happens that $S \cap T \neq 0$ yet still $[S, T]_{\mathbf{R}} = 0$, then the intersection $S \cap T$ is a nontrivial rectangular tolerance (by Theorem 5.2 (1)). We spend the rest of this section describing the structure of such objects.

DEFINITION 5.5. An operation $+$ on A is a **semilattice operation** on A if it is

i couldn't find line (9) after def. 5.1.

- (i) idempotent: $\forall x \in A(x + x \approx x)$;
- (ii) commutative: $\forall x, y \in A(x + y \approx y + x)$; and
- (iii) associative: $\forall x, y, z \in A(x + (y + z) \approx (x + y) + z)$.

If T is a tolerance on an algebra \mathbf{A} , then a partial operation $x + y$ on \mathbf{A} is a **compatible semilattice operation on T** if

- (i) $+$: $T \rightarrow \mathbf{A}$ is a homomorphism (with T considered as a subalgebra of \mathbf{A}^2); and
- (ii) $+$ is idempotent, commutative, and associative on each T -block.

THEOREM 5.6. *Let \mathbf{A} be an algebra and let θ be a congruence on \mathbf{A} . If θ supports a compatible semilattice operation, then θ is rectangular.*

PROOF. Let $+$ be a compatible semilattice operation on θ . Define a relation \sqsupseteq by

$$a \sqsupseteq b \iff (a, b) \in \theta \text{ and } a = a + b.$$

This is a partial order on A due to the facts that θ is reflexive and transitive and $+$ is a semilattice operation on θ -blocks. It is compatible because θ and $+$ are.

To show that $\mathbf{R}(\theta, \theta; \sqsupseteq)$ holds, let

$$\begin{bmatrix} p & q \\ r & s \end{bmatrix} = \begin{bmatrix} f(\mathbf{a}, \mathbf{u}) & f(\mathbf{a}, \mathbf{v}) \\ f(\mathbf{b}, \mathbf{u}) & f(\mathbf{b}, \mathbf{v}) \end{bmatrix}$$

be a θ, θ -matrix, and suppose that $u \sqsupseteq q$ and $u \sqsupseteq r$ for some $u \in A$. The definition of \sqsupseteq forces the elements u, q and r to belong to the same θ -block, and also forces $u = u + q = u + r$. This θ -block is closed under $+$, so $(u, q + r) \in \theta$. Moreover

$$u = u + u = (u + q) + (u + r) = u + (q + r),$$

proving that $u \sqsupseteq q + r$. But

$$\begin{aligned} q + r &= f(\mathbf{a}, \mathbf{v}) + f(\mathbf{b}, \mathbf{u}) \\ &= f(\mathbf{a} + \mathbf{b}, \mathbf{v} + \mathbf{u}) \\ &= f(\mathbf{a}, \mathbf{u}) + f(\mathbf{b}, \mathbf{v}) \\ &= p + s \sqsupseteq s, \end{aligned}$$

so $u \sqsupseteq s$. This completes the proof. \square

THEOREM 5.7. *Let \mathbf{A} be an algebra and let T be a tolerance on \mathbf{A} . The following conditions are equivalent*

- (1) T is rectangular.
- (2) There is an algebra $\widehat{\mathbf{A}} \in \mathcal{V}(\mathbf{A})$ and a tolerance \widehat{T} on $\widehat{\mathbf{A}}$ such that
 - (i) \mathbf{A} is a subalgebra of $\widehat{\mathbf{A}}$;

- (ii) $T \subseteq \widehat{T}|_A$;
- (iii) \widehat{T} is rectangular; and
- (iv) \widehat{T} supports a compatible semilattice operation.

PROOF. For (2) \implies (1), we leave it to the reader to check our definitions and verify that whenever $\mathbf{A} \leq \widehat{\mathbf{A}}$, \widehat{T} is a tolerance on $\widehat{\mathbf{A}}$ for which $T \subseteq \widehat{T}|_A$, and \widehat{T} rectangularizes itself with respect to the partial order \sqsupseteq , then T rectangularizes itself with respect to the restricted partial order $\sqsupseteq|_A$. (In particular, (2)(iv) is not needed to prove that (2) \implies (1).) The real content of the theorem is the claim that (1) \implies (2).

Assume that (1) holds, and that \sqsupseteq is a compatible partial order on \mathbf{A} for which $\mathbf{R}(T, T; \sqsupseteq)$ holds. We will construct $\widehat{\mathbf{A}}$ from an algebra \mathbf{B} , defined as follows. The universe of \mathbf{B} consists of the nonempty order ideals of the poset $\langle A; \sqsupseteq \rangle$. (A subset $I \subseteq A$ is an **order ideal** if $x \in I$ and $x \sqsupseteq y$ implies $y \in I$.) If f is a k -ary operation symbol in the language of \mathbf{A} , and $I_1, \dots, I_k \in B$, then $f^{\mathbf{B}}(I_1, \dots, I_k)$ is defined to be the order ideal I generated by the set of all elements of the form $f^{\mathbf{A}}(a_1, \dots, a_k)$ with $a_i \in I_i$. (Thus, $u \in I$ if and only if there exist $u_i \in I_i$ such that $f^{\mathbf{A}}(u_1, \dots, u_k) \sqsupseteq u$.) If $a \in A$, write $\langle a \rangle$ for the principal order ideal generated by a in $\langle A; \sqsupseteq \rangle$.

CLAIM 5.8. *The function $\mathbf{A} \rightarrow \mathbf{B}: x \mapsto \langle x \rangle$ is a 1-1 homomorphism.*

The fact that $\langle \rangle$ is a 1-1 function follows from the fact that $\langle A; \sqsupseteq \rangle$ is a partial order. To show that a typical k -ary basic operation f is preserved by $\langle \rangle$, choose $a_1, \dots, a_k \in A$. The ideal $f^{\mathbf{B}}(\langle a_1 \rangle, \dots, \langle a_k \rangle)$ consists of all elements $u \in A$ for which there exist elements $u_i \in \langle a_i \rangle$ such that $f^{\mathbf{A}}(u_1, \dots, u_k) \sqsupseteq u$. But since $a_i \sqsupseteq u_i$ and \sqsupseteq is compatible, this means that

$$f^{\mathbf{A}}(a_1, \dots, a_k) \sqsupseteq f^{\mathbf{A}}(u_1, \dots, u_k) \sqsupseteq u.$$

Thus $u \in f^{\mathbf{B}}(\langle a_1 \rangle, \langle a_2 \rangle, \dots, \langle a_k \rangle)$ if and only if $u \in \langle f^{\mathbf{A}}(a_1, \dots, a_k) \rangle$, which proves that $\langle \rangle$ is a homomorphism.

This claim shows that we may view \mathbf{B} as an extension of \mathbf{A} . In order to define a tolerance S on \mathbf{B} whose restriction to \mathbf{A} contains T , we first define

$$S' = \{(I, J) \in B^2 \mid (\forall u \in I)(\exists (v, w) \in (I \times J) \cap T)(v \sqsupseteq u)\}.$$

Now let S be the intersection of S' and its converse.

Let \mathcal{L} be the language of \mathbf{A} , and let $\mathcal{L}(\cup)$ be the expansion of the language to include the symbol \cup . By interpreting \cup in \mathbf{B} as union of ideals we get an $\mathcal{L}(\cup)$ -algebra $\langle \mathbf{B}; \cup \rangle$.

CLAIM 5.9. S is a tolerance on $\langle \mathbf{B}; \cup \rangle$. Moreover, if $(a, b) \in T$, then $(\langle a \rangle, \langle b \rangle) \in S$.

Since S is the intersection of S' and its converse, it will suffice to prove that S' is a compatible reflexive relation of $\langle \mathbf{B}; \cup \rangle$, and that if $(a, b) \in T$, then $(\langle a \rangle, \langle b \rangle) \in S'$.

Choose $I \in B$. For any $u \in I$ the pair $(u, u) \in (I \times I) \cap T$, so $(I, I) \in S'$. This shows that S' is reflexive.

Next, suppose that $(I_i, J_i) \in S'$, $1 \leq i \leq k$, and that f is a k -ary \mathcal{L} -operation. Let $I = f(I_1, \dots, I_k)$ and $J = f(J_1, \dots, J_k)$. To see that $(I, J) \in S'$ choose $u \in I$. There exist $u_i \in I_i$ such that $f(u_1, \dots, u_k) \sqsupseteq u$, by the definition of \mathbf{B} , and there exist $(v_i, w_i) \in (I_i \times J_i) \cap T$ with $v_i \sqsupseteq u_i$, since $(I_i, J_i) \in S'$. Since \sqsupseteq is compatible, this means that

$$v := f(v_1, \dots, v_k) \sqsupseteq f(u_1, \dots, u_k) \sqsupseteq u.$$

Since T is also compatible, for $w := f(w_1, \dots, w_k) \in J$ we have

$$(v, w) = f((v_1, w_1), \dots, (v_k, w_k)) \in T.$$

Thus, for our $u \in I$ (which we chose arbitrarily) there exist $(v, w) \in (I \times J) \cap T$ such that $v \sqsupseteq u$. Hence $(I, J) \in S'$, showing that S' is compatible with all \mathcal{L} -operations.

To show that S' is compatible with \cup , choose $(I', J'), (I'', J'') \in S'$ and let $(I, J) = (I' \cup I'', J' \cup J'')$. Choose any $u \in I = I' \cup I''$. If $u \in I'$, then since $(I', J') \in S'$ there is a pair

$$(v, w) \in ((I' \times J') \cap T) \subseteq ((I \times J) \cap T)$$

such that $v \sqsupseteq u$. The same conclusion holds if $u \in I''$, so we have that $(I, J) \in S'$. Hence S' is a compatible reflexive relation of $\langle \mathbf{B}; \cup \rangle$.

To finish the proof of the claim we must show that for $(a, b) \in T$ we have $(\langle a \rangle, \langle b \rangle) \in S'$. To see that this is so, observe that for any $u \in \langle a \rangle$ the pair $(v, w) = (a, b)$ satisfies the required properties: $(v, w) \in (\langle a \rangle \times \langle b \rangle) \cap T$ and $v \sqsupseteq u$.

We have constructed an algebra \mathbf{B} , a tolerance S , and we have a candidate \cup for the compatible semilattice operation on S . For \cup to be compatible on S it must commute with the \mathcal{L} -operations on S . This won't be true without some tinkering, so we factor by a congruence that collapses all failures of compatibility. Specifically, let θ be the $\mathcal{L}(\cup)$ -congruence on \mathbf{B} generated by all pairs of the form

$$\langle f(I_1, \dots, I_k) \cup f(J_1, \dots, J_k), f(I_1 \cup J_1, \dots, I_k \cup J_k) \rangle$$

where $f \in \mathcal{L}(\cup)$ and each $(I_i, J_i) \in S$. Since θ is defined to be a congruence of $\langle \mathbf{B}; \cup \rangle$, the quotient $\langle \mathbf{B}; \cup \rangle / \theta$ satisfies all equational properties satisfied by $\langle \mathbf{B}; \cup \rangle$. In particular, \cup is a semilattice operation of the

quotient and all operations are monotone with respect to the semilattice order.

CLAIM 5.10. *If $\langle U, V \rangle \in \theta$ and $u \in A$, then $U \subseteq \langle u \rangle \iff V \subseteq \langle u \rangle$. Hence if $\langle \langle u \rangle, \langle v \rangle \rangle \in \theta$, then $u = v$.*

To see that this is so, it suffices to check the claim only in the case where $\langle U, V \rangle$ is a polynomial image of a generating pair for θ . Here, by a polynomial, we mean a unary polynomial of $\langle \mathbf{B}; \cup \rangle$, although in our argument we will only consider special unary polynomials of $\langle \mathbf{B}; \cup \rangle$.

Let \mathbf{C} be any algebra. A **basic translation** of \mathbf{C} is a unary polynomial of the form $q(x, c_1, \dots, c_n)$ where each $c_i \in C$, and q is obtained from some basic operation by permuting variables. It is known (and not hard to see) that the congruence on \mathbf{C} generated by some set $G \subseteq C \times C$ can be described by first forming the closure of G under all basic translations, and then generating an equivalence relation. Hence $(x, y) \in \text{Cg}^{\mathbf{C}}(G)$ if and only if there is a sequence $x = x_0, \dots, x_m = y$ where each $\{x_i, x_{i+1}\} = \{p(g), p(h)\}$ for some $(g, h) \in G$ and some unary polynomial p which is a composition of basic translations.

For us, the fact that for any $f \in \mathcal{L}$ we have

$$\begin{aligned} f(x \cup W_1, W_2, \dots, W_k) &= f(x, W_2, \dots, W_k) \cup f(W_1, W_2, \dots, W_k) \\ &= f(x, W_2, \dots, W_k) \cup Z \end{aligned}$$

when all $W_i \in \langle \mathbf{B}; \cup \rangle$ and $Z := f(W_1, \dots, W_k)$ implies that if a composition of basic translations of $\langle \mathbf{B}; \cup \rangle$ involves \cup , then it can be rewritten so that it is a composition of basic translations not involving \cup , followed by a single basic translation of the form $p(x) = x \cup Z$. This observation shows that to prove the claim it suffices to show that if (H, K) is a generating pair for θ , then

$$p(H) \subseteq \langle u \rangle \iff p(K) \subseteq \langle u \rangle$$

when p has the form $p(x) = q(x, W_1, \dots, W_k) \cup Z$ where q is a composition of basic translations not involving \cup , and $W_i, Z \in B$.

Since the generating pairs for θ are comparable and the operations of $\langle \mathbf{B}; \cup \rangle$ are monotone, to prove that

$$p(H) \subseteq \langle u \rangle \iff p(K) \subseteq \langle u \rangle$$

for a generating pair (H, K) it suffices to show that

$$\begin{aligned} U := p(f(I_1, \dots, I_n) \cup f(J_1, \dots, J_n)) &\subseteq \langle u \rangle \implies \\ V := p(f(I_1 \cup J_1, \dots, I_n \cup J_n)) &\subseteq \langle u \rangle. \end{aligned}$$

Assume that $U \subseteq \langle u \rangle$ but that there is some $e \in V - \langle u \rangle$. Then, since $p(x) = q(x, W_1, \dots, W_k) \cup Z$ we get $Z \subseteq U \subseteq \langle u \rangle$, so since

$$e \in V - Z = q(f(I_1 \cup J_1, \dots, I_n \cup J_n), W_1, \dots, W_k) - Z,$$

we get that $e \in q(f(I_1 \cup J_1, \dots, I_n \cup J_n), W_1, \dots, W_k)$. We can reorder the variables of f so that there are elements $s_i \in I_i$ ($1 \leq i \leq j$), $t_i \in J_i$ ($j+1 \leq i \leq n$), $w_i \in W_i$ ($1 \leq i \leq k$) such that

$$e = q(f(s_1, \dots, s_j, t_{j+1}, \dots, t_n), w_1, \dots, w_k) = q(f(\mathbf{s}, \mathbf{t}), \mathbf{w}).$$

For each i between 1 and j we use the facts that $s_i \in I_i$ and $(I_i, J_i) \in S$ to choose $(s'_i, t'_i) \in (I_i \times J_i) \cap T$ with $s'_i \sqsupseteq s_i$, and similarly for i between $j+1$ and n we can choose $(t'_i, s'_i) \in (J_i \times I_i) \cap T$ with $t'_i \sqsupseteq t_i$. We use these elements to form tuples $\mathbf{s}' = (s'_1, \dots, s'_j)$, $\mathbf{s}'' = (s''_{j+1}, \dots, s''_n)$, \mathbf{t}' and \mathbf{t}'' . Since \mathbf{s}' and \mathbf{t}'' are T -related coordinatewise, and the same is true for \mathbf{s}'' and \mathbf{t}' , we have

$$\begin{bmatrix} q(f(\mathbf{t}'', \mathbf{s}''), \mathbf{w}) & q(f(\mathbf{t}'', \mathbf{t}'), \mathbf{w}) \\ q(f(\mathbf{s}', \mathbf{s}''), \mathbf{w}) & q(f(\mathbf{s}', \mathbf{t}'), \mathbf{w}) \end{bmatrix} = \begin{bmatrix} p & q \\ r & s \end{bmatrix} \in M(T, T)$$

with $q, r \in U \subseteq \langle u \rangle$. Since we have assumed that T rectangulates itself with respect to \sqsupseteq , and $u \sqsupseteq q, r$, we must have

$$u \sqsupseteq s = q(f(\mathbf{s}', \mathbf{t}'), \mathbf{w}) \sqsupseteq q(f(\mathbf{s}, \mathbf{t}), \mathbf{w}) = e,$$

which is contrary to our choice of e . This proves the first sentence of the claim. The second sentence follows from the first by taking $U = \langle u \rangle$, $V = \langle v \rangle$, and using the fact that \sqsupseteq is a partial order.

The relation $S/\theta = \{(I/\theta, J/\theta) \in (B/\theta)^2 \mid (I, J) \in S\}$ is a tolerance on \mathbf{B}/θ , hence a subalgebra of $\mathbf{B}/\theta \times \mathbf{B}/\theta$.

CLAIM 5.11. *The function $\cup: S/\theta \rightarrow \mathbf{B}/\theta$ is a homomorphism.*

We must show that if $(I_i/\theta, J_i/\theta) \in S/\theta$ and $f \in \mathcal{L}$, then in \mathbf{B}/θ we have

$$\begin{aligned} & f(I_1/\theta \cup J_1/\theta, \dots, I_k/\theta \cup J_k/\theta) \\ &= f(I_1/\theta, \dots, I_k/\theta) \cup f(J_1/\theta, \dots, J_k/\theta). \end{aligned}$$

Equivalently, we must show that if $(I_i, J_i) \in S$ and $f \in \mathcal{L}$, then in \mathbf{B} we have

$$f(I_1 \cup J_1, \dots, I_k \cup J_k) \equiv_{\theta} f(I_1, \dots, I_k) \cup f(J_1, \dots, J_k).$$

This is immediate from the definition of θ .

Claims 5.8 and 5.10 prove that the composite $x \mapsto \langle x \rangle \mapsto \langle x \rangle/\theta$ is an embedding: $\mathbf{A} \rightarrow \mathbf{B} \rightarrow \mathbf{B}/\theta$. We identify \mathbf{A} with its image under this homomorphism.

Let \mathcal{C} be the set of all algebras $\mathbf{C} \in \mathcal{V}(\mathbf{A})$ that are intermediate extensions: $\mathbf{A} \leq \mathbf{C} \leq \mathbf{B}/\theta$. This set contains \mathbf{A} , so it is nonempty, and it is closed under unions of chains. Therefore there is an algebra

$\widehat{\mathbf{A}} \in \mathcal{C}$ which is maximal among intermediate extensions that belong to $\mathcal{V}(\mathbf{A})$. Let $\widehat{T} = S/\theta|_{\widehat{A}}$ be the restriction of S/θ to \widehat{A} .

CLAIM 5.12. \widehat{T} is a tolerance on \widehat{A} , and conditions 2(i)–(iv) hold for \widehat{T} and \widehat{A} .

\widehat{T} is the restriction of the tolerance S/θ on \mathbf{B}/θ to \widehat{A} , so it is a tolerance on \widehat{A} . It is a consequence of Claim 5.9 that $T \subseteq \widehat{T}|_A$. What we have left to show is that \widehat{T} is rectangular, and supports a compatible semilattice operation. We show the latter first.

Claim 5.11 proves that $\cup: S/\theta \rightarrow \mathbf{B}/\theta$ is a homomorphism. This homomorphism can be restricted to \widehat{T} . The image $\cup(\widehat{T})$ of the restricted homomorphism contains \widehat{A} , since \widehat{T} is reflexive and $a \cup a = a$ for all $a \in \widehat{A}$. The image $\cup(\widehat{T})$ is a quotient of \widehat{T} , and \widehat{T} is a subalgebra of $\widehat{A} \times \widehat{A} \in \mathcal{V}(\mathbf{A})$, so the image is in $\mathcal{V}(\mathbf{A})$. The maximality of \widehat{A} prevents it from being properly contained in $\cup(\widehat{T})$, so $\cup(\widehat{T}) = \widehat{A}$. This shows that \cup is a homomorphism from \widehat{T} to \widehat{A} .

Let D be a \widehat{T} -block. If $a, b, c \in D$, then $(a, c), (b, c) \in \widehat{T} \subseteq S/\theta$ so, since S/θ is a tolerance of $\langle \mathbf{B}; \cup \rangle / \theta$,

$$(a \cup b, c) = (a, c) \cup (b, c) \in S/\theta|_{\widehat{A}} = \widehat{T}.$$

This shows that $a \cup b$ is \widehat{T} -related to any $c \in D$, so $a \cup b \in D$ by the fact that D is a block. Since \cup is a homomorphism from \widehat{T} to \widehat{A} that it is the restriction of a semilattice operation on \mathbf{B}/θ , and since each \widehat{T} -block is closed under \cup , it follows that \cup is a compatible semilattice operation on \widehat{T} .

Now we prove that \widehat{T} is rectangular. To check this we need a compatible partial order on \widehat{A} . We choose it to be the restriction of the \cup -semilattice order on \mathbf{B}/θ , which we denote by \gg and define by

$$a \gg b \iff a = a \cup b.$$

We have already observed that the operations of \mathbf{B}/θ are monotone with respect to the semilattice order, and this fact is inherited by \widehat{A} , hence \gg is a compatible partial order. To see that $\mathbf{R}(\widehat{T}, \widehat{T}; \gg)$ holds, choose $M \in M(\widehat{T}, \widehat{T})$,

$$M = \begin{bmatrix} p(\mathbf{a}, \mathbf{u}) & p(\mathbf{a}, \mathbf{v}) \\ p(\mathbf{b}, \mathbf{u}) & p(\mathbf{b}, \mathbf{v}) \end{bmatrix} = \begin{bmatrix} p & q \\ r & s \end{bmatrix},$$

with $u \gg q, r$ for some $u \in \widehat{A}$. All four elements of this matrix are \widehat{T} -related, so

$$u \gg u \cup u \gg q \cup r = p(\mathbf{a}, \mathbf{v}) \cup p(\mathbf{b}, \mathbf{u}) = p(\mathbf{a} \cup \mathbf{b}, \mathbf{v} \cup \mathbf{u}) \gg p(\mathbf{b}, \mathbf{v}) = s.$$

remove "it"

Thus, $u \gg q, r$ implies $u \gg s$, as required. This completes the proof of the claim, and also of the theorem. \square

It is worth recording that the proof produces from \mathbf{A} and T two algebras of interest: $\widehat{\mathbf{A}}$ and \mathbf{B}/θ . The properties of the two algebras are roughly the same with respect to rectangulation, but $\widehat{\mathbf{A}}$ is in the variety generated by \mathbf{A} while \mathbf{B}/θ may not be. On the other hand, the compatible semilattice operation that we constructed is only a partial operation on $\widehat{\mathbf{A}}$ (it is defined only on \widehat{T} -blocks), while it is a total operation on \mathbf{B}/θ and the semilattice order is a compatible partial order of \mathbf{B}/θ .

5.2. Rectangular Tolerances and Join Terms

We introduced join terms in Section 3.2, but did not investigate their properties. It is now important that we do so. In this section we explore the structure of rectangular tolerances on algebras in varieties with a join term. For such tolerances we prove that the representation given in Theorem 5.7 is simpler in that one may take $\widehat{\mathbf{A}} = \mathbf{A}$ and $\widehat{T} = T$. We then describe connections with the relations \geq_T and \sim_T that were introduced in Definition 3.14

THEOREM 5.13. *Let \mathcal{V} be a variety with a join term $+$, and let T be a tolerance on an algebra $\mathbf{A} \in \mathcal{V}$.*

- (1) *If B is a T -block and $w \geq_B u$ and $w \geq_B v$, then $w \geq_B u + v$. Similarly, if $w \geq_T u$ and $w \geq_T v$, then $w \geq_T u + v$.*
- (2) *If B is a T -block, then \geq_B is a join semilattice ordering on B modulo \sim_B , and $+$ is the join with respect to that ordering.*
- (3) *If B is a T -block containing all entries of $\begin{bmatrix} a & b \\ c & d \end{bmatrix} \in M(T, T)$, then $a + d \sim_B b + c$.*
- (4) *If $(q, r; s)$ is a (T, T) -triple, then $q + r \geq_T s$.*
- (5) *If p is a polynomial of \mathbf{A} and $\mathbf{u} T \mathbf{v}$, then*

$$p(\mathbf{u} + \mathbf{v}) \sim_T p(\mathbf{u}) + p(\mathbf{v}).$$

PROOF. From $w \geq_B u$ and $w \geq_B v$ it follows that $w = w + w \geq_B u + v$, since $+$ is idempotent and \geq_B is compatible with idempotent polynomials (Lemma 3.15 (1)). The statement about \geq_T follows in the same way from Lemma 3.15 (2).

If B is a T -block, then B is closed under $+$ since $+$ is idempotent. From Corollary 3.21, $u + v \geq_B u, v$, while item (1) proves conversely that if $w \geq_B u, v$, then $w \geq_B u + v$. Thus \geq_B is indeed a semilattice order modulo \sim_B on B , and $+$ is the join with respect to this order.

For item (3), Lemma 3.17 (3) proves that $b+c \geq_B (b+d)+(a+c)$. But from Corollary 3.21 we have $b+d \geq_B d$ and $a+c \geq_B a$, so

$$b+c \geq_B (b+d)+(a+c) \geq_B d+a.$$

Since \geq_B is transitive on B , we have $b+c \geq_B d+a$. If we interchange the two rows of this T, T -matrix, we derive conversely that $d+a \geq_B b+c$, hence that $b+c \sim_B d+a$. It follows from part (2) that $d+a \sim_B a+d$, so from the transitivity of \sim_B on B we get $a+d \sim_B b+c$, as claimed.

For (4), if $(q, r; s)$ is a (T, T) -triple, then its entries belong to a T, T -matrix. All four entries of any T, T -matrix are pairwise T -related, hence a block B exists that contains all four entries. Applying item (3) to this block we obtain $q+r \sim_B p+s$ for some $p \in B$, hence from item (1) it follows that $q+r \geq_B s$. This proves $q+r \geq_T s$.

We first prove item (5) in the case where $p = f$ is a basic operation of \mathbf{A} . Choose T -blocks B_i such that $u_i, v_i \in B_i$ for all i , and choose a T -block B such that $f(B_1, \dots, B_n) \subseteq B$. According to part (2) of this theorem we have

$$(5.4) \quad u_i + v_i \sim_{B_i} v_i + u_i$$

for all i . Now consider the T, T -matrix

$$\begin{bmatrix} f(\mathbf{u} + \mathbf{u}) & f(\mathbf{u} + \mathbf{v}) \\ f(\mathbf{v} + \mathbf{u}) & f(\mathbf{v} + \mathbf{v}) \end{bmatrix}.$$

All entries lie in B , so from part (3) we get that

$$f(\mathbf{u} + \mathbf{u}) + f(\mathbf{v} + \mathbf{v}) \sim_B f(\mathbf{u} + \mathbf{v}) + f(\mathbf{v} + \mathbf{u}).$$

The left hand side is $f(\mathbf{u}) + f(\mathbf{v})$, since $+$ is idempotent. The right side is \sim_B -related to $f(\mathbf{u} + \mathbf{v})$ since, by (5.4), Claim 3.16 and the idempotence of $+$ we have

$$f(\mathbf{u} + \mathbf{v}) + f(\mathbf{v} + \mathbf{u}) \sim_B f(\mathbf{u} + \mathbf{v}) + f(\mathbf{u} + \mathbf{v}) = f(\mathbf{u} + \mathbf{v}).$$

Since \sim_B is transitive on B we get $f(\mathbf{u}) + f(\mathbf{v}) \sim_B f(\mathbf{u} + \mathbf{v})$, hence $f(\mathbf{u}) + f(\mathbf{v}) \sim_T f(\mathbf{u} + \mathbf{v})$.

Now we prove (5) for polynomials. Consider the relation $\rho \subseteq A^3$ consisting of all triples (u, v, w) for which there is a T -block B such that $u, v, w \in B$ and $u+v \sim_B w$. It follows from the fact that item (5) is true when f is a basic operation, together with Claim 3.16 and the transitivity of $\sim_{B'}$ on B' for any block B' , that ρ is compatible with the basic operations of \mathbf{A} . It is compatible with all constant operations on A because $+$ is idempotent. Thus ρ is compatible with all polynomial operations. Since ρ contains all triples of the form $(u, v, u+v)$, it

therefore contains $(p(\mathbf{u}), p(\mathbf{v}), p(\mathbf{u} + \mathbf{v}))$, establishing that $p(\mathbf{u} + \mathbf{v}) \sim_T p(\mathbf{u}) + p(\mathbf{v})$. \square

We can prove more if T is rectangular.

THEOREM 5.14. *Let \mathcal{V} be a variety with a join term $+$, and let T be a tolerance on an algebra $\mathbf{A} \in \mathcal{V}$. If \sqsupseteq is a compatible partial order of \mathbf{A} for which $\mathbf{R}(T, T; \sqsupseteq)$ holds, then the following are true.*

- (1) *The relation \sim_T is the equality relation on \mathbf{A} .*
- (2) *\geq_T equals $\sqsupseteq \cap T$, and also equals $\{(u, v) \in T \mid u = u + v\}$.*
- (3) *If $\begin{bmatrix} a & b \\ c & d \end{bmatrix} \in M(T, T)$, then $a + d = b + c$.*
- (4) *If f is a polynomial of \mathbf{A} and $\mathbf{u} T \mathbf{v}$, then*

$$f(\mathbf{u} + \mathbf{v}) = f(\mathbf{u}) + f(\mathbf{v}).$$

- (5) *$+$: $T \rightarrow \mathbf{A}$ is the unique compatible semilattice operation supported by T .*

PROOF. If (r, s) is a $(T, T; 0)$ -pair, then there is a matrix

$$\begin{bmatrix} p & q \\ r & s \end{bmatrix} \in M(T, T)$$

with $q = r$. Since $r \sqsupseteq r$ and $r \sqsupseteq q (= r)$, we get $r \sqsupseteq s$ because $\mathbf{R}(T, T; \sqsupseteq)$ holds. This shows that the relation \sqsupseteq contains all $(T, T; 0)$ -pairs, hence contains \geq_T . Therefore, the reflexive relation \sim_T is contained in the intersection of \sqsupseteq and its converse, which forces \sim_T to be the equality relation. This proves item (1), and together with Theorem 5.13 (3) and (5) also proves items (3) and (4).

We will prove that \geq_T is contained in $\sqsupseteq \cap T$, which is contained in $\{(u, v) \in T \mid u = u + v\}$, which is contained in \geq_T . It follows from Definition 3.14 that \geq_T is contained in T , and it follows from the argument of the previous paragraph that \geq_T is contained in \sqsupseteq , so \geq_T is contained in $\sqsupseteq \cap T$. Next, choose any pair (u, v) in $\sqsupseteq \cap T$. Since \sqsupseteq is reflexive and compatible, and since $u \sqsupseteq v$, we get that $u = u + u \sqsupseteq u + v$. But $u + v \geq_T u$, according to the definition of a join term, so $u + v \sqsupseteq u$ by the argument of the first paragraph. This shows that both $(u, u + v)$ and $(u + v, u)$ belong to the partial order \sqsupseteq , so $u = u + v$. Hence $\sqsupseteq \cap T$ is contained in $\{(u, v) \in T \mid u = u + v\}$. Finally, choose any $(u, v) \in T$ such that $u = u + v$. If B is a T -block containing u and v , then $u = u + v \geq_B v$, since \geq_B is the join order for $+$ on B (Theorem 5.13 (2)). Therefore $u \geq_T v$, proving that $\{(u, v) \in T \mid u = u + v\}$ is contained in \geq_T . This completes the proof of item (2).

For item (5), it follows from Theorem 5.13 (2) and part (1) of this corollary that $+$ is a semilattice operation on the blocks of T . The property from item (4) of this corollary, in the cases where f is a basic operation of \mathbf{A} , is exactly the property that $+: T \rightarrow \mathbf{A}$ is a homomorphism. Thus $+$ is a compatible semilattice operation on T . If $s: T \rightarrow \mathbf{A}$ is another compatible semilattice operation, then using the commutativity and idempotence of s and $+$, and the fact that $s: T \rightarrow \mathbf{A}$ is a homomorphism, and that $+$ is a term, we have for any $(u, v) \in T$

$$\begin{aligned} s(u, v) &= s(u, v) + s(v, u) \\ &= s(u + v, v + u) \\ &= s(u + v, u + v) \\ &= u + v. \end{aligned}$$

Thus $+$ is unique. \square

COROLLARY 5.15. *Let \mathcal{V} be a variety. The following are equivalent.*

- (1) \mathcal{V} has no member with a nonzero congruence that is simultaneously rectangular and abelian.
- (2) \mathcal{V} has no member with a nonzero tolerance that is simultaneously rectangular and abelian.
- (3) \mathcal{V} satisfies a nontrivial idempotent Maltsev condition.

PROOF. Since tolerances are congruences, (2) \implies (1).

A strongly abelian congruence is both rectangular and abelian, so if (1) holds, then \mathcal{V} omits strongly abelian abelian congruences. Hence \mathcal{V} satisfies a nontrivial idempotent Maltsev condition, according to Theorem 3.12, showing that (1) \implies (3).

Finally, if \mathcal{V} satisfies a nontrivial idempotent Maltsev condition, then it has a join term, $+$. Suppose that $\mathbf{A} \in \mathcal{V}$ has a tolerance T that is rectangular and abelian. By Theorem 5.14 (5), $+$ is a semilattice operation on T -blocks. Therefore, if $(a, b) \in T$ and $a \geq b$ in the semilattice order, then

$$\begin{bmatrix} a & a \\ a & b \end{bmatrix} = \begin{bmatrix} a + a & a + b \\ b + a & b + b \end{bmatrix} \in M(T, T).$$

But since T is also abelian, the equality in the first row of this matrix implies equality in the second row. Hence $T = 0$, and (3) \implies (2). \square

One can use tame congruence theory, [34], to prove that a finite simple algebra has a nontrivial tolerance that is (a) simultaneously rectangular and abelian if and only if it has one that is (b) strongly rectangular if and only if it has one that is (c) strongly abelian. So perhaps it is not too surprising that the varieties that omit tolerances

of type (a) are the same as those that omit type (b) or (c). However, (a), (b) and (c) are truly different properties of tolerances, even for finite algebras, as we explain in the next example.

EXAMPLE 5.16. This example is an extension of Example 3.2. As in that example, let X be a set and let $\mathcal{M} \subseteq X^{2 \times 2}$ be a set of 2×2 matrices with entries in X . Define $\mathbf{A}(\mathcal{M})$ as in Example 3.2. Suppose that \gg is a quasiorder defined on the set X such that whenever $M = \begin{bmatrix} p & q \\ r & s \end{bmatrix}$ belongs to \mathcal{M} , then

$$u \gg q \ \& \ u \gg r \iff u \gg s \ \& \ u \gg p$$

for any $u \in X$. Let \sqsubseteq be the least extension of \gg to a quasiorder on $A = X \cup Y \cup \{0\}$ (namely, \sqsubseteq is the union of \gg with the equality relation on A). Then \sqsubseteq is a compatible quasiorder of $\mathbf{A}(\mathcal{M})$, and it is clear that $\mathbf{R}(R, S; \sqsubseteq)$ and $\mathbf{R}(T, T; \sqsubseteq)$ both hold (since for $u \neq v$ we have $u \sqsubseteq v$ if and only if $u \gg v$). Thus, we are again in the situation that we may ignore algebraic considerations and think set-theoretically when considering relationships between rectangulation and the term conditions introduced earlier.

To see, for example, that there is a finite algebra with a tolerance that is rectangular and abelian but not strongly rectangular (hence not strongly abelian), it is entirely sufficient to write down the single matrix $M = \begin{bmatrix} 1 & 3 \\ 3 & 2 \end{bmatrix}$. Then, for $\mathcal{M} = \{M\}$, and $\mathbf{A}(\mathcal{M})$ and T defined as in Example 3.2, the set of nontrivial T, T -matrices are just those that can be obtained from M by permuting its entries by some combination of row or column interchanges or transposes. As a T, T -matrix, M formally conflicts with $\mathbf{SR}(T, T; 0)$. But it does not formally conflict with $\mathbf{C}(T, T; 0)$ or $\mathbf{R}(T, T; \gg)$ where \gg is the partial order on X generated by the pairs $(3, 1)$ and $(3, 2)$. Therefore, on the algebra $\mathbf{A}(\mathcal{M})$, the tolerance T is rectangular and abelian but not strongly rectangular.

The construction of Example 5.16 can be used to show that it is possible for the transitive closure of the $(T, T; 0)$ -pairs to be a partial order even when T is not rectangular. For example, this occurs when

$$\mathcal{M} = \left\{ \begin{bmatrix} 1 & 3 \\ 3 & 2 \end{bmatrix}, \begin{bmatrix} 1 & 4 \\ 4 & 2 \end{bmatrix} \right\}.$$

Generalizations of this example show that the construction of the rectangular commutator of two tolerances can be arbitrarily complicated in general. The situation is better in varieties with join terms.

THEOREM 5.17. *Let \mathcal{V} be a variety with a join term $+$, and let T be a tolerance on an algebra $\mathbf{A} \in \mathcal{V}$. Let ρ be any congruence on \mathbf{A} that contains \sim_T , and let $\mathbf{C} = \mathbf{A}/\rho$ and $S = T/\rho$.*

- (1) *If $a_1 \geq_S a_2 \geq_S \cdots \geq_S a_n \geq_S a_1$ in \mathbf{C} , then $a_1 = a_2 = \cdots = a_n$.*
- (2) *If \sqsupseteq is the transitive closure of \geq_S , then \sqsupseteq is a partial order of \mathbf{C} whose restriction to any S -block agrees with \geq_S .*
- (3) *$\mathbf{R}(S, S; \sqsupseteq)$ holds.*
- (4) *The transitive closure of \sim_T is $[T, T]_{\mathbf{R}}$.*

PROOF. First we prove (1). We may assume that all (a_i, a_{i+1}) (with the understanding that $n+1 = 1$) are $(S, S; 0)$ -pairs. Now choose $n \geq 2$ minimal among such cycles for which the statement is false. We argue that $a_1 + a_2 = a_2 + a_1 = a_1$. As (a_1, a_2) is an $(S, S; 0)$ -pair, there exists a (T, T) -matrix

$$\begin{bmatrix} a & b \\ c & d \end{bmatrix}$$

such that $b/\rho = c/\rho = a_1$ and $d/\rho = a_2$. Choose a T -block B containing these four elements. By Theorem 5.13 we have that $b+c \sim_B a+d$, and that $+$ is a semilattice operation on B modulo \sim_B , which implies that $(b+c)+d \sim_B d+(b+c) \sim_B b+c$. The relation \sim_B is contained in ρ , and therefore in \mathbf{A}/ρ we have $(a_1+a_1)+a_2 = a_2+(a_1+a_1) = a_1+a_1$. But $+$ is idempotent, so $a_1+a_2 = a_2+a_1 = a_1$. A similar argument shows that $a_i+a_{i+1} = a_{i+1}+a_i = a_i$ for all i .

Consider

$$a_1 + a_2 \geq_S a_2 + a_3 \geq_S a_2 + a_4 \geq_S \cdots \geq_S a_2 + a_n \geq_S a_2 + a_1.$$

This is a cycle of length $n-1$. We show that every link in it is an $(S, S; 0)$ -pair. Indeed, the sum of two (S, S) -matrices is also an (S, S) -matrix, hence the sum of two (S, S) -triples is an (S, S) -triple as well. Now $(a_1, a_1; a_2)$ and $(a_2, a_2; a_3)$ are (S, S) -triples, and so their sum $(a_1+a_2, a_1+a_2; a_2+a_3)$ is an (S, S) -triple, proving that the first link is indeed an $(S, S; 0)$ -pair. Similarly, $(a_2, a_2; a_2)$ and $(a_i, a_i; a_{i+1})$ are (S, S) -triples, and their sum shows that all the other links in the chain are $(S, S; 0)$ -pairs, too.

The minimality of n then forces this chain to be trivial, that is, all links are equal. Thus $a_1+a_2 = a_2+a_3$. But we know that $a_1+a_2 = a_1$ and $a_2+a_3 = a_2$. Therefore $a_1 = a_2$. This contradicts the minimality of n in the original chain. This argument works for every $n > 1$, and so (1) is proved.

To prove (2) suppose that $u \sqsupseteq v$ and $v \sqsupseteq u$. Then there is a \geq_S -chain from u to v , and from v to u . Thus the statement in (1) shows

that $u = v$. Therefore \geq is indeed a partial order. To prove the other statement, suppose that

$$b_1 \geq_S b_2 \geq_S b_3 \geq_S \cdots \geq_S b_n,$$

where $(b_1, b_n) \in S$. Consider

$$(5.5) \quad b_1 + b_1 \geq_S b_2 + b_1 \geq_S b_3 + b_1 \geq_S \cdots \geq_S b_n + b_1.$$

Here $b_n + b_1 \geq_S b_1 = b_1 + b_1$ by the properties of $+$, since b_n and b_1 are S -related. Hence (5.5) is a cycle, which must be trivial by (1). Therefore $b_1 = b_n + b_1 \geq_S b_n$, proving (2).

Now we prove that $\mathbf{R}(S, S; \sqsupseteq)$ holds. Suppose that

$$\begin{bmatrix} a & b \\ c & d \end{bmatrix} \in M(S, S),$$

and $u \in C$ satisfies that $u \sqsupseteq b$ and $u \sqsupseteq c$. If B is an S -block that contains these four elements, then Theorem 5.13 (3) guarantees that $b + c \sim_B a + d$. But the relation \sim_B is the intersection of \geq_B and its converse, the relation \geq_B is contained in \geq_S , and \geq_S is antisymmetric by item (2). Thus \sim_B is equality on B , yielding $b + c = a + d$. As $+$ is idempotent and \sqsupseteq is compatible, we get that $u = u + u \sqsupseteq b + c = a + d$. But $a + d \geq_S d$ by the definition of a join term, so $u \sqsupseteq a + d \sqsupseteq d$, proving that $\mathbf{R}(S, S; \sqsupseteq)$ holds.

For item (4), let ρ be the transitive closure of \sim_T in \mathbf{A} . Since \sim_T is a tolerance, its transitive closure ρ is a congruence. Now suppose that \gg is the least compatible partial order on \mathbf{A} such that $\mathbf{R}(T, T; \gg)$ holds. If

$$\begin{bmatrix} p & q \\ r & s \end{bmatrix} \in M(T, T),$$

and $q = r$, then $r \gg r$ and $r \gg q (= r)$, so $r \gg s$. Thus \gg contains all $(T, T; 0)$ -pairs, and therefore contains \geq_T . According to Definition 5.3, $[T, T]_{\mathbf{R}}$ is the intersection of \gg and its converse, hence $[T, T]_{\mathbf{R}}$ contains the intersection of \geq_T and its converse, hence contains \sim_T . Since ρ is the least congruence containing \sim_T , we get $\rho \leq [T, T]_{\mathbf{R}}$.

