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For the complete proof, see NST Theorem 28.24, pages 609-611.

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- ① Well order $\text{dom}(\alpha) \in M$ with a bijection $f: \kappa^{\text{card}} \rightarrow \text{dom}(\alpha)$.

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- ⑦ Well order α_G by $\gamma_G \mapsto \text{least } \lambda \in \kappa \text{ such that } f(\lambda)_G = \gamma_G$.