

# Modules, III



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In particular, if the underlying set functor  $U : \mathcal{C} \rightarrow \mathbf{Set}$  is representable, then the underlying set of a product is the product of the underlying sets of the factors.

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We call  $(M \oplus N, p_M, p_N, \iota_M, \iota_N)$  the biproduct of  $M$  and  $N$ .