On the other hand, according to Theorem 5.4, $[T, T]_{\mathbf{R}}$ is the least congruence δ for which T/δ is a rectangular tolerance of \mathbf{A}/δ . By items (2) and (3) of this theorem, ρ is a congruence with this property. Therefore $[T, T]_{\mathbf{R}} = \rho$. \square

The following consequence of Theorem 5.17 is worth stating separately.

COROLLARY 5.18. *Let \mathcal{V} be a variety with a join term, and let T be a tolerance on some $\mathbf{A} \in \mathcal{V}$. T is rectangular if and only if \sim_T is the equality relation.*

PROOF. If T is rectangular, then \sim_T is the equality relation according to Theorem 5.14 (1). On the other hand, if \sim_T is the equality relation, then Theorem 5.17 (4) guarantees that $[T, T]_{\mathbf{R}} = 0$. \square

Corollary 5.18 implies the following characterization of rectangular tolerances in varieties with join terms, which is a refinement of Theorem 5.6.

COROLLARY 5.19. *Let \mathcal{V} be a variety with a join term $+$, and let T be a tolerance on an algebra $\mathbf{A} \in \mathcal{V}$. T is rectangular if and only if $+$ is a compatible semilattice operation on T .*

PROOF. If T is rectangular, then $+$ is a compatible semilattice operation on T by Theorem 5.14 (5). Conversely, suppose that $+$ is a compatible semilattice operation on T . If (r, s) is a $(T, T; 0)$ -pair, then there is a matrix

$$\begin{bmatrix} p & q \\ r & s \end{bmatrix} = \begin{bmatrix} f(\mathbf{a}, \mathbf{u}) & f(\mathbf{a}, \mathbf{v}) \\ f(\mathbf{b}, \mathbf{u}) & f(\mathbf{b}, \mathbf{v}) \end{bmatrix} \in M(T, T)$$

with $q = r$. In this situation

$$q + r = f(\mathbf{a} + \mathbf{b}, \mathbf{v} + \mathbf{u}) = p + s,$$

and all elements $p, q, r, s, q+r$ and $p+s$ lie in a single T -block. Therefore $r = q + r = p + s \geq s$ in the semilattice order on that block. This means that the equation $r + s = r$ holds, which is a property that is independent of the block that we are considering. Thus if (r, s) is a $(T, T; 0)$ -pair, then $r \geq s$ in the semilattice order on any block containing r and s . This conclusion implies that on any T -block the relation \geq_T is contained in the semilattice order, and therefore that \sim_T is the equality relation. Hence T is rectangular by Corollary 5.18. \square

The following corollary will be used in Section 9.2.

COROLLARY 5.20. *Let \mathcal{V} be a variety with a join term $+$, and let T be a rectangular tolerance on an algebra $\mathbf{A} \in \mathcal{V}$. Assume that a and b are distinct T -related elements of \mathbf{A} . If ρ is a congruence on \mathbf{A} that is maximal with respect to $(a, b) \notin \rho$, then T/ρ is a nontrivial rectangular tolerance on the subdirectly irreducible algebra \mathbf{A}/ρ .*

PROOF. If T is rectangular, then \sim_T is the equality relation according to Theorem 5.14 (1). Therefore every congruence ρ on \mathbf{A} contains \sim_T . From Theorem 5.17, T/ρ is a rectangular congruence of \mathbf{A}/ρ for any ρ . Choosing ρ as described in the statement of this corollary guarantees that T/ρ is nontrivial and that \mathbf{A}/ρ is subdirectly irreducible. \square

5.3. Varieties Omitting Rectangular Tolerances

The purposes of this section are to show that a variety omits rectangular tolerances if and only if it omits rectangular congruences, and then to give three different Maltsev characterizations of the class of varieties with these properties.

LEMMA 5.21. *Let \mathcal{V} be a variety with a join term and let $\mathbf{F} = \mathbf{F}_{\mathcal{V}}(x, y)$ be the free \mathcal{V} -algebra generated by $\{x, y\}$. Let $T = \text{Tg}^{\mathbf{F}}(x, y)$, let $\theta = \text{Cg}^{\mathbf{F}}(T)$, and let $B = x/\theta$. The relations \geq_{θ} and \geq_T have the same restriction to B .*

PROOF. It follows from Lemma 3.15 (4) that $\geq_T|_B = \text{tr.cl.}(S|_B)$, where S is the set of $(T, T; 0)$ -pairs of \mathbf{F} , and that $\geq_{\theta}|_B = \text{tr.cl.}(\overline{S}|_B)$, where \overline{S} is the (larger) set of $(\theta, \theta; 0)$ -pairs of \mathbf{F} . Thus, it is enough to show that if $(r, s) \in B^2$ is a $(\theta, \theta; 0)$ -pair, then (r, s) is in the transitive closure of the set of $(T, T; 0)$ -pairs. So choose a θ, θ -matrix of the form

$$\begin{bmatrix} p & r \\ r & s \end{bmatrix} = \begin{bmatrix} t(\mathbf{a}, \mathbf{c}) & t(\mathbf{a}, \mathbf{d}) \\ t(\mathbf{b}, \mathbf{c}) & t(\mathbf{b}, \mathbf{d}) \end{bmatrix}.$$

Let $h: \mathbf{F} \rightarrow \mathbf{F}$ be the endomorphism determined by $x \mapsto x, y \mapsto x$. The conditions $\mathbf{a} \equiv_{\theta} \mathbf{b}$ and $\mathbf{c} \equiv_{\theta} \mathbf{d}$ are equivalent to $h(\mathbf{a}) = h(\mathbf{b})$ and $h(\mathbf{c}) = h(\mathbf{d})$. As in the proof of Lemma 3.7, consider the following 4×4 matrix:

$$\begin{array}{c} \text{extra)} \\ \left[\begin{array}{cc|cc} t(\mathbf{a}, \mathbf{c}) & t(\mathbf{a}, h(\mathbf{c})) & = & t(\mathbf{a}, h(\mathbf{d})) & t(\mathbf{a}, \mathbf{d}) \\ t(h(\mathbf{a}), \mathbf{c}) & t(h(\mathbf{a}), h(\mathbf{c})) & = & t(h(\mathbf{a}), h(\mathbf{d})) & t(h(\mathbf{a}), \mathbf{d}) \\ \parallel & \parallel & & \parallel & \parallel \\ t(h(\mathbf{b}), \mathbf{c}) & t(h(\mathbf{b}), h(\mathbf{c})) & = & t(h(\mathbf{b}), h(\mathbf{d})) & t(h(\mathbf{b}), \mathbf{d}) \\ t(\mathbf{b}, \mathbf{c}) & t(\mathbf{b}, h(\mathbf{c})) & = & t(\mathbf{b}, h(\mathbf{d})) & t(\mathbf{b}, \mathbf{d}) \end{array} \right]. \end{array}$$

Recall from that earlier proof that the four elements in the middle are equal to x and that the four 2×2 matrices in the corners are T, T -matrices. All elements of this 4×4 matrix are contained in B , because all elements are θ -related, $r, s \in B$ by assumption, and B is a block of θ .

Proceeding as we did in the proof of Lemma 3.7, let $k: \mathbf{F} \rightarrow \mathbf{F}$ be the endomorphism of \mathbf{F} that sends x to y and y to y . Consider the analogous 4×4 matrix with respect to k , and apply $r(x, y)$ to these two matrices, the first matrix in the first argument, the second matrix in the second argument. As before, we get a 4×4 matrix which, after deleting one of the doubled middle columns and one of the doubled

middle rows, is a 3×3 matrix of the form:

$$J = \begin{bmatrix} p & i & r \\ j & r & \ell \\ r & m & s \end{bmatrix}.$$

The four 2×2 matrices in the corners of the 3×3 matrix J are T, T -matrices. From the matrix in the upper right corner we get that $r \geq_T \ell$. From the matrix in the lower left corner we get that $r \geq_T m$. From the matrix in the lower right corner we get, using Theorem 5.13 (4), that $\ell + m \geq_T s$. But $+$ is idempotent, \geq_T is compatible, and \geq_T is transitive on B , so we see that $r = r + r \geq_T \ell + m \geq_T s$. \square

THEOREM 5.22. *Let \mathcal{V} be a variety. The following are equivalent.*

- (1) *No member of \mathcal{V} has a nontrivial rectangular congruence.*
- (2) *No member of \mathcal{V} has a nontrivial rectangular tolerance.*

PROOF. Clearly (2) \implies (1), since congruences are tolerances. The same type of argument as that used in the proof of Lemma 3.5 shows that to prove the converse it is enough to show that

$$(x, y) \in [\theta, \theta]_{\mathbf{R}} \implies (x, y) \in [T, T]_{\mathbf{R}}$$

where $\theta = \text{Cg}^{\mathbf{F}}(x, y)$ and $T = \text{Tg}^{\mathbf{F}}(x, y)$ are the congruence and tolerance generated by (x, y) in the free algebra $\mathbf{F} = \mathbf{F}_{\mathcal{V}}(x, y)$. So assume that $(x, y) \in [\theta, \theta]_{\mathbf{R}}$. This assumption, which is equivalent to the assumption that \mathcal{V} has no member with a rectangular congruence, implies that \mathcal{V} has no algebra with a strongly rectangular congruence. Theorems 3.20 and 3.22 now guarantee that \mathcal{V} has a join term. This puts us into position to use Theorem 5.17. By part (4) of that theorem, the condition that $(x, y) \in [\theta, \theta]_{\mathbf{R}}$ means that (x, y) lies in the transitive closure of \sim_{θ} . But according to Lemma 3.15 (5), the fact that θ is a congruence implies that \sim_{θ} is transitive. Therefore $x \sim_{\theta} y$. This implies that $x \geq_{\theta} y$ and $y \geq_{\theta} x$, so by Lemma 5.21 we get that $x \geq_T y$ and $y \geq_T x$. Since $B = x/\theta$ is the unique T -block containing x and y we have $x \geq_B y$ and $y \geq_B x$; i.e., $x \sim_B y$ and hence $x \sim_T y$. Theorem 5.17 (4) now guarantees that $(x, y) \in [T, T]_{\mathbf{R}}$. \square

The proof of Theorem 5.22 shows that a variety \mathcal{V} contains no member with a rectangular tolerance or congruence if and only if $x \sim_T y$ in $\mathbf{F} = \mathbf{F}_{\mathcal{V}}(x, y)$ for $T = \text{Tg}^{\mathbf{F}}(x, y)$. This property is easily seen to be expressible by an idempotent Maltsev condition, which we record here.

THEOREM 5.23. *Let \mathcal{V} be a variety. The following are equivalent.*

- (1) *No member of \mathcal{V} has a nontrivial rectangular tolerance.*
- (2) *$x \sim_T y$ in $\mathbf{F} = \mathbf{F}_{\mathcal{V}}(x, y)$ for $T = \text{Tg}^{\mathbf{F}}(x, y)$.*

- (3) There exists an $m \geq 1$, and sixary terms f_1, \dots, f_m such that \mathcal{V} satisfies the following identities:
 - (i) $x \approx f_1(y, y, x, x, x, y)$,
 - (ii) $f_m(y, x, x, x, y, y) \approx y$,
 - (iii) $f_i(x, x, x, y, y, y) \approx f_i(y, y, x, x, x, y)$, $1 \leq i \leq m$
 - (iv) $f_i(y, x, x, x, y, y) \approx f_{i+1}(y, y, x, x, x, y)$, $1 \leq i \leq m - 1$.
- (4) For any algebra $\mathbf{A} \in \mathcal{V}$, tolerance S on \mathbf{A} , and $(a, b) \in S$, $a \geq_S b$ holds.
- (5) For any algebra $\mathbf{A} \in \mathcal{V}$, tolerance S on \mathbf{A} , and $(a, b) \in S$, $a \sim_S b$ holds.

PROOF. The proof of Theorem 5.22 shows that (1) and (2) are equivalent.

lower case "i"

If (2) holds, then $x \geq_T y$ in \mathbf{F} , so according to Lemma 3.19 there exists a sequence f_1, \dots, f_m of sixary terms satisfying the identities in (3).

Now suppose that (3) holds. Choose $\mathbf{A} \in \mathcal{V}$, a tolerance S on \mathbf{A} , an S -block B and $a, b \in B$. Since B is a block it is closed under all the idempotent terms that appear in item (3). By substituting a and b for x and y in the identities of (3) we obtain elements of S witnessing that $a \geq_B b$. Namely, the fact that $(a, b) \in S$ implies that for each i the matrix

$$\begin{bmatrix} p_i & q_i \\ r_i & s_i \end{bmatrix} = \begin{bmatrix} f_i(a, b, a, b, a, b) & f_i(a, a, a, b, b, b) \\ f_i(b, b, a, a, a, b) & f_i(b, a, a, a, b, b) \end{bmatrix}$$

remove "i"

"i" not "r"

is in $M(S, S)$. Identities of type (iii) imply that $q_i = r_i$, so each (r_i, s_i) is an $(S, S; 0)$ pair. The other identities imply that the sequence of these pairs is a connected chain from a to b . Thus $a \geq_S b$, and item (4) is established.

"b"

The argument of the previous paragraph showed that $a \geq_B b$ for some S -block B , and a similar one shows that $b \geq_B a$ for the same block B . Hence $a \sim_B b$, which implies that $a \sim_S b$. This establishes item (5). □

and clearly, (5) implies (2).

Our next task will be to connect the foregoing with the Hobby-McKenzie term. First we need a lemma.

LEMMA 5.24. *If \mathbf{C} is a nontrivial idempotent algebra with a compatible semilattice term, then there is an algebra \mathbf{D} that is a homomorphic image of a subalgebra of \mathbf{C} and is term equivalent to the 2-element semilattice.*

PROOF. Denote the semilattice operation by $+$, and call it join. Since the semilattice order is equationally definable by

$$x \geq y \iff x = x + y,$$

and $+$ is compatible, it follows that the semilattice order is a compatible relation. As C is not a singleton, there exist two distinct elements $0 < 1$ of \mathbf{C} . If c_1, \dots, c_k belong to the interval $[0, 1]$ (with respect to the semilattice ordering), and f is a term operation of \mathbf{C} , then

$$1 = f(1, \dots, 1) \geq f(c_1, \dots, c_k) \geq f(0, \dots, 0) = 0$$

by the idempotence of f and the compatibility of \geq . Thus $[0, 1]$ is a nontrivial subalgebra of \mathbf{C} . Replacing \mathbf{C} by this subalgebra, we may henceforth assume that \mathbf{C} has least and largest elements 0 and 1 .

If f is a k -ary term operation of \mathbf{C} , then since $+$ is compatible the operation f must commute with $x_1 + \dots + x_k$ on the matrix

$$\begin{bmatrix} x_1 & 0 & \cdots & 0 \\ 0 & x_2 & \cdots & 0 \\ \vdots & \vdots & & \vdots \\ 0 & 0 & \cdots & x_k \end{bmatrix}.$$

This says exactly that $f(x_1, \dots, x_k) = f_1(x_1) + \dots + f_k(x_k)$ where $f_i(x_i) := f(0, \dots, 0, x_i, 0, \dots, 0)$, with x_i in the i -th position. As f is idempotent, we have that

$$f_1(x) + \dots + f_n(x) = x$$

for every x . In particular, $f_i(x) \leq x$ for every x . The functions f_i are of course $+$ -endomorphisms, too, since f is compatible with $+$ and $0 + 0 = 0$.

Choose a maximal semilattice ideal I of \mathbf{C} that does not contain 1 . Let $F = C - I$ be the complementary filter. Let θ be the equivalence relation on C whose classes are I and F . We claim that θ is a congruence of \mathbf{C} , and that \mathbf{C}/θ is term equivalent to a semilattice. To show this, represent an arbitrarily chosen term f as a sum of unary polynomials as described in the previous paragraph. The condition $f_i(x) \leq x$ implies that $f_i(I) \subseteq I$. We prove that either $f_i(F) \subseteq F$, or $f_i(F) \subseteq I$. Assume that $f_i(F) \not\subseteq F$, and choose $u \in F$ so that $f_i(u) \in I$. By the maximality of I , there exists an element $a \in I$ such that $a + u = 1$. Since f_i is a $+$ -endomorphism we have $f_i(1) = f_i(a) + f_i(u) \in I$. This shows that if $f_i(u) \in I$, then $f_i(1) \in I$, hence $f_i(C) \subseteq I$. Consequently each f_i preserves θ . Since $+$ also preserves θ , it follows that f preserves it, so θ is a congruence of \mathbf{C} .

The quotient $\mathbf{D} = \mathbf{C}/\theta$ is a 2-element idempotent algebra with a compatible semilattice term $+$. Any such algebra is term equivalent to the 2-element semilattice, as we now show. Suppose that the elements of D are denoted 0 and 1, and that $0 < 1$ in the $+$ -order. As above, each term operation of \mathbf{D} may be represented as $f(x_1, \dots, x_k) = f_1(x_1) + \dots + f_k(x_k)$ where each f_i is a decreasing $+$ -endomorphism. But since $D = \{0, 1\}$ this forces each f_i to be the identity function or the constant function with value 0. If $J \subseteq \{1, \dots, k\}$ is the set of indices j such that f_j is the identity, then $J \neq \emptyset$, since f is idempotent, and $f(x_1, \dots, x_k) = \sum_J x_j$. Hence f equals a semilattice term operation. \square

Now we are in position to give a second Maltsev characterization of the class of varieties that omit rectangular tolerances.

THEOREM 5.25. *Let \mathcal{V} be a variety. The following are equivalent.*

- (1) *No member of \mathcal{V} has a nontrivial rectangular tolerance.*
- (2) *\mathcal{V} satisfies an idempotent (linear) Maltsev condition that fails in the variety of semilattices.*
- (3) *\mathcal{V} has a Hobby–McKenzie term.*

PROOF. The equivalence of (2) and (3) is Theorem 2.16.

To show that (1) \implies (2) we argue that the idempotent linear Maltsev condition from Theorem 5.23 (3) fails in the variety of semilattices. We give two proofs of this.

First, the Maltsev condition from Theorem 5.23 holds only in varieties omitting rectangular congruences. A calculation shows that the largest congruence on the 2-element semilattice is rectangular. Therefore the Maltsev condition fails in the variety of semilattices.

For the second proof, observe that if $f_1(x_1, \dots, x_6)$ is a semilattice term, then identity (i) of Theorem 5.23 (3), $x \approx f_1(y, y, x, x, x, y)$, implies that f_1 does not involve the variables x_1, x_2 or x_6 . (Otherwise, by putting the absorbing element of the 2-element semilattice in for these variables and the nonabsorbing element for the other variables one obtains the contradiction that these two elements are equal.) This, together with the first instance of identity (iii), $f_1(x, x, x, y, y, y) \approx f_1(y, y, x, x, x, y) (\approx x)$, shows that f_1 does not involve x_4, x_5 or x_6 either. Therefore, by idempotence, $f_1(x_1, \dots, x_6) \approx x_3$. This argument propagates down the chain of f_i 's via identity (iv), showing that each term satisfies $f_i(x_1, \dots, x_6) \approx x_3$. Identity (ii) then reduces to $x \approx y$, which fails in the variety of semilattices.

Now we prove (2) \implies (1). Assume that \mathcal{V} satisfies a (not necessarily linear) idempotent Maltsev condition, called W , that fails in the variety of semilattices. Let $T = \text{Tg}(x, y)$ be the tolerance generated

remove

by (x, y) on the free algebra $\mathbf{F} = \mathbf{F}_{\mathcal{V}}(x, y)$, and let B be the T -block of idempotent binary terms. B is closed under all idempotent terms of \mathcal{V} , so let \mathbf{B} be the algebra with universe B and basic operations all idempotent terms of \mathcal{V} . \mathbf{B} is an idempotent algebra that satisfies the idempotent Maltsev condition W . The relation \sim_T restricted to B is a congruence θ of \mathbf{B} .

Since the Maltsev condition W is nontrivial, the variety generated by \mathbf{B} has a join term $+$. By Theorem 5.13 (2) and (5), $+$ is a compatible semilattice term of the idempotent algebra $\mathbf{C} = \mathbf{B}/\theta$. If \mathbf{C} is nontrivial, then Lemma 5.24 guarantees that some homomorphic image \mathbf{D} of some subalgebra \mathbf{C} is term equivalent to the 2-element semilattice. The Maltsev condition W is satisfied by \mathbf{C} , hence by \mathbf{D} , hence by the variety of semilattices. But we have assumed the contrary, so we conclude that \mathbf{C} is trivial. This proves that \mathbf{B} is a single \sim_T class, and therefore that $x \sim_T y$ in $\mathbf{F}_{\mathcal{V}}(x, y)$. According to Theorem 5.23, this means that \mathcal{V} contains no algebra with a nontrivial rectangular tolerance. \square

In the preceding proof the implication (3) \implies (1) is verified indirectly. Before proceeding, we would like to describe a direct proof of this implication. Suppose that $f(x_1, \dots, x_n)$ is a Hobby–McKenzie term for \mathcal{V} . Recall from Chapter 2 that for a subset $U \subseteq N := \{1, \dots, n\}$ the expression $f_U(x, y)$ denotes the term obtained from f by substituting x for x_i if $i \in U$ and y for x_i otherwise. The statement that f is a Hobby–McKenzie term for \mathcal{V} is precisely the statement that f is idempotent and for any nonempty subset $U \subseteq N$ there exist subsets $V, W \subseteq N$ with $U \subseteq V$, $U \not\subseteq W$ and $\mathcal{V} \models f_V(x, y) \approx f_W(x, y)$.

As usual, let $\mathbf{F} = \mathbf{F}_{\mathcal{V}}(x, y)$, $T = \text{Tg}^{\mathbf{F}}(x, y)$, and B be the T -block of idempotent binary terms. Our aim is to prove that $x \sim_T y$ using the equations of the Hobby–McKenzie term. Recall that the restriction of \sim_T to B is transitive. Let $U \subseteq N$ be minimal under inclusion with respect to the property that $f_U(x, x+y) \sim_T x$. (Such a set U exists, since $f_N(x, x+y) = x \sim_T x$.) If U is not empty, then there exist $V, W \subseteq N$ such that $U \subseteq V$ but $U \not\subseteq W$, and $f_V(x, x+y) = f_W(x, x+y)$. Hence by Theorem 5.13,

$$(5.6) \quad f_U(x, x+y) + f_W(x, x+y) \sim_T f_{U \cap W}(x, x+y),$$

and

$$(5.7) \quad f_U(x, x+y) + f_V(x, x+y) \sim_T f_{U \cap V}(x, x+y) = f_U(x, x+y).$$

Since $f_V(x, x+y) = f_W(x, x+y)$, the left hand sides of lines (5.6) and (5.7) are the same. Therefore

$$f_{U \cap W}(x, x+y) \sim_T f_U(x, x+y) \sim_T x.$$

This contradicts the minimality of U unless $U = \emptyset$, which leads to the conclusion that $x \sim_T f_{\emptyset}(x, x+y) = x+y$. A symmetric argument shows that $y \sim_T x+y$, and therefore $x \sim_T y$. At this point, Theorem 5.23 may be invoked to deduce that no algebra in \mathcal{V} has a rectangular tolerance.

We turn to our third Maltsev characterization of the class of varieties that omit rectangular tolerances. This new Maltsev condition involves four variable terms, and somewhat resembles Alan Day's Maltsev condition characterizing the class of congruence modular varieties (cf. [8]).

DEFINITION 5.26. Let \mathbf{A} be an algebra. If $\mathbf{Con}(\mathbf{A})$ has a sublattice isomorphic to the lattice \mathbf{D}_2 with $(\theta \circ \mu) \cap (\nu \circ \delta) \not\subseteq \alpha$, then we call this

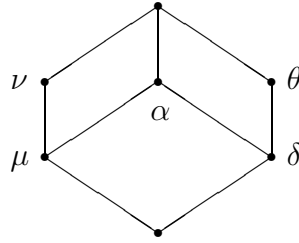


FIGURE 5.1. \mathbf{D}_2

sublattice a **special \mathbf{D}_2** in $\mathbf{Con}(\mathbf{A})$. We say that \mathcal{V} **omits special \mathbf{D}_2 's** if for no $\mathbf{A} \in \mathcal{V}$ does $\mathbf{Con}(\mathbf{A})$ have a sublattice that is a special \mathbf{D}_2 .

As an aid to recognizing copies of \mathbf{D}_2 in congruence lattices, we record the following.

LEMMA 5.27. *A presentation of \mathbf{D}_2 relative to the variety of all lattices is $\langle G \mid R \rangle$ where $G = \{x, y, z\}$ and R consists of the relations:*

- (i) $x \wedge z \leq y$,
- (ii) $x \leq z \vee (y \wedge x)$,
- (iii) $z \leq x \vee (y \wedge z)$, and
- (iv) $y \leq (x \wedge y) \vee (y \wedge z)$.

Moreover any lattice generated by G and satisfying the relations in R and also satisfying $x \not\leq y$ is isomorphic to \mathbf{D}_2 .

PROOF. The reader can easily verify that the copy of \mathbf{D}_2 in Figure 5.2 is generated by x, y and z , and that the relations in the statement of the lemma hold for these elements. To show that it is a presentation of \mathbf{D}_2 we must prove from the relations (and the laws of

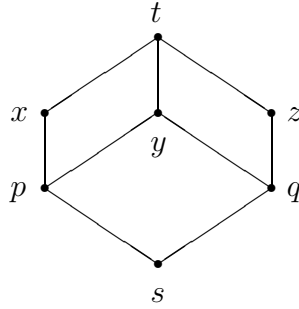


FIGURE 5.2.

lattice theory) that the seven elements $x, y, z, p := x \wedge y, q := z \wedge y, s := x \wedge y \wedge z$, and $t := x \vee y \vee z$ are closed under meet and join.

It is clear that the seven elements are ordered as they should be, and consequently s and t are the bottom and top elements of the sublattice they generate. The only incomparable doubletons are $\{x, y\}, \{x, z\}, \{x, q\}, \{y, z\}, \{p, z\}$ and $\{p, q\}$. By symmetry, we do not need to consider the meets or joins of $\{y, z\}$ or $\{p, z\}$. For the remaining cases:

- $\{x, y\}$: $x \wedge y = p$ by the definition of p . $x \vee y$ is above z by relation (iii). Since $x \vee y$ is formally below $x \vee y \vee z = t$, it must equal it.
- $\{x, z\}$: $x \wedge z$ is below y by relation (i). This implies $x \wedge z = x \wedge y \wedge z = s$, as desired. $x \vee z$ is above $(x \wedge y) \vee (y \wedge z)$, since corresponding joinands are larger, and $(x \wedge y) \vee (y \wedge z) \geq y$ by (iv). Therefore $x \vee z = x \vee y \vee z = t$.
- $\{x, q\}$: $x \wedge q = x \wedge (y \wedge z) = s$ by the definition of q . $x \vee q = x \vee (y \wedge z)$ is above z by (iii), therefore it is above $x \vee z = t$ and hence equal to t .
- $\{p, q\}$: $p \wedge q$ is below $x \wedge z$, which has already been shown to be the bottom element, so $p \wedge q = s$. By (iv) we have $y \leq (x \wedge y) \vee (y \vee z) = p \vee q \leq y$, so $p \vee q = y$.

This proves that $\langle G \mid R \rangle$ presents \mathbf{D}_2 . For the last claim of the lemma, any lattice generated by G satisfying the relations R which is not isomorphic to \mathbf{D}_2 must be a proper quotient of \mathbf{D}_2 . Since \mathbf{D}_2 is a subdirectly irreducible lattice whose least congruence collapses p and x , any such lattice will satisfy $x = p \leq y$. Therefore, if a lattice is generated by G and satisfies R but does not satisfy $x \leq y$, then the lattice is isomorphic to \mathbf{D}_2 . \square

THEOREM 5.28. *The following are equivalent for a variety \mathcal{V} .*

- (1) \mathcal{V} satisfies an idempotent Maltsev condition that fails in the variety of semilattices.
- (2) Whenever θ, μ, ν and δ are congruences on some $\mathbf{A} \in \mathcal{V}$, then $(\theta \circ \mu) \cap (\nu \circ \delta) \subseteq [(\theta \vee \delta) \wedge (\nu \vee \mu)] \vee \delta \vee \mu$.
- (3) \mathcal{V} has a sequence of terms $f_i(x, y, u, v)$, $0 \leq i \leq 2m + 1$, such that
 - (i) $\mathcal{V} \models f_0(x, y, u, v) \approx x$;
 - (ii) $\mathcal{V} \models f_i(x, y, y, y) \approx f_{i+1}(x, y, y, y)$ for even i ;
 - (iii) $\mathcal{V} \models f_i(x, x, y, y) \approx f_{i+1}(x, x, y, y)$ and $\mathcal{V} \models f_i(x, y, x, y) \approx f_{i+1}(x, y, x, y)$ for odd i ;
 - (iv) $\mathcal{V} \models f_{2m+1}(x, y, u, v) \approx v$.
- (4) \mathcal{V} omits special \mathbf{D}_2 's.

PROOF. We first establish the equivalence of (2), (3) and (4), and connect them to item (1) afterward.

Assume that (2) holds. On $\mathbf{F} = \mathbf{F}_{\mathcal{V}}(x, y, u, v)$ let $\theta = \text{Cg}^{\mathbf{F}}(x, u)$, $\mu = \text{Cg}^{\mathbf{F}}(u, v)$, $\nu = \text{Cg}^{\mathbf{F}}(x, y)$, and $\delta = \text{Cg}^{\mathbf{F}}(y, v)$. Since $(x, v) \in (\theta \circ \mu) \cap (\nu \circ \delta)$, and (2) holds:

$$(\theta \circ \mu) \cap (\nu \circ \delta) \subseteq [(\theta \vee \delta) \wedge (\nu \vee \mu)] \vee \delta \vee \mu,$$

we have $(x, v) \in [(\theta \vee \delta) \wedge (\nu \vee \mu)] \vee \delta \vee \mu$. Therefore, there is a sequence $f_0, f_1, \dots, f_{2m+1}$ of elements of \mathbf{F} such that $f_0 = x, f_{2m+1} = v, (f_i, f_{i+1}) \in \delta \vee \mu$ for even i while $(f_i, f_{i+1}) \in (\theta \vee \delta) \wedge (\nu \vee \mu)$ for odd i . If one chooses 4-ary terms representing the f_i 's, it is straightforward to see that the equations of (3) hold.

Assume that (3) holds and that some $\mathbf{A} \in \mathcal{V}$ has congruences μ, ν, α, δ and θ that generate a sublattice of $\mathbf{Con}(\mathbf{A})$ isomorphic to \mathbf{D}_2 labeled as in Figure 5.1. Choose any $(a, d) \in (\theta \circ \mu) \cap (\nu \circ \delta)$. There exist elements $b, c \in A$ such that $a \nu b \delta d$ and $a \theta c \mu d$. The pairs $(f_i(a, b, c, d), f_{i+1}(a, b, c, d))$ belong to $\delta \vee \mu = \alpha$ for even i and $(\theta \vee \delta) \wedge (\nu \vee \mu) \leq \alpha$ for odd i . Hence $a = f_0(a, b, c, d)$ is α -related to $f_{2m+1}(a, b, c, d) = d$. Since $(a, d) \in (\theta \circ \mu) \cap (\nu \circ \delta)$ was chosen arbitrarily, this implies that $(\theta \circ \mu) \cap (\nu \circ \delta) \subseteq \alpha$. Therefore this copy of \mathbf{D}_2 is not a special \mathbf{D}_2 . The argument works for every \mathbf{D}_2 that appears in a congruence lattice of a member of \mathcal{V} , so \mathcal{V} omits special \mathbf{D}_2 's.

Next we argue the contrapositive of (4) \implies (2), so assume that some $\mathbf{A} \in \mathcal{V}$ has congruences μ, ν, δ and θ such that

$$(\theta \circ \mu) \cap (\nu \circ \delta) \not\subseteq [(\theta \vee \delta) \wedge (\nu \vee \mu)] \vee \delta \vee \mu.$$

Choose $(a, d) \in (\theta \circ \mu) \cap (\nu \circ \delta)$ satisfying $(a, d) \notin [(\theta \vee \delta) \wedge (\nu \vee \mu)] \vee \delta \vee \mu$, as well as elements $b, c \in A$ such that $a \equiv_{\theta} b \not\equiv_{\mu} d$ and $a \equiv_{\nu} c \equiv_{\delta} d$. Let $\theta^* := \text{Cg}^{\mathbf{A}}((a, b), (c, d))$, $\nu^* := \text{Cg}^{\mathbf{A}}((a, c), (b, d))$, and $\alpha^* := \text{Cg}^{\mathbf{A}}((b, d), (c, d)) \vee (\theta^* \wedge \nu^*)$. We will show that the congruences

slight mismatch
in presentation
here and on the
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$x := \theta^*$, $y := \alpha^*$, and $z := \nu^*$ satisfy the presentation for \mathbf{D}_2 given in Lemma 5.27.

Item (i) of Lemma 5.27 requires that $\theta^* \wedge \nu^* \leq \alpha^*$, which is an immediate consequence of the definition of α^* .

Item (ii) requires that $\theta^* \leq \nu^* \vee (\alpha^* \wedge \theta^*)$. Since ν^* contains (a, c) and (b, d) while $\alpha^* \wedge \theta^*$ contains (c, d) , the join is the congruence generated by $X \times X$ where $X = \{a, b, c, d\}$. Hence the join contains the generators (a, b) and (c, d) of θ^* . Item (iii) is symmetric to item (ii).

For item (iv), the congruence $\theta^* \wedge \alpha^*$ contains $\{(c, d)\} \cup (\theta^* \wedge \nu^*)$, since this is true of each of θ^* and α^* . Similarly $\nu^* \wedge \alpha^*$ contains $\{(b, d)\} \cup (\theta^* \wedge \nu^*)$. This shows that $(\theta^* \wedge \alpha^*) \vee (\alpha^* \wedge \nu^*)$ contains all the generators of α^* . Hence $\alpha^* \leq (\theta^* \wedge \alpha^*) \vee (\alpha^* \wedge \nu^*)$.

Finally we must show that $\theta^* \not\leq \alpha^*$. Assume that this is not the case. Since $(a, b) \in \theta^*$, this assumption forces $(a, b) \in \alpha^*$. By definition, $\theta^* \leq (\theta \vee \delta)$, $\nu^* \leq (\nu \vee \mu)$, $\text{Cg}^{\mathbf{A}}(c, d) \leq \delta$ and $\text{Cg}^{\mathbf{A}}(b, d) \leq \mu$, so

$$\alpha^* = (\theta^* \wedge \nu^*) \vee \text{Cg}^{\mathbf{A}}(c, d) \vee \text{Cg}^{\mathbf{A}}(b, d) \leq [(\theta \vee \delta) \wedge (\nu \vee \mu)] \vee \delta \vee \mu.$$

Combining the facts that $(a, b) \in \alpha^*$ and $(b, d) \in \mu$, we obtain that the right hand side of the previous displayed line contains (a, d) . This is contrary to the choice of (a, d) , so we are forced to conclude that $\theta^* \not\leq \alpha^*$. This completes the proof that $\{\theta^*, \alpha^*, \nu^*\}$ generates a copy of \mathbf{D}_2 in $\mathbf{Con}(\mathbf{A})$. Indeed, this is a *special* \mathbf{D}_2 since $a \equiv_{\theta^*} b \equiv_{\alpha^* \wedge \nu^*} d$ and $a \equiv_{\nu^*} c \equiv_{\alpha^* \wedge \theta^*} d$, so

$$(a, d) \in (\theta^* \circ (\alpha^* \wedge \nu^*)) \cap (\nu^* \circ (\alpha^* \wedge \theta^*)) - \alpha^*.$$

This proves that item (4) of the theorem fails to hold.

Now we turn our attention to item (1). If $\mathbf{2} = \langle \{0, 1\}; + \rangle$ is the 2-element join semilattice, $\theta, \nu \in \mathbf{Con}(\mathbf{2} \times \mathbf{2})$ are the projection kernels, and $\alpha \in \mathbf{Con}(\mathbf{2} \times \mathbf{2})$ is the kernel of the homomorphism $+: \mathbf{2}^2 \rightarrow \mathbf{2}$, then θ, ν and α generate a special \mathbf{D}_2 . Therefore the variety of semilattices does not satisfy the idempotent Maltsev condition in (3). This shows that (3) \implies (1). To complete the proof, we must prove that (1) \implies (3).

Both (1) and (3) are statements about the idempotent Maltsev conditions satisfied by \mathcal{V} . The idempotent reduct $\text{Id}(\mathcal{V})$ satisfies the same idempotent Maltsev conditions as \mathcal{V} , so to prove (1) \implies (3) we may (and do) assume that \mathcal{V} is idempotent.

Suppose that (3) fails. Then \mathcal{V} does not have terms satisfying the identities listed in (3), so the method that we used to construct those terms in our proof that (2) \implies (3) must fail. Recall from that part of the proof that we defined $\mathbf{F} = \mathbf{F}_{\mathcal{V}}(x, y, u, v)$, $\theta = \text{Cg}^{\mathbf{F}}(x, u)$, $\mu = \text{Cg}^{\mathbf{F}}(u, v)$, $\nu = \text{Cg}^{\mathbf{F}}(x, y)$, and $\delta = \text{Cg}^{\mathbf{F}}(y, v)$. We argued that if $(x, v) \in [(\theta \vee \delta) \wedge (\nu \vee \mu)] \vee \delta \vee \mu$, then (3) holds. Therefore, in our current situation, the

isn't τ_2 a transversal for E_1 ?
 τ_1 isn't a transversal for E_2 .
 E_2 should be (x,y) and (u,v)

congruence $\alpha := [(\theta \vee \delta) \wedge (\nu \vee \mu)] \vee \delta \vee \mu$ does not contain the pair (x, v) . Our plan is to use this fact to construct a nontrivial algebra in \mathcal{V} that has a compatible semilattice operation.

Let $X = \{x, y, u, v\}$, so that $\mathbf{F} = \mathbf{F}_{\mathcal{V}}(X)$. Let $\mathbf{G} = \mathbf{F}_{\mathcal{V}}(x, v)$ be the subalgebra of \mathbf{F} generated by $\{x, v\}$. The functions $\tau_i: X \rightarrow X$, $i = 1, 2$, defined by $\tau_1(x) = \tau_1(y) = x = \tau_2(x) = \tau_2(u)$ and $\tau_1(u) = \tau_1(v) = v = \tau_2(y) = \tau_2(v)$ are transversals for the equivalence relations $E_1 = \text{Eq}^X((x, u), (y, v))$ and $E_2 = \text{Eq}^X((x, v), (y, u))$ respectively. These transversals induce homomorphisms $h_i: \mathbf{F} \rightarrow \mathbf{F}$ whose kernels are $\theta \vee \delta$ and $\nu \vee \mu$ respectively (according to Lemma 3.6 (1)), and whose images are both \mathbf{G} . Therefore the kernel of

$$h := h_1 \times h_2: \mathbf{F} \rightarrow \mathbf{G} \times \mathbf{G}: t(x, y, u, v) \mapsto (t(x, x, v, v), t(x, v, x, v))$$

is $[(\theta \vee \delta) \wedge (\nu \vee \mu)]$. The idempotence of \mathcal{V} implies that this homomorphism is surjective. Namely, if $(p, q) = (p(x, v), q(x, v)) \in G \times G$, then $(p \circ q)(x, y, u, v) := p(q(x, y), q(u, v)) \in F$ is a preimage of (p, q) under h . This shows that \mathbf{G} is generated by the four elements $h(x) = (x, x)$, $h(y) = (x, v)$, $h(u) = (v, x)$, and $h(v) = (v, v)$. Since the kernel of h is contained in $\alpha = [(\theta \vee \delta) \wedge (\nu \vee \mu)] \vee \delta \vee \mu$, we may use h to recast our earlier assumption that $(x, v) \notin \alpha$ as a statement about $\mathbf{G} \times \mathbf{G}$. That assumption is equivalent to the assumption that $h(x, v) = ((x, x), (v, v))$ does not belong to the congruence

$$\begin{aligned} h(\alpha) &= h([(\theta \vee \delta) \wedge (\nu \vee \mu)] \vee \delta \vee \mu) \\ &= \text{Cg}^{\mathbf{G} \times \mathbf{G}}(h(\delta \vee \mu)) \\ &= \text{Cg}^{\mathbf{G} \times \mathbf{G}}(h(y, v), h(u, v)) \\ &= \text{Cg}^{\mathbf{G} \times \mathbf{G}}(((x, v), (v, v)), ((v, x), (v, v))). \end{aligned}$$

Therefore, if we let $\beta = h(\alpha)$, then $\beta = \text{Cg}^{\mathbf{G} \times \mathbf{G}}(Y \times Y)$ where $Y = \{(x, v), (v, x), (v, v)\}$, and our earlier assumption that $(x, v) \notin \alpha$ in \mathbf{F} is equivalent to the assumption that $((x, x), (v, v)) \notin \beta$ in $\mathbf{G} \times \mathbf{G}$.

For any $(p, q) \in G^2$ there is an $r \in G$ such that $(p, q) \beta (r, r)$. Indeed, if $p = p(x, v)$ and $q = q(x, v)$, then for $r = p(q(x, v), v)$ we have

$$\begin{aligned} (p, q) &= (p(x, v), q(x, v)) \\ &= (p(q(x, x), q(v, v)), p(q(x, v), q(x, v))) \\ &= (p \circ q)((x, x), (x, v), (v, x), (v, v)) \\ (5.8) \quad &\equiv_{\beta} (p \circ q)((x, x), (v, v), (v, v), (v, v)) \\ &= (p(q(x, v), q(v, v)), p(q(x, v), q(v, v))) \\ &= (p(q(x, v), v), p(q(x, v), v)) \\ &= (r, r). \end{aligned}$$

this needs to be fixed too and references to h in the proof need to be adjusted, depending on how E_1 and E_2 are defined.

Of course, the element r need not be unique, since it may happen that $(p, q) \beta (r, r) \beta (s, s)$ for some $s \neq r$. To account for this nonuniqueness, let γ be the congruence on \mathbf{G} defined by

$$r \gamma s \iff (r, r) \beta (s, s).$$

This is indeed a congruence, as it is $\iota^{-1}(\beta)$ for $\iota: \mathbf{G} \rightarrow \mathbf{G}^2: p \mapsto (p, p)$ equal to the diagonal embedding. Write \overline{G} for G/γ and \overline{p} for p/γ if $p \in G$. In this notation, we have shown in (5.8) that there is a well defined function

$$*: G \times G \rightarrow \overline{G}: (p, q) \mapsto \overline{r} \text{ if } (p, q) \beta (r, r).$$

Explicitly, the value of $p * q$ is $\overline{p(q(x, v), v)}$, as we have shown.

If $q \gamma q'$, then

$$p * q = p(q(x, v), v) \gamma p(q'(x, v), v) = p * q',$$

since p is a term and γ is a congruence. Thus, left $*$ -multiplication by any $p \in G$ is compatible with γ . Since the generators of β are invariant under switching coordinates, the property $(p, q) \beta (r, r)$ is equivalent to $(q, p) \beta (r, r)$, so $p * q = q * p$. This commutativity implies that right $*$ -multiplication is also compatible with γ . Consequently $*$ induces a well defined binary operation (also called $*$) on \overline{G} :

$$*: \overline{G} \times \overline{G} \rightarrow \overline{G}: (\overline{p}, \overline{q}) \mapsto \overline{r} \text{ if } (p, q) \beta (r, r).$$

We have already established the commutativity of this operation. The idempotence of $*$ follows from the fact that $(p, p) \beta (p, p)$ for any $p \in G$. Associativity is clear, too, since

$$(\overline{p} * \overline{q}) * \overline{r} = \overline{p(q(r(x, v), v), v)} = \overline{p} * (\overline{q} * \overline{r}).$$

Thus $*$ is a semilattice operation on the set \overline{G} . Moreover, $*$ is a compatible semilattice operation of the algebra $\overline{\mathbf{G}} := \mathbf{G}/\gamma$. For suppose that f is an n -ary basic operation of \overline{G} , and that $\overline{p}_i, \overline{q}_i \in \overline{G}$ are chosen arbitrarily. Choose r_i such that $(p_i, q_i) \beta (r_i, r_i)$ for all i . Then, since β is a congruence, $(f(\mathbf{p}), f(\mathbf{q})) \beta (f(\mathbf{r}), f(\mathbf{r}))$. This shows that

$$f(\overline{p}) * f(\overline{q}) = f(\overline{r}) = f(\overline{p} * \overline{q}).$$

Finally (and critically), $\overline{\mathbf{G}}$ is nontrivial. Indeed, since $((x, x), (v, v)) \notin \beta$, we get that $(x, v) \notin \gamma$, so \overline{x} and \overline{v} are distinct elements of $\overline{\mathbf{G}}$.

This proves that $\overline{\mathbf{G}} \in \mathcal{V}$ is a nontrivial algebra with a compatible semilattice operation. Theorem 5.6 applies to show that the total relation on $\overline{\mathbf{G}}$ is rectangular. From Theorem 5.25 we derive that any idempotent Maltsev condition satisfied by \mathcal{V} is also satisfied by the variety of semilattices. Thus (1) fails, and the proof is complete. \square

strictly speaking,
 $p * q$ is not a
 member of \mathbf{G} , but
 is rather a γ -
 class.

In Theorem 8.11 we will learn that if \mathcal{V} satisfies an idempotent Maltsev condition that fails in the variety of semilattices, then \mathcal{V} omits all \mathbf{D}_2 's, not just the special ones.

CHAPTER 6

A Theory of Solvability

A normal subgroup N of a group G is **solvable** if there is a finite chain of normal subgroups of G ,

$$\{1\} = N_0 \subseteq N_1 \subseteq \cdots \subseteq N_k = N,$$

such that each factor N_{i+1}/N_i is abelian, or equivalently such that $[N_{i+1}, N_{i+1}] \subseteq N_i$ for all i . An interval $I[M, N]$ in the normal subgroup lattice of G is a **solvable interval** if N/M is a solvable normal subgroup of G/M . Normal subgroups M and N of G are **solvably related**, written $M \overset{s}{\sim} N$, if they belong to a solvable interval. We will call the relation $\overset{s}{\sim}$ “the solvability relation”, and refer to the study of this solvability relation as the “theory of solvability” for groups. Among the important features of the theory of solvability for groups are:

- The solvability relation is a congruence on the normal subgroup lattice of any group.
- The solvability relation is preserved by homomorphisms in the sense that if M, N and P are normal subgroups of G with $P \leq M \cap N$, then M is solvably related to N if and only if M/P is solvably related to N/P in G/P .
- If N is a minimal normal subgroup of G , then N is solvable if and only if it is abelian.

The ring-theoretic concept that corresponds directly to that of a solvable normal subgroup is that of a nilpotent ideal. In the early days of ring theory the concept of solvability was developed in analogy with the group theoretical concept via the Wedderburn radical (= the largest nilpotent ideal, if it exists). Other radicals soon made their appearance, which permitted new approaches to solvability not bound to *finite* chains of congruences with abelian factors. Moreover, in the case of some radicals, the resulting “solvability theory” was no longer bound to the abelian property. For example, if \mathbf{R} is a subdirectly irreducible ring that is not simple, then the Brown-McCoy radical (= the intersection of maximal two-sided ideals) contains the monolith *whether or*

not the monolith is abelian. Thus, it can happen that a minimal non-abelian congruence is contained in the Brown-McCoy radical, hence is “solvable” in the Brown-McCoy sense.

A third approach to solvability, valid for arbitrary finite algebras, is developed by D. Hobby and R. McKenzie in Chapter 7 of [34]. This approach does not use chains of congruences with abelian factors nor does it use radicals to define the solvability concept. Rather, one starts by defining a pair $(0, 1) \in A^2$ from an algebra \mathbf{A} to be a **2-snag** if there is a polynomial $p(x, y) \in \text{Pol}_2(\mathbf{A})$ such that $p(0, 0) = p(0, 1) = p(1, 0) = 0$ and $p(1, 1) = 1$. Hobby and McKenzie observe that if \mathbf{A} is a finite algebra and $\alpha \prec \beta$ are congruences on \mathbf{A} , then $\mathbf{C}(\beta, \beta; \alpha)$ holds if and only if $\beta - \alpha$ contains no 2-snag. From this it is clear that a chain of congruences:

$$\alpha = \alpha_0 \prec \alpha_1 \prec \cdots \prec \alpha_k = \beta$$

has abelian factors α_{i+1}/α_i if and only if $\beta - \alpha$ contains no 2-snag. Starting from this observation, it can be argued that the relation $\alpha \stackrel{s}{\sim} \beta$, defined by the condition that α and β contain the same 2-snags, is a complete congruence on $\mathbf{Con}(\mathbf{A})$ for any locally finite algebra.

Our aim in this chapter is to synthesize ideas from the solvability theories for groups, rings, and finite algebras in order to develop a solvability theory for a wide class of varieties. The theory resembles the theory for groups in that it is defined in terms of chains of congruences with abelian factors. However, as we will show in Section 6.3, the theory may be viewed as an extension of the theory of solvability for rings that is based on the prime radical ($=$ the intersection of all prime congruences), since for rings our theory defines an ideal to be solvable precisely when it is contained in the prime radical. Finally, anyone familiar with the theory developed by Hobby and McKenzie in [34] will recognize that our “solvability obstructions” (Definition 6.5) play the role of their 2-snags, while Theorems 6.17, 6.23 and 6.25 of this chapter are analogues of Theorem 7.9 (3), Corollary 7.13 and Lemma 6.5 of [34].

6.1. Varieties With A Weak Difference Term

There is an obstacle to the development of a very general solvability theory that we must deal with first. Any absolutely free algebra in any given signature is abelian in the sense of Chapter 2. If there were a solvability theory that applied to all such algebras, where the solvability relation is a congruence preserved by homomorphisms, then such a theory must identify every congruence on every algebra as solvable. Such a theory would, of course, be trivial. We cannot avoid this type of problem even if we restrict our attention to varieties satisfying

consider replacing "observe" with "prove", since a fair amount of work is needed to establish this connection.

a nontrivial idempotent Maltsev condition, since there exist algebras (such as $\mathbf{A} = \langle [0, 1]; \frac{x+y}{2} \rangle$, the unit interval of the real line under the averaging operation) that are abelian in the sense of Chapter 2, have a Taylor term ($f(x, y) = \frac{x+y}{2}$), but which have simple nonabelian homomorphic images (in this case, \mathbf{A}/θ where $\theta = [0, 1) \times [0, 1) \cup \{(1, 1)\}$). This phenomenon, that solvable algebras have nonsolvable homomorphic images, prevents the development of a powerful solvability theory.

As it happens, there is a weakest idempotent Maltsev condition that guarantees that homomorphic images of solvable algebras are solvable, and this Maltsev condition is a strong enough hypothesis for the theorems that we have in mind. We start this chapter by discussing the idempotent Maltsev condition on which the solvability theory is based, and then we discuss the theory.

DEFINITION 6.1. Let \mathcal{V} be a variety. A ternary term $d(x, y, z)$ is a **weak difference term** for \mathcal{V} if whenever $\mathbf{A} \in \mathcal{V}$, $\theta \in \text{Con}(\mathbf{A})$ and $(a, b) \in \theta$ it is the case that

$$d(a, a, b) \equiv b \equiv d(b, a, a) \pmod{[\theta, \theta]}.$$

The following theorem is included to show that our development of a solvability theory only for varieties with a weak difference term is reasonable.

THEOREM 6.2. *The class of varieties with a weak difference term is definable by a nontrivial idempotent Maltsev condition. This idempotent Maltsev condition is the weakest conjunction of idempotent Maltsev conditions that implies that homomorphic images of solvable algebras are solvable. It is also the weakest conjunction of idempotent Maltsev conditions that implies that abelian algebras are affine.*

PROOF. The first and third statements are proved in Theorem 4.8 of [52]. We prove the second only.

If \mathcal{V} has a weak difference term, then it can be shown (by using Lemma 6.8 and following the same argument that works for groups) that homomorphic images of solvable algebras are solvable.

Conversely assume that \mathcal{V} satisfies some conjunction of idempotent Maltsev conditions that implies that homomorphic images of solvable algebras are solvable. At least one of these Maltsev conditions must be nontrivial, since there exist solvable algebras with nonsolvable homomorphic images (e.g. $\langle [0, 1]; \frac{x+y}{2} \rangle$). The idempotent reduct $\text{Id}(\mathcal{V})$ of \mathcal{V} satisfies the same idempotent Maltsev conditions as \mathcal{V} , so homomorphic images of solvable algebras in $\text{Id}(\mathcal{V})$ are again solvable. Let \mathcal{A} be the subvariety of $\text{Id}(\mathcal{V})$ that is generated by the abelian algebras in

$\text{Id}(\mathcal{V})$. Since the class of abelian algebras in any language is axiomatizable by the universal Horn sentences of the term condition, this class is closed under the formation of subalgebras and products. From our assumption that homomorphic images of solvable algebras are solvable we get that \mathcal{A} consists of solvable algebras. In particular, \mathcal{A} contains no subvariety term equivalent to the variety of semilattices. Since \mathcal{A} also satisfies the nontrivial idempotent Maltsev conditions satisfied by \mathcal{V} , it follows from Lemma 2.5 of [47] that \mathcal{A} must satisfy an idempotent Maltsev condition that fails in the variety of semilattices. From Theorem 4.10 of [52], it follows that the variety \mathcal{A} is affine, and in particular congruence permutable. If $d(x, y, z)$ is a Maltsev operation for \mathcal{A} , then it will be a weak difference term for \mathcal{V} , as is explained in the proof of Theorem 4.8 of [52]. \square

This theorem has a corollary which will prove useful in what follows.

COROLLARY 6.3. *If a variety has a Hobby–McKenzie term, then it has a weak difference term.*

PROOF. Theorem 4.10 of [52] proves that if \mathcal{V} satisfies an idempotent Maltsev condition that fails in the variety of semilattices, then abelian algebras in \mathcal{V} are affine. Theorem 5.25 proves that the hypothesis of this statement is equivalent to the existence of a Hobby–McKenzie term, while Theorem 6.2 proves that the conclusion of the statement implies the existence of a weak difference term. \square

The concept of a weak difference term entered mathematics during the development of commutator theory. Recall from the introduction that A. I. Maltsev showed in [62] that a variety \mathcal{V} has permuting congruences if and only if it has a term $d(x, y, z)$ such that $\mathcal{V} \models d(x, x, y) \approx y \approx d(y, x, x)$. In [31], C. Herrmann proved that abelian algebras in a congruence modular variety \mathcal{V} are affine by constructing a term $d(x, y, z)$ from the Day terms for congruence modularity and showing that $d(x, y, z)$ is a Maltsev term for any abelian algebra in \mathcal{V} .¹ Then H. P. Gumm observed in [27] that the term constructed by Herrmann has the stronger properties that the identity $d(x, x, y) \approx y$ holds in \mathcal{V} , and $d(b, a, a) \equiv b \pmod{[\theta, \theta]}$ for any $\mathbf{A} \in \mathcal{V}$, $\theta \in \text{Con}(\mathbf{A})$, and $(a, b) \in \theta$. Such a term is now called a **difference term**. Difference terms play a central role in most developments of modular commutator theory and in some extensions of the theory to nonmodular varieties (see [45, 56, 58]). Weak difference terms were

explain why the term "difference" is used for this kind of term.

¹This fact, together with the fact that the class of congruence modular varieties is definable by an idempotent Maltsev condition, already implies that $d(x, y, z)$ is a weak difference term for \mathcal{V} , as the proof of Theorem 6.2 shows.

first identified in [34]. They occur far more frequently than ordinary difference terms, and are equally useful in most circumstances.

EXAMPLES 6.4. As noted in Theorem 6.2, any variety with a weak difference term satisfies a nontrivial idempotent Maltsev condition, and therefore has a Taylor term. The converse is not true in general, as we explained in the first paragraph of this section, but the converse does hold for certain classes of varieties. We describe some of these classes now, thereby providing a number of examples of varieties with weak difference terms.

Theorem 9.6 of [34] proves the converse for locally finite varieties: that is, any locally finite variety that has a Taylor term has a weak difference term.

Any idempotent abelian (or solvable) variety that has a Taylor term has a weak difference term. (The proof of this is contained in the proof of Theorem 6.2.)

A variety of semigroups that has a Taylor term has a weak difference term. For if $f(x_1, \dots, x_n) = x_{i_0} \cdots x_{i_k}$ is a Taylor term for \mathcal{V} , then necessarily $k \geq 1$. The idempotence of f implies that $\mathcal{V} \models x^{k+1} \approx x$. Any semigroup \mathbf{S} satisfying this equation has the property that $(s^k)^2 = s^{k-1}s^{k+1} = s^{k-1}s = s^k$ for any $s \in S$, so k -th powers of elements of \mathbf{S} are idempotent. If $e, f \in S$ are idempotent, and $ef \neq fe$, then it can be shown that the subsemigroup of \mathbf{S} generated by $\{e, f\}$ has a homomorphism onto a nontrivial left zero or right zero semigroup. But a nontrivial left zero semigroup fails the i_0 -th Taylor identity for $f(x_1, \dots, x_n) = x_{i_0} \cdots x_{i_k}$, and a nontrivial right zero semigroup fails the i_k -th Taylor identity. It follows that \mathcal{V} has no nontrivial left or right zero semigroups, and hence that no $\mathbf{S} \in \mathcal{V}$ has a pair of non-commuting idempotents. This fact may be expressed by saying that the identity $x^k y^k \approx y^k x^k$ holds in \mathcal{V} . Now we argue that the term $d(x, y, z) = xy^{2k-1}z$ is a weak difference term for the variety of semigroups axiomatized by $x^{k+1} \approx x$ and $x^k y^k \approx y^k x^k$. To justify this, it suffices to show that if \mathbf{A} is a semigroup in \mathcal{V} and $a, b \in A$ are elements for which $\theta := \text{Cg}^{\mathbf{A}}(a, b)$ is abelian, then $d(a, a, b) = b = d(b, a, a)$. To show that $b = d(a, a, b) = a^{2k}b = a^k b$ and $b = ba^k$ it is enough to prove that $a^k = b^k$. Since $e = a^k$ and $f = b^k$ are idempotent and $(e, f) = (a^k, b^k) \in \theta$, it suffices to prove that idempotents related by an abelian congruence are equal. The θ, θ -term condition applied to $\underline{e}e\underline{f} = e\underline{f} = e\underline{f}f$ allows us to replace the underlined f with e and obtain $\underline{e}ee = e\underline{f}e$. Since idempotents commute, this is equivalent to $e = ef$. The same argument shows that $f = ef$, so $e = f$.

It is easily seen that any congruence n -permutable variety satisfies a nontrivial idempotent Maltsev condition, hence has a Taylor term. Less obvious is the theorem of P. Lipparini in [57] that any congruence n -permutable variety has a weak difference term.

6.2. ∞ -Solvability

We now begin the development of a solvability theory for varieties with a weak difference term. Any successor ordinal, considered as a lattice under its usual ordering, is a meet continuous lattice. By a **continuous sequence of congruences** on an algebra \mathbf{A} we mean a meet continuous lattice homomorphism $\sigma : \kappa + 1 \rightarrow \mathbf{Con}(\mathbf{A})$. We write such a sequence as $(\sigma_\lambda)_{\lambda < \kappa+1}$ where $\sigma_\lambda := \sigma(\lambda)$. It is not hard to verify that $(\sigma_\lambda)_{\lambda < \kappa+1}$ is a continuous sequence of congruences if and only if $\sigma_\mu \leq \sigma_\lambda$ for $\mu < \lambda$ and $\sigma_\lambda = \bigcup_{\mu < \lambda} \sigma_\mu$ for all limit ordinals λ .

DEFINITION 6.5. Let \mathbf{A} be an algebra with congruences $\alpha, \beta, \theta, \theta_\lambda, \delta$ and δ' .

If $\delta \leq \theta$ and $\mathbf{C}(\theta, \theta; \delta)$ holds then we say that θ is **abelian over δ** , and write $\delta \triangleleft \theta$.

If κ is an ordinal, and $(\theta_\lambda)_{\lambda < \kappa+1}$ is a continuous sequence of congruences such that $\theta_\lambda \triangleleft \theta_{\lambda+1}$ for all $\lambda < \kappa$, then we say that θ_κ is **κ -step solvable over θ_0** .

We say that θ is **∞ -solvable over δ** if θ is κ -step solvable over δ for some κ . We may denote this by $\delta \lll \theta$. If $0 \lll \theta$, then we say that θ is **∞ -solvable**. If $0 \lll 1$, then say that \mathbf{A} is **∞ -solvable**.

If $\alpha \wedge \beta \lll \alpha \vee \beta$, then we say that α is **solvably related to β** and denote this by $\alpha \stackrel{s}{\sim} \beta$.

If $\delta < \theta$, and $\delta \triangleleft \delta' \leq \theta$ implies $\delta = \delta'$, then we call the congruence interval $I[\delta, \theta]$ a **solvability obstruction** and we denote this by $\delta \blacktriangleleft \theta$.

Note that for comparable congruences $\delta \leq \theta$ the relations $\delta \lll \theta$ and $\delta \stackrel{s}{\sim} \theta$ are the same.

EXAMPLE 6.6. If \mathcal{V} is a variety whose members have the property that their principal congruences are abelian, then every member of \mathcal{V} is ∞ -solvable. For if $\mathbf{A} \in \mathcal{V}$, $(a_\lambda)_{\lambda < \kappa}$ is an enumeration of \mathbf{A} , and θ_λ is the congruence generated by $\{(a_0, a_\mu) \mid \mu < \lambda\}$, then it is not hard to see that

- (i) $\theta_0 = 0$ and $\theta_\kappa = 1$,
- (ii) $\theta_\lambda \triangleleft \theta_{\lambda+1}$ for all $\lambda < \kappa$, and
- (iii) $\theta_\lambda = \bigcup_{\mu < \lambda} \theta_\mu$ for all limit ordinals.

Items (i) and (iii) follow immediately from the definition of θ_λ . For item (ii), note that $\theta_{\lambda+1} = \text{Cg}^{\mathbf{A}}(a_0, a_\lambda) \vee \theta_\lambda$. Thus $\theta_{\lambda+1}/\theta_\lambda$ is a principal (hence abelian) congruence on $\mathbf{A}/\theta_\lambda$. This implies that $\theta_\lambda \triangleleft \theta_{\lambda+1}$ for all $\lambda < \kappa$.

An example of a variety in which principal congruences are abelian is the variety \mathcal{V} of nonunital rings satisfying $x^2 \approx 0$. This identity implies that $0 \approx (x+y)^2 \approx x^2 + xy + yx + y^2 \approx xy + yx$, so the anticommutative law $xy \approx -yx$ holds in \mathcal{V} . This law implies that, for any $\mathbf{R} \in \mathcal{V}$, all principal ideals have the form $(a) = R \cdot a$. Hence $(a)^2 = RaRa = \sum \{ra \cdot sa = -rsa^2 \mid r, s \in R\} = (0)$. Since congruences correspond to ideals, and the commutator operation corresponds to ideal multiplication, this is indeed a variety in which principal congruences are abelian. Hence \mathcal{V} consists entirely of ∞ -solvable algebras.

However, there is no cardinal bound on solvability degree in \mathcal{V} . In fact, if κ is any infinite cardinal, \mathcal{V} contains rings of solvability degree exactly κ . One such ring, \mathbf{R}_κ , may be presented by generators and relations relative to \mathcal{V} as follows. Let S be an infinite set of cardinality κ . Take as generators for \mathbf{R}_κ all symbols x_U where U is a nonempty subset of S . Impose the relations

- $2x_U = 0$ for all nonempty $U \subseteq S$,
- $x_U \cdot x_V = 0$ if $U \cap V \neq \emptyset$, and
- $x_U \cdot x_V = x_{U \cup V}$ if $U \cap V = \emptyset$.

These relations force each element of \mathbf{R}_κ to be expressible as a finite sum of generators $x_{U_1} + \cdots + x_{U_n}$ where $U_i \neq U_j$ if $i \neq j$. We leave it as an exercise to prove that the ring \mathbf{R}_κ is κ -step solvable, but not λ -step solvable for any ordinal $\lambda < \kappa$.

LEMMA 6.7. *Let \mathcal{V} be a variety with a weak difference term. If $\mathbf{A} \in \mathcal{V}$ has a tolerance T and a congruence δ such that $\mathbf{C}(T, T; \delta)$ holds, then $T' := \delta \circ T \circ \delta$ is the smallest congruence on \mathbf{A} containing $T \cup \delta$. Furthermore, $\mathbf{C}(T', T'; \delta)$ holds.*

PROOF. We know from the basic properties of the centralizer relation (enumerated in Theorem 2.19) that if $\mathbf{C}(T, T; \delta)$ holds and $T' := \delta \circ T \circ \delta$, then $\mathbf{C}(T', T'; \delta)$ holds. Since $(T \cup \delta) \subseteq T' = \delta \circ T \circ \delta \subseteq \text{Cg}^{\mathbf{A}}(T \cup \delta)$, the only thing left to show is that T' is a congruence. For this we only need to establish that $T' \circ T' \subseteq T'$.

Let $\theta = \text{Cg}^{\mathbf{A}}(T')$. Since $\mathbf{C}(T', T'; \delta)$ holds, it follows from Theorems 3.23 and 6.2 that $\mathbf{C}(\theta, \theta; \delta)$ holds. If $(a, c) \in T' \circ T'$, then there is an element b such that $(a, b), (b, c) \in T' \subseteq \theta$. Since $a \equiv_{[\theta, \theta]} d(a, b, b) \equiv_{T'} d(b, b, c) \equiv_{[\theta, \theta]} c$ and $[\theta, \theta] \leq \delta$, we get that $(a, c) \in \delta \circ T' \circ \delta = T'$. Since $(a, c) \in T' \circ T'$ was chosen arbitrarily, $T' \circ T' \subseteq T'$. \square

LEMMA 6.8. *Let \mathcal{V} be a variety with a weak difference term. If $\mathbf{A} \in \mathcal{V}$ has congruences θ and δ such that $[\theta, \theta] \leq \delta$, then $\mathbf{C}(\theta, \theta; \delta)$ holds.*

PROOF. Assume that $[\theta, \theta] \leq \delta$. If $\mathbf{C}(\theta, \theta; \delta)$ does not hold, then there are a polynomial $p(\mathbf{x}, \mathbf{y})$ and tuples $\mathbf{a} \theta \mathbf{b}$ and $\mathbf{u} \theta \mathbf{v}$ such that

$$p(\mathbf{a}, \mathbf{u}) \equiv_{\delta} p(\mathbf{a}, \mathbf{v})$$

and

$$(6.1) \quad p(\mathbf{b}, \mathbf{u}) \not\equiv_{\delta} p(\mathbf{b}, \mathbf{v}).$$

Let $p'(\mathbf{x}, \mathbf{y}) = d(p(\mathbf{x}, \mathbf{y}), p(\mathbf{x}, \mathbf{u}), p(\mathbf{b}, \mathbf{u}))$. Using that d is an idempotent term that is Maltsev on θ -classes modulo $[\theta, \theta]$, we get that

$$p'(\mathbf{a}, \underline{\mathbf{u}}) = d(p(\mathbf{a}, \mathbf{u}), p(\mathbf{a}, \mathbf{u}), p(\mathbf{b}, \mathbf{u})) \equiv_{[\theta, \theta]} p(\mathbf{b}, \mathbf{u}) = p'(\mathbf{b}, \underline{\mathbf{u}}).$$

Since $\mathbf{C}(\theta, \theta; [\theta, \theta])$ holds, we can change the underlined \mathbf{u} to \mathbf{v} without changing the fact that the left hand side is $[\theta, \theta]$ -related to the right hand side:

$$(6.2) \quad p'(\mathbf{a}, \underline{\mathbf{v}}) \equiv_{[\theta, \theta]} p'(\mathbf{b}, \underline{\mathbf{v}}).$$

Moreover we have that

$$(6.3) \quad \begin{aligned} p'(\mathbf{a}, \mathbf{v}) &= d(p(\mathbf{a}, \mathbf{v}), p(\mathbf{a}, \mathbf{u}), p(\mathbf{b}, \mathbf{u})) \\ &\equiv_{\delta} d(p(\mathbf{a}, \mathbf{u}), p(\mathbf{a}, \mathbf{u}), p(\mathbf{b}, \mathbf{u})) \equiv_{[\theta, \theta]} p(\mathbf{b}, \mathbf{u}). \end{aligned}$$

Here

$$d(p(\mathbf{a}, \mathbf{v}), p(\mathbf{a}, \mathbf{u}), p(\mathbf{b}, \mathbf{u})) \equiv_{\delta} d(p(\mathbf{a}, \mathbf{u}), p(\mathbf{a}, \mathbf{u}), p(\mathbf{b}, \mathbf{u}))$$

is a consequence of the fact that $p(\mathbf{a}, \mathbf{v}) \equiv_{\delta} p(\mathbf{a}, \mathbf{u})$, and

$$d(p(\mathbf{a}, \mathbf{u}), p(\mathbf{a}, \mathbf{u}), p(\mathbf{b}, \mathbf{u})) \equiv_{[\theta, \theta]} p(\mathbf{b}, \mathbf{u})$$

is a consequence of the facts that $\mathbf{a} \theta \mathbf{b}$ and that d is Maltsev on θ -classes modulo $[\theta, \theta]$. We also have that

$$(6.4) \quad p'(\mathbf{b}, \mathbf{v}) = d(p(\mathbf{b}, \mathbf{v}), p(\mathbf{b}, \mathbf{u}), p(\mathbf{b}, \mathbf{u})) \equiv_{[\theta, \theta]} p(\mathbf{b}, \mathbf{v}).$$

since $\mathbf{u} \theta \mathbf{v}$ and d is Maltsev on θ -classes modulo $[\theta, \theta]$. Combining lines (6.2), (6.3), and (6.4) yields

$$p(\mathbf{b}, \mathbf{u}) \equiv_{\delta} p'(\mathbf{a}, \mathbf{v}) \equiv_{[\theta, \theta]} p'(\mathbf{b}, \mathbf{v}) \equiv_{[\theta, \theta]} p(\mathbf{b}, \mathbf{v}),$$

or simply $p(\mathbf{b}, \mathbf{u}) \equiv_{\delta} p(\mathbf{b}, \mathbf{v})$ (since $[\theta, \theta] \leq \delta$). This contradicts line (6.1), thereby completing the proof. \square

The following corollary will be used in Section 9.2.

COROLLARY 6.9. *Let \mathcal{V} be a variety with a weak difference term. Assume that T is an abelian tolerance on some $\mathbf{A} \in \mathcal{V}$, and that a and b are distinct T -related elements. If ρ is a congruence on \mathbf{A} that is maximal with respect to $(a, b) \notin \rho$, then T/ρ is a nontrivial abelian congruence on the subdirectly irreducible algebra \mathbf{A}/ρ .*

PROOF. It follows from Lemma 6.7 (using $\delta = 0$) that T is an abelian congruence on \mathbf{A} . Then, using Lemma 6.8 in the case where $\theta = T$ and $\delta = \rho$, we obtain that $\mathbf{C}(T, T; \rho)$ holds, so T/ρ is a nontrivial abelian tolerance on the subdirectly irreducible algebra \mathbf{A}/ρ . Using Lemma 6.7, we find that T/ρ is a congruence. \square

LEMMA 6.10. *Let \mathcal{V} be a variety with a weak difference term. Assume that $\mathbf{A} \in \mathcal{V}$ has congruences α, β , and γ .*

- (1) *If $\alpha \triangleleft \beta$, then $(\alpha \wedge \gamma) \triangleleft (\beta \wedge \gamma)$ and $(\alpha \vee \gamma) \triangleleft (\beta \vee \gamma)$.*
- (2) *If $\alpha \triangleleft\triangleleft \beta$, then $(\alpha \wedge \gamma) \triangleleft\triangleleft (\beta \wedge \gamma)$ and $(\alpha \vee \gamma) \triangleleft\triangleleft (\beta \vee \gamma)$.*

PROOF. To prove the part of item (1) that concerns the compatibility of \triangleleft with meet, recall that $\alpha \triangleleft \beta$ means that $\mathbf{C}(\beta, \beta; \alpha)$ holds. By item (1) of Theorem 2.19 we have that $\mathbf{C}(\beta \wedge \gamma, \beta \wedge \gamma; \alpha)$ holds. By item (8) of Theorem 2.19 we have that $\mathbf{C}(\beta \wedge \gamma, \beta \wedge \gamma; \gamma)$ holds. Item (6) of Theorem 2.19 now yields that $\mathbf{C}(\beta \wedge \gamma, \beta \wedge \gamma; \alpha \wedge \gamma)$ holds, which means that $(\alpha \wedge \gamma) \triangleleft (\beta \wedge \gamma)$.

To prove the compatibility of \triangleleft with join, assume that $\mathbf{C}(\beta, \beta; \alpha)$ holds. Then $[\beta, \beta] \leq \alpha \leq \alpha \vee \gamma$, so by Lemma 6.8 we get that $\mathbf{C}(\beta, \beta; \alpha \vee \gamma)$ holds. Using Lemma 6.7 with $T = \beta$ and $\delta = \alpha \vee \gamma$ we deduce that $(\alpha \vee \gamma) \circ \beta \circ (\alpha \vee \gamma) = \beta \vee (\alpha \vee \gamma) = \beta \vee \gamma$ and that $\mathbf{C}(\beta \vee \gamma, \beta \vee \gamma; \alpha \vee \gamma)$ holds. This shows that $(\alpha \vee \gamma) \triangleleft (\beta \vee \gamma)$, so item (1) holds.

For (2), assume that $(\theta_\lambda)_{\lambda < \kappa+1}$ is a continuous sequence of congruences for which $\alpha = \theta_0, \beta = \theta_\kappa, \theta_\lambda \triangleleft \theta_{\lambda+1}$ for all λ . The sequences $(\theta_\lambda \wedge \gamma)_{\lambda < \kappa+1}$ and $(\theta_\lambda \vee \gamma)_{\lambda < \kappa+1}$ are also continuous, since meeting or joining with the fixed element γ preserves order and unions of chains. Therefore it follows from item (1) of this lemma that $(\theta_\lambda \wedge \gamma)_{\lambda < \kappa+1}$ and $(\theta_\lambda \vee \gamma)_{\lambda < \kappa+1}$ are sequences witnessing $(\alpha \wedge \gamma) \triangleleft\triangleleft (\beta \wedge \gamma)$ and $(\alpha \vee \gamma) \triangleleft\triangleleft (\beta \vee \gamma)$, respectively. \square

THEOREM 6.11. *Let \mathcal{V} be a variety with a weak difference term. If $\mathbf{A} \in \mathcal{V}$, then $\overset{\text{s}}{\sim}$ is a congruence on $\mathbf{Con}(\mathbf{A})$ that is compatible with the complete join operation.*

PROOF. Recall that $\alpha \overset{\text{s}}{\sim} \beta$ holds if and only if $(\alpha \wedge \beta) \triangleleft\triangleleft (\alpha \vee \beta)$. It is easy to see that this relation is reflexive and symmetric. To show that it is transitive, assume that $\alpha \overset{\text{s}}{\sim} \beta \overset{\text{s}}{\sim} \gamma$. Then $(\alpha \wedge \beta) \triangleleft\triangleleft (\alpha \vee \beta)$ and $(\beta \wedge \gamma) \triangleleft\triangleleft (\beta \vee \gamma)$, so by the second statement of Lemma 6.10 we

get that $(\alpha \wedge \beta \wedge \gamma) = (\alpha \wedge \beta) \wedge (\beta \wedge \gamma) \lll (\alpha \wedge \beta) \wedge (\beta \vee \gamma) = (\alpha \wedge \beta)$
and $(\alpha \vee \beta) = (\alpha \vee \beta) \vee (\beta \wedge \gamma) \lll (\alpha \vee \beta) \vee (\beta \vee \gamma) = (\alpha \vee \beta \vee \gamma)$.
Thus we have

$$(\alpha \wedge \beta \wedge \gamma) \lll (\alpha \wedge \beta) \lll (\alpha \vee \beta) \lll (\alpha \vee \beta \vee \gamma).$$

The well-ordered chains of congruences that witness these three instances of ∞ -solvability may be concatenated to produce a single such chain from $\alpha \wedge \beta \wedge \gamma$ to $\alpha \vee \beta \vee \gamma$, thus proving that $(\alpha \wedge \beta \wedge \gamma) \lll (\alpha \vee \beta \vee \gamma)$. Since $(\alpha \wedge \beta \wedge \gamma) \leq \alpha \wedge \gamma \leq \alpha \vee \gamma \leq (\alpha \vee \beta \vee \gamma)$, we can deduce that $(\alpha \wedge \gamma) \lll (\alpha \vee \gamma)$ if we show first that $\delta \leq \delta' \leq \theta' \leq \theta$ and $\delta \lll \theta$ imply $\delta' \lll \theta'$. We show this by using Lemma 6.10 again: joining $\delta \lll \theta$ with δ' we obtain

$$\delta' = (\delta' \vee \delta) \lll (\delta' \vee \theta) = \theta.$$

Meeting with θ' we find that

$$\delta' = (\delta' \wedge \theta') \lll (\theta \wedge \theta') = \theta'.$$

This reasoning completes the argument that $\overset{s}{\sim}$ is transitive, hence an equivalence relation, and also shows that subintervals of ∞ -solvable intervals are again ∞ -solvable. Hence $\overset{s}{\sim}$ -classes are convex sublattices of $\mathbf{Con}(\mathbf{A})$. This implies that $\alpha \overset{s}{\sim} \beta$ if and only if α and β belong to some ∞ -solvable interval. This reformulation of the definition of the $\overset{s}{\sim}$ -relation makes it easy to see that it is a congruence: $\alpha \overset{s}{\sim} \beta$ if and only if $\delta := \alpha \wedge \beta \lll \alpha \vee \beta =: \theta$. But $\delta \lll \theta$ implies $\delta \wedge \gamma \lll \theta \wedge \gamma$ and $\delta \vee \gamma \lll \theta \vee \gamma$ by Lemma 6.10. Since $\alpha \wedge \gamma, \beta \wedge \gamma \in I[\delta \wedge \gamma, \theta \wedge \gamma]$ and $\alpha \vee \gamma, \beta \vee \gamma \in I[\delta \vee \gamma, \theta \vee \gamma]$, we conclude that $\alpha \overset{s}{\sim} \beta$ implies that $(\alpha \wedge \gamma) \overset{s}{\sim} (\beta \wedge \gamma)$ and $(\alpha \vee \gamma) \overset{s}{\sim} (\beta \vee \gamma)$.

It remains to show that $\overset{s}{\sim}$ is compatible with the complete join operation of $\mathbf{Con}(\mathbf{A})$. Suppose that $\alpha_i \overset{s}{\sim} \beta_i$ for all $i \in I$. Let $\delta_i = \alpha_i \wedge \beta_i$ and $\theta_i = \alpha_i \vee \beta_i$. Then $\delta_i \lll \theta_i$ for all i . If $\Delta := \bigvee_{i \in I} \delta_i$ and $\Theta := \bigvee_{i \in I} \theta_i$, then both $\bigvee_{i \in I} \alpha_i$ and $\bigvee_{i \in I} \beta_i$ belong to the interval $I[\Delta, \Theta]$. We will be done if we show that Θ is ∞ -solvable over Δ .

Consider the collection \mathcal{P} of all sequences of the form $\sigma = (\sigma_\lambda)_{\lambda < \kappa+1}$ where $\sigma_\lambda \in \mathbf{Con}(\mathbf{A})$ and

- (0) $\sigma_0 = \Delta, \sigma_\kappa \in I[\Delta, \Theta]$,
- (1) $\sigma_\lambda \triangleleft \sigma_{\lambda+1}$ for all $\lambda < \kappa$, and
- (2) $\sigma_\lambda = \bigcup_{\mu < \lambda} \sigma_\mu$ for limit ordinals λ .

The elements of \mathcal{P} are the continuous sequences of congruences that witness the ∞ -solvability over Δ of some elements $\sigma_\kappa \in I[\Delta, \Theta]$. \mathcal{P} is nonempty, since it contains the one-term sequence (Δ) . \mathcal{P} can be partially ordered by the rule: $\sigma \leq \sigma'$ if σ' is an extension of σ . That

is, $\sigma \leq \sigma'$ if the domain of σ is contained in the domain of σ' and the functions agree on their common domain. This is an inductive ordering of \mathcal{P} , as we now explain. Assume that $(\sigma^\mu)_{\mu < \nu}$ is an increasing sequence of sequences in \mathcal{P} . Define Σ' to be the union of the functions σ^μ . Σ' is a function from the ordinal $\rho := \bigcup_{\mu < \nu} \text{dom}(\sigma^\mu)$ into $\mathbf{Con}(\mathbf{A})$ that preserves all existing joins. Σ' may not be in \mathcal{P} because elements of \mathcal{P} are functions from successor ordinals into $\mathbf{Con}(\mathbf{A})$ and ρ may be a limit ordinal. But if it is, then there is a unique meet continuous extension of Σ' to $\rho + 1$: it is the sequence $\Sigma = (\Sigma_\lambda)_{\lambda < \rho+1}$ where $(\Sigma_\lambda)_{\lambda < \rho} = \Sigma'$ and Σ_ρ is defined to be $\bigcup_{\lambda < \rho} \Sigma_\lambda$. Since Σ is a meet continuous homomorphism from a successor ordinal into $\mathbf{Con}(\mathbf{A})$, Σ extends each σ^μ , and $\Sigma_\lambda = \sigma_\lambda^\mu \triangleleft \sigma_{\lambda+1}^\mu = \Sigma_{\lambda+1}$ for some μ for each $\lambda < \rho$, Σ is an upper bound in \mathcal{P} for all the sequences σ^μ .

Since \mathcal{P} is nonempty and inductively ordered, Zorn's lemma applies to show that \mathcal{P} has a maximal element $(\psi_\lambda)_{\lambda < \kappa+1}$. The final congruence in this sequence, ψ_κ , is ∞ -solvable over Δ and is a maximal congruence in the interval $I[\Delta, \Theta]$ among those that are ∞ -solvable over Δ . For each $i \in I$, $\delta_i \triangleleft \theta_i$ and $\delta_i \leq \Delta \leq \psi_\kappa$, so by Lemma 6.10 we have

$$\psi_\kappa = (\psi_\kappa \vee \delta_i) \triangleleft (\psi_\kappa \vee \theta_i) \leq \Theta.$$

The maximality condition on ψ_κ implies that $\psi_\kappa = \psi_\kappa \vee \theta_i$, or $\psi_\kappa \geq \theta_i$. Since this is true for all $i \in I$, $\psi_\kappa \geq \bigvee_{i \in I} \theta_i = \Theta$. This proves that $\Delta \triangleleft \psi_\kappa = \Theta$, as desired. \square

DEFINITION 6.12. Assume that \mathcal{V} is a variety with a weak difference term, and $\mathbf{A} \in \mathcal{V}$. If $\alpha \in \mathbf{Con}(\mathbf{A})$, then

$$\text{rad}(\alpha) := \bigvee \{ \beta \in \mathbf{Con}(\mathbf{A}) \mid \alpha \stackrel{s}{\sim} \beta \}$$

is the ∞ -solvable radical of α . The ∞ -solvable radical of \mathbf{A} is $\text{rad}(0)$. A **radical congruence** is a congruence α such that $\text{rad}(\alpha) = \alpha$.

THEOREM 6.13. *Let \mathcal{V} be a variety with a weak difference term. Assume that $\mathbf{A} \in \mathcal{V}$ and that $\alpha, \beta \in \mathbf{Con}(\mathbf{A})$. Then*

- (1) $\text{rad}(\alpha)$ is the largest congruence $\stackrel{s}{\sim}$ -related to α .
- (2) $\text{rad}(\text{rad}(\alpha)) = \text{rad}(\alpha)$.
- (3) $\text{rad}(\alpha \wedge \beta) = \text{rad}(\alpha) \wedge \text{rad}(\beta)$.
- (4) $\alpha \triangleleft \beta$ if and only if $\alpha < \beta$ and $\alpha = \beta \wedge \text{rad}(\alpha)$.
- (5) If $\text{Rad}(\mathbf{A})$ denotes the meet subsemilattice of $\mathbf{Con}(\mathbf{A})$ consisting of the radical congruences, then the restriction to $\text{Rad}(\mathbf{A})$ of the natural homomorphism $\nu : \mathbf{Con}(\mathbf{A}) \rightarrow \mathbf{Con}(\mathbf{A}) / \stackrel{s}{\sim} : \alpha \mapsto \alpha / \stackrel{s}{\sim}$ to $\text{Rad}(\mathbf{A})$ is an isomorphism onto $\mathbf{Con}(\mathbf{A}) / \stackrel{s}{\sim}$.

(6) *The natural homomorphism ν from item (5) is meet continuous, and its image $\mathbf{Con}(\mathbf{A})/\overset{s}{\sim}$ meet semidistributive.*

PROOF. For item (1), the fact that $\overset{s}{\sim}$ is compatible with the complete join operation implies that the join of all congruence $\overset{s}{\sim}$ -related to α is $\overset{s}{\sim}$ -related to α , and it is necessarily the largest congruence $\overset{s}{\sim}$ -related to α .

For (2), $\alpha \overset{s}{\sim} \text{rad}(\alpha)$ and $\text{rad}(\alpha) \overset{s}{\sim} \text{rad}(\text{rad}(\alpha))$ hold by (1), so $\alpha \overset{s}{\sim} \text{rad}(\text{rad}(\alpha))$. Since $\text{rad}(\alpha)$ is the largest congruence $\overset{s}{\sim}$ -related to α , it follows that $\text{rad}(\text{rad}(\alpha)) \leq \text{rad}(\alpha)$. The reverse inclusion follows from the facts that $\text{rad}(\alpha) \overset{s}{\sim} \text{rad}(\alpha)$ and $\text{rad}(\text{rad}(\alpha))$ is the largest congruence $\overset{s}{\sim}$ -related to $\text{rad}(\alpha)$.

For (3), $\alpha \overset{s}{\sim} \text{rad}(\alpha)$ and $\beta \overset{s}{\sim} \text{rad}(\beta)$, so $\alpha \wedge \beta \overset{s}{\sim} \text{rad}(\alpha) \wedge \text{rad}(\beta)$ since $\overset{s}{\sim}$ is a congruence. Since $\text{rad}(\alpha \wedge \beta)$ is the largest congruence $\overset{s}{\sim}$ -related to $\alpha \wedge \beta$ we get that $\text{rad}(\alpha) \wedge \text{rad}(\beta) \leq \text{rad}(\alpha \wedge \beta)$. On the other hand, $\alpha \wedge \beta \overset{s}{\sim} \text{rad}(\alpha \wedge \beta)$, so

$$\alpha = \alpha \vee (\alpha \wedge \beta) \overset{s}{\sim} \alpha \vee \text{rad}(\alpha \wedge \beta) \leq \text{rad}(\alpha)$$

from the definition of $\text{rad}(\alpha)$. Thus, $\text{rad}(\alpha \wedge \beta) \leq \text{rad}(\alpha)$, and similarly $\text{rad}(\alpha \wedge \beta) \leq \text{rad}(\beta)$, so $\text{rad}(\alpha \wedge \beta) \leq \text{rad}(\alpha) \wedge \text{rad}(\beta)$. This completes the proof of (3).

For (4), assume that $\alpha \blacktriangleleft \beta$. Then $\alpha < \beta$ according to the definition of \blacktriangleleft . The congruence $\alpha' := \beta \wedge \text{rad}(\alpha)$ satisfies $\alpha \leq \alpha' \leq \text{rad}(\alpha)$ and is in $I[\alpha, \beta]$, so it is a congruence in $I[\alpha, \beta]$ that is ∞ -solvable over α . If this fact is witnessed by a chain of congruences $(\theta_\lambda)_{\lambda < \kappa+1}$ starting at $\theta_0 = \alpha$ and ending at $\theta_\kappa = \alpha'$, and satisfying the conditions of Definition 6.5, then using the fact that $\alpha \blacktriangleleft \beta$ and induction it follows that $\theta_\lambda = \alpha$ for all λ , so $\alpha = \alpha' = \beta \wedge \text{rad}(\alpha)$. Conversely, if $\alpha \blacktriangleleft \beta$ fails to hold and $\alpha < \beta$, then there is a congruence $\alpha' \in I[\alpha, \beta]$ such that $\alpha \triangleleft \alpha'$ and $\alpha < \alpha'$. Since $\alpha \overset{s}{\sim} \alpha'$ we get that $\alpha < \alpha' \leq \beta \wedge \text{rad}(\alpha)$, so $\alpha = \beta \wedge \text{rad}(\alpha)$ fails to hold.

For (5), it is clear that $\nu : \text{Rad}(\mathbf{A}) \rightarrow \mathbf{Con}(\mathbf{A})/\overset{s}{\sim}$ is an order-preserving function between meet semilattices. To prove that it is an isomorphism it suffices to show that it is bijective, or equivalently that the elements of $\text{Rad}(\mathbf{A})$ form a transversal for the congruence $\overset{s}{\sim}$. The fact that $\alpha \overset{s}{\sim} \text{rad}(\alpha)$ (proved in (1)) together with the fact that $\text{rad}(\alpha) \in \text{Rad}(\mathbf{A})$ (an equivalent form of (2)) shows that every $\overset{s}{\sim}$ -class contains a radical congruence. But the fact proved in (1), that every radical element is the largest element of its $\overset{s}{\sim}$ -class, implies that each $\overset{s}{\sim}$ -class has at most one radical congruence. This shows that the elements of $\text{Rad}(\mathbf{A})$ form a transversal for $\overset{s}{\sim}$.

For (6), the fact that $\overset{s}{\sim}$ is compatible with the complete join operation implies that it is a meet continuous congruence, hence the quotient modulo this congruence is meet continuous and the natural homomorphism $\nu : \mathbf{Con}(\mathbf{A}) \rightarrow \mathbf{Con}(\mathbf{A})/\overset{s}{\sim}$ is a meet continuous homomorphism. To prove that $\mathbf{Con}(\mathbf{A})/\overset{s}{\sim}$ is meet semidistributive, suppose that $(\alpha/\overset{s}{\sim}) \wedge (\beta/\overset{s}{\sim}) = (\alpha/\overset{s}{\sim}) \wedge (\gamma/\overset{s}{\sim}) = (\delta/\overset{s}{\sim})$. From item (5) we get that $\text{rad}(\alpha) \wedge \text{rad}(\beta) = \text{rad}(\alpha) \wedge \text{rad}(\gamma) = \text{rad}(\delta)$ in $\text{Rad}(\mathbf{A})$. Since $\text{Rad}(\mathbf{A})$ is a meet subsemilattice of $\mathbf{Con}(\mathbf{A})$, these equalities hold in $\mathbf{Con}(\mathbf{A})$. Theorem 2.19 (8) implies that $\mathbf{C}(\text{rad}(\beta), \text{rad}(\alpha); \text{rad}(\delta))$ and $\mathbf{C}(\text{rad}(\gamma), \text{rad}(\alpha); \text{rad}(\delta))$ hold. Theorem 2.19 (5) now yields that $\mathbf{C}(\text{rad}(\beta) \vee \text{rad}(\gamma), \text{rad}(\alpha); \text{rad}(\delta))$ holds. If

$$\theta := \text{rad}(\alpha) \wedge (\text{rad}(\beta) \vee \text{rad}(\gamma)),$$

then $\theta \geq \text{rad}(\delta)$ and $\mathbf{C}(\theta, \theta; \text{rad}(\delta))$ holds according to Theorem 2.19 (1). This shows that $\text{rad}(\delta) \triangleleft \theta$. But $\text{rad}(\delta) \blacktriangleleft 1$ or else $\text{rad}(\delta) = 1$, by items (2) and (4) of this theorem, and in either case we get

$$\text{rad}(\delta) = \theta = \text{rad}(\alpha) \wedge (\text{rad}(\beta) \vee \text{rad}(\gamma)).$$

Using the isomorphism from (5) once again, we conclude that $(\alpha/\overset{s}{\sim}) \wedge ((\beta/\overset{s}{\sim}) \vee (\gamma/\overset{s}{\sim})) = (\delta/\overset{s}{\sim})$. This proves that $\mathbf{Con}(\mathbf{A})/\overset{s}{\sim}$ is meet semidistributive. \square

THEOREM 6.14. *Let \mathcal{V} be a variety with a weak difference term. If $\mathbf{A} \in \mathcal{V}$ has congruences $\delta \leq \theta$, then θ is ∞ -solvable over δ if and only if the interval $I[\delta, \theta]$ contains no solvability obstruction.*

PROOF. Assume that θ is ∞ -solvable over δ , and that $I[\alpha, \beta]$ is an arbitrarily chosen but nontrivial subinterval of $I[\delta, \theta]$. Then $\delta = \delta \wedge \alpha \overset{s}{\sim} \theta \wedge \alpha = \alpha$, so $\text{rad}(\alpha) = \text{rad}(\delta) \geq \theta \geq \beta$. This shows that $\alpha < \beta = \beta \wedge \text{rad}(\alpha)$. By item (4) of Theorem 6.13, $I[\alpha, \beta]$ is not a solvability obstruction.

Conversely, assume that $I[\delta, \theta]$ contains no solvability obstruction. Let $\alpha = \theta \wedge \text{rad}(\delta)$. Then since $\delta \leq \alpha \leq \text{rad}(\delta)$ we get that $\delta \overset{s}{\sim} \alpha$, and hence that $\text{rad}(\alpha) = \text{rad}(\delta)$. This means that $\alpha = \theta \wedge \text{rad}(\alpha)$. According to item (4) of Theorem 6.13, either $\alpha \blacktriangleleft \theta$ or else $\alpha = \theta$. The first possibility is ruled out by the assumption that $I[\delta, \theta]$ contains no solvability obstruction, so it must be that $\alpha = \theta$. Since $\delta \overset{s}{\sim} \alpha = \theta$, the proof is complete. \square

LEMMA 6.15. *Let \mathcal{V} be a variety with a weak difference term. Assume that $\mathbf{A} \in \mathcal{V}$ has congruences γ, γ', δ and δ' such that $\gamma' \triangleleft \gamma$ and $\delta' \triangleleft \delta$. Then $\gamma \circ \delta \subseteq \gamma' \circ \delta \circ \gamma \circ \delta'$.*

PROOF. Choose any $(a, c) \in \gamma \circ \delta$. There is a $b \in A$ such that $a \equiv_\gamma b \equiv_\delta c$, so

$$a \equiv_{\gamma'} d(a, b, b) \equiv_\delta d(a, b, c) \equiv_\gamma d(b, b, c) \equiv_{\delta'} c$$

shows that $(a, c) \in \gamma' \circ \delta \circ \gamma \circ \delta'$. □

THEOREM 6.16. *Let \mathcal{V} be a variety with a weak difference term. If $\mathbf{A} \in \mathcal{V}$, then $\overset{s}{\sim}$ -classes of $\mathbf{Con}(\mathbf{A})$ consist of permuting congruences. In particular, these classes are modular sublattices.*

PROOF. We assume that the theorem is false and argue to a contradiction. If some $\overset{s}{\sim}$ -class of $\mathbf{Con}(\mathbf{A})$ contains a pair of nonpermuting congruences, then there is a subinterval $I[\sigma, \tau]$ of that $\overset{s}{\sim}$ -class that contains the two congruences. (For example, we could take σ to be the meet and τ to be the join of the two nonpermuting congruences.) Let A denote the ordered set of all congruences $\theta \in I[\sigma, \tau]$ such that every congruence in $I[\sigma, \theta]$ permutes with every congruence in $I[\sigma, \tau]$:

$$A = \{\theta \in I[\sigma, \tau] \mid \forall \theta' \in I[\sigma, \theta], \forall \psi \in I[\sigma, \tau] (\theta' \circ \psi = \psi \circ \theta')\}.$$

A is nonempty (since $\sigma \in A$) and it is closed under updirected unions, therefore A has maximal elements. Let α be a maximal element of A . It must be that $\alpha < \tau$ since otherwise all congruences in $I[\sigma, \tau]$ permute, contrary to assumption. Since $I[\sigma, \tau]$ is ∞ -solvable and $\sigma \leq \alpha < \tau$, there is a congruence $\hat{\alpha}$ such that $\alpha < \hat{\alpha} \leq \tau$ and $\alpha \triangleleft \hat{\alpha}$. The set

$$B = \{\theta \in I[\sigma, \tau] \mid \forall \theta' \in I[\sigma, \theta], \forall \psi \in I[\sigma, \hat{\alpha}] (\theta' \circ \psi = \psi \circ \theta')\}$$

remove

of all $\theta \in I[\sigma, \tau]$ such that every congruence in $I[\sigma, \theta]$ permutes with every congruence in $I[\sigma, \hat{\alpha}]$ is also nonempty and closed under updirected unions, so B has a maximal element β . We cannot have $\beta = \tau$ since this would force $\hat{\alpha} \in A$, contradicting the the maximality of α . Since $\beta < \tau$, there is a congruence $\hat{\beta}$ such that $\beta < \hat{\beta} \leq \tau$ and $\beta \triangleleft \hat{\beta}$.

By the maximality of β and the fact that $\beta < \hat{\beta}$, there is some $\delta \leq \hat{\beta}$ which fails to permute with some $\gamma \leq \hat{\alpha}$. Every congruence below β permutes with every congruence below $\hat{\alpha}$, so $\delta \not\leq \beta$. Every congruence below α permutes with every congruence in $I[\sigma, \tau]$, so $\gamma \not\leq \alpha$. Therefore $\gamma' = \alpha \wedge \gamma$ is strictly less than γ and $\delta' = \beta \wedge \delta$ is strictly less than δ . The four congruences γ, γ', δ and δ' have the following properties:

- (i) $\gamma' < \gamma \leq \hat{\alpha}$, $\gamma' \leq \alpha$;
- (ii) $\delta' < \delta \leq \hat{\beta}$, $\delta' \leq \beta$;
- (iii) $\gamma' \triangleleft \gamma$, $\delta' \triangleleft \delta$; and
- (iv) any two congruences in $\{\gamma, \gamma', \delta, \delta'\}$ permute, except $\gamma \circ \delta \not\leq \delta \circ \gamma$.

(The first, second and fourth items follow from the way we chose $\alpha, \widehat{\alpha}, \beta, \widehat{\beta}, \gamma, \gamma', \delta$, and δ' . The third item follows from Lemma 6.10 since, for example, $\alpha \triangleleft \widehat{\alpha}$ so $\gamma' = \alpha \wedge \gamma \triangleleft \widehat{\alpha} \wedge \gamma = \gamma$.) By item (iii) and Lemma 6.15 we have $\gamma \circ \delta \subseteq \gamma' \circ \delta \circ \gamma \circ \delta'$. But by items (i), (ii) and (iv) we have

$$\gamma' \circ \delta \circ \gamma \circ \delta' = \delta \circ \gamma' \circ \gamma \circ \delta' = \delta \circ \gamma \circ \delta' = \delta \circ \delta' \circ \gamma = \delta \circ \gamma.$$

Thus $\gamma \circ \delta \subseteq \delta \circ \gamma$, in contradiction with (iv). This is a contradiction to the assumption that $I[\sigma, \tau]$ contains a pair of nonpermuting congruences.

For the final statement of the theorem, recall that any lattice of permuting equivalence relations is modular. \square

This result may be used to establish a new characterization of varieties with a weak difference term. In the following theorem, a lattice \mathbf{L} has an $\text{SD}_\wedge/\text{Modular}$ **factorization** in a category \mathcal{C} of lattices if there is a \mathcal{C} -morphism of \mathbf{L} onto a meet semidistributive lattice whose kernel classes are modular.

THEOREM 6.17. *The following are equivalent for a variety \mathcal{V} .*

- (1) \mathcal{V} has a weak difference term.
- (2) For every $\mathbf{A} \in \mathcal{V}$, $\mathbf{Con}(\mathbf{A})$ has an $\text{SD}_\wedge/\text{Modular}$ factorization in the category of meet continuous lattices.
- (3) For every $\mathbf{A} \in \mathcal{V}$, $\mathbf{Con}(\mathbf{A})$ has an $\text{SD}_\wedge/\text{Modular}$ factorization in the category of ordinary lattices.
- (4) For every $\mathbf{A} \in \mathcal{V}$, each SD_\wedge -failure in $\mathbf{Con}(\mathbf{A})$ is modular.
- (5) If \mathbf{L} belongs to the meet continuous congruence variety of \mathcal{V} , then each SD_\wedge -failure in \mathbf{L} is modular.

PROOF. If \mathcal{V} has a weak difference term, then the natural homomorphism $\nu: \mathbf{Con}(\mathbf{A}) \rightarrow \mathbf{Con}(\mathbf{A})/\overset{s}{\sim}$ is a homomorphism of meet continuous lattices whose image is meet semidistributive according to Theorem 6.13 (6). The kernel classes are modular according to Theorems 6.16. This proves that (1) \implies (2).

The implications (2) \implies (3) \implies (4) are straightforward; the first since any factorization in the category of meet continuous lattices is a factorization in the category of lattices, and the second because SD_\wedge -failures must be collapsed by any homomorphism into a meet semidistributive lattice.

Assume that SD_\wedge -failures of $\mathbf{Con}(\mathbf{A})$ are modular for any $\mathbf{A} \in \mathcal{V}$. As explained in the third paragraph of Remark 2.24, the class of meet continuous lattices whose SD_\wedge -failures are modular is a definable by a meet continuous identity. If (4) holds, then the congruence lattices of

remove "a"

algebras in \mathcal{V} satisfy this identity, so the meet continuous congruence variety of \mathcal{V} also satisfies it, hence (5) holds.

According to Theorem 4.7 the class of varieties \mathcal{U} that satisfy the the meet continuous congruence identity expressing the fact that SD_\wedge -failures are modular is definable by idempotent Maltsev conditions. To show that (5) \implies (1) we must show that these Maltsev conditions are strong enough to force the existence of a weak difference term. By Theorem 6.2, it is enough to show that if \mathcal{V} is a variety satisfying (5) then the abelian algebras in \mathcal{V} are affine. So let \mathcal{V} be any variety satisfying condition (5), and let \mathcal{A} be the subvariety of \mathcal{V} that is generated by the abelian algebras of \mathcal{V} . \mathcal{A} must also satisfy item (5), and in particular (since the variety of sets does not satisfy (5)) this implies that \mathcal{A} has a Taylor term. Let $\mathbf{F} = \mathbf{F}_{\mathcal{A}}(4)$ be the 4-generated free algebra of this variety. \mathbf{F} is abelian since \mathcal{A} is generated by abelian algebras, the abelian property is preserved by subalgebras and products, and free algebras may be constructed as subalgebras of products of generating algebras. Since \mathbf{F} is abelian, there is a congruence Δ on \mathbf{F}^2 that has the diagonal $D = \{(f, f) \mid f \in F\}$ as a class. It follows from Lemma 4.4 of [52] (in the case $\gamma = \delta = 1$ of that result) that since \mathcal{A} has a Taylor term the congruence Δ is a lattice complement of the coordinate projection kernels $\eta_1, \eta_2 \in \text{Con}(\mathbf{F}^2)$. Since $\Delta \wedge \eta_1 = \Delta \wedge \eta_2 = 0$, the interval $I[0, \Delta \vee (\eta_1 \vee \eta_2)] = I[0, \Delta]$ is an SD_\wedge -failure. By item (5), this interval is modular. Since the diagonal D is a Δ -class, it follows from Lemma 2.4 of [34] that restriction is a lattice homomorphism from the interval $I[0, \Delta]$ to $\text{Con}(\mathbf{F}^2|_D)$, where $\mathbf{F}^2|_D$ is the algebra induced on D by \mathbf{F}^2 , so $\text{Con}(\mathbf{F}^2|_D)$ is modular. The operations of $\mathbf{F}^2|_D$ are the polynomial operations of \mathbf{F}^2 under which D is closed, each of which agrees with some polynomial operation of \mathbf{F} acting coordinatewise on D . Thus, $\mathbf{F}^2|_D$ is isomorphic as a nonindexed algebra to the algebra $\mathbf{F}|_F$ that \mathbf{F} induces on itself. But the congruence lattices of \mathbf{F} and $\mathbf{F}|_F$ are the same, so $\mathbf{F} = \mathbf{F}_{\mathcal{A}}(4)$ has a modular congruence lattice. By Day's Theorem, [8], this implies that \mathcal{A} is congruence modular. Since \mathcal{A} is generated by abelian algebras, modular commutator theory implies that all algebras in \mathcal{A} are affine. \square

this should be "meet".

EXAMPLE 6.18. There exists a purely lattice-theoretic proof that (2) \implies (3) \implies (4) in Theorem 6.17, but not for the reverse implications. Those implications are false even for congruence lattices of individual algebras.

For example, if \mathbf{L} is a lattice of height ≤ 3 , then any SD_\wedge -failure has height ≤ 2 , hence is modular. Thus, any finite, nonmodular, simple lattice of height ≤ 3 has the properties that it is algebraic (hence

isomorphic to the congruence lattice of some algebra), has modular SD_\wedge -failures, but has no $\text{SD}_\wedge/\text{Modular}$ factorization. There are many graph lattices satisfying the required conditions (cf. Section 4.4).

There also exist algebraic lattices with $\text{SD}_\wedge/\text{Modular}$ factorizations in the category of ordinary lattices and no such factorization in the category of meet continuous lattices, such as the one in Figure 6.1.

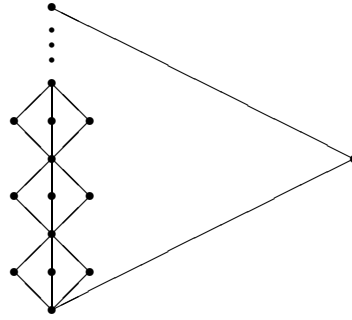


FIGURE 6.1. No $\text{SD}_\wedge/\text{Modular}$ factorization in \mathcal{L}_{MC}

We next show that the relation $\overset{s}{\sim}$ behaves well with respect to the class operators \mathbf{H} , \mathbf{S} and \mathbf{P}_{fin} .

THEOREM 6.19. *Let \mathcal{V} be a variety with a weak difference term.*

- (1) *If $\mathbf{A} \in \mathcal{V}$ and $\alpha, \beta \geq \gamma$ are congruences on \mathbf{A} , then $\alpha \overset{s}{\sim} \beta$ in $\mathbf{Con}(\mathbf{A})$ if and only if $\alpha/\gamma \overset{s}{\sim} \beta/\gamma$ in $\mathbf{Con}(\mathbf{A}/\gamma)$.*
- (2) *If \mathbf{B} is a subalgebra of $\mathbf{A} \in \mathcal{V}$, then $\alpha \overset{s}{\sim} \beta$ in $\mathbf{Con}(\mathbf{A})$ implies that $\alpha|_B \overset{s}{\sim} \beta|_B$ in $\mathbf{Con}(\mathbf{B})$.*
- (3) *If $\mathbf{A} \in \mathcal{V}$ has congruences α and β and $\mathbf{B} \in \mathcal{V}$ has congruence γ and δ , then $\alpha \times \gamma \overset{s}{\sim} \beta \times \delta$ in $\mathbf{Con}(\mathbf{A} \times \mathbf{B})$ if and only if $\alpha \overset{s}{\sim} \beta$ in $\mathbf{Con}(\mathbf{A})$ and $\gamma \overset{s}{\sim} \delta$ in $\mathbf{Con}(\mathbf{B})$.*

In particular, the class of ∞ -solvable algebras in \mathcal{V} is closed under the formation of homomorphic images, subalgebras and finite products.

PROOF. It follows from Theorem 2.19 (10) that if $\gamma \leq \delta \leq \theta$, then $\delta \triangleleft \theta$ if and only if $\delta/\gamma \triangleleft \theta/\gamma$. It is also true that a sequence $(\theta_\lambda)_{\lambda < \kappa}$ of congruences on \mathbf{A} , each containing γ , is continuous if and only if $(\theta_\lambda/\gamma)_{\lambda < \kappa}$ is. The property $\alpha \overset{s}{\sim} \beta$ in $\mathbf{Con}(\mathbf{A})$ is the property that there exists a continuous sequence of congruences witnessing ∞ -solvability from $\alpha \wedge \beta$ to $\alpha \vee \beta$, which is therefore equivalent to the existence of a continuous sequence of congruences witnessing ∞ -solvability from

$(\alpha \wedge \beta)/\gamma = (\alpha/\gamma) \wedge (\beta/\gamma)$ to $(\alpha \vee \beta)/\gamma = (\alpha/\gamma) \vee (\beta/\gamma)$. Hence $\alpha \overset{s}{\sim} \beta$ is equivalent to $\alpha/\gamma \overset{s}{\sim} \beta/\gamma$ when $\gamma \leq \alpha \wedge \beta$. This proves (1).

It follows from Theorem 2.19 (9) that restriction to \mathbf{B} preserves the relation \triangleleft (i.e., $\delta \triangleleft \theta$ in $\mathbf{Con}(\mathbf{A})$ implies $\delta|_{\mathbf{B}} \triangleleft \theta|_{\mathbf{B}}$). Since restriction to \mathbf{B} also preserves unions of chains, it must preserve $\triangleleft\!\!\triangleleft$. Hence ∞ -solvable intervals in $\mathbf{Con}(\mathbf{A})$ restrict to ∞ -solvable intervals in $\mathbf{Con}(\mathbf{B})$. Since $\alpha \overset{s}{\sim} \beta$ if and only if α and β lie in a single ∞ -solvable interval, $\overset{s}{\sim}$ is also preserved under restriction to \mathbf{B} . This proves (2).

Applying item (1) to $\mathbf{A} \times \mathbf{B}$ (with the coordinate projection kernels η_1 and η_2 playing the role of γ in (1)) we obtain that $\alpha \overset{s}{\sim} \beta$ in $\mathbf{Con}(\mathbf{A})$ if and only if $\alpha_1 \overset{s}{\sim} \beta_1$ in $\mathbf{Con}(\mathbf{A} \times \mathbf{B})$ and similarly $\gamma \overset{s}{\sim} \delta$ in $\mathbf{Con}(\mathbf{B})$ if and only if $\gamma_2 \overset{s}{\sim} \delta_2$ in $\mathbf{Con}(\mathbf{A} \times \mathbf{B})$. Since ∞ -solvability is a congruence, the conjunction of $\alpha \overset{s}{\sim} \beta$ and $\gamma \overset{s}{\sim} \delta$ implies that

$$\alpha \times \gamma := \alpha_1 \wedge \gamma_2 \overset{s}{\sim} \beta_1 \wedge \delta_2 =: \beta \times \delta.$$

Conversely, if $\alpha \times \gamma \overset{s}{\sim} \beta \times \delta$, then again because ∞ -solvability is a congruence we get

$$\alpha_1 = \eta_1 \vee (\alpha \times \gamma) \overset{s}{\sim} \eta_1 \vee (\beta \times \delta) = \beta_1,$$

and similarly $\gamma_2 \overset{s}{\sim} \delta_2$. As proved above, these instances of $\overset{s}{\sim}$ hold if and only if $\alpha \overset{s}{\sim} \beta$ in $\mathbf{Con}(\mathbf{A})$ and $\gamma \overset{s}{\sim} \delta$ in $\mathbf{Con}(\mathbf{B})$. \square

By a **partial Maltsev operation** on a set A we mean a function $m: D \rightarrow A$, for some domain $D \subseteq A^3$ containing all triples (x, y, y) and (y, y, x) , $x, y \in A$, which satisfies $m(x, y, y) = x = m(y, y, x)$ for all $x, y \in A$.

THEOREM 6.20. *Let \mathcal{V} be a variety with a weak difference term. Assume that $\mathbf{A} \in \mathcal{V}$, $\theta \in \mathbf{Con}(\mathbf{A})$ is ∞ -solvable, and $S \subseteq A^n[\theta]$ is a subset. If S is compatible with the operations of \mathbf{A} , then S is also compatible with some partial Maltsev operation on A*

PROOF. Recall from the opening paragraph of Chapter 2 that $A^n[\theta]$ is defined to be $\{\mathbf{x} \in A^n \mid x_i \equiv_{\theta} x_j \text{ for all } i, j\}$. Assume that $S \subseteq A^n[\theta]$ is compatible with the operations of \mathbf{A} . We will show that S is compatible with the partial Maltsev operation on A with smallest domain, which is

$$D = \{(a, b, b) \mid a, b \in A\} \cup \{(a, a, b) \mid a, b \in A\}.$$

We must show that if $\mathbf{a}, \mathbf{b}, \mathbf{c} \in S$ and the tuple

$$m(\mathbf{a}, \mathbf{b}, \mathbf{c}) = (m(a_1, b_1, c_1), m(a_2, b_2, c_2), \dots, m(a_n, b_n, c_n)) =: \mathbf{d}$$

is defined in all coordinates, then $\mathbf{d} \in S$. To be defined in all coordinates, it must be the case that $a_i = b_i$ or $b_i = c_i$ (or both) for all i . Assume that this is so. Let $I \subseteq N := \{1, \dots, n\}$ be the set of coordinates where $a_i = b_i$. If $I = \emptyset$, then $\mathbf{b} = \mathbf{c}$, so $\mathbf{d} = \mathbf{a} \in S$. If $I = N$, then $\mathbf{a} = \mathbf{b}$, so $\mathbf{d} = \mathbf{c} \in S$. Thus the only nontrivial cases are those where $\emptyset \neq I \neq N$. Restrict attention to those, and let $J = N - I$.

Let \mathbf{S} be the subalgebra of \mathbf{A}^n that is supported by $S = A^n[\theta]$. Elements $\mathbf{u}, \mathbf{v} \in S$ are related by θ_i ($= \pi_i^{-1}(\theta)$) if and only if $u_i \equiv_{\theta} v_i$. Since $S \subseteq A^n[\theta]$ this implies that $u_j \equiv_{\theta} u_i \equiv_{\theta} v_i \equiv_{\theta} v_j$ for any j , hence that $\theta_i \subseteq \theta_j$ for any i and j . Let $\widehat{\theta}$ denote the congruence $\theta_1 = \theta_2 = \dots = \theta_n$ of \mathbf{S} . It follows from Theorem 6.19 (1) and (2) that $\theta_i \stackrel{s}{\sim} \eta_i$ for each i , so $\widehat{\theta} = \bigwedge \theta_i \stackrel{s}{\sim} \bigwedge \eta_i = 0$, proving that $\widehat{\theta}$ is ∞ -solvable.

If $i \in I$, then $\eta_I \leq \eta_i \leq \theta_i = \widehat{\theta}$, and a similar argument proves that $\eta_J \leq \widehat{\theta}$. Since $\widehat{\theta}$ is ∞ -solvable, $\eta_I \stackrel{s}{\sim} \widehat{\theta} \stackrel{s}{\sim} \eta_J$. This and Theorem 6.16 prove that η_I and η_J are permuting congruences of \mathbf{S} . According to the definition of I and J , we have $\mathbf{a} \eta_I \mathbf{b} \eta_J \mathbf{c}$, so there must exist a tuple $\mathbf{x} \in S$ such that $\mathbf{a} \eta_J \mathbf{x} \eta_I \mathbf{c}$. The tuple \mathbf{x} satisfies the condition that $x_i = c_i$ if $i \in I$ and $x_j = a_j$ if $j \in N - I = J$. That is, $x_i = c_i$ if $m(a_i, b_i, c_i) = m(a_i, a_i, c_i) = c_i$, and $x_j = a_j$ if $m(a_j, b_j, c_j) = m(a_j, c_j, c_j) = a_j$. This proves that $\mathbf{x} = \mathbf{d}$, hence $\mathbf{d} = \mathbf{x} \in S$ as desired. \square

It is a general fact that if \mathbf{A} is an algebra, m is an operation on A , and for every cardinal κ any relation $S \subseteq A^\kappa$ compatible with the operations of \mathbf{A} is also compatible with m , then m is a term operation of \mathbf{A} . If instead one only has that any *finitary* relation S on A that is compatible with the operations of \mathbf{A} is also compatible with m , then m is only guaranteed to be a **local term operation** of \mathbf{A} . This means that for any finite subset $F \subseteq A$ there is a term operation t_F that agrees with m on F . Theorem 6.20 suggests the possibility that if \mathcal{V} has a weak difference term and $\mathbf{A} \in \mathcal{V}$ is ∞ -solvable, then \mathbf{A} has a local Maltsev term operation. Admittedly Theorem 6.20 speaks only about partial Maltsev operations, but that is exactly what one would expect of a result like Theorem 6.20 if the local Maltsev term operation were not unique. These observations lead to speculations recorded in the following problem.

PROBLEM 6.21. Assume that \mathcal{V} has a weak difference term, and that $\mathbf{A} \in \mathcal{V}$ is ∞ -solvable. Does \mathbf{A} have a local Maltsev term operation? Does \mathbf{A} have a Maltsev term operation? If \mathcal{V} is generated by ∞ -solvable algebras, then must \mathcal{V} be congruence permutable?

is a term...

Although affirmative answers to these questions seem unlikely, it is true that for each finite n there is term that is Maltsev on classes of n -step solvable congruences of algebras in \mathcal{V} . It would be interesting if this property extended to congruences with higher degrees of solvability.

COROLLARY 6.22. *Let \mathcal{V} be a variety with a weak difference term. If $\mathbf{A} \in \mathcal{V}$, $\rho \stackrel{s}{\sim} 0$ in $\mathbf{Con}(\mathbf{A})$ and $S \subseteq \rho (= A^2[\rho])$ is a reflexive compatible binary relation on \mathbf{A} , then S is a congruence.*

PROOF. This is proved in the same way that one proves that a reflexive compatible binary relation on a Maltsev algebra is a congruence, since that argument works equally well with a partial Maltsev operation.

For the record, the details are as follows. If $S \subseteq \rho$ is reflexive and compatible, then S is compatible with a partial Maltsev operation according to Theorem 6.20. Thus, if $(a, b) \in S$, then since S is reflexive $(a, a), (b, b) \in S$ also, so $(b, a) = m((a, a), (a, b), (b, b)) \in S$. This proves that S is symmetric. If $(a, b), (b, c) \in S$, then $(a, c) = m((a, b), (b, b), (b, c)) \in S$. This proves that S is transitive, hence a congruence. \square

THEOREM 6.23. *Let \mathcal{V} be a variety with a weak difference term.*

(1) *If $\mathbf{A} \in \mathcal{V}$ has congruences $\alpha' \triangleleft \alpha$ and $\beta' \triangleleft \beta$, then*

$$\alpha \vee \beta = (\alpha' \vee \beta') \circ \alpha \circ \beta \circ (\alpha' \vee \beta').$$

(2) *If $\mathbf{A} \in \mathcal{V}$ has congruences α and β with $\beta \stackrel{s}{\sim} \alpha \wedge \beta$, then*

$$\alpha \vee \beta = \alpha \circ \beta \circ \alpha.$$

PROOF. For item (1), the relation $R = (\alpha' \vee \beta') \circ \alpha \circ \beta \circ (\alpha' \vee \beta')$ is an $(\alpha' \vee \beta')$ -closed compatible reflexive relation on \mathbf{A} whose transitive closure is $\alpha \vee \beta$. To prove this part of the theorem it suffices to show that R is a congruence.

Since $\alpha' \triangleleft \alpha$ and $\beta' \triangleleft \beta$, it follows from Lemma 6.10 and the transitivity of \triangleleft that $(\alpha' \vee \beta') \triangleleft (\alpha \vee \beta)$. In the quotient algebra $\mathbf{A}/(\alpha' \vee \beta')$ we have that

$$\rho := (\alpha \vee \beta)/(\alpha' \vee \beta') \stackrel{s}{\sim} (\alpha' \vee \beta')/(\alpha' \vee \beta') = 0$$

by Theorem 6.19. The relation $S := R/(\alpha' \vee \beta')$ is a compatible reflexive relation of $\mathbf{A}/(\alpha' \vee \beta')$, and $S \subseteq \rho$, hence S is a congruence on $\mathbf{A}/(\alpha' \vee \beta')$ according to Corollary 6.22. The fact that R is $(\alpha' \vee \beta')$ -closed means that for the natural homomorphism $\nu : \mathbf{A} \rightarrow \mathbf{A}/(\alpha' \vee \beta') : a \mapsto a/(\alpha' \vee \beta')$ we have $R = \nu^{-1}(\nu(R))$. But $\nu(R) = R/(\alpha' \vee \beta') = S$ is a congruence, and the inverse image of a congruence is a congruence. Hence R is a congruence.

Item (2) is the special case of item (1) in which $\alpha = \alpha' \geq \beta' = \alpha \wedge \beta$. \square

Theorem 6.16 follows from the case of Theorem 6.23 (1) where $\alpha' = \beta' = \alpha \wedge \beta$.

THEOREM 6.24. *Let \mathcal{V} be a variety with a weak difference term. If $\mathbf{A} \in \mathcal{V}$ has congruences α and β with $\beta \stackrel{s}{\sim} \alpha \wedge \beta$ and $\alpha \leq \gamma \leq \alpha \vee \beta$, then $\gamma = \alpha \vee (\beta \wedge \gamma)$. Hence the map $J : I[\alpha \wedge \beta, \beta] \rightarrow I[\alpha, \alpha \vee \beta] : x \mapsto x \vee \alpha$ is surjective.*

PROOF. The hypotheses of Theorem 6.23 (2) are met, so $\alpha \vee \beta = \alpha \circ \beta \circ \alpha$. Arbitrarily choose $(a, b) \in \gamma$. Since $\gamma \leq \alpha \vee \beta = \alpha \circ \beta \circ \alpha$, there are elements $u, v \in A$ such that

$$(6.5) \quad a \equiv_{\alpha} u \equiv_{\beta} v \equiv_{\alpha} b.$$

This means that $u \equiv_{\alpha} a \equiv_{\gamma} b \equiv_{\alpha} v$, so $(u, v) \in \beta \wedge (\alpha \vee \gamma) = \beta \wedge \gamma$. This allows us to strengthen the β -relation in line (6.5) to a $(\beta \wedge \gamma)$ -relation, and in the strengthened form that line indicates that $(a, b) \in \alpha \vee (\beta \wedge \gamma)$. Since (a, b) was chosen arbitrarily, $\gamma \subseteq \alpha \vee (\beta \wedge \gamma)$. The reverse inclusion follows from the fact that $\alpha \leq \gamma$.

For the final statement of the theorem, note that for any $\gamma \in I[\alpha, \alpha \vee \beta]$ the element $\beta \wedge \gamma$ is in $I[\alpha \wedge \beta, \beta]$ and we have proved above that $\gamma = J(\beta \wedge \gamma)$. \square

THEOREM 6.25. *Let \mathcal{V} be a variety with a weak difference term. Assume that $\mathbf{A} \in \mathcal{V}$.*

- (1) *If \mathbf{N}_5 , with the labeling indicated in Figure 6.2, appears as a sublattice of $\mathbf{Con}(\mathbf{A})$, then $\beta \not\stackrel{s}{\sim} \beta \wedge \alpha$.*

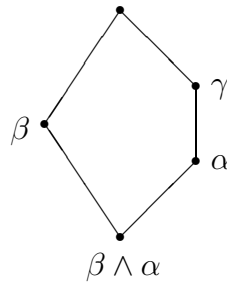


FIGURE 6.2. $\beta \not\stackrel{s}{\sim} \beta \wedge \alpha$

- (2) *If \mathbf{D}_2 appears as a sublattice of $\mathbf{Con}(\mathbf{A})$, then no two distinct congruences in the sublattice are ∞ -solvably related.*

remove

PROOF. In item (1), we have $\alpha \leq \gamma \leq \alpha \vee \beta$. Therefore, if we had $\beta \overset{s}{\sim} \beta \wedge \alpha$, then all of the the hypotheses of Theorem 6.24 would be met. This would force $\gamma = \alpha \vee (\beta \wedge \gamma)$. But $\alpha \vee (\beta \wedge \gamma) = \alpha \neq \gamma$, so we cannot have $\beta \overset{s}{\sim} \beta \wedge \alpha$.

For item (2), the relation $\overset{s}{\sim}$ restricts to a congruence on \mathbf{D}_2 , which is a subdirectly irreducible lattice. If the copy of \mathbf{D}_2 under consideration is labeled as in Figure 6.3, then the critical intervals are $I[\mu, \nu]$, $I[\alpha, \tau]$, and $I[\delta, \theta]$. Therefore, if any two distinct elements of this sublattice are $\overset{s}{\sim}$ -related then $\mu \overset{s}{\sim} \nu$, $\alpha \overset{s}{\sim} \tau$, and $\delta \overset{s}{\sim} \theta$. Assume that this is so

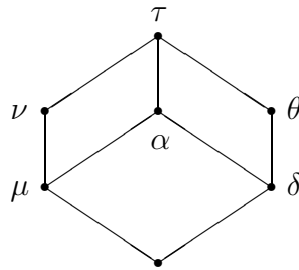


FIGURE 6.3.

le

and choose α' so that $\alpha < \alpha' \leq \tau$ and $\alpha \triangleleft \alpha'$. By Theorem 6.24 the maps $I[\mu, \nu] \rightarrow I[\alpha, \tau]: x \mapsto x \vee \alpha$ and $I[\delta, \theta] \rightarrow I[\alpha, \tau]: x \mapsto x \vee \alpha$ are surjective, so there exist μ' and δ' such that $\mu < \mu' \leq \nu$, $\delta < \delta' < \theta$ and $\mu' \vee \alpha = \alpha' = \delta' \vee \alpha$. Since $I[\mu, \mu']$ and $I[\delta, \delta']$ are perspective with $I[\alpha, \alpha']$ we get that $\mu \triangleleft \mu'$ and $\delta \triangleleft \delta'$ from Theorem 3.26 (or Lemma 6.10).

z, not y

CLAIM 6.26. μ', α , and δ' and generate a sublattice isomorphic to \mathbf{D}_2

Our argument uses Lemma 5.27 with $x = \mu'$, $y = \alpha$ and $z = \delta'$, as well as the fact that ν, α and θ generate a copy of \mathbf{D}_2 .

For item (i) of Lemma 5.27,

$$x \wedge y = \mu' \wedge \delta' \leq \nu \wedge \theta \leq \alpha = y.$$

For item (ii),

$$z \vee (y \wedge x) = \delta' \vee (\alpha \wedge \mu') = \delta' \vee \mu = \alpha' = \alpha \vee \mu' \geq \mu' = x.$$

Here we have used the facts that $I[\mu, \mu']$ and $I[\delta, \delta']$ are perspective up to $I[\alpha, \alpha']$, and that $\delta' \vee \mu = \delta' \vee (\delta \vee \mu) = \delta' \vee \alpha = \alpha'$.

Item (iii) is similar to (ii). For (iv),

$$(x \wedge y) \vee (y \wedge z) = (\mu' \wedge \alpha) \vee (\alpha \wedge \delta') = \mu \vee \delta = \alpha.$$

proves

intervals

Since $x = \mu' \not\leq \alpha = y$, the sublattice generated by $\{\mu', \alpha, \delta'\}$ is isomorphic to \mathbf{D}_2 . This prove the claim.

From the assumption that $\mathbf{Con}(\mathbf{A})$ has a sublattice isomorphic to \mathbf{D}_2 with ∞ -solvable critical interval we have shown that $\mathbf{Con}(\mathbf{A})$ has a sublattice isomorphic to \mathbf{D}_2 with abelian critical interval. This contradicts Theorem 4.16 (2). \square

The rest of the results in this section provide some information about failures of join semidistributivity in the congruence lattice of an algebra with a weak difference term.

LEMMA 6.27. *Assume that \mathcal{V} has a weak difference term. If $\mathbf{A} \in \mathcal{V}$ has congruences α, β and γ satisfying*

- (1) $\alpha \vee \beta \stackrel{s}{\sim} \alpha \stackrel{s}{\sim} \alpha \vee \gamma$, and
- (2) $\beta \wedge \gamma \leq \alpha$,

then $(\beta \vee (\alpha \wedge \gamma)) \wedge (\gamma \vee (\alpha \wedge \beta)) = (\alpha \wedge \beta) \vee (\alpha \wedge \gamma)$.

PROOF. Since $\beta \wedge \gamma$ lies below all congruences of interest, we may factor by it and assume that $\beta \wedge \gamma = 0$.

In the next claim we label some of the congruences in the sublattice of $\mathbf{Con}(\mathbf{A})$ generated by $\{\alpha, \beta, \gamma\}$, and prove some of the more obvious relationships between them.

CLAIM 6.28. *Let $\delta = (\alpha \wedge \beta) \vee (\alpha \wedge \gamma)$, $\widehat{\beta} = \beta \vee \delta$, $\widehat{\gamma} = \gamma \vee \delta$, $\theta = \widehat{\beta} \wedge \widehat{\gamma}$, $\beta' = \beta \wedge \theta$, and $\gamma' = \gamma \wedge \theta$. Then*

- (i) $\delta \leq \alpha$,
- (ii) $\beta \wedge \delta = \beta \wedge \alpha$ and $\gamma \wedge \delta = \gamma \wedge \alpha$,
- (iii) $\beta \stackrel{s}{\sim} \beta \wedge \delta$ and $\gamma \stackrel{s}{\sim} \gamma \wedge \delta$,
- (iv) $\delta \leq \theta \leq \widehat{\beta}$ and $\delta \leq \theta \leq \widehat{\gamma}$,
- (v) $\widehat{\beta} \stackrel{s}{\sim} \delta \stackrel{s}{\sim} \widehat{\gamma}$,
- (vi) $\widehat{\beta} \vee \alpha = \beta \vee \alpha$ and $\widehat{\gamma} \vee \alpha = \gamma \vee \alpha$, and
- (vii) $\beta' \vee \delta = \theta = \gamma' \vee \delta$.

Some of these congruences and relationships are indicated in Figure 6.4.

We omit the proofs of the parts of the claim that concern γ if those parts are symmetric to the arguments we give for β .

lemma

For item (i) of Claim 6.28, note that $\delta = (\alpha \wedge \beta) \vee (\alpha \wedge \gamma)$ is a join of two elements below α , so $\delta \leq \alpha$.

For item (ii), since $\delta \leq \alpha$ we get that $\beta \wedge \delta \leq \beta \wedge \alpha$. On the other hand, both β and δ are above $\beta \wedge \alpha$, so $\beta \wedge \delta \geq \beta \wedge \alpha$.

For item (iii), start with $\alpha \stackrel{s}{\sim} \alpha \vee \beta$, which holds by assumption (1) of the theorem. Using the fact that $\stackrel{s}{\sim}$ is a congruence, meet both

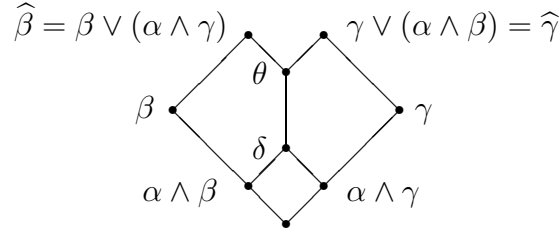


FIGURE 6.4.

sides of $\alpha \stackrel{s}{\sim} \alpha \vee \beta$ with β to obtain $\beta \wedge \alpha \stackrel{s}{\sim} \beta$. Applying the result from (ii) of this claim yields $\beta \wedge \delta \stackrel{s}{\sim} \beta$.

For item (iv), the fact that $\theta = \widehat{\beta} \wedge \widehat{\gamma} = (\beta \vee \delta) \wedge (\gamma \vee \delta)$ is a meet of elements above δ is enough to prove that $\delta \leq \theta$, while $\theta \leq \widehat{\beta}$ is immediate from the definition $\theta = \widehat{\beta} \wedge \widehat{\gamma}$.

For item (v), join the relation $\beta \wedge \delta \stackrel{s}{\sim} \beta$ from (iii) with δ to obtain $\delta \stackrel{s}{\sim} \beta \vee \delta = \widehat{\beta}$.

For item (vi), $\widehat{\beta} \vee \alpha = (\beta \vee \delta) \vee \alpha = \beta \vee (\delta \vee \alpha) = \beta \vee \alpha$, where the last equality follows from item (i).

For item (vii), use the fact established in (iii) that $\beta \stackrel{s}{\sim} \beta \wedge \delta$, the fact established in (iv) that $\delta \leq \theta \leq \widehat{\beta} = \beta \vee \delta$, and Theorem 6.24 to deduce that $\theta = \delta \vee (\beta \wedge \theta)$. Since $\beta' = \beta \wedge \theta$, this yields $\theta = \beta' \vee \delta$. This completes the proof of Claim 6.28.

By examining the definitions of δ and θ , the reader will verify that the statement of the next claim is precisely what is asserted in the statement of this lemma.

CLAIM 6.29. $\delta = \theta$.

We assume instead that $\delta \neq \theta$ and show that the sublattice of $\mathbf{Con}(\mathbf{A})$ generated by $\{\beta', \gamma', \delta\}$ is isomorphic to \mathbf{D}_2 . In order to show this, we verify that the congruences $x = \beta'$, $y = \delta$, and $z = \gamma'$ satisfy the relations from the presentation of \mathbf{D}_2 given in Lemma 5.27. Throughout the argument we will make free use of Claim 6.28.

To establish relation (i) of Lemma 5.27, we have

$$x \wedge z = \beta' \wedge \gamma' \leq \beta \wedge \gamma = 0 \leq \delta = y.$$

For relation (ii),

$$\begin{aligned}
z \vee (y \wedge x) &= \gamma' \vee (\delta \wedge \beta') \\
&= (\gamma \wedge \theta) \vee (\delta \wedge (\beta \wedge \theta)) \\
&= (\gamma \wedge \theta) \vee (\beta \wedge \delta) \\
&= (\gamma \wedge \theta) \vee (\gamma \wedge \delta) \vee (\beta \wedge \delta) \\
&= \gamma' \vee ((\gamma \wedge \alpha) \vee (\beta \wedge \alpha)) \\
&= \gamma' \vee \delta \\
&= \theta \geq \beta \wedge \theta = \beta' = x.
\end{aligned}$$

The proof that $x \vee (y \wedge z) \geq z$ is similar to the one just given. For relation (iv),

$$\begin{aligned}
(x \wedge y) \vee (y \wedge z) &= (\beta' \wedge \delta) \vee (\delta \wedge \gamma') \\
&= ((\beta \wedge \theta) \wedge \delta) \vee (\delta \wedge (\gamma \wedge \theta)) \\
&= (\beta \wedge \delta) \vee (\gamma \wedge \delta) \\
&= (\beta \wedge \alpha) \vee (\gamma \wedge \alpha) \\
&= \delta = y.
\end{aligned}$$

Finally, $x = \beta' \not\leq \delta = y$ since $\beta' \vee \delta = \theta$ by Claim 6.28 (vii), and we are assuming that $\delta \neq \theta$.

Thus, our assumption that $\delta \neq \theta$ and Lemma 5.27 together imply that the sublattice generated by $\{\beta', \gamma', \delta\}$ is isomorphic to \mathbf{D}_2 . It follows from Claim 6.28 (iv) and (v) that $I[\delta, \theta]$ is an ∞ -solvable interval in this sublattice. This is forbidden by Theorem 6.25, hence Claim 6.29 and the lemma are proved. \square

THEOREM 6.30. *Assume that \mathcal{V} has a weak difference term. If $\mathbf{A} \in \mathcal{V}$ has congruences α, β and γ satisfying*

- (1) $\alpha \vee \beta \stackrel{s}{\sim} \alpha \stackrel{s}{\sim} \alpha \vee \gamma$, and
- (2) $\beta \wedge \gamma \leq \alpha$,

then $\mathbf{C}(\alpha \vee \beta, \alpha \vee \gamma; \alpha)$ holds.

PROOF. As in the proof of Lemma 6.27, we may assume that $\beta \wedge \gamma = 0$. As in that proof we let $\delta = (\alpha \wedge \beta) \vee (\alpha \wedge \gamma)$, $\widehat{\beta} = \beta \vee \delta$, and $\widehat{\gamma} = \gamma \vee \delta$.

$$\text{CLAIM 6.31. } \alpha \wedge \widehat{\beta} = \alpha \wedge \widehat{\gamma} = \widehat{\beta} \wedge \widehat{\gamma}.$$

We proved in Lemma 6.27 that $\delta = \widehat{\beta} \wedge \widehat{\gamma}$, so it suffices (by symmetry) to prove only that $\alpha \wedge \widehat{\beta} = \delta$. Both α and $\widehat{\beta}$ are above δ , according to Claim 6.28 (i) and (iv) of Lemma 6.27, so we certainly have $\alpha \wedge \widehat{\beta} \geq \delta$. This forces $\alpha \wedge \widehat{\beta} \in I[\delta, \widehat{\beta}] = I[\delta, \beta \vee \delta]$. The perspective interval $I[\beta \wedge \delta, \beta] = I[\beta \wedge \alpha, \beta]$ is ∞ -solvable by assumption, so Theorem 6.24 guarantees that

$$\alpha \wedge \widehat{\beta} = \delta \vee (\beta \wedge (\alpha \wedge \widehat{\beta})).$$

this is pretty much immediate, no need to refer to earlier claims.

But we can simplify the expression on the right as follows: $\delta \vee (\beta \wedge (\alpha \wedge \widehat{\beta})) = \delta \vee (\alpha \wedge (\beta \wedge \widehat{\beta})) = \delta \vee (\alpha \wedge \beta) = \delta$. This proves the claim.

is a ...

CLAIM 6.32. $\mathbf{C}(\widehat{\beta}, \widehat{\gamma}; \alpha)$ holds.

It follows from Claim 6.31 that $\widehat{\gamma} \wedge (\widehat{\beta} \vee (\widehat{\gamma} \wedge \alpha)) \leq \alpha$, so Claim 6.32 is a direct application of Theorem 2.19 (8).

why distinguish this as a separate claim? this is what is supposed to be proved in the theorem.

CLAIM 6.33. $\mathbf{C}(\alpha \vee \beta, \alpha \vee \gamma; \alpha)$ holds.

It follows from Theorem 2.19 (7) or (8) that $\mathbf{C}(\alpha, \widehat{\gamma}; \alpha)$ holds. From this, Claim 6.32, and Theorem 2.19 (5) we get that $\mathbf{C}(\alpha \vee \widehat{\beta}, \widehat{\gamma}; \alpha)$ holds. We established that $\alpha \vee \widehat{\beta} = \alpha \vee \beta$ in Claim 6.28 (vi) of Lemma 6.27, so $\mathbf{C}(\alpha \vee \beta, \widehat{\gamma}; \alpha)$ holds. Using Theorem 2.19 (3) we finally get that $\mathbf{C}(\alpha \vee \beta, \alpha \circ \widehat{\gamma} \circ \alpha; \alpha)$ holds. This claim, and this theorem, will be established if we show that $\alpha \circ \widehat{\gamma} \circ \alpha = \alpha \vee \widehat{\gamma}$ ($= \alpha \vee \gamma$). For this, meet both sides of the relation $\alpha \vee \gamma \stackrel{s}{\sim} \alpha$, which is part of assumption (1) of the theorem, with the congruence $\widehat{\gamma}$. The result is that $\widehat{\gamma} \stackrel{s}{\sim} \alpha \wedge \widehat{\gamma}$. From Theorem 6.23 (2) we conclude that $\alpha \circ \widehat{\gamma} \circ \alpha = \alpha \vee \widehat{\gamma}$, as desired. \square

There is a special case of the previous theorem that is worth recording.

COROLLARY 6.34. *Assume that \mathcal{V} has a weak difference term, that $\mathbf{A} \in \mathcal{V}$, and that I is an $\text{SD}_{\mathcal{V}}$ -failure in $\mathbf{Con}(\mathbf{A})$. If I is ∞ -solvable, then I is abelian.*

PROOF. An $\text{SD}_{\mathcal{V}}$ -failure is a nontrivial interval $I = I[\alpha, \alpha \vee \beta]$ determined by congruences α, β and γ satisfying $\alpha \vee \beta = \alpha \vee \gamma$ and $\beta \wedge \gamma \leq \alpha$. The assumption that I is ∞ -solvable means that $\alpha \stackrel{s}{\sim} \alpha \vee \beta$ ($= \alpha \vee \gamma$). This establishes hypothesis (1) of Theorem 6.30. Hypothesis (2) holds simply because I is an $\text{SD}_{\mathcal{V}}$ -failure. Hence we have the conclusion of Theorem 6.30, that $\mathbf{C}(\alpha \vee \beta, \alpha \vee \gamma; \alpha)$ (equivalently $\mathbf{C}(\alpha \vee \beta, \alpha \vee \beta; \alpha)$) holds. This is precisely what it means for I to be abelian. \square

develop

6.3. An Alternative Development

we show that it is possible to ...

In the preceding section we developed the theory of solvability in analogy with the theory of solvability for groups, using chains of congruences with abelian factors. In this short section we show that it would have been possible to develop the theory of solvability in analogy with the theory of solvability for rings. Namely, we show that the basic notions of the preceding section, rad , $\stackrel{s}{\sim}$ and \lll , can be defined in terms of prime congruences.

DEFINITION 6.35. Let \mathbf{A} be an algebra. A congruence θ on \mathbf{A} is **semiprime** if $\theta \triangleleft 1$; it is **prime** if it is semiprime and meet irreducible in $\mathbf{Con}(\mathbf{A})$.

An m -**sequence** is a sequence $\sigma = (\theta_\lambda)_{\lambda < \omega}$ of principal congruences where $[\theta_\lambda, \theta_\lambda] \geq \theta_{\lambda+1}$ holds for all λ . The m -sequence σ **avoids** a congruence θ if $\theta_\lambda \not\leq \theta$ for all λ , otherwise it **intersects** θ .

Our definitions for “prime” and “semiprime” agree with those from ring theory. The concept of an m -sequence generalizes the concept of a multiplicatively closed subset of a commutative ring, or an m -system in a noncommutative ring. All of these notions are definable in terms of the centralizer relation.

Observe that if $(\theta_\lambda)_{\lambda < \omega}$ is an m -sequence, then it is descending, since $\theta_\lambda \geq [\theta_\lambda, \theta_\lambda] \geq \theta_{\lambda+1}$ for all λ .

THEOREM 6.36. Let \mathcal{V} be a variety with a weak difference term. Assume that $\mathbf{A} \in \mathcal{V}$, and $\theta \in \mathbf{Con}(\mathbf{A})$. The following conditions concerning θ are equivalent.

- (1) θ is semiprime.
- (2) θ is an intersection of prime congruences.
- (3) For every principal congruence $\theta_0 \not\leq \theta$ there is an m -sequence $\sigma = (\theta_0, \theta_1, \dots)$ starting at θ_0 that avoids θ .
- (4) θ is a radical congruence.

PROOF. It follows immediately from the definitions that (1) and (4) are equivalent, so we prove only that (1) \implies (3) \implies (2) \implies (1).

Assume that θ is semiprime and that $\theta_0 \not\leq \theta$. If $[\theta_0, \theta_0] \leq \theta$, then since $[\theta_0, \theta_0] \triangleleft \theta_0$ we get from Lemma 6.10 that

$$\theta = \theta \vee [\theta_0, \theta_0] \triangleleft \theta \vee \theta_0.$$

But since $\theta < \theta \vee \theta_0$, this contradicts the assumption that $\theta \triangleleft 1$ (i.e., that θ is semiprime). Thus it must be that $[\theta_0, \theta_0] \not\leq \theta$, and this means that there is a principal congruence $\theta_1 \leq [\theta_0, \theta_0]$ such that $\theta_1 \not\leq \theta$. The same argument applied to θ_1 produces a principal congruence θ_2 such that $\theta_2 \leq [\theta_1, \theta_1]$ and $\theta_2 \not\leq \theta$. Repeating the argument indefinitely produces an m -sequence $(\theta_0, \theta_1, \dots)$ starting at θ_0 that avoids θ .

Now suppose that Condition (3) holds for θ . Let θ' be the intersection of the prime congruences containing θ . If $\theta < \theta'$, then there is a principal congruence $\theta_0 \leq \theta'$ such that $\theta_0 \not\leq \theta$. By Condition (3), there is an m -sequence $\sigma = (\theta_0, \theta_1, \dots)$ that avoids θ . Extend θ to a congruence π that is maximal for the property that σ avoids π . (The fact that such a congruence exists uses Zorn's Lemma, and relies on the fact that each θ_i in σ is compact.) We claim that π is a prime congruence.

it isn't obvious why there are any prime congruences above θ (unless i've missed something). so, up front, θ' could be the intersection of the empty set. (which is this equal to $1_{\mathbf{A}}$). in any case, the construction of π shows that there must be prime congruences $\geq \theta$. you might want to address this issue in some manner, since others may have the same problem at this point. you could introduce a lemma that shows that if σ avoids a congruence then it can be extended to one that is prime and that is still avoided by σ . this is used in the next theorem as well.

To see that it is meet irreducible, assume instead that $\pi = \alpha \wedge \beta$ where $\pi < \alpha, \beta$. Then by the maximality of π it must be that σ intersects both α and β , so $\theta_i \leq \alpha$ and $\theta_j \leq \beta$ for some i and j . But then for $k = \max(i, j)$ we have $\theta_k \leq \alpha \wedge \beta = \pi$, since σ is a descending chain of congruences, and this contradicts the fact that σ avoids π . To show that π is semiprime, assume instead that there is a congruence δ such that $\pi \triangleleft \delta$ and $\pi < \delta$. Then the maximality condition on π forces σ to intersect δ , so for some i we have $\theta_i \leq \delta$. But now the monotonicity of the commutator operation implies that

$$\theta_{i+1} \leq [\theta_i, \theta_i] \leq [\delta, \delta] \leq \pi.$$

This contradicts the fact that σ avoids π . This concludes the proof that π is prime. Since π extends θ , and θ' is the intersection of prime congruences containing θ , it follows that $\theta' \leq \pi$. But now we have the contradiction that $\theta_0 \leq \theta' \leq \pi$ while π avoids $\sigma = (\theta_\lambda)_{\lambda < \omega}$. This final contradiction proves that $\theta' = \theta$, and hence that Condition (2) holds.

Finally, assume that Condition (2) holds. Then $\theta = \bigcap_{i \in I} \pi_i$ where each π_i is a prime congruence. If $\theta \triangleleft \delta$, then (since $\theta \leq \pi_i$) we get $\pi_i \triangleleft \pi_i \vee \delta$ for each i . Since each π_i is semiprime we deduce that $\pi_i = \pi_i \vee \delta$, or that $\delta \leq \pi_i$ for all i . Hence $\delta \leq \bigcap_{i \in I} \pi_i = \theta$. This shows that $\theta \triangleleft \delta$ implies that $\theta = \delta$, so $\theta \blacktriangleleft 1$. Hence Condition (1) holds. \square

COROLLARY 6.37. *Let \mathcal{V} be a variety with a weak difference term. Assume that $\mathbf{A} \in \mathcal{V}$ has congruences α and β .*

- (1) $\text{rad}(\alpha)$ equals the intersection of the prime congruences containing α .
- (2) $\alpha \overset{s}{\sim} \beta$ if and only if α and β are contained in the same prime congruences.
- (3) $\alpha \blacktriangleleft \beta$ if and only if $\alpha \leq \beta$ and every m -sequence that intersects β also intersects α .

PROOF. Item (1) follows from the fact, proved in Theorem 6.36, that the radical congruences are exactly the intersections of prime congruences. Item (2) follows from (1) and the fact that $\alpha \overset{s}{\sim} \beta$ if and only if $\text{rad}(\alpha) = \text{rad}(\beta)$.

For item (3), assume that $\alpha \blacktriangleleft \beta$ and that the m -sequence σ intersects β . If σ avoids α , then α can be extended to a prime π that σ avoids, as in the proof of Theorem 6.36 (3) \implies (2). The prime π cannot contain β since σ avoids π but intersects β . Therefore π contains α and not β . Since $\alpha \blacktriangleleft \beta$, this contradicts Item (2) of this corollary. For the other direction, assume that $\alpha \not\blacktriangleleft \beta$. Then, since $\alpha \leq \beta$, Item (2) guarantees that there is a prime π containing α and not containing β . If θ_0 is a principal congruence satisfying $\theta_0 \leq \beta$ and $\theta_0 \not\leq \pi$, then

\sigma avoids \pi

Theorem 6.36 guarantees the existence of an m -sequence starting at θ_0 and avoiding π since the prime π is semiprime. This m -sequence also avoids α since $\alpha \leq \pi$, but it does not avoid β since $\theta_0 \leq \beta$. Thus (3) holds. \square

CHAPTER 7

Ordinary Congruence Identities

In this chapter we will prove that a variety satisfies a nontrivial lattice identity as a congruence identity if and only if it satisfies an idempotent Maltsev condition that fails in the variety of semilattices. In view of the results of Chapter 5, this shows that the class of varieties that satisfy a congruence identity is definable by a Maltsev condition (for example, the Maltsev condition of Theorem 5.28).

7.1. A Rank for Solvability Obstructions

The congruence identity that we describe depends on the number of variables in a Hobby–McKenzie term for the variety, so let us now fix that arity. We assume in this chapter that \mathcal{V} is a variety with a fixed Hobby–McKenzie term $F(x_1, \dots, x_n)$, and we let $N = \{1, 2, \dots, n\}$. By Corollary 6.3, \mathcal{V} also has a weak difference term, so we are free to use the results of Chapter 6 in this chapter.

DEFINITION 7.1. If $I[\delta, \theta]$ is a solvability obstruction and T is a tolerance, then T **separates** $I := I[\delta, \theta]$ if $T \subseteq \theta$ and $T \not\subseteq \delta$. We write $I[\delta, \theta]_T$ or I_T to refer to the obstruction I in a manner that indicates that T separates I .

If f and g are m -ary terms, then I_T **supports the local equation** $f(x_1, \dots, x_m) \approx_{I_T} g(x_1, \dots, x_m)$ if $f(x_1, \dots, x_m) \equiv_{\delta} g(x_1, \dots, x_m)$ holds whenever all $x_i \in \{a, b\}$ for each $(a, b) \in T$.

DEFINITION 7.2. Assume that $\mathbf{A} \in \mathcal{V}$, that $I = I[\delta, \theta]$ is a solvability obstruction, and that T is a tolerance that separates I . The **rank of I_T** is the set

$$\text{Rank}(I_T) := \{U \subseteq N \mid F_U(x, y) \approx_{I_T} x\}.$$

LEMMA 7.3. Assume that $\mathbf{A} \in \mathcal{V}$ and that $I_T := I[\delta, \theta]_T$ is a separated solvability obstruction in $\mathbf{Con}(\mathbf{A})$. Then

- (1) $N \in \text{Rank}(I_T)$.
- (2) $\emptyset \notin \text{Rank}(I_T)$.
- (3) $\text{Rank}(I_T)$ is a closed subset of $\mathcal{B}(F)$ that is not a lattice filter.

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it would be helpful to point the reader to the definition of this. it has been many pages since we last saw this.

PROOF. For (1), $N \in \text{Rank}(I_T)$ if and only if

$$F(x, x, \dots, x) = F_N(x, y) \approx_{I_T} x.$$

That this is true follows from the fact that F is idempotent.

For (2), if $\emptyset \in \text{Rank}(I_T)$, then

$$y = F(y, y, \dots, y) = F_\emptyset(x, y) \approx_{I_T} x.$$

This statement means that $y \equiv_\delta x$ for all $(x, y) \in T$, contradicting the fact that T separates $I[\delta, \theta]$. Thus $\emptyset \notin \text{Rank}(I_T)$,

To prove (3), notice that if $\mathcal{V} \models F_U(x, y) \approx F_V(x, y)$, then $U \in \text{Rank}(I_T)$ if and only if $V \in \text{Rank}(I_T)$. Therefore $\text{Rank}(I_T)$ is a closed subset of $\mathcal{B}(F)$. It is not a lattice filter since, by parts (2) and (3), it is proper and nonempty, and $\mathcal{B}(F)$ has no closed, proper, nonempty lattice filter. \square

A certain configuration involving a pair of separated solvability obstructions forces one obstruction to have rank that properly contains the rank of the other obstruction. We identify this configuration now, and introduce a shorthand notation for it.

DEFINITION 7.4. $I[\mu, \nu]_S \rightsquigarrow I[\delta, \theta]_T$ means that

- (a) $S \cap (\delta : T) \subseteq \mu$,
- (b) $T \cap (S \circ (T \cap \delta) \circ (S \cap \mu)) \subseteq \delta$, and
- (c) $S \cap ((T \cap \delta) \circ (S \cap \mu) \circ (T \cap \delta)) \subseteq \mu$.

The three conditions of this definition describe a property of the tolerance

$$\bar{T} := T \cap (S \circ (T \cap \delta) \circ S).$$

This tolerance consists of all pairs (c, d) for which there exist elements a and b related as in the quadrangle in Figure 7.1. Condition (a) implies

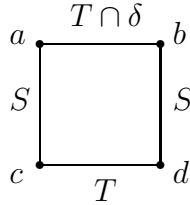


FIGURE 7.1.

that $\mathbf{C}(S, T; \delta)$ fails (since $S \not\subseteq \mu$). An S, T -matrix $\begin{bmatrix} a & b \\ c & d \end{bmatrix}$ witnessing this failure produces a pair $(c, d) \in \bar{T}$ with $(c, d) \notin \delta$. Thus, if Condition (a) holds, the tolerance $\bar{T} \subseteq T$ also separates $I[\delta, \theta]$. Condition (b)

implies that if a, b, c and d are any four elements related as in the above quadrangle, and $(b, d) \in \mu$, then $(c, d) \in \delta$. If this happens, and Condition (c) holds, then $(a, c) \in \mu$. In particular, if $I[\mu, \nu]_S \rightsquigarrow I[\delta, \theta]_T$, then there will exist quadrangles related as above with $(c, d) \notin \delta$, but any such quadrangle must have $(a, c), (b, d) \notin \mu$.

LEMMA 7.5. *Assume that $\mathbf{A} \in \mathcal{V}$, that $I_S = I[\mu, \nu]_S$ and $J_T = I[\delta, \theta]_T$ are separated solvability obstructions in $\mathbf{Con}(\mathbf{A})$, and that $I_S \rightsquigarrow J_T$. If $\bar{S} \subseteq S$ also separates $I[\mu, \nu]$ and $\bar{T} = T \cap (\bar{S} \circ (T \cap \delta) \circ \bar{S})$, then*

- (1) $I_{\bar{S}} \rightsquigarrow J_T$, and
- (2) \bar{T} separates $I[\delta, \theta]$.

PROOF. For item (1), suppose that $\bar{S} \subseteq S$ and that \bar{S} separates I . Then the three conditions of the definition of \rightsquigarrow hold with \bar{S} in place of S , because the symbol S appears only on the left side of each inclusion and the left side of each inclusion is an expression that is monotone in S with respect to inclusion.

Item (2) follows from $I_{\bar{S}} \rightsquigarrow J_T$ as we explained in the remarks preceding the statement of this lemma (using S instead of \bar{S}). \square

LEMMA 7.6. *Assume that $\mathbf{A} \in \mathcal{V}$, that $I_S = I[\mu, \nu]_S$ and $J_T = I[\delta, \theta]_T$ are separated solvability obstructions in $\mathbf{Con}(\mathbf{A})$, and that $I_S \rightsquigarrow J_T$. Let $\bar{T} = T \cap (S \circ (T \cap \delta) \circ S)$. If $f(x, y, z)$ is an idempotent term of \mathbf{A} , and I_S supports the local equation $f(x, y, y) \approx_{I_S} x$, then both of the local equations*

$$f(x, x, y) \approx_{J_{\bar{T}}} x \quad \text{and} \quad f(x, y, x) \approx_{J_{\bar{T}}} x$$

are supported by $J_{\bar{T}}$.

PROOF. The fact that the tolerance $\bar{T} := T \cap (S \circ (T \cap \delta) \circ S)$ separates J follows from Condition (a) defining the relation $I_S \rightsquigarrow J_T$, as we remarked prior to the statement of Lemma 7.5 (or, it follows from Lemma 7.5 (2)). Now, to prove this lemma, it suffices to show that if I_S supports the local equation $f(x, y, y) \approx_{I_S} x$, then $J_{\bar{T}}$ supports the local equation $f(x, x, y) \approx_{J_{\bar{T}}} x$. The same argument applied to $f'(x, y, z) = f(x, z, y)$ shows that if the local equation $f(x, y, y) \approx_{I_S} x$ is supported by I_S , then $f(x, y, x) \approx_{J_{\bar{T}}} x$ is supported by $J_{\bar{T}}$.

To show that $J_{\bar{T}}$ supports the local equation $f(x, x, y) \approx_{J_{\bar{T}}} x$, choose a pair $(0, 1) \in \bar{T} = T \cap (S \circ (T \cap \delta) \circ S)$. There exist elements u and v such that $0 \equiv_S u \equiv_{T \cap \delta} v \equiv_S 1$. The relationships between $0, 1, u$ and v are described in the left hand square in Figure 7.2. The right hand square of this figure describes the relationships between four other elements.

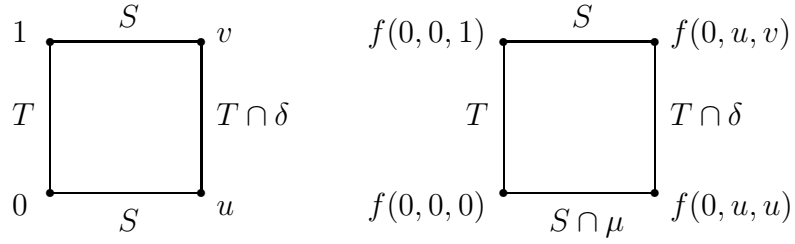


FIGURE 7.2.

Let's convince ourselves that the information indicated in the right hand square of Figure 7.2 is correct. First,

$$f(0, 0, 0) \equiv_T f(0, 0, 1)$$

follows from the fact that $(0, 0), (0, 1) \in T$. Second,

$$f(0, u, u) \equiv_{T \cap \delta} f(0, u, v)$$

follows from the fact that $(0, 0), (u, u), (u, v) \in T \cap \delta$. Third,

$$f(0, 0, 1) \equiv_S f(0, u, v) \text{ and } f(0, 0, 0) \equiv_S f(0, u, u)$$

follow from $(0, 0), (0, u), (1, v) \in S$. Finally,

$$f(0, 0, 0) \equiv_\mu f(0, u, u)$$

since both sides are μ -related to 0: the left hand side equals 0 since f is idempotent and the right hand side is μ -related to 0 since $(0, u) \in S$ and $I[\mu, \nu]_S$ supports $f(x, y, y) \approx_{I[\mu, \nu]_S} x$.

The relations that we have just verified imply that

$$(f(0, 0, 1), f(0, 0, 0)) \in T \cap (S \circ (T \cap \delta) \circ (S \cap \mu)).$$

By Condition (b) of the definition of $I_S \rightsquigarrow J_T$ (Definition 7.4) this implies that

$$f(0, 0, 1) \equiv_\delta f(0, 0, 0) (= 0)$$

holds for any pair $(0, 1) \in \bar{T}$. Hence $J_{\bar{T}}$ supports the local equation $f(x, x, y) \approx_{J_{\bar{T}}} x$, and the proof is complete. \square

LEMMA 7.7. *Assume that $\mathbf{A} \in \mathcal{V}$, that $I_S = I[\mu, \nu]_S$ and $J_T = I[\delta, \theta]_T$ are separated solvability obstructions in $\mathbf{Con}(\mathbf{A})$, and that $I_S \rightsquigarrow J_T$. Let $\bar{T} = T \cap (S \circ (T \cap \delta) \circ S)$. If $f(x, y, z)$ is an idempotent term of \mathbf{A} , and I_S supports the local equations*

$$f(x, x, y) \approx_{I_S} x \text{ and } f(x, y, x) \approx_{I_S} x,$$

then $f(x, y, y) \approx_{J_{\bar{T}}} x$ is supported by $J_{\bar{T}}$.

or just I_S.

PROOF. As in the last proof, choose $(0, 1) \in \bar{T} = T \cap (S \circ (T \cap \delta) \circ S)$, and let u and v denote elements of A for which $0 \equiv_S u \equiv_{T \cap \delta} v \equiv_S 1$. As before, relationships between $0, 1, u$ and v are expressed in Figure 7.3. In this proof we will be interested in the relationships between some

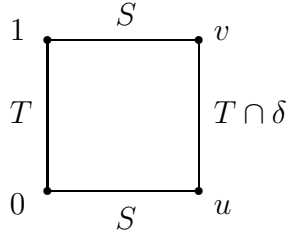


FIGURE 7.3.

of the elements depicted in Figure 7.4. In this figure each element has

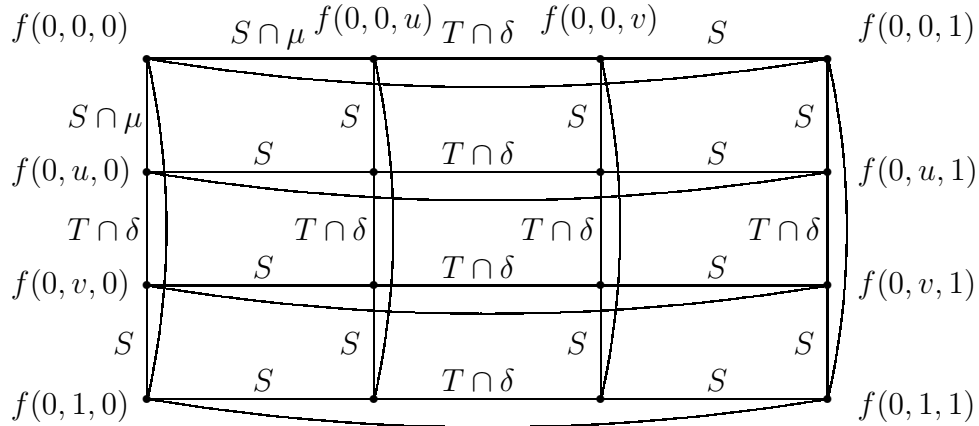


FIGURE 7.4.

the form $f(0, a, b)$. As one moves row-by-row from top to bottom the value of a changes from 0 to u to v to 1 . As one moves column-by-column from left to right the value of b changes from 0 to u to v to 1 . All but two of the indicated relations follow from $(0, u), (1, v) \in S$, $(u, v) \in T \cap \delta$, and that S and $T \cap \delta$ are compatible. The exceptions are the two instances of $S \cap \mu$ in the upper left hand corner of Figure 7.4. The relevant pairs are clearly S -related, since $(0, u) \in S$ and S is compatible, so we must justify the claim that

$$f(0, u, 0) \equiv_{\mu} f(0, 0, 0) \equiv_{\mu} f(0, 0, u).$$

For this one can use the facts that $I_{[\mu, \nu]_S}$ supports

$$f(x, y, x) \approx_{I_{[\mu, \nu]_S}} x = f(x, x, x) \approx_{I_{[\mu, \nu]_S}} f(x, x, y),$$

that $(0, u) \in S$, and that S is a tolerance. The curved lines in the figure represent T -relations, and all follow from $(0, 1) \in T$ and the fact that T is a tolerance. We must prove that $f(0, 1, 1) \equiv_\delta 0$. We will accomplish this by showing that the curved lines of Figure xxx that appear across the top and along the right hand side represent $T \cap \delta$ -relations, and not merely T -relations.

We begin our work by examining the top row of Figure 7.4. Following the straight line from right to left we see that $(f(0, 0, 1), f(0, 0, 0)) \in S \circ (T \cap \delta) \circ (S \cap \mu)$, while following the curved path we see that these elements are related by T . By Condition (b) of the definition of $I_S \rightsquigarrow J_T$ we get that $(f(0, 0, 1), f(0, 0, 0)) \in \delta$. We can apply this information to the rectangle formed from the leftmost and rightmost elements of the first two rows, as shown in Figure 7.5. Condition (b) of the definition

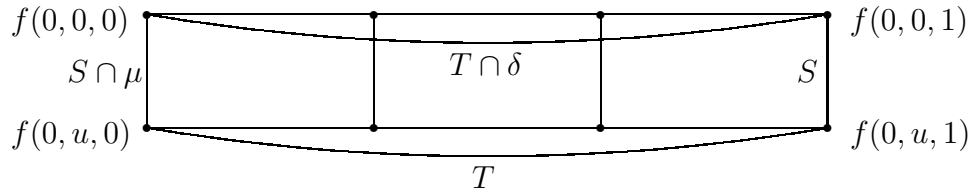


FIGURE 7.5.

of $I_S \rightsquigarrow J_T$ can be applied again to show that

$$(f(0, u, 1), f(0, u, 0)) \in T \cap (S \circ (T \cap \delta) \circ (S \cap \mu)) \subseteq \delta,$$

so the bottom curved line in this figure represents a δ -relation. Now we see that the two elements on the right hand side are related by $S \cap ((T \cap \delta) \circ (S \cap \mu) \circ (T \cap \delta))$. Condition (c) of the definition of $I_S \rightsquigarrow J_T$ asserts that $S \cap ((T \cap \delta) \circ (S \cap \mu) \circ (T \cap \delta)) \subseteq \mu$, so $((f(0, 0, 1), f(0, u, 1)) \in \mu$.

Returning to Figure 7.4, we examine the four elements of the right hand side:

$$f(0, 0, 1) \equiv_{S \cap \mu} f(0, u, 1) \equiv_{T \cap \delta} f(0, v, 1) \equiv_S f(0, 1, 1),$$

and $f(0, 0, 1) \equiv_T f(0, 1, 1)$. Hence

$$(f(0, 1, 1), f(0, 0, 1)) \in T \cap (S \circ (T \cap \delta) \circ (S \cap \mu)) \subseteq \delta.$$

We have shown that the curved lines across the top and down the right side are δ -relations, so we have

$$f(0, 0, 0) \equiv_{\delta} f(0, 0, 1) \equiv_{\delta} f(0, 1, 1).$$

Since f is idempotent, this proves that $f(0, 1, 1) \equiv_{\delta} f(0, 0, 0) = 0$, as desired. \square

Next is the promised result showing that instances of \rightsquigarrow are related to proper increases in rank.

THEOREM 7.8. *Assume that $\mathbf{A} \in \mathcal{V}$, that $I_S = I[\mu, \nu]_S$ and $J_T = I[\delta, \theta]_T$ are separated solvability obstructions in $\mathbf{Con}(\mathbf{A})$, and that $I_S \rightsquigarrow J_T$. If $\overline{T} = T \cap (S \circ (T \cap \delta) \circ S)$, then $\text{Rank}(I_S) \subsetneq \text{Rank}(J_{\overline{T}})$. In fact,*

- (1) *if $U \in \text{Rank}(I_S)$ and $U \subseteq V \subseteq N$, then $V \in \text{Rank}(J_{\overline{T}})$, and*
- (2) *if $U, V \in \text{Rank}(I_S)$ and $U \cup V = N$, then $U \cap V \in \text{Rank}(J_{\overline{T}})$.*

PROOF. It follows from Lemma 7.5 that \overline{T} separates J , so $\text{Rank}(J_{\overline{T}})$ is defined. We will prove that (1) and (2) hold first, and then we will deduce that $\text{Rank}(I_S) \subsetneq \text{Rank}(J_{\overline{T}})$. The following notation will be helpful: Given a partition of $N = \{1, \dots, n\}$ into three sets P, Q and R , one or two of which may be empty, we write $F_{P/Q/R}(x, y, z)$ to denote the term that results from substituting x for x_i in the Hobby–McKenzie term $F(x_1, \dots, x_n)$ when $i \in P$, y for x_i when $i \in Q$, and z for x_i when $i \in R$. Any $F_{P/Q/R}(x, y, z)$ is idempotent, since F is.

We begin the proof of (1). Assume that $U \in \text{Rank}(I_S)$ and $U \subseteq V \subseteq N$. Partition N as $P/Q/R$ with $P = U, Q = V - U$ and $R = N - V$.

The fact that $U \in \text{Rank}(I_S)$ means precisely that the local equation

$$\rightarrow F_{U/(V-U)/(N-V)}(x, y, y) = F_U(x, y) \approx_{I_S} x$$

is supported by I_S . From Lemma 7.6 we get that $J_{\overline{T}}$ supports

$$F_V(x, y) = \rightarrow F_{U/(V-U)/(N-V)}(x, x, y) \approx_{J_{\overline{T}}} x.$$

Hence $V \in \text{Rank}(J_{\overline{T}})$.

Now we prove (2). Assume that $U, V \in \text{Rank}(I_S)$ and $U \cup V = N$, Partition N as $P/Q/R$ with $P = U \cap V, Q = U - V$ and $R = V - U$. The fact that $U, V \in \text{Rank}(I_S)$ translates into

$$F_{(U \cap V)/(U - V)/(V - U)}(x, x, y) = F_U(x, y) \approx_{I_S} x$$

and

$$F_{(U \cap V)/(U - V)/(V - U)}(x, y, x) = F_V(x, y) \approx_{I_S} x,$$

so Lemma 7.7 proves that

$$F_{U \cap V}(x, y) = F_{(U \cap V)/(U - V)/(V - U)}(x, y, y) \approx_{J_{\overline{T}}} x$$

is supported by $J_{\overline{T}}$. Hence $U \cap V \in \text{Rank}(J_{\overline{T}})$.

you could use P/Q/R here.

Now we use (1) and (2) to prove that $\text{Rank}(I_S) \subsetneq \text{Rank}(J_{\overline{T}})$. It follows from (1) that $U \in \text{Rank}(I_S)$ implies $U \in \text{Rank}(J_{\overline{T}})$, so we certainly have $\text{Rank}(I_S) \subseteq \text{Rank}(J_{\overline{T}})$. Suppose that $\text{Rank}(I_S) = \text{Rank}(J_{\overline{T}})$. Then from (1) we get that $\text{Rank}(I_S)$ is an order filter in $\mathcal{B}(F)$. If $U, V \in \text{Rank}(I_S)$ and $U \not\subseteq V$, then since V is the intersection of the $(n-1)$ -element sets that contain it there must be an $(n-1)$ -element set $W \supseteq V$ such that $U \not\subseteq W$. For this set W we have $W \in \text{Rank}(I_S)$ and $U \cup W = N$, so by (2) we get that $U \cap W \in \text{Rank}(J_{\overline{T}}) = \text{Rank}(I_S)$. Since $U \cap W$ is a proper subset of U , and is in $\text{Rank}(I_S)$, this shows that if U is not the least element of $\text{Rank}(I_S)$, then U is not minimal under inclusion in $\text{Rank}(I_S)$. It follows that $\text{Rank}(I_S)$ is a principal order filter, hence a lattice filter in $\mathcal{B}(F)$. But by Lemma 7.3 (3) the set $\text{Rank}(I_S)$ is not a lattice filter, hence $\text{Rank}(I_S) = \text{Rank}(J_{\overline{T}})$ is impossible. This completes the proof that $\text{Rank}(I_S) \subsetneq \text{Rank}(J_{\overline{T}})$. \square

Since an instance of the \rightsquigarrow -relation leads to a proper increase in rank, and there are only finitely many ranks, one expects a bound on the length of \rightsquigarrow -sequences. That bound is provided by the next theorem.

THEOREM 7.9. *Assume that $\mathbf{A} \in \mathcal{V}$. $\mathbf{Con}(\mathbf{A})$ does not contain a \rightsquigarrow -sequence of separated solvability obstructions*

$$I_{S^0}^0 \rightsquigarrow I_{S^1}^1 \rightsquigarrow I_{S^2}^2 \rightsquigarrow \cdots \rightsquigarrow I_{S^k}^k$$

of length $k \geq 2n - 4$.

PROOF. Suppose otherwise, and that $I^j = I[\mu^j, \nu^j]$. Define tolerances $\overline{S}^j \subseteq S^j$ according to the rules

- (i) $\overline{S}^0 := S^0$, and
- (ii) $\overline{S}^{j+1} := S^{j+1} \cap (\overline{S}^j \circ (S^{j+1} \cap \mu^{j+1}) \circ \overline{S}^j)$.

Now consider two sequences of claims:

- (a)^j $I_{\overline{S}^j}^j \rightsquigarrow I_{S^{j+1}}^{j+1}$.
- (b)^j \overline{S}^j separates I^j .

Claims (a)⁰ and (b)⁰ hold since $I_{S^0}^0 \rightsquigarrow I_{S^1}^1$ and $\overline{S}^0 = S^0$ separates I^0 . Claim (a)^j implies Claim (b)^{j+1} by Lemma 7.5 (2). Claim (b)^j implies Claim (a)^j by Lemma 7.5 (1) and the facts that $\overline{S}^j \subseteq S^j$ and $I_{S^j}^j \rightsquigarrow I_{S^{j+1}}^{j+1}$. By induction, Claims (a)^j and (b)^j hold for all j . Applying Theorem 7.8 to Claim (a)^j yields that $\text{Rank}(I_{\overline{S}^j}^j) \subsetneq \text{Rank}(I_{S^{j+1}}^{j+1})$ for each $j < k$. This shows that

$$I_{S^0}^0 \rightsquigarrow I_{S^1}^1 \rightsquigarrow I_{S^2}^2 \rightsquigarrow \cdots \rightsquigarrow I_{S^k}^k$$

implies that

$$\text{Rank}(I_{\overline{S}^0}^0) \subsetneq \text{Rank}(I_{\overline{S}^1}^1) \subsetneq \text{Rank}(I_{\overline{S}^2}^2) \subsetneq \cdots \subsetneq \text{Rank}(I_{\overline{S}^k}^k)$$

where the tolerance \overline{S}^j is contained in S^j yet still separates I^j .

Let $\mathcal{U}^j := \text{Rank}(I_{\overline{S}^j}^j)$. In the last paragraph we showed that

$$\mathcal{U}^0 \subsetneq \mathcal{U}^1 \subsetneq \mathcal{U}^2 \subsetneq \cdots \subsetneq \mathcal{U}^k.$$

Other properties of the sets \mathcal{U}^j include:

- Each \mathcal{U}^j is a closed subset of $\mathcal{B}(F)$ that is not a lattice filter (by Lemma 7.3 (3)).
- $N \in \mathcal{U}^j$ and $\emptyset \notin \mathcal{U}^j$ for all j (by Lemma 7.3 (1) and (2)).
- \mathcal{U}^{j+1} contains the order filter generated by \mathcal{U}^j (by Lemma 7.8 (1)).
- \mathcal{U}^{j+1} contains the intersections of comaximal pairs from \mathcal{U}^j (by Lemma 7.8 (2)).

We now argue that such a sequence of sets cannot exist if $k \geq 2n - 4$. First, \mathcal{U}^0 contains N , is closed, and is not a lattice filter, so either \mathcal{U}^0 contains a set of size $\leq n - 2$ or else at least 2 sets of size $n - 1$. In either case, since \mathcal{U}^1 contains the order filter generated by \mathcal{U}^0 and contains the intersections of comaximal pairs in \mathcal{U}^0 , the set \mathcal{U}^1 properly contains an order filter generated by a set U of size $n - 2$. Now suppose that for some $r \geq 0$ the set \mathcal{U}^{2r+1} contains a principal order filter generated by a set U of size $\leq n - 2 - r$. If U is not empty, then there is some $W \in \mathcal{U}^{2r+1}$ such that $W \not\supseteq U$ (since \mathcal{U}^{2r+1} is closed). The set $V := U \cup (N - W)$ is in \mathcal{U}^{2r+1} since $V \supseteq U$, and $V \cup W = N$ since $V \supseteq N - W$. Thus $V \cap W = U \cap W \in \mathcal{U}^{2r+2} = \mathcal{U}^{2(r+1)}$. Since $W \not\supseteq U$, the set $U \cap W$ has size $\leq n - 2 - (r + 1)$. The set $\mathcal{U}^{2(r+1)+1}$ contains the order filter generated by $U \cap W$. By induction we see that for all r the set \mathcal{U}^{2r+1} contains a principal order filter generated by a set of size $\leq n - 2 - r$. For $r = n - 3$ we get that \mathcal{U}^{2n-5} contains a principal order filter generated by a set U' of size 1. The containment is proper, since \mathcal{U}^{2n-5} is not a lattice filter. If $W' \in \mathcal{U}^{2n-5}$ does not contain U' , then as above we get that $V' := U' \cup (N - W') \in \mathcal{U}^{2n-5}$ and $V' \cap W' = \emptyset \in \mathcal{U}^{2n-4} \subseteq \mathcal{U}^k$. This is a contradiction, since no \mathcal{U}^j contains the empty set. This completes the proof that there is no sequence

$$I_{S^0}^0 \rightsquigarrow I_{S^1}^1 \rightsquigarrow I_{S^2}^2 \rightsquigarrow \cdots \rightsquigarrow I_{S^k}^k$$

with $k \geq 2n - 4$. □

Now that we have shown that there is a bound on the length of \rightsquigarrow -sequences in congruence lattices of algebras in \mathcal{V} , our next goal is to determine how to recognize \rightsquigarrow -related intervals lattice-theoretically.

LEMMA 7.10. *Assume that $\mathbf{A} \in \mathcal{V}$, and that $I_S = I[\mu, \nu]_S$ and $J_T = I[\delta, \theta]_T$ are separated solvability obstructions in $\mathbf{Con}(\mathbf{A})$. Choose and fix congruences δ' and μ' satisfying $T \cap \delta \subseteq \delta' \subseteq \text{Cg}^{\mathbf{A}}(T)$ and $S \cap \mu \subseteq \mu' \subseteq \text{Cg}^{\mathbf{A}}(S)$. If the conditions*

- (d) $\text{Cg}^{\mathbf{A}}(T) \cap \text{Cg}^{\mathbf{A}}(S) \subseteq \delta' \vee \mu'$,
- (e) $T \cap (\delta' \vee \mu') \subseteq \delta$, and
- (f) $S \cap (\delta' \vee \mu') \subseteq \mu$

are satisfied, then the conditions

- (b) $T \cap (S \circ (T \cap \delta) \circ (S \cap \mu)) \subseteq \delta$, and
- (c) $S \cap ((T \cap \delta) \circ (S \cap \mu) \circ (T \cap \delta)) \subseteq \mu$

from Definition 7.4 are also satisfied.

PROOF. Condition (c) follows immediately from Condition (f), since $T \cap \delta \subseteq \delta'$ and $S \cap \mu \subseteq \mu'$.

We will prove that (b) follows from (d) and (e) by showing that $T \cap (S \circ (T \cap \delta) \circ (S \cap \mu)) \subseteq T \cap (\delta' \vee \mu')$. Choose $(c, d) \in T \cap (S \circ (T \cap \delta) \circ (S \cap \mu))$. There exist a and b such that $c \equiv_S a \equiv_{T \cap \delta} b \equiv_{S \cap \mu} d$, as shown in Figure 7.6. The pair (c, b) belongs to $S \circ (T \cap \delta) \subseteq \text{Cg}^{\mathbf{A}}(S) \circ \delta'$

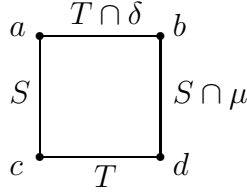


FIGURE 7.6.

and also to $T \circ (S \cap \mu) \subseteq \text{Cg}^{\mathbf{A}}(T) \circ \mu'$. Using Theorem 5.28 (2) (to go from the first line below to the second) and Condition (d) of this theorem we get that

$$\begin{aligned}
 (c, b) &\in (\text{Cg}^{\mathbf{A}}(T) \circ \mu') \cap (\text{Cg}^{\mathbf{A}}(S) \circ \delta') \\
 &\subseteq ((\text{Cg}^{\mathbf{A}}(T) \vee \delta') \wedge (\text{Cg}^{\mathbf{A}}(S) \vee \mu')) \vee \delta' \vee \mu' \\
 &\leq (\text{Cg}^{\mathbf{A}}(T) \wedge \text{Cg}^{\mathbf{A}}(S)) \vee \delta' \vee \mu' \\
 &\leq \delta' \vee \mu'.
 \end{aligned}$$

(Here we used the facts that $\text{Cg}^{\mathbf{A}}(T) \supseteq \delta'$ and $\text{Cg}^{\mathbf{A}}(S) \supseteq \mu'$ to replace $(\text{Cg}^{\mathbf{A}}(T) \vee \delta') \wedge (\text{Cg}^{\mathbf{A}}(S) \vee \mu')$ by $\text{Cg}^{\mathbf{A}}(T) \wedge \text{Cg}^{\mathbf{A}}(S)$.) Hence $c \equiv_{\delta' \vee \mu'} b \equiv_{\mu'} d$, so $(c, d) \in T \cap (\delta' \vee \mu')$. Since $(c, d) \in T \cap (S \circ (T \cap \delta) \circ (S \cap \mu))$ was chosen arbitrarily, we get that

$$T \cap (S \circ (T \cap \delta) \circ (S \cap \mu)) \subseteq T \cap (\delta' \vee \mu').$$

is it the case that up until this point, the fact that the intervals are solvability obstructions has not been used? if so, it might be worth mentioning this.

With this and Condition (e) we get Condition (b). This completes the proof. \square

The following lemma concerns the behavior of the \blacktriangleleft -relation in varieties with a weak difference term, so this lemma could have appeared in Chapter 6. It appears here instead because its primary function is to aid in the proof of the succeeding lemma, which further simplifies the task of recognizing some instances of the \rightsquigarrow relation.

LEMMA 7.11. *Let \mathcal{U} be a variety with a weak difference term, $\mathbf{A} \in \mathcal{U}$, and $I[\delta, \theta]_T$ a separated solvability obstruction in $\mathbf{Con}(\mathbf{A})$.*

- (1) $(\delta : T)$ is the largest congruence $\xi \in \mathbf{Con}(\mathbf{A})$ such that $\xi \geq \delta$ and $\xi \cap T \subseteq \delta$.
- (2) $(\delta : T)$ is a radical congruence.
- (3) If ρ is a radical congruence and $\nu \not\leq \rho$, then $\nu \wedge \rho \blacktriangleleft \nu$.

ν is any congruence?

PROOF. Recall that $(\delta : T)$ is defined to be the largest congruence ξ such that $\mathbf{C}(\xi, T; \delta)$ holds, so in particular $\mathbf{C}((\delta : T), T; \delta)$ holds. If $T' := (\delta : T) \cap T$, then the monotonicity of the centralizer relation guarantees that $\mathbf{C}(T', T'; \delta)$ holds. If $T'' := \delta \circ T' \circ \delta$ is the δ -closure of T' , then Lemma 6.7 proves that T'' is a congruence containing δ for which $\mathbf{C}(T'', T''; \delta)$ holds. Therefore $\delta \triangleleft T'' \subseteq \text{Cg}^{\mathbf{A}}(T \cup \delta) \subseteq \theta$, which (since $I[\delta, \theta]$ is a solvability obstruction) forces $T'' = \delta$. Hence $(\delta : T) \cap T = T' \subseteq T'' = \delta$, proving $(\delta : T)$ is one of the congruences above δ whose intersection with T is contained in δ .

Now suppose that ξ is any congruence above δ whose intersection with T is contained in δ . These properties of ξ may be used to show that

$$T \cap (\xi \circ (T \cap \delta) \circ \xi) \subseteq T \cap \xi \subseteq \delta.$$

By Theorem 2.19 (7) it follows that $\mathbf{C}(\xi, T; \delta)$, or $\xi \leq (\delta : T)$. This proves item (1).

For item (2) assume instead that $(\delta : T) \triangleleft \alpha \leq 1$ for some congruence α . Let $T' = \alpha \cap T$. Since $\mathbf{C}(\alpha, \alpha; (\delta : T))$ holds, we also have $\mathbf{C}(T', T'; (\delta : T))$ by monotonicity of the centralizer. We will derive that $\mathbf{C}(T', T'; \delta)$ holds using Theorem 2.19 (4). To do this we need to prove that $T' \cap (\delta : T) = T' \cap \delta$. The inclusion \supseteq follows from $(\delta : T) \supseteq \delta$, while the inclusion \subseteq follows from $T' \cap (\delta : T) \subseteq T \cap (\delta : T) \subseteq \delta$. Hence $\mathbf{C}(T', T'; \delta)$ holds indeed. Now, using an argument like the one in the first paragraph of this proof, this contradicts the fact that $I[\delta, \theta]$ is a solvability obstruction unless $T' \subseteq \delta$. But then α is a congruence above δ for which $\alpha \cap T = T' \subseteq \delta$, so by item (1) we get that $\alpha \leq (\delta : T)$. This proves that $(\delta : T) \blacktriangleleft 1$, and therefore that $(\delta : T)$ is a radical congruence.

adjust the spacing here.

For item (3), if $\nu \not\leq \rho$, then we certainly have $(\nu \wedge \rho) < \nu$. Given any λ satisfying $(\nu \wedge \rho) \triangleleft \lambda \leq \nu$ we may join with ρ to obtain $\rho \triangleleft (\rho \vee \lambda) \leq (\rho \vee \nu)$, according to Lemma 6.10 (1). Since $\rho \blacktriangleleft 1$, it follows that $\lambda \leq \rho$. Since $\lambda \leq \nu$ as well, we must have $\lambda \leq (\nu \wedge \rho)$. This proves that $(\nu \wedge \rho) \blacktriangleleft \nu$. \square

Because of Lemma 7.11 (1), we may refer to $(\delta : T)$ as the **pseudocomplement of T over δ** when $I[\delta, \theta]_T$ is a separated solvability obstruction.

The following set of criteria for producing \rightsquigarrow -related intervals is simpler than the definition of \rightsquigarrow or the criteria of Lemma 7.10, but is still strong enough for all of our intended applications.

LEMMA 7.12. *Suppose that $I[\mu, \nu]_S$ and $I[\delta, \theta]_T$ are separated solvability obstructions in $\mathbf{Con}(\mathbf{A})$. If*

- (d)' $\text{Cg}^{\mathbf{A}}(S) \cap \text{Cg}^{\mathbf{A}}(T) \subseteq (\text{Cg}^{\mathbf{A}}(S) \wedge (\delta : T)) \vee (\text{Cg}^{\mathbf{A}}(T) \wedge (\delta : T))$
and
(g) $S \cap \mu = S \cap (\delta : T)$,

then $I[\mu, \nu]_S \rightsquigarrow I[\delta, \theta]_T$.

In particular, assume that $I[\delta, \theta]_T$ is any separated solvability obstruction, $\nu \not\leq (\delta : T)$, and $\mu := \nu \wedge (\delta : T)$. Then $I[\mu, \nu]_\nu$ is a separated solvability obstruction, and if $\nu \cap \text{Cg}^{\mathbf{A}}(T) \subseteq \mu$ then $I[\mu, \nu]_\nu \rightsquigarrow I[\delta, \theta]_T$.

PROOF. Assume that (d)' and (g) hold. We will argue that Condition (a) of Definition 7.4 holds, and Conditions (d), (e) and (f) of Lemma 7.10 all hold in the case where $\delta' = \text{Cg}^{\mathbf{A}}(T) \cap (\delta : T)$ and $\mu' = \text{Cg}^{\mathbf{A}}(S) \cap (\delta : T)$. This choice for δ' satisfies the required condition $T \cap \delta \subseteq \delta' \subseteq \text{Cg}^{\mathbf{A}}(T)$ because $T \subseteq \text{Cg}^{\mathbf{A}}(T)$ and $\delta \subseteq (\delta : T)$. The choice for μ' satisfies the condition $S \cap \mu \subseteq \mu' \subseteq \text{Cg}^{\mathbf{A}}(S)$ because $S \subseteq \text{Cg}^{\mathbf{A}}(S)$, while $S \cap \mu \subseteq (\delta : T)$ is part of assumption (g).

From Condition (g) we get $S \cap (\delta : T) = S \cap \mu \subseteq \mu$ as Condition (a) requires.

Condition (d)' is simply Condition (d) of Lemma 7.10 restricted to the case where $\delta' = \text{Cg}^{\mathbf{A}}(T) \cap (\delta : T)$ and $\mu' = \text{Cg}^{\mathbf{A}}(S) \cap (\delta : T)$.

Condition (e) asserts that $T \cap (\delta' \vee \mu') \subseteq \delta$. For this, observe that $\delta' \vee \mu' \leq (\delta : T)$ from our choice of primed congruences, so from Lemma 7.11 (1) we get that $T \cap (\delta' \vee \mu') \subseteq T \cap (\delta : T) \subseteq \delta$.

Condition (f) asserts that $S \cap (\delta' \vee \mu') \subseteq \mu$. According to the previous paragraph we have $\delta' \vee \mu' \leq (\delta : T)$. Together with (g), this yields $S \cap (\delta' \vee \mu') \subseteq S \cap (\delta : T) = S \cap \mu \subseteq \mu$.

For the second paragraph of the lemma statement, assume that $I[\delta, \theta]_T$ is any separated solvability obstruction. If $\nu \not\leq (\delta : T)$ and $\mu := \nu \wedge (\delta : T)$, then Lemma 7.11 (2) and (3) prove that $I[\mu, \nu]$ is a

at first i thought that this was a typo. consider using the \prime symbol, or some other symbol.

solvability obstruction. It is certainly separated by $S := \nu$. To see that (d)' holds under the assumption $\nu \cap \text{Cg}^{\mathbf{A}}(T) \subseteq \mu$, observe that since $S = \nu$ we have

$$\begin{aligned} \text{Cg}^{\mathbf{A}}(S) \cap \text{Cg}^{\mathbf{A}}(T) &= \nu \cap \text{Cg}^{\mathbf{A}}(T) \\ &\subseteq \mu \quad (\text{by the assumption}) \\ &:= \nu \cap (\delta : T) \\ &= \text{Cg}^{\mathbf{A}}(S) \cap (\delta : T) \\ &\subseteq (\text{Cg}^{\mathbf{A}}(S) \cap (\delta : T)) \vee (\text{Cg}^{\mathbf{A}}(T) \cap (\delta : T)). \end{aligned}$$

For (g), $S \cap \mu = \nu \cap \mu = \mu := \nu \cap (\delta : T) = S \cap (\delta : T)$. \square

7.2. Congruence Identities

in a failure

The first theorem of this section shows that solvability obstructions of minimal rank cannot be contained in failure of the distributive law. It is pretty easy to combine this with our bound on the length of \rightsquigarrow -sequences to derive a nontrivial congruence identity, which is something that we will do in Theorem 7.15.

THEOREM 7.13. *Assume that $\mathbf{A} \in \mathcal{V}$ and that $J_T := I[\delta, \theta]_T$ is a separated solvability obstruction in $\mathbf{Con}(\mathbf{A})$. If α, β , and γ are congruences on \mathbf{A} such that*

$$(\alpha \wedge \beta) \vee (\alpha \wedge \gamma) \leq \delta < \theta \leq \alpha \wedge (\beta \vee \gamma),$$

then there is a separated solvability obstruction I_S that is a subinterval of either $I[\alpha \wedge \beta, \beta]$ or $I[\alpha \wedge \gamma, \gamma]$ such that $I_S \rightsquigarrow J_T$.

PROOF. We will apply Lemma 7.12. The pseudocomplement of T over δ cannot lie above both β and γ , for if it did we would have

$$T \subseteq \theta \leq \beta \vee \gamma \leq (\delta : T),$$

which is in conflict with the properties $T \cap (\delta : T) \subseteq \delta$ and $T \not\subseteq \delta$. We may therefore assume that $\gamma \not\leq (\delta : T)$. If $\mu := \gamma \wedge (\delta : T)$, then the second paragraph of the statement of Lemma 7.12 proves that $I[\mu, \gamma]$ is a solvability obstruction. Since $\alpha \wedge \gamma \leq \delta \leq (\delta : T)$, we have $\alpha \wedge \gamma \leq \gamma \wedge (\delta : T) = \mu$, so $I[\mu, \nu]$ is contained in $I[\alpha \wedge \gamma, \gamma]$. Using the second paragraph of the statement of Lemma 7.12 again we obtain from

$$\gamma \cap \text{Cg}^{\mathbf{A}}(T) \leq \gamma \cap (\alpha \wedge (\beta \vee \gamma)) = \alpha \wedge \gamma \leq \mu$$

that $I[\mu, \nu]_{\nu} \rightsquigarrow I[\delta, \theta]_T$. \square

COROLLARY 7.14. *Assume that $\mathbf{A} \in \mathcal{V}$, and that $\mathbf{Con}(\mathbf{A})$ has a sublattice of congruences labeled as in Figure 7.7. If J_T is a separated solvability obstruction in $I[\beta, \alpha]$, then there is a separated solvability obstruction I_S contained in $I[\sigma, \gamma]$ such that $I_S \rightsquigarrow J_T$.*

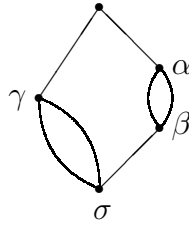


FIGURE 7.7.

PROOF. The hypotheses of Theorem 7.13 are met, since $(\alpha \wedge \beta) \vee (\alpha \wedge \gamma) = \beta \vee \sigma = \beta \leq \delta < \theta \leq \alpha = \alpha \wedge (\beta \vee \gamma)$. Therefore there is a separated solvability obstruction I_S contained in either $I[\alpha \wedge \beta, \beta] = I[\beta, \beta]$ or $I[\alpha \wedge \gamma, \gamma] = I[\sigma, \gamma]$ such that $I_S \rightsquigarrow J_T$. The first of these intervals is trivial, so I_S must be a subinterval of the second one. \square

THEOREM 7.15. *Let \mathcal{V} be a variety. The following conditions are equivalent.*

- (1) \mathcal{V} satisfies a nontrivial congruence identity.
- (2) \mathcal{V} satisfies an idempotent Maltsev condition that fails in the variety of semilattices.
- (3) There exists $\ell < \omega$ such that the lattice $\mathbf{N}_{\ell+5}$, depicted in Figure 7.8, cannot be embedded into $\mathbf{Con}(\mathbf{A})$ for any $\mathbf{A} \in \mathcal{V}$.

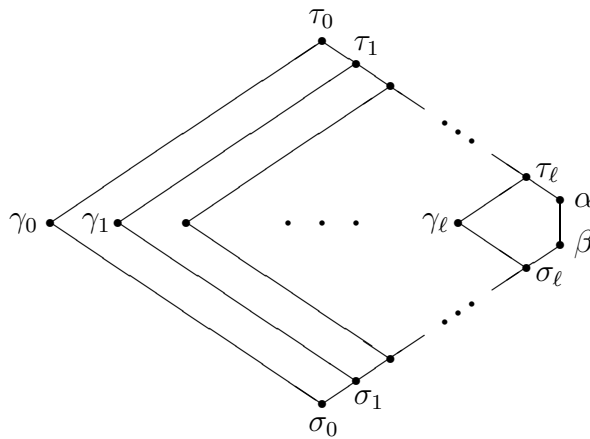


FIGURE 7.8. The lattice $\mathbf{N}_{\ell+5}$

PROOF. If \mathcal{V} satisfies a nontrivial congruence identity, ε , then any variety satisfying the same idempotent Maltsev conditions as \mathcal{V} must also satisfy ε , by the Pixley-Wille algorithm. But Ralph Freese and J. B.Nation proved in [20] that the variety of semilattices satisfies no

nontrivial congruence identity. Therefore \mathcal{V} satisfies an idempotent Maltsev condition that fails in the variety of semilattices.

Assume that (2) holds. Then \mathcal{V} has a Hobby–McKenzie term F of arity n for some $n < \omega$. We shall argue that for $\ell = 2n - 4$ the lattice $\mathbf{N}_{\ell+5}$ cannot be embedded into $\mathbf{Con}(\mathbf{A})$ for any $\mathbf{A} \in \mathcal{V}$. The interval $I[\sigma_\ell, \tau_\ell]$ contains a sublattice $\{\sigma_\ell, \alpha, \beta, \gamma_\ell, \tau_\ell\}$ that is isomorphic to \mathbf{N}_5 , so this interval is nonmodular. By Theorem 6.16, it follows that $I[\sigma_\ell, \tau_\ell]$ is not ∞ -solvable. Therefore, by Theorem 6.14, this interval contains a solvability obstruction $I[\mu_\ell, \nu_\ell]$. Choose a tolerance S_ℓ that separates this interval (e.g., $S_\ell = \nu_\ell$). By Corollary 7.14, applied to the sublattice $\{\sigma_\ell, \sigma_{\ell-1}, \gamma_{\ell-1}, \tau_{\ell-1}, \tau_\ell\}$, there is a separated solvability obstruction $I[\mu_{\ell-1}, \nu_{\ell-1}]_{S_{\ell-1}}$ contained in the interval $I[\sigma_{\ell-1}, \gamma_{\ell-1}]$ such that

$$I[\mu_{\ell-1}, \nu_{\ell-1}]_{S_{\ell-1}} \rightsquigarrow I[\mu_\ell, \nu_\ell]_{S_\ell}.$$

Similarly, for each $i = 1, 2, \dots, \ell$ successively we can apply Corollary 7.14 to the sublattice $\{\sigma_{\ell-i+1}, \sigma_{\ell-i}, \gamma_{\ell-i}, \tau_{\ell-i}, \tau_{\ell-i+1}\}$ to obtain a separated solvability obstruction $I[\mu_{\ell-i}, \nu_{\ell-i}]_{S_{\ell-i}}$ contained in the interval $I[\sigma_{\ell-i}, \gamma_{\ell-i}]$ such that

$$I[\mu_{\ell-i}, \nu_{\ell-i}]_{S_{\ell-i}} \rightsquigarrow I[\mu_{\ell-i+1}, \nu_{\ell-i+1}]_{S_{\ell-i+1}}.$$

When we reach $i = \ell$ we will have a sequence

$$I[\mu_0, \nu_0]_{S_0} \rightsquigarrow I[\mu_1, \nu_1]_{S_1} \rightsquigarrow I[\mu_2, \nu_2]_{S_2} \rightsquigarrow \dots \rightsquigarrow I[\mu_\ell, \nu_\ell]_{S_\ell}.$$

According to Theorem 7.9, there is no such sequence if $\ell \geq 2n - 4$. This completes the proof of (2) \implies (3).

To prove that (3) \implies (1), we observe that:

- (i) Each $\mathbf{N}_{\ell+5}$ is subdirectly irreducible.
- (ii) Each $\mathbf{N}_{\ell+5}$ is projective in the variety of lattices. (See [64].)

By Theorem 2.7, there is an identity $\varepsilon_{\ell+5}$ that holds in a lattice \mathbf{L} if and only if \mathbf{L} has no sublattice isomorphic to $\mathbf{N}_{\ell+5}$. Since (3) asserts that $\mathbf{N}_{\ell+5}$ is not embeddable in $\mathbf{Con}(\mathbf{A})$ for any $\mathbf{A} \in \mathcal{V}$, $\varepsilon_{\ell+5}$ is a congruence identity of \mathcal{V} . Since $\varepsilon_{\ell+5}$ fails in $\mathbf{N}_{\ell+5}$, it is a nontrivial congruence identity. \square

COROLLARY 7.16. (Cf. [57]) *Any congruence n -permutable variety satisfies a nontrivial congruence identity.*

PROOF. It suffices to observe that congruence n -permutability can be characterized by an idempotent Maltsev condition, for example by the Maltsev condition of J. Hagemann and A. Mitschke from [29], and that the variety of semilattices is not congruence n -permutable for any n . \square

COROLLARY 7.17. *Any congruence join semidistributive variety satisfies a nontrivial congruence identity.*

PROOF. Congruence join semidistributivity can be characterized by a set of idempotent Maltsev conditions, as is shown in [3], and the variety of semilattices is not congruence join semidistributive. \square

See Theorem 8.14 (7) for explicit congruence identities for congruence join semidistributive varieties.

CHAPTER 8

Congruence Meet and Join Semidistributivity

In this chapter we will characterize the satisfaction of a nontrivial congruence identity in several new ways (Theorems 8.3, 8.10, 8.11, 8.12, and 8.13). Our goal is to understand congruence join semidistributivity and its relationship to the satisfaction of a nontrivial congruence identity. Although most of the results of this chapter concern varieties satisfying congruence identities, we start the chapter by recalling known results about congruence meet semidistributivity and we end the chapter with new results (Theorem 8.14) that provide several characterizations of congruence join semidistributivity for varieties. In the process we solve an old problem of B. Jónsson (Problem 2.18 of [40]) by showing that the class of congruence join semidistributive varieties is definable by a Maltsev condition.

It should be noted that any congruence join semidistributive variety is also congruence meet semidistributive (see Exercise 7.14 (10) of [34] or (1) \implies (6) of Theorem 8.14 of this chapter). Therefore *congruence join semidistributivity* and *congruence semidistributivity* are equivalent properties for varieties.

8.1. Congruence Meet Semidistributivity

Using straightforward arguments, it follows from a combination of the statements in Theorem 2.19 that if α and β are congruences on an algebra \mathbf{A} , then $[\alpha, \beta] \leq \alpha \wedge \beta$ and $(\alpha \wedge \beta)/[\alpha, \beta]$ is an abelian congruence on $\mathbf{A}/[\alpha, \beta]$. Therefore, if \mathcal{V} is a variety that omits abelian congruences, it follows that $[\alpha, \beta] = \alpha \wedge \beta$ for any $\alpha, \beta \in \text{Con}(\mathbf{A})$ for any $\mathbf{A} \in \mathcal{V}$. Theorem 2.19 (5) implies that the commutator operation is semidistributive in its first variable i.e.,

$$[\alpha, \gamma] = \delta \ \& \ [\beta, \gamma] = \delta \ \rightarrow \ [\alpha \vee \beta, \gamma] = \delta.$$

Hence if \mathcal{V} omits abelian congruences, then the congruence lattices of members of \mathcal{V} must satisfy the ordinary meet semidistributive law:

$$\alpha \wedge \gamma = \delta \ \& \ \beta \wedge \gamma = \delta \ \rightarrow \ (\alpha \vee \beta) \wedge \gamma = \delta.$$

The converse of this statement is true, but is much more difficult to prove. It was established for locally finite varieties in Theorem 9.10

of [34] using tame congruence theory, and in full generality in Corollary 4.7 of [52] using combinatorial arguments. This result links many different properties, as we will show in the next theorem.

Recall the type of meet continuous lattice words defined before Corollary 2.25: if x, y, z are variables let $y_0 = y$, $z_0 = z$, $y_{n+1} = y \vee (x \wedge z_n)$, and $z_{n+1} = z \vee (x \wedge y_n)$. Let $y_\omega = \bigvee_{n < \omega} y_n$ and let $z_\omega = \bigvee_{n < \omega} z_n$. When it is necessary to indicate the variables involved and their order we will write $y_\omega(x, y, z)$ for y_ω and $z_\omega(x, y, z)$ for z_ω .

THEOREM 8.1. *Let \mathcal{V} be a variety. The following conditions are equivalent.*

- (1) \mathcal{V} is congruence meet semidistributive.
- (2) There is a positive integer n such that if $\mathbf{A} \in \mathcal{V}$ and $\alpha, \beta, \gamma \in \text{Con}(\mathbf{A})$, then

$$\alpha \cap (\beta \circ \gamma) \subseteq \beta_n,$$

where $\beta_n := y_n(\alpha, \beta, \gamma)$.

- (3) $[\alpha, \beta] = \alpha \wedge \beta$. for all $\alpha, \beta \in \text{Con}(\mathbf{A})$ and all $\mathbf{A} \in \mathcal{V}$.
- (4) No member of \mathcal{V} has a nontrivial abelian congruence.
- (5) No member of \mathcal{V} has a nontrivial abelian tolerance.
- (6) $\mathbf{C}(\alpha, \beta; \delta) \iff \beta \wedge (\alpha \vee (\beta \wedge \delta)) \leq \delta$ for all $\mathbf{A} \in \mathcal{V}$ and all $\alpha, \beta, \delta \in \text{Con}(\mathbf{A})$.
- (7) \mathcal{V} satisfies the meet continuous congruence identity

$$x \wedge (y \vee z) \approx x \wedge y_\omega,$$

where $y_\omega = y_\omega(x, y, z)$.

- (8) The meet continuous congruence variety of \mathcal{V} is meet semidistributive.
- (9) \mathbf{M}_3 is not embeddable in $\text{Con}(\mathbf{A})$ for any $\mathbf{A} \in \mathcal{V}$.
- (10) \mathcal{V} satisfies a family of idempotent Maltsev conditions that, taken together, fail in any nontrivial variety of modules. (Equivalently, \mathcal{V} satisfies a single idempotent Maltsev condition that fails in any nontrivial variety of modules.)

PROOF. The equivalence of items (1), (2) and (3) is Corollary 4.7 of [52].

If (3) holds, and $\mathbf{A} \in \mathcal{V}$ has an abelian congruence γ , then $0 = [\gamma, \gamma] = \gamma \wedge \gamma = \gamma$. Hence γ is trivial, and (4) holds. Conversely, assume that (4) holds. It follows from Theorem 2.19 (8) that $\mathbf{C}(\alpha, \beta; \alpha \wedge \beta)$ holds, so $[\alpha, \beta] \leq \alpha \wedge \beta$ holds for any pair of congruences on any algebra \mathbf{A} . Moreover $\mathbf{C}(\alpha, \beta; [\alpha, \beta])$ holds, according to the definition of the commutator, and

$$\mathbf{C}(\alpha, \beta; [\alpha, \beta]) \implies \mathbf{C}(\alpha \wedge \beta, \alpha \wedge \beta; [\alpha, \beta])$$

by Theorem 2.19 (1), so $(\alpha \wedge \beta)/[\alpha, \beta]$ is an abelian congruence on $\mathbf{A}/[\alpha, \beta]$ by Theorem 2.19 (10). Therefore, from (4) we derive (3).

Assume that (4) holds. Since \mathcal{V} omits nontrivial abelian congruences, it omits nontrivial strongly abelian congruences. Therefore, by Theorem 3.12, \mathcal{V} satisfies a nontrivial idempotent Maltsev condition. In this situation, Theorem 3.23 guarantees that an abelian tolerance generates an abelian congruence. Since \mathcal{V} omits nontrivial abelian congruences, \mathcal{V} omits nontrivial abelian tolerances. This shows that (4) \implies (5). The reverse implication is trivial.

To connect item (6) with the preceding conditions, assume that (3) holds and that

(a) $\mathbf{C}(\alpha, \beta; \delta)$ holds.

By Theorem 2.19 (8),

(b) $\mathbf{C}(\alpha, \beta; \beta)$ holds,

so by Theorem 2.19 (6) we derive from (a) and (b) that

(c) $\mathbf{C}(\alpha, \beta; \beta \wedge \delta)$ holds.

Again by Theorem 2.19 (8) we get that

(d) $\mathbf{C}(\beta \wedge \delta, \beta; \beta \wedge \delta)$ holds,

so by Theorem 2.19 (5) we get from (c) and (d) that

(e) $\mathbf{C}(\alpha \vee (\beta \wedge \delta), \beta; \beta \wedge \delta)$ holds.

From this and the definition of the commutator we get that

(f) $[(\alpha \vee (\beta \wedge \delta)), \beta] \leq \beta \wedge \delta$ holds.

Now from (f) and item (3) we derive that

$$\beta \wedge (\alpha \vee (\beta \wedge \delta)) = [(\alpha \vee (\beta \wedge \delta)), \beta] \leq \beta \wedge \delta \leq \delta.$$

Conversely, $\beta \wedge (\alpha \vee (\beta \wedge \delta)) \leq \delta$ implies $\mathbf{C}(\alpha, \beta; \delta)$ for any three congruences on any algebra, as Theorem 2.19 (8) proves. This shows that (3) implies (6).

On the other hand, to see that (6) \implies (4), take $\alpha = \beta = \gamma$ and $\delta = 0$ in (6). The result is the assertion that γ is abelian if and only if $\gamma = 0$.

The equivalence of (1), (7) and (8) follows from Theorem 2.23 and the fact that congruence lattices are meet continuous.

Item (8) implies item (9), since \mathbf{M}_3 is not meet semidistributive.

Assume that (9) holds. Then, since \mathbf{M}_3 satisfies Whitman's condition (W), it follows from Theorem 4.18 that \mathcal{V} satisfies a nontrivial idempotent Maltsev condition. In this situation, a result from [52] explains how to construct from a nontrivial abelian congruence on an algebra in \mathcal{V} a copy of \mathbf{M}_3 in the congruence lattice of some other algebra in \mathcal{V} . Namely, suppose that α is a nonzero abelian congruence on some $\mathbf{A} \in \mathcal{V}$, that $\mathbf{A} \times_{\alpha} \mathbf{A}$ is the subalgebra of \mathbf{A}^2 whose universe

is α , that $\eta_1, \eta_2 \in \text{Con}(\mathbf{A} \times_\alpha \mathbf{A})$ are the coordinate projection kernels, and that Δ is the congruence on $\mathbf{A} \times_\alpha \mathbf{A}$ generated by the pairs $\langle (a, a), (b, b) \rangle$ with $(a, b) \in \alpha$. It is a consequence of Theorem 3.5 of [52] that η_1, η_2 , and Δ generate a sublattice of $\text{Con}(\mathbf{A} \times_\alpha \mathbf{A})$ isomorphic to \mathbf{M}_3 . Thus, if (9) holds, then (4) must hold.

The Pixley–Wille algorithm shows that the class of varieties satisfying a congruence condition like the one in item (2) can be defined by a single idempotent Maltsev condition. Since (2) is equivalent to (4), this Maltsev condition cannot be satisfied by any variety with nontrivial abelian algebras, hence by any nontrivial variety of modules. Thus (7) implies the stronger of the two statements in (10).

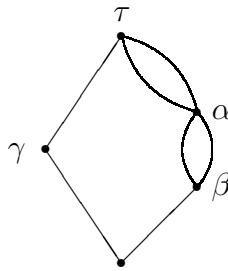
Finally, we show that the weaker of the two statements in (10) implies (4). Assume instead that \mathcal{V} satisfies a family \mathcal{F} of idempotent Maltsev conditions that fail in every nontrivial variety of modules, but that some $\mathbf{A} \in \mathcal{V}$ has a nontrivial abelian congruence α . We may assume, without loss of generality, that \mathcal{V} is an idempotent variety. For if we replace \mathcal{V} and \mathbf{A} by their idempotent reducts, then we do not affect the idempotent Maltsev conditions satisfied by \mathcal{V} , nor that α is an abelian congruence on \mathbf{A} . Thus, we make this assumption. We may (and do) further assume that \mathbf{A} itself is abelian, since if it is not we may simply replace it by one of its subalgebras that is supported by nontrivial α -block.

At least one of the Maltsev conditions in \mathcal{F} is nontrivial, since these conditions taken together fail in every nontrivial variety of modules. Corollary 4.5 of [52] proves that an abelian algebra \mathbf{A} in a variety satisfying some nontrivial idempotent Maltsev condition is **quasi-affine**, which by definition means that \mathbf{A} is a subalgebra of a reduct of an algebra that has the same universe and polynomial operations as some module. Since \mathbf{A} is both idempotent and quasi-affine, it is a subalgebra of an algebra \mathbf{M}^* that is a reduct of some module \mathbf{M} over some ring \mathbf{R} . There is no harm in renaming the zero element of \mathbf{M} so that it lies in A . If p is the module term $x - y + z$, then $p: \mathbf{M}^3 \rightarrow \mathbf{M}$ is a homomorphism, and therefore $p: \mathbf{A}^3 \rightarrow \mathbf{M}^*$ is also a homomorphism. The image $\mathbf{A}' := p(\mathbf{A}^3)$ contains A , since $p(a, a, a) = a$. Moreover, \mathbf{A}' is a quotient of $\mathbf{A}^3 \in \mathcal{V}$, hence $\mathbf{A}' \in \mathcal{V}$. This produces a possibly larger subalgebra $\mathbf{A} \leq \mathbf{A}' \leq \mathbf{M}^*$. Iterating this, $\mathbf{A} \leq \mathbf{A}' \leq \mathbf{A}'' \leq \dots$, and taking a union yields an algebra $\mathbf{B} \in \mathcal{V}$ that is a subalgebra of \mathbf{M}^* and is closed under p . Since $0 \in A \subseteq B$, the set B is closed under the abelian group operations $x + y = p(x, 0, y)$, $-x = p(0, x, 0)$, and 0 . Thus $\widehat{\mathbf{B}} := \langle \mathbf{B}; p, 0 \rangle$ is an additive subgroup of \mathbf{M} that is also a subalgebra of a reduct of \mathbf{M} . This implies that $\widehat{\mathbf{B}}$ is term equivalent to a module over

the subring of \mathbf{R} consisting of all elements $r \in R$ such that $rB \subseteq B$. Clearly $\widehat{\mathbf{B}}$ is nontrivial, since \mathbf{B} contains \mathbf{A} as a subalgebra. Moreover, $\widehat{\mathbf{B}}$ satisfies the conditions in \mathcal{F} , since its reduct \mathbf{B} does. Therefore the variety generated by $\widehat{\mathbf{B}}$ is (up to term equivalence) a nontrivial variety of modules that satisfies the conditions in \mathcal{F} , contrary to (10). This contradicts our assumption that (4) fails to hold, so we are done. \square

8.2. More on Congruence Identities

THEOREM 8.2. *Assume that $\mathbf{A} \in \mathcal{V}$, and that $\mathbf{Con}(\mathbf{A})$ has a sublattice of congruences labeled as in Figure 8.1. If J_T is a separated*



what assumptions are you making about \mathcal{V} in this theorem? that is has a hobby-mckenzie term?

FIGURE 8.1.

solvability obstruction in $I[\beta, \alpha]$, then there is a separated solvability obstruction I_S contained in $I[\alpha, \tau]$ and a \rightsquigarrow -sequence from I_S to J_T .

PROOF. We will show first that if J_T is a separated solvability obstruction in $I[\beta, \alpha]$, then there is a separated solvability obstruction $I'_{S'}$ contained in either $I[\beta, \alpha]$ or $I[\alpha, \tau]$ and a \rightsquigarrow -sequence of length 1 or 2 from $I'_{S'}$ to J_T . If it happens that $I'_{S'}$ is contained in $I[\alpha, \tau]$, then we are done. If instead $I'_{S'}$ is contained in $I[\beta, \alpha]$, then we can repeat the argument to find $I''_{S''}$ contained in either $I[\beta, \alpha]$ or $I[\alpha, \tau]$ and a \rightsquigarrow -sequence from $I''_{S''}$ to $I'_{S'}$. The argument cannot be repeated indefinitely since concatenating \rightsquigarrow -sequences forms longer \rightsquigarrow -sequences and there is a bound on the length of \rightsquigarrow -sequences. Therefore the argument will produce a separated solvability obstruction in $I[\alpha, \tau]$ that initiates a \rightsquigarrow -sequence terminating at J_T .

Recall from the proof of Theorem 7.13, applied to the setting of Corollary 7.14, that if $\mu := \gamma \cap (\delta : T)$, then $I[\mu, \gamma]_\gamma \rightsquigarrow I[\delta, \theta]_T$. Since $\mu \leq (\delta : T)$ and $\delta \leq (\delta : T)$, we have

$$\gamma \cap (\delta \vee \mu) \leq \gamma \cap (\delta : T) = \mu,$$

and therefore $\gamma \wedge (\delta \vee (\gamma \wedge \mu)) \leq \mu$. By Theorem 2.19 (8) we get that $\mathbf{C}(\delta, \gamma; \mu)$ holds, or equivalently $\delta \leq (\mu : \gamma)$. If $\alpha \not\leq (\mu : \gamma)$, then for

$\lambda := \alpha \wedge (\mu : \gamma)$ ($\geq \delta \geq \beta$) we have that $I[\lambda, \alpha]$ is a nontrivial subinterval of $I[\beta, \alpha]$. According to the second paragraph of Lemma 7.12, the interval $I[\lambda, \alpha]$ is a solvability obstruction. That paragraph further guarantees that, since $\alpha \cap \text{Cg}^{\mathbf{A}}(\gamma) \leq \lambda$, we have $I[\lambda, \alpha]_{\alpha} \rightsquigarrow I[\mu, \gamma]_{\gamma}$. Therefore, in the case where $\alpha \not\leq (\mu : \gamma)$ we have located a separated solvability obstruction $I[\lambda, \alpha]_{\alpha}$ in $I[\beta, \alpha]$ and a \rightsquigarrow -sequence

$$I[\lambda, \alpha]_{\alpha} \rightsquigarrow I[\mu, \gamma]_{\gamma} \rightsquigarrow I[\delta, \theta]_T = J_T$$

of length two.

Now consider the case where $\alpha \leq (\mu : \gamma)$. Since $\gamma \not\leq (\mu : \gamma)$ and $\gamma \leq \tau$ we have $\tau \not\leq (\mu : \gamma)$, so $\kappa := \tau \cap (\mu : \gamma)$ satisfies $\alpha \leq \kappa < \tau$ in this case. According to the second paragraph of Lemma 7.12, the interval $I[\kappa, \tau]$ is a solvability obstruction. Since $\gamma \not\leq \kappa$, $I[\kappa, \tau]_{\gamma}$ is a separated solvability obstruction. We prove that $I[\kappa, \tau]_{\gamma} \rightsquigarrow I[\delta, \tau]_T$ by verifying criteria (d)' and (g) from Lemma 7.12. In our setting these statements are

- (d)' $\gamma \cap \text{Cg}^{\mathbf{A}}(T) \subseteq (\gamma \wedge (\delta : T)) \vee (\text{Cg}^{\mathbf{A}}(T) \wedge (\delta : T))$ and
- (g) $\gamma \cap \kappa = \gamma \cap (\delta : T)$.

Item (d)' follows from the observation that

$$\gamma \cap \text{Cg}^{\mathbf{A}}(T) \subseteq \gamma \cap \alpha \subseteq \mu := \gamma \cap (\delta : T),$$

since $\gamma \cap (\delta : T)$ is the first joinand in the right hand side of (d)'. For item (g), observe that

$$\gamma \cap \kappa = \gamma \cap \tau \cap (\mu : \gamma) = \gamma \cap (\mu : \gamma) = \mu = \gamma \cap (\delta : T).$$

This proves that $I[\kappa, \tau]_{\gamma} \rightsquigarrow I[\delta, \tau]_T$. □

Let \mathbf{L} be a lattice. A **herringbone** in \mathbf{L} is the union of three descending chains $\{\alpha^i\} \cup \{\beta^{2i}\} \cup \{\gamma^{2i+1}\}$ in \mathbf{L} where $\{\alpha^i\} \cup \{\beta^{2i}\}$ and $\{\alpha^i\} \cup \{\gamma^{2i+1}\}$ are sublattices of \mathbf{L} ordered as in Figure 8.2. In other

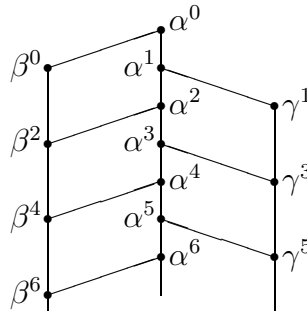


FIGURE 8.2. A herringbone

words, a herringbone is a partial sublattice of \mathbf{L} , ordered as in Figure 8.2, with

$$\begin{aligned} (\beta) \quad & \alpha^{2i+1} \wedge \beta^{2i} = \beta^{2i+2}, \alpha^{2i+2} \vee \beta^{2i} = \alpha^{2i}, \text{ and} \\ (\gamma) \quad & \alpha^{2i+2} \wedge \gamma^{2i+1} = \gamma^{2i+3}, \alpha^{2i+3} \vee \gamma^{2i+1} = \alpha^{2i+1}. \end{aligned}$$

The **length** of the herringbone is ∞ if there are infinitely many distinct α 's. Otherwise, the **length** of the herringbone is the supremum of the superscripts k such that $\alpha^0 > \alpha^1 > \dots > \alpha^k$. (It is easy to see from (β) and (γ) that if some $\alpha^k = \alpha^{k+1}$ then $\alpha^k = \alpha^{k+1} = \alpha^{k+2} = \dots$.) Our first goal is to prove that if a variety \mathcal{V} satisfies a nontrivial congruence identity, then there is a fixed bound on the length of any herringbone that appears in the congruence lattice of a member of \mathcal{V} .

Let x, y and z be lattice variables. Define lattice words by $y^0 = y$, $z^0 = z$, $y^{n+1} = y \wedge (x \vee z^n)$, and $z^{n+1} = z \wedge (x \vee y^n)$. The construction of these words is dual to the words y_n and z_n , which preceded the statement of Theorem 8.1. As there, we will write $y^n(x, y, z)$ and $z^n(x, y, z)$ when it is necessary to indicate the order of the variables.

THEOREM 8.3. *Let \mathcal{V} be a variety. The following conditions are equivalent.*

- (1) \mathcal{V} satisfies a nontrivial congruence identity.
- (2) There is a positive integer N such that no algebra in \mathcal{V} has a herringbone of length greater than N in its congruence lattice.
- (3) \mathcal{V} satisfies the congruence identity

$$z^M \approx z^{M+1}$$

for some M .

PROOF. If \mathcal{V} satisfies a nontrivial congruence identity, then by Theorems 7.15 and 5.25 the variety \mathcal{V} has a Hobby–McKenzie term of arity n for some $n < \omega$. We argue that no algebra in \mathcal{V} has a herringbone of length greater than $2n - 3$ in its congruence lattice. If this is not the case, then some $\mathbf{A} \in \mathcal{V}$ has a herringbone $\{\alpha^i\} \cup \{\beta^{2i}\} \cup \{\gamma^{2i+1}\}$ with $\alpha^0 > \alpha^1 > \dots > \alpha^{2n-3} > \alpha^{2n-2}$. Figure 8.3 isolates five congruences near the bottom of the herringbone that form a sublattice of $\mathbf{Con}(\mathbf{A})$ isomorphic to \mathbf{N}_5 . For the congruences in this figure, if $\alpha^{2n-3} \lll \alpha^{2n-4}$, then by Lemma 6.10 (2) we derive that

$$\beta^{2n-4} = \beta^{2n-4} \wedge \alpha^{2n-4} \lll \beta^{2n-4} \wedge \alpha^{2n-3} = \beta^{2n-2}.$$

But this was proved to be impossible in Theorem 6.25 (1). Therefore we cannot have $\alpha^{2n-3} \lll \alpha^{2n-4}$. By Theorem 6.14 there is a solvability obstruction $I[\mu_{2n-4}, \nu_{2n-4}]$ contained in the interval $I[\alpha^{2n-3}, \alpha^{2n-4}]$. This obstruction is separated by $S_{2n-4} := \nu_{2n-4}$. We are now in a position to apply Theorem 8.2 to a copy of \mathbf{N}_5 that is slightly higher up

the ordering is reversed here:
 $\beta^{2n-2} \leq \beta^{2n-4}$

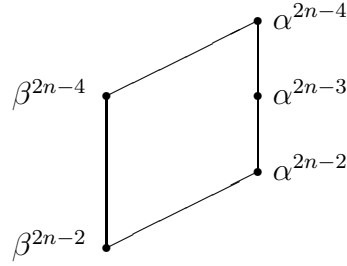


FIGURE 8.3. A piece near the bottom of the herringbone

in this herringbone (Figure 8.4). The separated solvability obstruction

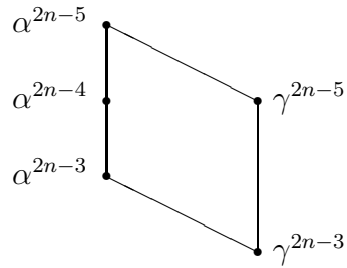


FIGURE 8.4. A piece higher up

2n-4

$I[\mu_{2n-4}, \nu_{2n-4}]_{S_{2n-4}}$ that is contained in $I[\alpha^{2n-3}, \alpha^{2n-4}]$ plays the role of J_T of that theorem. The theorem guarantees the existence of a separated solvability obstruction $I[\mu_{2n-5}, \nu_{2n-5}]_{S_{2n-5}}$ in $I[\alpha^{2n-5}, \alpha^{2n-5}]$ and a \rightsquigarrow -sequence

$$I[\mu_{2n-5}, \nu_{2n-5}]_{S_{2n-5}} \rightsquigarrow \cdots \rightsquigarrow I[\mu_{2n-4}, \nu_{2n-4}]_{S_{2n-4}}.$$

We can work our way up the herringbone, alternately using copies of \mathbf{N}_5 on the right and left and using Theorem 8.2 at each step, to derive that for each $i = 4, 5, \dots, 2n - 1$, if $I[\alpha^{2n-i+1}, \alpha^{2n-i}]$ contains a separated solvability obstruction $I[\mu_{2n-i}, \nu_{2n-i}]_{S_{2n-i}}$, then $I[\alpha^{2n-i}, \alpha^{2n-i-1}]$ contains a separated solvability obstruction $I[\mu_{2n-i-1}, \nu_{2n-i-1}]_{S_{2n-i-1}}$ for which there exists a \rightsquigarrow -sequence

$$I[\mu_{2n-i-1}, \nu_{2n-i-1}]_{S_{2n-i-1}} \rightsquigarrow \cdots \rightsquigarrow I[\mu_{2n-i}, \nu_{2n-i}]_{S_{2n-i}}.$$

When we reach the top of the herringbone we will have located a \rightsquigarrow -sequence

$$I[\mu_0, \nu_0]_{S_0} \rightsquigarrow \cdots \rightsquigarrow I[\mu_1, \nu_1]_{S_1} \rightsquigarrow \cdots \rightsquigarrow I[\mu_{2n-4}, \nu_{2n-4}]_{S_{2n-4}}$$

of length at least $2n - 4$. But we proved in Theorem 7.9 that there is no \rightsquigarrow -sequence of this length. This contradiction proves that there is no herringbone of length greater than $2n - 3$ in $\mathbf{Con}(\mathbf{A})$.

Now we prove that (2) \implies (3). Start with three congruences α, β and γ , and define $\beta^n := y^n(\alpha, \beta, \gamma)$ and $\gamma^n := z^n(\alpha, \beta, \gamma)$ where $y^n(x, y, z)$ and $z^n(x, y, z)$ are the terms defined before the statement of this theorem. Let $\alpha^n := \alpha \vee \beta^n$ if n is even, and $\alpha^n := \alpha \vee \gamma^n$ if n is odd. Since $\beta^0 = \beta \geq \beta \wedge (\alpha \vee \gamma^0) = \beta^1$, and $\gamma^0 \geq \gamma^1$, it is easy to see inductively that

$$\beta^{n+1} = \beta \wedge (\alpha \vee \gamma^n) \geq \beta \wedge (\alpha \vee \gamma^{n+1}) = \beta^{n+2},$$

and $\gamma^{n+1} \geq \gamma^{n+2}$. Thus the β and γ -sequences are descending chains, which forces the α -sequence to be a descending chain. We claim that

$$\{\alpha^n \mid \text{all } n\} \cup \{\beta^n \mid \text{even } n\} \cup \{\gamma^n \mid \text{odd } n\}$$

is a herringbone. To see this, we must verify conditions (β) and (γ) from the definition of a herringbone:

$$\begin{aligned} (\beta) \quad & \alpha^{2i+1} \wedge \beta^{2i} = \beta^{2i+2}, \quad \alpha^{2i+2} \vee \beta^{2i} = \alpha^{2i} \\ (\gamma) \quad & \alpha^{2i+2} \wedge \gamma^{2i+1} = \gamma^{2i+3}, \quad \alpha^{2i+3} \vee \gamma^{2i+1} = \alpha^{2i+1}. \end{aligned}$$

To show that $\alpha^{2i+1} \wedge \beta^{2i} = \beta^{2i+2}$, observe that

$$\beta^{2i+2} = \beta \wedge (\alpha \vee \gamma^{2i+1}) = \beta \wedge \alpha^{2i+1} \leq \alpha^{2i+1},$$

and (as observed earlier) $\beta^{2i+2} \leq \beta^{2i}$. Thus $\beta^{2i+2} \leq \alpha^{2i+1} \wedge \beta^{2i}$. Conversely,

$$\alpha^{2i+1} \wedge \beta^{2i} = (\alpha \vee \gamma^{2i+1}) \wedge \beta^{2i} \leq (\alpha \vee \gamma^{2i+1}) \wedge \beta = \beta^{2i+2}.$$

To show that $\alpha^{2i+2} \vee \beta^{2i} = \alpha^{2i}$, observe that

$$\alpha^{2i+2} \vee \beta^{2i} \geq \alpha \vee \beta^{2i} = \alpha^{2i},$$

while $\alpha^{2i+2} \leq \alpha^{2i}$ and $\beta^{2i} \leq \alpha \vee \beta^{2i} = \alpha^{2i}$, so $\alpha^{2i+2} \vee \beta^{2i} \leq \alpha^{2i}$. This establishes (β) , and (γ) can be established the same way.

It follows that if N bounds the length of any herringbone in any congruence lattice of a member of \mathcal{V} , then however we choose our original three congruences α, β and γ , the sequences defined above must satisfy $\alpha^N = \alpha^{N+1} = \alpha^{N+2} = \dots$, and therefore $\gamma^M = \gamma^{M+2}$ holds for each $M \geq N$. Since the γ 's form a descending chain, we conclude that $\gamma^M = \gamma^{M+1}$ holds for each $M \geq N$. Therefore $z^M \approx z^{M+1}$ holds in all congruence lattices of members of \mathcal{V} .

Item (1) follows from (3) and the fact that the identity $z^M \approx z^{M+1}$ is nontrivial. To see that it fails in some lattice, take any lattice containing a herringbone of length at least $M + 1$. By setting $\alpha := \alpha^{M+1}$, $\beta := \beta^0$ and $\gamma := \gamma^1$, it is easy to see that $z^M(\alpha, \beta, \gamma) \neq z^{M+1}(\alpha, \beta, \gamma)$. \square

This argument shows that if \mathcal{V} has a Hobby–McKenzie term of arity n , then \mathcal{V} satisfies the congruence identity $z^{2n-3} \approx z^{2n-2}$. These exponents are not optimal, since if $n = 3$ then our proof shows that \mathcal{V} satisfies the congruence identity $z^3 \approx z^4$. But it can be argued (by considering all Hobby–McKenzie terms of arity 3) that when $n = 3$ the variety \mathcal{V} is congruence modular, and consequently satisfies $z^1 \approx z^2$ as a congruence identity.

Recall from Section 2.2 that the SD_{\vee} -configuration is the pair

$$(\mathbf{P}(\text{SD}_{\vee}), p \approx j)$$

where $\mathbf{P}(\text{SD}_{\vee})$ is the partial lattice presented by $\langle G \mid R \rangle$ where $G = \{p, q, r, j, m\}$ and R consists of the relations $p \vee q = j = p \vee r$ and $q \wedge r = m \leq p$.

LEMMA 8.4. *$\mathbf{P}(\text{SD}_{\vee})$ is projective relative to any variety of lattices satisfying an identity of the form $z^M \approx z^{M+1}$.*

In fact, if $\sigma: \mathbf{L} \rightarrow \mathbf{K}$ is a surjective homomorphism in this variety and $\varphi: \mathbf{P}(\text{SD}_{\vee}) \rightarrow \mathbf{K}$ is a partial lattice homomorphism, then for any function $f: \{p, q, r, s\} \rightarrow \mathbf{L}$ satisfying $\sigma \circ f = \varphi$ on the domain of f there is a homomorphism of partial lattices $\bar{\varphi}: \mathbf{P}(\text{SD}_{\vee}) \rightarrow \mathbf{L}$ such that $\sigma \circ \bar{\varphi} = \varphi$ and $\bar{\varphi}(p) \geq f(p)$, $\bar{\varphi}(q) \leq f(q)$, and $\bar{\varphi}(r) \leq f(r)$.

PROOF. Define $\mathbf{P} := \mathbf{P}(\text{SD}_{\vee})$, and suppose that $\varphi: \mathbf{P} \rightarrow \mathbf{K}$ is a homomorphism of partial lattices and that $f: \{p, q, r, s\} \rightarrow \mathbf{L}$ is a function satisfying $\sigma \circ f = \varphi$ on the domain of f . Let $a = f(p)$, $b = f(q)$ and $c = f(r)$. Recursively define $b^0 := b$, $c^0 := c$, $b^{n+1} := b \wedge (a \vee c^n)$, and $c^{n+1} := c \wedge (a \vee b^n)$, so that $b^n = y^n(a, b, c)$ and $c^n = z^n(a, b, c)$ where $y^n(x, y, z)$ and $z^n(x, y, z)$ are the terms defined for the proof of Theorem 8.3. As we showed in the proof of Theorem 8.3 (1) \implies (2), $b = b^0 \geq b^1 \geq \dots$ and $c = c^0 \geq c^1 \geq \dots$. From the identity $z^M \approx z^{M+1}$ we get that these chains terminate, in which case $b^M = b^{M+1} = b \wedge (a \vee c^M) \leq a \vee c^M$ and similarly $c^M \leq a \vee b^M$. Thus $b^M \leq b$, $c^M \leq c$ and $a \vee b^M = a \vee c^M$. Define $\bar{\varphi}: \mathbf{P} \rightarrow \mathbf{L}$ by $\bar{\varphi}(p) = a \vee (b^M \wedge c^M)$ ($\geq f(p)$), $\bar{\varphi}(q) = b^M$ ($\leq f(q)$), $\bar{\varphi}(r) = c^M$ ($\leq f(r)$), $\bar{\varphi}(j) = a \vee b^M$, and $\bar{\varphi}(m) = b^M \wedge c^M$. It is clear that

$$\bar{\varphi}(j) = a \vee b^M = \bar{\varphi}(p) \vee \bar{\varphi}(q) = \bar{\varphi}(p) \vee \bar{\varphi}(r)$$

and that

$$\bar{\varphi}(m) = b^M \wedge c^M = \bar{\varphi}(q) \wedge \bar{\varphi}(r) \leq \bar{\varphi}(p).$$

Therefore $\bar{\varphi}: \mathbf{P} \rightarrow \mathbf{L}$ is a homomorphism of partial lattices. To finish the proof we must show that $\sigma \circ \bar{\varphi} = \varphi$. It suffices to prove that $\sigma \circ \bar{\varphi}(x) = \sigma \circ f(x)$ ($= \varphi(x)$) when $x \in \{p, q, r\}$ is one of the generators of \mathbf{P} .

By simultaneous induction, $\sigma(b^n) = \sigma(b)$ and $\sigma(c^n) = \sigma(c)$ for all n , where one of the induction steps is proved by

$$\begin{aligned}\sigma(b^{n+1}) &= \sigma(b) \wedge (\sigma(a) \vee \sigma(c^n)) \\ &= \sigma(b) \wedge (\sigma(a) \vee \sigma(c)) \\ &= \sigma(b) \wedge (\sigma(a) \vee \sigma(b)) = \sigma(b).\end{aligned}$$

Thus $\sigma \circ f(q) = \sigma(b) = \sigma(b^M) = \sigma \circ \overline{\varphi}(q)$, and similarly $\sigma \circ f(r) = \sigma(c) = \sigma(c^M) = \sigma \circ \overline{\varphi}(r)$. Since $q \wedge r \leq p$ and φ is a homomorphism we have $\varphi(q) \wedge \varphi(r) \leq \varphi(p)$. This yields the second equality in

$$\begin{aligned}\sigma \circ f(p) &= \sigma(a) \\ &= \sigma(a) \vee (\sigma(b) \wedge \sigma(c)) \\ &= \sigma(a) \vee (\sigma(b^M) \wedge \sigma(c^M)) \\ &= \sigma \circ \overline{\varphi}(p).\end{aligned}$$

The lemma is proved. \square

Our first use of Lemma 8.4 will require the following consequence of Theorem 8.3, which seems to be new even for locally finite varieties.

THEOREM 8.5. *Assume that \mathcal{V} satisfies a nontrivial congruence identity. If I is an SD_{\vee} -failure in $\mathbf{Con}(\mathbf{A})$ for some $\mathbf{A} \in \mathcal{V}$, then I is abelian.*

PROOF. We proved in Corollary 6.34 that if \mathcal{V} has a weak difference term, then ∞ -solvable SD_{\vee} -failures are abelian. Since the assumption that \mathcal{V} satisfies a nontrivial congruence identity implies that \mathcal{V} has a Hobby–McKenzie term, and therefore a weak difference term (Corollary 6.3), to prove the present theorem it will suffice to show that all SD_{\vee} -failures are ∞ -solvable. To set up notation for this, assume that $\mathbf{A} \in \mathcal{V}$ has congruences α, β and γ satisfying $\alpha \vee \beta = \alpha \vee \gamma$ and $\beta \wedge \gamma \leq \alpha$. Our goal is to prove that $\alpha \lll \alpha \vee \beta$. Factoring by $\beta \wedge \gamma$ if necessary, we also assume that $\beta \wedge \gamma = 0$.

Our strategy will be to show that if it is not true that $\alpha \lll \alpha \vee \beta$, then $\mathbf{Con}(\mathbf{A})$ contains an infinite \rightsquigarrow -sequence. This will suffice to prove the theorem, since Theorem 7.9 provides a finite bound on the length of \rightsquigarrow -sequences. Therefore, assume that α is *not* ∞ -solvably related to $\alpha \vee \beta = \alpha \vee \gamma$. Define new congruences as follows:

- $\beta^0 := \beta \wedge \text{rad}(\alpha)$, and $\gamma^0 := \gamma \wedge \text{rad}(\alpha)$.
- $\beta^{k+1} := \beta \wedge (\gamma^k : \gamma)$ and $\gamma^{k+1} := \gamma \wedge (\beta^k : \beta)$.

Our first goal will be to prove that $\beta^k \triangleleft \beta$ and $\gamma^k \triangleleft \gamma$ for all k .

CLAIM 8.6. $\beta^0 \triangleleft \beta$ and $\gamma^0 \triangleleft \gamma$.

We have assumed that α is not ∞ -solvably related to $\alpha \vee \beta$, so $\alpha \vee \beta \not\leq \text{rad}(\alpha)$. Since $\alpha \leq \text{rad}(\alpha)$, we conclude that $\beta \not\leq \text{rad}(\alpha)$.

Since $\text{rad}(\alpha)$ is a radical congruence, Lemma 7.11 (3) guarantees that $\beta^0 := \beta \wedge \text{rad}(\alpha) \triangleleft \beta$ and similarly $\gamma^0 \triangleleft \gamma$.

CLAIM 8.7. $\alpha \leq (\beta^k : \beta)$ and $\alpha \leq (\gamma^k : \gamma)$ for all k .

From $\beta^0 := \beta \wedge \text{rad}(\alpha)$ we derive $\mathbf{C}(\text{rad}(\alpha), \beta; \beta^0)$ using Theorem 2.19 (8). Hence $\alpha \leq \text{rad}(\alpha) \leq (\beta^0 : \beta)$. An identical argument proves that $\alpha \leq (\gamma^0 : \gamma)$.

Now assume that $\alpha \leq (\beta^k : \beta)$ for some k . Since $\gamma^{k+1} = \gamma \wedge (\beta^k : \beta)$ we compute that

$$\gamma^{k+1} \leq \gamma \wedge (\alpha \vee \gamma^{k+1}) \leq \gamma \wedge ((\beta^k : \beta) \vee \gamma^{k+1}) = \gamma \wedge (\beta^k : \beta) = \gamma^{k+1},$$

so $\gamma^{k+1} = \gamma \wedge (\alpha \vee \gamma^{k+1})$. By Theorem 2.19 (8) we conclude that $\mathbf{C}(\alpha, \gamma; \gamma^{k+1})$ holds, or $\alpha \leq (\gamma^{k+1} : \gamma)$. This proof that $\alpha \leq (\beta^k : \beta)$ implies $\alpha \leq (\gamma^{k+1} : \gamma)$ works for every k , and for β and γ interchanged, so the claim follows by induction.

CLAIM 8.8. If $\beta^k \triangleleft \beta$, then $\gamma^{k+1} \triangleleft \gamma$. Similarly, if $\gamma^k \triangleleft \gamma$, then $\beta^{k+1} \triangleleft \beta$.

If $\beta^k \triangleleft \beta$, then $I[\beta^k, \beta]_\beta$ is a separated solvability obstruction, so according to Lemma 7.11 (2) the pseudocomplement $(\beta^k : \beta)$ is a radical congruence. We proved in Claim 8.7 that $\alpha \leq (\beta^k : \beta)$. If it is also the case that $\gamma \leq (\beta^k : \beta)$, then

$$\beta \leq \alpha \vee \beta = \alpha \vee \gamma \leq (\beta^k : \beta),$$

forcing $\mathbf{C}(\beta, \beta; \beta^k)$, or equivalently $\beta^k \triangleleft \beta$. This is contrary to $\beta^k \triangleleft \beta$, so it must be that $\gamma \not\leq (\beta^k : \beta)$. Now Lemma 7.11 (3) guarantees that

$$\gamma^{k+1} := \gamma \wedge (\beta^k : \beta) \triangleleft \gamma.$$

The proof that $\gamma^k \triangleleft \gamma$ implies $\beta^{k+1} \triangleleft \beta$ is the same.

CLAIM 8.9. $I[\gamma^{k+1}, \gamma]_\gamma \rightsquigarrow I[\beta^k, \beta]_\beta$ and $I[\beta^{k+1}, \beta]_\beta \rightsquigarrow I[\gamma^k, \gamma]_\gamma$ for all k .

The second paragraph of Lemma 7.12 describes a method to produce \rightsquigarrow -related intervals, which we employ now. Since $I[\beta^k, \beta]_\beta$ is a separated solvability obstruction and $\gamma \not\leq (\beta^k : \beta)$, then for $\gamma^{k+1} := \gamma \wedge (\beta^k : \beta)$ we have that $I[\gamma^{k+1}, \gamma]_\gamma$ is also a separated solvability obstruction. Moreover, as is shown in Lemma 7.12, to establish that $I[\gamma^{k+1}, \gamma]_\gamma \rightsquigarrow I[\beta^k, \beta]_\beta$ it is sufficient to verify that $\gamma \cap \text{Cg}^{\mathbf{A}}(\beta) \subseteq \gamma^{k+1}$. But $\gamma \cap \text{Cg}^{\mathbf{A}}(\beta) = \gamma \cap \beta = 0$, so this certainly holds.

Under the assumption that α is not ∞ -solvably related to $\alpha \vee \beta$ we have produced an infinite \rightsquigarrow -sequence

$$\dots \rightsquigarrow I[\beta^4, \beta]_\beta \rightsquigarrow I[\gamma^3, \gamma]_\gamma \rightsquigarrow I[\beta^2, \beta]_\beta \rightsquigarrow I[\gamma^1, \gamma]_\gamma \rightsquigarrow I[\beta^0, \beta]_\beta.$$

This completes the proof of the theorem, as we explained in the second paragraph of this proof. \square

Let \mathcal{V} be a variety satisfying a nontrivial congruence identity. Theorem 8.5 shows that if $\mathbf{A} \in \mathcal{V}$, then the SD_\vee -failures in $\mathbf{Con}(\mathbf{A})$ are abelian. Since \mathcal{V} has a weak difference term, these abelian intervals consist of permuting equivalence relations, hence these intervals are modular. It turns out that this property characterizes the class of varieties satisfying nontrivial congruence identities.

THEOREM 8.10. *The following are equivalent for a variety \mathcal{V} .*

- (1) \mathcal{V} satisfies a nontrivial congruence identity.
- (2) For every $\mathbf{A} \in \mathcal{V}$, $\mathbf{Con}(\mathbf{A})$ has an $\text{SD}_\vee/\text{Modular}$ factorization in the category of meet continuous lattices.
- (3) For every $\mathbf{A} \in \mathcal{V}$, $\mathbf{Con}(\mathbf{A})$ has an $\text{SD}_\vee/\text{Modular}$ factorization in the category of ordinary lattices.
- (4) For every $\mathbf{A} \in \mathcal{V}$, each SD_\vee -failure in $\mathbf{Con}(\mathbf{A})$ is modular.
- (5) If \mathbf{L} belongs to $\text{CON}(\mathcal{V})$, then each SD_\vee -failure in \mathbf{L} is modular.

PROOF. If (1) holds, then from Theorem 7.15 (1) \implies (2), Theorem 5.25 (2) \implies (3), and Corollary 6.3 it can be deduced that \mathcal{V} has a weak difference term. Therefore the natural homomorphism

$$\nu: \mathbf{Con}(\mathbf{A}) \rightarrow \mathbf{Con}(\mathbf{A})/\overset{\text{s}}{\sim}$$

provides a factorization of $\mathbf{Con}(\mathbf{A})$ of the form $\text{SD}_\wedge/\text{Modular}$ in the category of meet continuous lattices. We need to argue that the quotient $\mathbf{Con}(\mathbf{A})/\overset{\text{s}}{\sim}$ is not only meet semidistributive, but is also join semidistributive.

SD_\vee -failures in $\mathbf{Con}(\mathbf{A})/\overset{\text{s}}{\sim}$ are each of the form $I = I[\varphi(p), \varphi(j)]$ for some homomorphism of partial lattices $\varphi: \mathbf{P} \rightarrow \mathbf{Con}(\mathbf{A})/\overset{\text{s}}{\sim}$ from the partial lattice $\mathbf{P} := \mathbf{P}(\text{SD}_\vee)$ of the SD_\vee -configuration. By the projectivity of \mathbf{P} , proved in Lemma 8.4, to each such homomorphism there is a homomorphism $\bar{\varphi}: \mathbf{P} \rightarrow \mathbf{Con}(\mathbf{A})$ such that $\nu \circ \bar{\varphi} = \varphi$ where $\nu: \mathbf{Con}(\mathbf{A}) \rightarrow \mathbf{Con}(\mathbf{A})/\overset{\text{s}}{\sim}$ is the natural map. The homomorphism $\bar{\varphi}$ identifies an SD_\vee -failure $I[\bar{\varphi}(p), \bar{\varphi}(j)]$ in $\mathbf{Con}(\mathbf{A})$, which by Theorem 8.5 must be abelian. But then $\bar{\varphi}(p) \overset{\text{s}}{\sim} \bar{\varphi}(j)$, so since $\overset{\text{s}}{\sim}$ is contained in $\ker(\nu)$ we derive that

$$\varphi(p) = \nu \circ \bar{\varphi}(p) = \nu \circ \bar{\varphi}(j) = \varphi(j).$$

This shows that our original SD_{\vee} -failure $I[\varphi(p), \varphi(j)]$ in $\mathbf{Con}(\mathbf{A})/\overset{s}{\sim}$ is trivial. Since it was chosen arbitrarily, $\mathbf{Con}(\mathbf{A})/\overset{s}{\sim}$ is join semidistributive.

As noted when proving the dual result (Theorem 6.17) the implications $(2) \implies (3) \implies (4)$ are straightforward.

We prove the contrapositive of $(4) \implies (5)$. If there is an $\mathbf{K} \in \mathbf{CON}(\mathcal{V})$ that has a nonmodular SD_{\vee} -failure, then since \mathcal{V} consists of the homomorphic images of lattices in the congruence prevariety of \mathcal{V} there is an algebra $\mathbf{A} \in \mathcal{V}$, a sublattice $\mathbf{L} \leq \mathbf{Con}(\mathbf{A})$, and surjective homomorphism $\sigma: \mathbf{L} \rightarrow \mathbf{K}$. The nonmodular SD_{\vee} -failure I in \mathbf{K} can be lifted to an SD_{\vee} -failure \bar{I} in \mathbf{L} using the projectivity of the SD_{\vee} -configuration. The fact that I is nonmodular means that it contains a sublattice isomorphic to the pentagon, \mathbf{N}_5 . The pentagon is a projective lattice, so it can be lifted to a pentagon in \bar{I} , establishing that \bar{I} is a nonmodular SD_{\vee} -failure in \mathbf{L} , hence the interval in $\mathbf{Con}(\mathbf{A})$ generated by \bar{I} is also a nonmodular SD_{\vee} -failure. This proves that $(4) \implies (5)$.

The implication $(5) \implies (4)$ is trivial, so we complete the proof by showing that $(4) \implies (1)$. Assume that (4) holds. This fact can be expressed by a quasi-identity valid in congruence lattices of \mathcal{V} . We have already discussed how to write down a quasi-identity for the dual property that SD_{\wedge} -failures are modular in Remark 2.24, so it is easy to write down the one we need now. Namely, set $t := p \vee (q \wedge r)$, and for each i set $x_i^* := (x_i \wedge s) \vee t$. A quasi-identity expressing the fact that SD_{\vee} -failures are modular is

$$(8.1) \quad ((p \vee q) \approx s) \ \& \ ((p \vee r) \approx s) \rightarrow u(x_1^*, x_2^*, x_3^*) \approx v(x_1^*, x_2^*, x_3^*)$$

where $u \approx v$ is the modular law. As explained in the paragraph after Definition 2.4, the fact that the premises of (8.1) are meet-free implies that (8.1) satisfies (W). By Theorem 2.22 this means that the class of varieties for which item (4) is true is definable by a set of idempotent Maltsev conditions. At least one of these must fail in the variety of semilattices, since, as we now prove, the variety of semilattices does not satisfy item (4). Let $\mathbf{2} = \langle \{0, 1\}; + \rangle$ be the 2-element join semilattice, and let $\mathbf{B} = \mathbf{2}^2$. We claim that $\mathbf{Con}(\mathbf{B}^2)$ has a nonmodular SD_{\vee} -failure. To see this, let θ be the kernel of the homomorphism $+: \mathbf{B}^2 \rightarrow \mathbf{B}$. That is, $\theta = \{ \langle (a, b), (c, d) \rangle \mid a + b = c + d \}$. Let $\eta_1 = \{ \langle (a, b), (c, d) \rangle \mid a = c \}$ and $\eta_2 = \{ \langle (a, b), (c, d) \rangle \mid b = d \}$ be the coordinate projection kernels. For any two pairs $(p, q), (r, s) \in B^2$ it is the case that

$$(p, q) \ \eta_1 \ (p, p + q + r + s) \ \theta \ (p + q + r + s, p + q + r + s),$$

and similarly $(p + q + r + s, p + q + r + s) \ \theta \circ \eta_1 \ (r, s)$, so (p, q) is related to (r, s) by $\theta \vee \eta_1$. The pairs were arbitrary, so $\theta \vee \eta_1 = 1$; similarly

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$\theta \vee \eta_2 = 1$. This proves that $I[\theta \vee (\eta_1 \wedge \eta_2), 1] = I[\theta, 1]$ is an SD_\vee -failure. But \mathbf{B}^2/θ is isomorphic to \mathbf{B} , since $+$ is a surjective homomorphism, so $I[\theta, 1] \cong \mathbf{Con}(\mathbf{B}) = \mathbf{Con}(\mathbf{2}^2) \cong \mathbf{D}_2$, which is a nonmodular lattice. This proves that any set of idempotent Maltsev conditions defining the class of varieties satisfying item (4) includes one that fails in the variety of semilattices, hence from Theorem 7.15 any variety satisfying (4) must satisfy (1). \square

The next result is an analogue of the \mathbf{D}_1 -version of Theorem 4.23.

THEOREM 8.11. *Let \mathcal{V} be a variety. The following conditions are equivalent.*

- (1) \mathcal{V} satisfies an idempotent Maltsev condition that fails in the variety of semilattices.
- (2) \mathcal{V} satisfies a nontrivial congruence identity.
- (3) \mathbf{D}_2 does not appear as a sublattice of $\mathbf{Con}(\mathbf{A})$ for any $\mathbf{A} \in \mathcal{V}$.
- (4) $\mathbf{D}_2 \notin \text{CON}(\mathcal{V})$.

PROOF. We proved the equivalence of (1) and (2) in Theorem 7.15.

If item (3) holds, then \mathcal{V} omits special \mathbf{D}_2 's. According to Theorems 5.28, this implies that \mathcal{V} satisfies an idempotent Maltsev condition that fails in the variety of semilattices, so (3) \implies (1).

Now suppose that item (2) holds and that item (3) fails to hold. The failure of (3) implies that some $\mathbf{A} \in \mathcal{V}$ had a copy of \mathbf{D}_2 in its congruence lattice. Choose a copy and label it as in Figure 8.5. Then

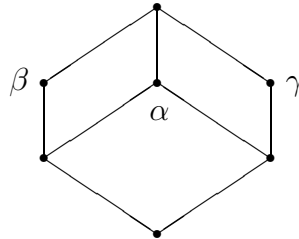


FIGURE 8.5.

$\alpha \vee \beta = \alpha \vee \gamma$, and $\alpha \vee (\beta \wedge \gamma) = \alpha$. Therefore, since item (2) holds, Theorem 8.5 guarantees that $\alpha \triangleleft \alpha \vee \beta$. But this is forbidden by Theorem 4.16 (2). Thus (1), (2) and (3) are equivalent.

Both implications (4) \implies (2) and (4) \implies (3) are trivial, so to complete the proof we will show that items (2) and (3) jointly imply (4). Assume that items (2) and (3) are true for \mathcal{V} . From (2) and Theorem 8.3 we obtain that there is some positive integer M such that the identity $z^M \approx z^{M+1}$ holds in $\text{CON}(\mathcal{V})$. Let \mathcal{Z}^M denote the variety

of lattices axiomatized by $z^M \approx z^{M+1}$. It is proved in Theorem 2.4 (5) \implies (6) of [48] that \mathbf{D}_2 is projective in any \mathcal{Z}^M that contains \mathbf{D}_2 . Since \mathbf{D}_2 is projective and subdirectly irreducible in \mathcal{Z}^M , the subclass $\mathcal{U} \subseteq \mathcal{Z}^M$ of lattices that have no sublattice isomorphic to \mathbf{D}_2 is a variety. By item (3), \mathcal{U} contains the congruence lattices of members of \mathcal{V} , hence \mathcal{U} contains the variety generated by these lattices. Since $\text{CON}(\mathcal{V}) \subseteq \mathcal{U}$ and $\mathbf{D}_2 \notin \mathcal{U}$, (4) holds. \square

The statement of the next theorem, which characterizes the class of varieties satisfying congruence identities in yet another way, should be compared to the statement of Theorem 4.12, which characterizes the class of varieties satisfying meet continuous congruence identities.

THEOREM 8.12. *The following conditions on a variety \mathcal{V} are equivalent.*

- (1) \mathcal{V} satisfies an idempotent Maltsev condition that fails in the variety of semilattices.
- (2) The quasi-identity

$$(8.2) \quad ((p \vee q) \approx s) \ \& \ ((p \vee r) \approx s) \ \& \ (p^{[2]} \approx s) \ \rightarrow \ ((p \vee (q \wedge r)) \approx s).$$

holds in the congruence lattices of algebras in \mathcal{V} .

- (3) Quasi-identity (8.2) holds in $\text{CON}(\mathcal{V})$.
- (4) \mathcal{V} satisfies a nontrivial congruence identity.

Quasi-identity (8.2) in item (2) is the dual of quasi-identity (4.4) of Theorem 4.12 (2).

PROOF. We proved the equivalence of (1) and (4) in Theorem 7.15. The implication (3) \implies (2) is obvious. We have previously noted that \mathbf{D}_1 fails quasi-identity (4.4) of Theorem 4.12 (2), so of course \mathbf{D}_2 fails the dual quasi-identity of the present theorem. So, if the congruence lattices of algebras in \mathcal{V} satisfy quasi-identity (8.2), then \mathbf{D}_2 is not embeddable in the congruence lattice of any algebra in \mathcal{V} . From Theorem 8.11 it follows that \mathcal{V} satisfies a nontrivial congruence identity. This shows that (2) \implies (4). In summary, these observations show that (3) \implies (2) \implies (4) \iff (1). The rest of the proof will be devoted to showing that (4) \implies (3). We argue by contradiction, so assume that \mathcal{V} satisfies a nontrivial congruence identity, but $\text{CON}(\mathcal{V})$ fails to satisfy quasi-identity (8.2).

Quasi-identity (8.2) is equivalent modulo the axioms of lattice theory to the following quasi-identity:

$$(8.3) \quad ((p \vee q) \approx s) \ \& \ ((p \vee r) \approx s) \ \& \ (p^{[2]} \approx s) \ \& \ ((q \wedge r) \leq p) \ \rightarrow \ (p \approx s).$$

This is because any assignment of variables in some lattice that fails (8.3) also fails (8.2), while, if $p \mapsto a, q \mapsto b, r \mapsto c, s \mapsto d$ is an assignment that fails (8.2), then the reassignment $p \mapsto a' := a \vee (b \wedge c), q \mapsto b, r \mapsto c, s \mapsto d$ clearly satisfies the first, second and fourth premise of (8.3) and fails the conclusion. To see that this reassignment also satisfies the third premise of (8.3) we must verify that $(a')^{[2]}(a', b, c) = d$. The fact that the first two premises are satisfied implies that d is the largest element of the sublattice generated by $\{a, b, c\}$. This shows that $(a')^{[2]}(a', b, c) \leq d$. But since the lattice operations are monotone, and $a \leq a'$, we get $d = a^{[2]}(a, b, c) \leq (a')^{[2]}(a', b, c)$. Thus $(a')^{[2]}(a', b, c) = d$, and the third premise is satisfied.

Let Q denote quasi-identity (8.3), and let $(\mathbf{P}(Q), p \approx s)$ be the Q -configuration. Since we have assumed that some $\mathbf{K} \in \text{CON}(\mathcal{V})$ fails Q , there is a lattice $\mathbf{K} \in \text{CON}(\mathcal{V})$ and a homomorphism of partial lattices $\varphi: \mathbf{P}(Q) \rightarrow \mathbf{K}$ such that $\varphi(p) < \varphi(s)$. For concreteness, let $a := \varphi(p), b := \varphi(q), c := \varphi(r)$ and $d := \varphi(s)$. Since $\mathbf{K} \in \text{CON}(\mathcal{V})$ there is an algebra $\mathbf{A} \in \mathcal{V}$, a sublattice $\mathbf{L} \leq \text{Con}(\mathbf{A})$ and a surjective homomorphism $\sigma: \mathbf{L} \rightarrow \mathbf{K}$. Choose a function $f: \{p, q, r, s\} \rightarrow \mathbf{L}$ such that $\sigma \circ f = \varphi$ on the domain of f . Since φ is a homomorphism of partial lattices and $q \wedge r \leq p, q, r, s$ in $\mathbf{P}(Q)$, the element $\varphi(q) \wedge \varphi(r) = b \wedge c$ is the smallest element of the sublattice generated by $\{a, b, c, d\}$. Thus, we may assume that f is chosen so that $f(q) \wedge f(r) \leq f(p), f(q), f(r), f(s)$.

The function f is an assignment of the variables of Q in \mathbf{L} . We wish to modify this assignment to a better one. Since $\mathbf{P}(Q)$ may be obtained (up to isomorphism) from $\mathbf{P}(\text{SD}_\vee)$ by adding more relations, there is a (surjective) homomorphism $\psi: \mathbf{P}(\text{SD}_\vee) \rightarrow \mathbf{P}(Q)$ that is the identity on the generators. Applying Lemma 8.4 to the assignment $f \circ \psi: \{p, q, r, s\} \rightarrow \mathbf{L}$ we obtain a homomorphism of partial lattices $\bar{\psi}: \mathbf{P}(\text{SD}_\vee) \rightarrow \mathbf{L}$ such that $\sigma \circ \bar{\psi} = \varphi \circ \psi$ (Figure 8.6). The second

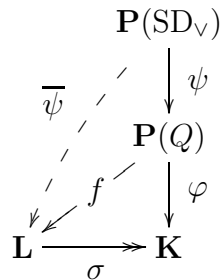


FIGURE 8.6.

clause of Lemma 8.4 guarantees that

$$(8.4) \quad \begin{aligned} \overline{\psi}(q) \wedge \overline{\psi}(r) &\leq (f \circ \psi(q)) \wedge (f \circ \psi(r)) \\ &= f(q) \wedge f(r) \leq f(p) \\ &= f \circ \psi(p) \\ &\leq \overline{\psi}(p), \end{aligned}$$

so if $\overline{\psi}$ is taken to define an assignment of the variables of Q in \mathbf{L} , $p \mapsto \alpha, q \mapsto \beta, r \mapsto \gamma, s \mapsto \delta$, then

- (i) $\alpha \vee \beta = \delta$ and $\alpha \vee \gamma = \delta$, (by the projectivity of $\mathbf{P}(\text{SD}_\vee)$)
- (ii) $\beta \wedge \gamma \leq \alpha$ (by (8.4)), and
- (iii) $\sigma(\alpha) = a < d = \sigma(\delta)$ (by the projectivity of $\mathbf{P}(\text{SD}_\vee)$).

This assignment is not necessarily a failure of Q in \mathbf{L} , since the third premise $\alpha^{[2]} = \delta$ of Q may not be satisfied. But since the third premise was satisfied by $\varphi: \mathbf{P}(Q) \rightarrow \mathbf{K}$ we at least have

$$(iv) \sigma(\alpha^{[2]}) = a^{[2]} = d = \sigma(\delta).$$

The rest of the argument produces a contradiction from properties (i)—(iv). From items (i) and (ii) we derive that $I[\alpha, \alpha \vee \beta]$ is an SD_\vee -failure, hence is an abelian interval according to Theorem 8.5. By this and item (i) we get $\alpha \vee \beta \overset{s}{\sim} \alpha \overset{s}{\sim} \alpha \vee \gamma$, which is hypothesis (1) of Lemma 6.27. Item (ii) above is the same as hypothesis (2) of Lemma 6.27, and our current assumption that \mathcal{V} satisfies a congruence identity is strong enough to guarantee that \mathcal{V} has a weak difference term. Since all the hypotheses of Lemma 6.27 currently hold, its conclusion also holds, which is the equality

$$(\beta \vee (\alpha \wedge \gamma)) \wedge (\gamma \vee (\alpha \wedge \beta)) = (\alpha \wedge \beta) \vee (\alpha \wedge \gamma).$$

Using this to move from the first line to the second in

$$\begin{aligned} \alpha^{[2]}(\alpha, \beta, \gamma) &:= \alpha \vee ((\beta \vee (\alpha \wedge \gamma)) \wedge (\gamma \vee (\alpha \wedge \beta))) \\ &= \alpha \vee ((\alpha \wedge \beta) \vee (\alpha \wedge \gamma)) \\ &= \alpha, \end{aligned}$$

we derive that $\alpha^{[2]} = \alpha$. But from item (iii) we have $\sigma(\alpha) < \sigma(\delta)$, while from item (iv) we have $\sigma(\alpha^{[2]}) = \sigma(\delta)$, which together conflict with $\alpha^{[2]} = \alpha$. This contradiction completes the proof. \square

The following result defines the class of varieties satisfying a non-trivial congruence identity by a 3-variable Maltsev condition.

THEOREM 8.13. *Let \mathcal{V} be a variety. The following conditions are equivalent.*

- (1) \mathcal{V} satisfies a nontrivial congruence identity.
- (2) $\mathcal{V} \models_{\text{con}} \alpha \cap (\beta \circ \gamma) \subseteq (\gamma \vee (\alpha \wedge \beta)) \circ (\beta \vee (\alpha \wedge \gamma))$.

- (3) *There is a positive integer k and ternary terms $d_0, \dots, d_{2k+1}, e_0, \dots, e_{2k+1}, p$ such that \mathcal{V} satisfies the following equations:*
- (i) $d_0(x, y, z) \approx p(x, y, z) \approx e_0(x, y, z)$;
 - (ii) $d_i(x, y, y) \approx d_{i+1}(x, y, y)$ and $e_i(x, x, y) \approx e_{i+1}(x, x, y)$ for even i ;
 - (iii) $d_i(x, x, y) \approx d_{i+1}(x, x, y)$, $d_i(x, y, x) \approx d_{i+1}(x, y, x)$, $e_i(x, y, y) \approx e_{i+1}(x, y, y)$ and $e_i(x, y, x) \approx e_{i+1}(x, y, x)$ for odd i ;
 - (iv) $d_{2k+1}(x, y, z) \approx x$ and $e_{2k+1}(x, y, z) \approx z$.

PROOF. For locally finite varieties, this result is part of Theorem 9.8 of [34]. The proof given there shows that items (2) and (3) are equivalent for any variety, and that the idempotent Maltsev condition in item (3) fails in the variety of semilattices. Therefore (3) \implies (1), according to Theorem 7.15.

What remains to show is that (1) \implies (2). Let $\mathbf{F} = \mathbf{F}_{\mathcal{V}}(x, y, z)$ and set $\alpha := \text{Cg}^{\mathbf{F}}(x, z)$, $\beta := \text{Cg}^{\mathbf{F}}(x, y)$, and $\gamma := \text{Cg}^{\mathbf{F}}(y, z)$. Since $(x, z) \in (\gamma \vee (\alpha \wedge \beta)) \circ (\beta \vee (\alpha \wedge \gamma))$ it will suffice to prove that $(x, z) \in (\gamma \vee (\alpha \wedge \beta)) \circ (\beta \vee (\alpha \wedge \gamma))$. Since $\alpha \vee \beta = \alpha \vee \gamma = \beta \vee \gamma =: \delta$, the intervals $I[\beta \vee (\alpha \wedge \gamma), \delta]$ and $I[\gamma \vee (\alpha \wedge \beta), \delta]$ are SD_{\vee} -failures, and therefore $\beta \vee (\alpha \wedge \gamma) \stackrel{s}{\sim} \delta \stackrel{s}{\sim} \gamma \vee (\alpha \wedge \beta)$ (Theorem 8.5). Since these congruences are solvably related, they permute (Theorem 6.16). Since $(x, y) \in \beta \vee (\alpha \wedge \gamma)$ and $(y, z) \in \gamma \vee (\alpha \wedge \beta)$, we get that

$$(x, z) \in (\beta \vee (\alpha \wedge \gamma)) \circ (\gamma \vee (\alpha \wedge \beta)) = (\gamma \vee (\alpha \wedge \beta)) \circ (\beta \vee (\alpha \wedge \gamma)).$$

□

8.3. Congruence Join Semidistributivity

The main purpose of this section is to prove that a variety is congruence join semidistributive if and only if it is both congruence meet semidistributive and satisfies a nontrivial congruence identity. Thus, congruence join semidistributive varieties are exactly those varieties whose members have no congruences that are ‘abelian’ in any of the senses that have been introduced.

THEOREM 8.14. *Let \mathcal{V} be a variety. The following conditions are equivalent.*

- (1) \mathcal{V} is congruence join semidistributive.
- (2) If $\mathbf{A} \in \mathcal{V}$ and $\alpha, \beta, \gamma \in \text{Con}(\mathbf{A})$, then

$$\alpha \cap (\beta \circ \gamma) \subseteq \beta \vee (\alpha \wedge \gamma).$$

- (3) *There is a positive integer k and ternary terms d_0, \dots, d_k such that \mathcal{V} satisfies the following equations:*

consider pointing out the similarity of this theorem and theorem 9.11 from hobby-mckenzie.

- (i) $d_0(x, y, z) \approx x$;
 - (ii) $d_i(x, y, y) \approx d_{i+1}(x, y, y)$ and $d_i(x, y, x) \approx d_{i+1}(x, y, x)$ for even $i < k$;
 - (iii) $d_i(x, x, y) \approx d_{i+1}(x, x, y)$, for odd $i < k$;
 - (iv) $d_k(x, y, z) \approx z$.
- (4) \mathcal{V} satisfies an idempotent Maltsev condition that fails in any nontrivial variety of modules and in the variety of semilattices.
- (5) No member of \mathcal{V} has a nontrivial abelian or rectangular congruence.
- (6) \mathcal{V} is congruence meet semidistributive and satisfies a nontrivial congruence identity.
- (7) There is a positive integer N such that if $\mathbf{A} \in \mathcal{V}$ and $\alpha, \beta, \gamma \in \text{Con}(\mathbf{A})$, then

$$\alpha \vee (\beta \wedge \gamma) \approx (\alpha \vee \beta^N) \wedge (\alpha \vee \gamma^N).$$

- (8) $\text{CON}(\mathcal{V})$ is a join semidistributive variety of lattices.
- (9) $\mathbf{M}_3, \mathbf{D}_2 \notin \text{CON}(\mathcal{V})$.
- (10) \mathbf{M}_3 and \mathbf{D}_2 are not embeddable in $\text{Con}(\mathbf{A})$ for any $\mathbf{A} \in \mathcal{V}$.

PROOF. We prove (1)–(6) cyclically, and then complete the proof by showing that ((1) & (6)) \implies (7) \implies (8) \implies (9) \implies (10) \implies (6).

Assume that (1) holds and that $\mathbf{A} \in \mathcal{V}$ has congruences α, β and γ . Choose any $(a, c) \in \alpha \cap (\beta \circ \gamma)$. There exists $b \in A$ such that $a \equiv_\beta b \equiv_\gamma c$. Let $\alpha' = \text{Cg}^{\mathbf{A}}(a, c)$, $\beta' = \text{Cg}^{\mathbf{A}}(a, b)$ and $\gamma' = \text{Cg}^{\mathbf{A}}(b, c)$. Since the generating pairs are related by α, β and γ respectively, we get that $\alpha' \leq \alpha$, $\beta' \leq \beta$ and $\gamma' \leq \gamma$. The congruence $\beta' \vee \alpha'$ is generated by (a, c) and (a, b) , so it is the congruence on \mathbf{A} generated by $X \times X$ where $X = \{a, b, c\}$. Similarly $\beta' \vee \gamma'$ is the congruence generated by $X \times X$, so $\beta' \vee \alpha' = \beta' \vee \gamma'$. By (1), we get that this congruence is also $\beta' \vee (\alpha' \wedge \gamma')$. Since this congruence contains (a, c) , we get that

$$(a, c) \in \beta' \vee (\alpha' \wedge \gamma') \leq \beta \vee (\alpha \wedge \gamma),$$

where the inequality follows from the monotonicity of the lattice operations. This argument shows that $(a, c) \in \beta \vee (\alpha \wedge \gamma)$ for every $(a, c) \in \alpha \cap (\beta \circ \gamma)$, so $\alpha \cap (\beta \circ \gamma) \subseteq \beta \vee (\alpha \wedge \gamma)$ and (2) holds.

For (2) \implies (3), let $p = \alpha \cap (\beta \circ \gamma)$, $r = \beta \vee (\alpha \wedge \gamma)$, and let r_i be the $\{\circ, \wedge\}$ -word obtained from r by replacing \vee with i -fold relational product. Then the congruence inclusion of (2) has the form $p \subseteq \bigcup_{i \in \omega} r_i$ with the r_i 's increasing. By the remark after Theorem 4.7, such a congruence inclusion is associated to a Maltsev condition. The method described in the proof of Theorem 4.7 for producing the Maltsev condition yields the one in item (3).

To prove that (3) \implies (4), we will show that the idempotent Maltsev condition in (3) fails in every nontrivial variety of modules and in the variety of semilattices. Assume that \mathcal{M} is the variety of all modules over the ring \mathbf{R} , and that \mathcal{M} satisfies the Maltsev condition in item (3). Since each d_i is a module term, it must be that $d_i(x, y, z) = a_i x + b_i y + c_i z$ for some $a_i, b_i, c_i \in R$. Since $d_i(x, x, x) \approx x$, we get that $a_i x + b_i x + c_i x \approx (a_i + b_i + c_i)x \approx x$ holds in any $\mathbf{M} \in \mathcal{M}$, hence $a_i + b_i + c_i = 1$ in \mathbf{R} . The first equation in (3)(ii) implies that $a_i = a_{i+1}$ for even i , the second equation in (3)(ii) implies that $b_i = b_{i+1}$ for even i , and with the equations of the type $a_i + b_i + c_i = 1$ we derive that $c_i = c_{i+1}$ holds for all even i . The equation from (3)(iii) implies that $c_i = c_{i+1}$ holds for odd i , so $c_i = c_{i+1}$ for all i . But the equation from (3)(i) implies that $c_0 = 0$, while the equation from (3)(iv) implies that $c_k = 1$. Altogether we derive that $0 = 1$ in \mathbf{R} , so \mathcal{M} is a trivial variety of modules.

Now suppose that \mathcal{S} is the variety of semilattices, and that \mathcal{S} satisfies the Maltsev condition in (3). Choose k minimal for the condition in (3) to be satisfied in \mathcal{S} . By equations (3)(i) and (3)(iv), it must be that $k > 0$. If $k > 0$ is even, then by equation (3)(iii) we find that $d_{k-1}(x, x, y) \approx d_k(x, x, y) \approx y$, which for semilattices implies that $d_{k-1}(x, y, z) \approx z$. Thus we can delete the term d_k and have a shorter sequence of terms satisfying the conditions in (3). If $k > 0$ is odd, then by the equations in (3)(ii) that d_{k-1} does not depend on its first variable (since $d_{k-1}(x, y, y) \approx d_k(x, y, y) \approx y$), and does not depend on its second variable (since $d_{k-1}(x, y, x) \approx d_k(x, y, x) \approx x$). Therefore $d_{k-1}(x, y, z) \approx d_{k-1}(z, z, z) \approx z$, and we can again shorten the sequence. It follows that the equations in (3) cannot be satisfied by semilattice terms. Thus (3) \implies (4).

It follows from Theorem 8.1 (10) \implies (4) that if \mathcal{V} satisfies an idempotent Maltsev condition that fails in every nontrivial variety of modules, then \mathcal{V} omits abelian congruences. It follows from Theorem 5.25 that if \mathcal{V} satisfies an idempotent Maltsev condition that fails in the variety of semilattices, then \mathcal{V} omits rectangular tolerances (hence \mathcal{V} omits rectangular congruences by Theorem 5.22). Thus (4) \implies (5).

If (5) holds, then \mathcal{V} has no abelian congruence, so from (4) \implies (1) of Theorem 8.1 we derive that \mathcal{V} is congruence meet semidistributive. \mathcal{V} also has no rectangular congruence, or tolerance, so by Theorems 5.25 and 7.15 we have that \mathcal{V} satisfies a nontrivial congruence identity. Hence (6) holds.

Now suppose that \mathcal{V} is congruence meet semidistributive and satisfies a nontrivial congruence identity. The latter supposition implies that the $\text{SD}_{\mathcal{V}}$ -failures in congruence lattices of members of \mathcal{V} are abelian,

according to Theorem 8.5. But the former supposition implies that there are no abelian intervals in congruence lattices of members of \mathcal{V} , according to (1) \implies (4) of Theorem 8.1. This shows that there are no $\text{SD}_{\mathcal{V}}$ -failures at all, which means of course that \mathcal{V} is congruence join semidistributive.

Now that the first six items of the theorem have been shown to be equivalent, we turn to the seventh. Assume that (1) and (the equivalent property) (6) hold. From (6) and Theorem 8.3 we know that there is a positive integer M such that the identity

$$z^M(x, y, z) \approx z^{M+1}(x, y, z)$$

holds in all congruence lattices of members of \mathcal{V} . Thus, if α, β and γ are as in the statement of item (7), and $\beta^k := y^k(\alpha, \beta, \gamma)$ and $\gamma^k := z^k(\alpha, \beta, \gamma)$, then $\beta^M = \beta^{M+1}$ and $\gamma^M = \gamma^{M+1}$. This implies that $\alpha \vee \beta^M \geq \alpha \vee \gamma^{M+1} = \alpha \vee \gamma^M$, and similarly that $\alpha \vee \gamma^M \geq \alpha \vee \beta^M$, so $\alpha \vee \beta^M = \alpha \vee \gamma^M$. Since (1) holds, we have

$$\alpha \vee (\beta^M \wedge \gamma^M) = \alpha \vee \beta^M.$$

But $\beta \wedge \gamma \leq \beta^M \wedge \gamma^M \leq \beta^M \leq \beta$ and $\beta \wedge \gamma \leq \beta^M \wedge \gamma^M \leq \gamma^M \leq \gamma$, so $\beta^M \wedge \gamma^M = \beta \wedge \gamma$. Therefore

$$\alpha \vee (\beta \wedge \gamma) = \alpha \vee \beta^M,$$

and similarly

$$\alpha \vee (\beta \wedge \gamma) = \alpha \vee \gamma^M.$$

This proves that $\alpha \vee (\beta \wedge \gamma) = (\alpha \vee \beta^M) \wedge (\alpha \vee \gamma^M)$ for the fixed value of M provided by Theorem 8.3.

If item (7) holds, then $\text{CON}(\mathcal{V})$ satisfies the weakened distributive law $x \vee (y \wedge z) \approx (x \vee y^N) \wedge (x \vee z^N)$ for some N . We claim that any lattice satisfying this law is join semidistributive. Indeed, if \mathbf{L} is a lattice, $a, b, c \in L$, and $b^k := y^k(a, b, c)$ and $c^k := z^k(a, b, c)$, then it may be shown by induction that if $a \vee b = a \vee c$, then $b^k = b$ and $c^k = c$ for all k . Thus, if $d := a \vee b = a \vee c$, the identity in (7) implies that

$$a \vee (b \wedge c) = (a \vee b^N) \wedge (a \vee c^N) = (a \vee b) \wedge (a \vee c) = d.$$

Since $a, b, c \in L$ were chosen arbitrarily, \mathbf{L} is join semidistributive.

Item (8) implies item (9) because \mathbf{M}_3 and \mathbf{D}_2 are not join semidistributive.

Item (9) implies item (10) because sublattices of $\mathbf{Con}(\mathbf{A})$ belong to $\text{CON}(\mathcal{V})$.

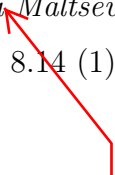
Assume that item (10) holds. Since \mathbf{M}_3 is not embeddable in the congruence lattice of any member of \mathcal{V} , it follows from Theorem 8.1 that \mathcal{V} is congruence meet semidistributive. Since \mathbf{D}_2 is not embeddable in

the congruence lattice of any member of \mathcal{V} , it follows from Theorem 8.11 that \mathcal{V} satisfies a nontrivial congruence identity. Thus item (6) holds. The proof is complete. \square

This theorem provides a positive solution to Problem 2.18 of [40] from this theorem, which we extract as the following corollary.

COROLLARY 8.15. *The class of congruence join semidistributive varieties is definable by a Maltsev condition.*

PROOF. By Theorem 8.14 (1) \iff (3). \square



an idempotent
maltsev condition.

CHAPTER 9

Residually Small Varieties

A variety is **residually large** if it has a proper class of subdirectly irreducible members, otherwise it is **residually small**. Our primary goal in this chapter is to prove that a residually small variety satisfies a congruence identity if and only if it is congruence modular. (David Hobby and Ralph McKenzie proved essentially the same result for locally finite varieties in Chapter 10 of [34].) Our result nearly completes the classification of congruence varieties associated to residually small varieties. In Section 9.2, we apply the result to show that almost congruence distributive varieties cannot have a Taylor term.

9.1. Residual Smallness and Congruence Modularity

LEMMA 9.1. *Let \mathcal{V} be a variety that satisfies a nontrivial congruence identity. If $\mathbf{A} \in \mathcal{V}$ has congruences α, β and γ that generate a sublattice isomorphic to \mathbf{N}_5 in $\text{Con}(\mathbf{A})$, as depicted in Figure 9.1, then*

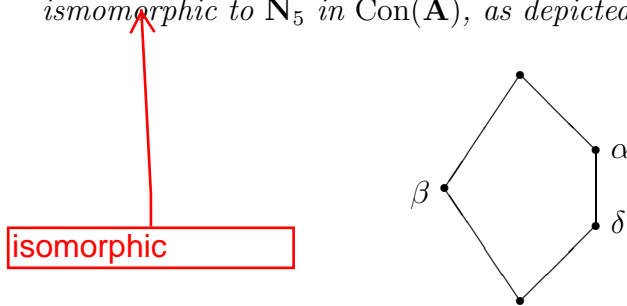


FIGURE 9.1.

- (1) $\mathbf{C}(\alpha, \beta; \delta)$ holds, and
- (2) $\mathbf{C}(\beta, \alpha; \delta)$ fails in the following strong way: For some $(\beta, \beta; 0)$ -pair (a, b) (cf. Definition 3.13) there is a polynomial $f(x, \mathbf{y})$ and tuples $\mathbf{c} \alpha \mathbf{d}$ such that

$$f(b, \mathbf{c}) \equiv_{\delta} f(b, \mathbf{d}),$$

but

$$f(a, \mathbf{c}) \not\equiv_{\delta} f(a, \mathbf{d}).$$

PROOF. Item (1) follows from Theorem 2.19 (8), since

$$\beta \wedge (\alpha \vee (\beta \wedge \delta)) \leq \delta.$$

For (2), assume instead that

$$f(b, \mathbf{c}) \equiv_{\delta} f(b, \mathbf{d}) \implies f(a, \mathbf{c}) \equiv_{\delta} f(a, \mathbf{d})$$

holds for every $(\beta, \beta; 0)$ -pair (a, b) . Then by transitivity this implication holds whenever $a \geq_{\beta} b$. But since \mathcal{V} has no algebra with a rectangular tolerance, Theorem 5.23 (1) \iff (4) guarantees that the set of pairs (a, b) satisfying $a \geq_{\beta} b$ contains β . Thus, if this implication holds for every f and every choice of tuples $\mathbf{c} \alpha \mathbf{d}$ whenever $(a, b) \in \beta$, then $\mathbf{C}(\beta, \alpha; \delta)$ holds. But $\mathbf{C}(\beta, \alpha; \delta)$ does not hold, as we proved in Theorem 4.16 (2). Therefore f, \mathbf{c} and \mathbf{d} exist. \square

As it happens, the properties described in conditions (1) and (2) of Lemma 9.1 are all that are needed to establish the existence of a proper class of subdirectly irreducibles in \mathcal{V} . (We no longer need the assumption that \mathcal{V} satisfies a nontrivial congruence identity.)

is this similar to any of the configurations used in chapter 10 of Hobby-McKenzie? If so, please indicate which one.

THEOREM 9.2. *Let \mathcal{V} be a variety. Assume that $\mathbf{A} \in \mathcal{V}$ has congruences $\delta < \alpha$ and a tolerance T such that*

- (1) $\mathbf{C}(\alpha, T; \delta)$ holds, and
- (2) For some $(T, T; 0)$ -pair (a, b) there is a polynomial $f(x, \mathbf{y})$ and tuples $\mathbf{c} \alpha \mathbf{d}$ such that

$$v := f(b, \mathbf{c}) \equiv_{\delta} f(b, \mathbf{d}) =: v',$$

but

$$u := f(a, \mathbf{c}) \not\equiv_{\delta} f(a, \mathbf{d}) =: u'.$$

Then \mathcal{V} is residually large.

PROOF. Both (1) and (2) continue to hold if we factor by δ , so there is no loss in generality in assuming that $\delta = 0$. (Under this assumption $v = v'$.) We now explain how to construct a subdirectly irreducible algebra in \mathcal{V} of cardinality $\geq \kappa$ for any infinite cardinal κ .

Let \mathbf{B} be the subalgebra of \mathbf{A}^{κ} consisting of those κ -tuples \mathbf{a} where $(a_i, a_j) \in T$ for all coordinates $i < j < \kappa$. For any $a \in A$ let \hat{a} denote the constant κ -tuple satisfying $(\hat{a})_i = a$ for all i . Clearly $\hat{a} \in B$ for all $a \in A$. Let $N = \{u, v\}$ and let $N' = \{u', v'\}$. Since $(a, b) \in T$ it follows that $(u, v) = (f(a, \mathbf{c}), f(b, \mathbf{c})) \in T$ and $(u', v') \in T$, so $N^{\kappa} \cup (N')^{\kappa} \subseteq B$. For an element \mathbf{z} of N^{κ} denote by \mathbf{z}' the element of $(N')^{\kappa}$ obtained from \mathbf{z} by priming the entry in each coordinate. Furthermore, let

$$G = \{(\mathbf{z}, \mathbf{z}') \in B^2 \mid \mathbf{z} \in N^{\kappa} \setminus \{\hat{u}\}\},$$

and let $\psi = \text{Cg}^{\mathbf{B}}(G)$.

CLAIM 9.3. *The pair (\hat{u}, \hat{u}') is not contained in ψ .*

As $u \neq u'$, it is sufficient to prove that $\{\hat{u}\}$ is a singleton ψ -class. For this we must show that if $(\mathbf{z}, \mathbf{z}') \in G$ and p is a unary polynomial of \mathbf{B} , then

$$p(\mathbf{z}) = \hat{u} \iff p(\mathbf{z}') = \hat{u}.$$

Since $\mathbf{z} \in N^\kappa \setminus \{\hat{u}\}$, at least one coordinate of \mathbf{z} is v . If $z_i = v$, then $z'_i = v' = v$, so \mathbf{z} and \mathbf{z}' agree in the i -th coordinate. This implies that $p(\mathbf{z})$ and $p(\mathbf{z}')$ agree in the i -th coordinate. We now argue that these tuples also agree in the j -th coordinate for every other $j < \kappa$. There is nothing to prove for those coordinates where $z_j = v = v' = z'_j$, by the above argument, so we consider only those j where $z_j = u$ and $z'_j = u'$.

The polynomial p of \mathbf{B} comes from a term $g(x, \mathbf{y})$, where appropriate parameters from \mathbf{B} are substituted in place of \mathbf{y} in every coordinate. Denote by \mathbf{r} and \mathbf{s} the parameters occurring at the i -th and the j -th coordinate, respectively. Then $\mathbf{r} T \mathbf{s}$ holds by the definition of \mathbf{B} . We have to prove that $g(u, \mathbf{s}) = g(u', \mathbf{s})$, and we have either

$$(9.1) \quad g(u, \mathbf{s}) = u = g(v, \mathbf{r})$$

(in case $p(\mathbf{z}) = \hat{u}$), or

$$(9.2) \quad g(u', \mathbf{s}) = u = g(v', \mathbf{r})$$

(in case $p(\mathbf{z}') = \hat{v}$). However $C(\alpha, T; 0)$ holds and $\mathbf{c} \alpha \mathbf{d}$, so

$$(g(u, \mathbf{s}) =) \quad g(f(a, \mathbf{c}), \mathbf{s}) = g(f(b, \mathbf{c}), \mathbf{r}) \quad (= g(v, \mathbf{r}))$$

is equivalent to

$$(g(u', \mathbf{s}) =) \quad g(f(a, \mathbf{d}), \mathbf{s}) = g(f(b, \mathbf{d}), \mathbf{r}) \quad (= g(v', \mathbf{r})).$$

Since at least one of (9.1) or (9.2) holds, this implies that $g(u, \mathbf{s}) = g(v, \mathbf{r})$ and $g(u', \mathbf{s}) = g(v', \mathbf{r})$ both hold. Moreover, $v = v'$ implies that $g(v, \mathbf{r}) = g(v', \mathbf{r})$, so in fact

$$g(u, \mathbf{s}) = g(v, \mathbf{r}) = g(v', \mathbf{r}) = g(u', \mathbf{s}).$$

Thus $g(u, \mathbf{s}) = g(u', \mathbf{s})$ indeed holds, completing the argument that $p(\mathbf{z}) = \hat{u} \iff p(\mathbf{z}') = \hat{u}$ holds. This proves the claim.

Since $(\hat{u}, \hat{u}') \notin \psi$ we may extend ψ to a congruence $\psi_0 \geq \psi$ that is maximal with respect to $(\hat{u}, \hat{u}') \notin \psi_0$. Then \mathbf{B}/ψ_0 is a subdirectly irreducible algebra in \mathcal{V} . We claim that $|B/\psi_0| \geq \kappa$.

In order to prove that \mathbf{B}/ψ_0 is as large as this, we need only to exhibit κ elements of \mathbf{B} that are pairwise incongruent modulo ψ_0 . We introduce the following notation for certain elements of A^κ : for $i < j < \kappa$, and for arbitrary elements $a, b, c \in A$ let $[a, b, c]^{i,j}$ denote the tuple whose k -th coordinate is a for $0 \leq k < i$, is b for $i \leq k < j$, and is

c for $j \leq k < \kappa$. (Note that $[a, b, c]^{i,j} \in B$ if a, b and c are pairwise T -related.) Similarly, if $\mathbf{a}, \mathbf{b}, \mathbf{c} \in A^n$, then let $[\mathbf{a}, \mathbf{b}, \mathbf{c}]^{i,j}$ denote the member of B^n whose ℓ -th coordinate is $[a_\ell, b_\ell, c_\ell]^{i,j}$ for $1 \leq \ell \leq n$. We shall omit the superscripts i, j if we think it is clear from the context what they are.

As (a, b) is a $(T, T; 0)$ -pair, there exists a T, T -matrix

$$\begin{bmatrix} f(\mathbf{p}, \mathbf{r}) & f(\mathbf{p}, \mathbf{s}) \\ f(\mathbf{q}, \mathbf{r}) & f(\mathbf{q}, \mathbf{s}) \end{bmatrix} = \begin{bmatrix} * & a \\ a & b \end{bmatrix}.$$

CLAIM 9.4. *For each $0 < i < j < \kappa$ the tuples $[\mathbf{q}, \mathbf{p}, \mathbf{p}]$ and $[\mathbf{q}, \mathbf{q}, \mathbf{p}]$ are not ψ_0 -congruent in every coordinate.*

Suppose that they are. Then in the algebra \mathbf{B} we have that

$$f([\mathbf{q}, \mathbf{p}, \mathbf{p}], [\mathbf{r}, \mathbf{s}, \mathbf{s}]) \psi_0 f([\mathbf{q}, \mathbf{q}, \mathbf{p}], [\mathbf{r}, \mathbf{s}, \mathbf{s}]),$$

that is,

$$[a, a, a] \psi_0 [a, b, a].$$

Therefore

$$[u, u, u] = f([a, a, a], [\mathbf{c}, \mathbf{c}, \mathbf{c}]) \psi_0 f([a, b, a], [\mathbf{c}, \mathbf{c}, \mathbf{c}]) = [u, v, u],$$

and

$$[u', u', u'] = f([a, a, a], [\mathbf{d}, \mathbf{d}, \mathbf{d}]) \psi_0 f([a, b, a], [\mathbf{d}, \mathbf{d}, \mathbf{d}]) = [u', v', u'].$$

But then

$$\widehat{u} = [u, u, u] \psi_0 [u, v, u] \psi [u', v', u'] \psi_0 [u', u', u'] = \widehat{u}'$$

By transitivity and the fact that $\psi_0 \geq \psi$, this yields $(\widehat{u}, \widehat{u}') \in \psi_0$, contradicting the choice of ψ_0 . So the claim is proved.

To finish the proof, let m be the length of the tuples \mathbf{p} and \mathbf{q} , and define a mapping $h: \kappa \rightarrow (B/\psi_0)^m$ in the following way. Let $i < k < \kappa$, and

$$h(i) = [\mathbf{q}, \mathbf{p}, \mathbf{p}]^{i,k} / \psi_0^m.$$

The value of $h(i)$ does not depend on k because the last two blocks of coordinates of the tuple $[\mathbf{q}, \mathbf{p}, \mathbf{p}]$ are both \mathbf{p} . We claim that h is injective. Indeed, suppose that $i < j < \kappa$. Then

$$h(i) = [\mathbf{q}, \mathbf{p}, \mathbf{p}]^{i,j} / \psi_0^m$$

and

$$h(j) = [\mathbf{q}, \mathbf{q}, \mathbf{p}]^{i,j} / \psi_0^m,$$

and these tuples are different by Claim 9.4. Therefore h is indeed injective, so $\kappa \leq |B/\psi_0|^m = |B/\psi_0|$, proving the theorem. \square

consider remarking
that this
generalizes
Theorem 10.4 of
Hobby-McKenzie.

We obtain the main theorem of this chapter by combining Lemma 9.1 with Theorem 9.2.

THEOREM 9.5. *A residually small variety satisfies a congruence identity if and only if it is congruence modular.*

REMARK 9.6. We describe a three-step plan for classifying all congruence varieties associated to residually small varieties.

- (1) Prove that if \mathcal{V} is residually small and satisfies a nontrivial congruence identity, then \mathcal{V} is congruence modular.
- (2) Prove that if \mathcal{V} is residually small and congruence modular, then \mathcal{V} is congruence distributive or $\text{CON}(\mathcal{V}) = \text{CON}(\mathcal{M})$ for some variety \mathcal{M} of modules.
- (3) Classify the congruence varieties of varieties of modules.

Theorem 9.5 accomplishes Step (1). Step (3) was accomplished by Gábor Czédli and George Hutchinson in [36]. Partial results on Step (2) were obtained by Alan Day and Emil W. Kiss in [12]. Namely, Day and Kiss accomplished Step (2) under the further assumption that \mathcal{V} is either locally finite or locally solvable. Their ideas extend a little further than these two cases, but fall short of fully accomplishing Step (2). Therefore, we pose the problem of accomplishing Step (2) in full generality.

PROBLEM 9.7. Show that if \mathcal{V} is residually small and congruence modular, then \mathcal{V} is congruence distributive or $\text{CON}(\mathcal{V}) = \text{CON}(\mathcal{M})$ for some variety \mathcal{M} of modules.

If this is proved, then we will know that the congruence varieties $\text{CON}(\mathcal{V})$ of residually small varieties \mathcal{V} are precisely

- the variety of all lattices — e.g. when \mathcal{V} is the (residually small) variety of sets,
- $\text{CON}(\mathcal{M})$ where \mathcal{M} is a variety of modules,
- the variety of distributive lattices — e.g. when \mathcal{V} is the (residually small) variety of distributive lattices, and
- the variety of trivial lattices (e.g. when \mathcal{V} is a trivial variety).

9.2. Almost Congruence Distributive Varieties

A variety is **affine complete** if every congruence preserving function on every member is a polynomial. In [43], Kalle Kaarli and Alden Pixley announced that any affine complete variety is residually finite and has the property that every finitely generated subvariety is congruence distributive. At the time it was not known whether the entire

variety had to be congruence distributive. Pixley then introduced the following concept.

DEFINITION 9.8. A variety \mathcal{V} is **almost congruence distributive**, or **ACD**, if it is residually finite, every finitely generated subvariety is congruence distributive, but \mathcal{V} itself is not congruence distributive.

Pixley asked if there exists an ACD variety. A negative answer would, of course, show that every affine complete variety is congruence distributive.

Kaarli and Ralph McKenzie proved in [42] that every affine complete variety is congruence distributive, but it remains open whether there is an ACD variety. This is still an interesting question, since we know very little about how local properties of a variety, like the structure of the subdirectly irreducible algebras, imply global properties, such as the satisfaction of a Maltsev condition.

In this section we prove that an ACD variety satisfies no nontrivial idempotent Maltsev condition. This implies that if an ACD variety exists, then some member has a nontrivial strongly abelian congruence, but for some reason such congruences do not show up in any subdirectly irreducible algebra or even in any subvariety generated by a finite set of subdirectly irreducible algebras.

THEOREM 9.9. *Let \mathcal{V} be a variety satisfying a nontrivial idempotent Maltsev condition.*

- (1) *If every subdirectly irreducible algebra in \mathcal{V} generates a subvariety satisfying a nontrivial congruence identity, then \mathcal{V} satisfies a nontrivial congruence identity.*
- (2) *If every subdirectly irreducible algebra in \mathcal{V} generates a congruence join semidistributive subvariety, then \mathcal{V} is congruence join semidistributive.*

Now assume also that \mathcal{V} is residually small.

- (1)' *If every subdirectly irreducible algebra in \mathcal{V} generates a subvariety satisfying a nontrivial congruence identity, then \mathcal{V} is congruence modular.*
- (2)' *If every subdirectly irreducible algebra in \mathcal{V} generates a congruence join semidistributive subvariety, then \mathcal{V} is congruence distributive.*

use a different symbol here.

PROOF. If each subdirectly irreducible algebra in \mathcal{V} generates a subvariety satisfying a nontrivial congruence identity, then by Theorem 7.15 each subdirectly irreducible algebra in \mathcal{V} has a Hobby–McKenzie term. By Theorem 5.25 the rectangular tolerances on subdirectly irreducible members of \mathcal{V} are trivial. Corollary 5.20 guarantees that no algebra in \mathcal{V} has a rectangular tolerance. Reusing Theorems 5.25 and 7.15, we get that \mathcal{V} satisfies a nontrivial congruence identity. This proves (1).

For (2), if every subdirectly irreducible in \mathcal{V} generates a congruence join semidistributive subvariety, then no subdirectly irreducible algebra has a nontrivial rectangular or abelian tolerance. By part (1), this implies that \mathcal{V} satisfies a nontrivial congruence identity, hence \mathcal{V} has a weak difference term according to Corollary 6.3. This shows that the hypotheses of Corollary 6.9 are met, so from the fact that no subdirectly irreducible algebra has an abelian tolerance we may derive that no algebra in \mathcal{V} has an abelian tolerance. By Theorem 8.14 (5) \implies (1) it follows that \mathcal{V} is congruence join semidistributive.

Items (1)' and (2)' follow from items (1) and (2) using Theorem 9.5. \square

COROLLARY 9.10. *An ACD variety satisfies no nontrivial idempotent Maltsev condition.*

PROOF. This follows from Theorem 9.9 (2)' and the definition of an ACD variety. \square

Theorem 9.9 suggests the following problem.

PROBLEM 9.11. Suppose that \mathcal{V} satisfies a nontrivial idempotent Maltsev condition, and every subdirectly irreducible member of \mathcal{V} generates a congruence modular subvariety. Must \mathcal{V} be congruence modular? (We are interested in the case where \mathcal{V} is residually large.)

The answer to Problem 9.11 is affirmative if \mathcal{V} is locally finite. Indeed, it follows from tame congruence theory that the following conditions are equivalent for a locally finite variety:

- (1) \mathcal{V} is congruence modular.
- (2) The minimal sets of the finite algebras are of type **2**, **3** or **4** and have empty tails.
- (3) The minimal sets of the finite subdirectly irreducible algebras are of type **2**, **3** or **4** and have empty tails.
- (4) Each finite subdirectly irreducible algebra generates a congruence modular subvariety.

However, the modular law is special in this regard, because of its link with the empty tails condition. We expect that for identities weaker than modularity the problem analogous to Problem 9.11 will have a negative answer.

Another problem is suggested by the proof of Theorem 9.9.

PROBLEM 9.12. Suppose that \mathcal{V} has a Taylor term. If some algebra in \mathcal{V} has a nontrivial abelian tolerance, must some subdirectly irreducible algebra in \mathcal{V} have a nontrivial abelian tolerance?

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