Definitions and Laws of Arithmetic on $\mathbb{N}$. With Hints!

Addition

$$\begin{aligned} m+0 &:= m && \text{(IC)} \\ m+S(n) &:= S(m+n) && \text{(RR)} \end{aligned}$$

Multiplication

$$\begin{aligned} m\cdot 0 &:= 0 && \text{(IC)} \\ m\cdot S(n) &:= m\cdot n+m && \text{(RR)} \end{aligned}$$

Exponentiation

$$\begin{aligned} m^0 &:= 1 && \text{(IC)} \\ m^{S(n)} &:= m^n\cdot m && \text{(RR)} \end{aligned}$$

(Each of these operations is defined by recursion on its *second* variable.)

Laws of successor. (These should be proved first.)

(a) 0 is not a successor. Every nonzero natural number is the successor of some natural number.

For the first part, $0=\emptyset$ has no elements, while any successor has at least one element ($x\in x\cup\{x\}=S(x)$).

For the second part, the set of natural numbers that are successors of natural numbers, together with 0, namely the set

$$\{n\in\mathbb{N}\mid \exists k((k\in\mathbb{N})\wedge(n=S(k)))\}\cup\{0\},$$

is an inductive subset of $\mathbb{N}$, hence equals $\mathbb{N}$. This implies that every nonzero element $n\in\mathbb{N}$ is the successor of some element $k\in\mathbb{N}$.

(b) Successor is injective. ($S(m)=S(n)$ implies $m=n$.)

If $S(x)=S(y)$, then $x\cup\{x\}=y\cup\{y\}$. Our goal is to prove $x=y$, so let's assume that this is not the case and derive a contradiction.

We have $x\in x\cup\{x\}$, and $x\cup\{x\}=y\cup\{y\}$, so $x\in y\cup\{y\}$. We have assumed that $x\neq y$, so we must have $x\in y$. A similar argument shows that $y\in x$. This contradicts the Axiom of Foundation. (Specifically, the unordered pair $\{x,y\}$ has no $\in$-minimal element.)

Laws of addition.

(a) $S(n) = n + 1$

$$\begin{aligned} n+1 &= n + S(0) && \text{(Defn of 1)} \\ &= S(n+0) && ((\text{RR}), +) \\ &= S(n) && ((\text{IC}), +) \end{aligned}$$

(b) (Associative Law) $m + (n + k) = (m + n) + k$

We prove this by induction on k.
(Base Case: $k = 0$)

$$\begin{aligned} m + (n+0) &= m + n && ((\text{IC}), +) \\ &= (m+n) + 0 && ((\text{IC}), +) \end{aligned}$$

(Inductive Step: Assume true for k, prove true for $S(k)$)

$$\begin{aligned} m + (n + S(k)) &= m + S(n+k) && ((\text{RR}), +) \\ &= S(m + (n+k)) && ((\text{RR}), +) \\ &= S((m+n)+k) && (\text{IH}) \\ &= (m+n) + S(k) && ((\text{RR}), +) \end{aligned}$$

(c) (Unit Law for 0) $m + 0 = 0 + m = m$

The fact that $m + 0 = m$ is part of the definition of addition, so we only need to prove that $0 + m = m$. We argue this by induction on m.
(Base Case: $m = 0$)

$$0 + 0 = 0 \qquad ((\text{IC}), +)$$

(Inductive Step: Assume true for m, prove true for $S(m)$)

$$\begin{aligned} 0 + S(m) &= S(0+m) && ((\text{RR}), +) \\ &= S(m) && (\text{IH}) \end{aligned}$$

(d) (Commutative Law) $m + n = n + m$

We argue this by induction on n.
(Base Case: $n = 0$)

$$m + 0 = 0 + m \qquad \text{(Part (c), +)}$$

Before proceeding to the inductive step, we prove a lemma. It is the "$n = 1$ case" of the Commutative Law.

Lemma. $m + 1 = 1 + m$.

Proof of Lemma.
(Base Case: $m = 0$)

$$\begin{aligned}
m + 1 = 0 + 1 &= 0 + S(0) && \text{(Defn of 1)}\\
&= S(0 + 0) && \text{((RR), +)}\\
&= S(0) && \text{((IC), +)}\\
&= 1 && \text{(Defn of 1)}\\
&= 1 + 0 = 1 + m && \text{((IC), +)}
\end{aligned}$$

(Inductive Step: Assume $m + 1 = 1 + m$ for some m, prove $S(m) + 1 = 1 + S(m)$)

$$\begin{aligned}
1 + S(m) &= S(1 + m) && \text{((RR), +)}\\
&= S(m + 1) && \text{(IH)}\\
&= S(S(m)) && \text{(Part (a), } S\text{)}\\
&= S(m) + 1 && \text{(Part (a), } S\text{)}
\end{aligned}$$

Now we give the Inductive Step for the proof of (d). We assume that $m+n = n+m$ holds and derive that $m + S(n) = S(n) + m$.

$$\begin{aligned}
m + S(n) &= S(m + n) && \text{((RR), +)}\\
&= S(n + m) && \text{(IH)}\\
&= n + S(m) && \text{((RR), +)}\\
&= n + (m + 1) && \text{(Part (a), } S\text{)}\\
&= n + (1 + m) && \text{(Lemma)}\\
&= (n + 1) + m && \text{(Part (b), +)}\\
&= S(n) + m && \text{((RR), +)}
\end{aligned}$$

(e) (+-Irreducibility of 0) $m + n = 0$ implies $m = n = 0$.

If $n \neq 0$, then $n = S(k)$ by Part (a) of the Laws of Successor. Then $0 = m + n = m + S(k) = S(m + k)$, contradicting that 0 is not a successor. Hence $0 = m + n$ forces $n = 0$. But now $0 = m + n = m + 0 = m$, so $m = 0$ too.

(f) (Cancellation) $m + k = n + k$ implies $m = n$.

(Base Case: $k = 0$)

$$\begin{aligned} m &= m + 0 && ((\mathrm{IC}), +) \\ &= n + 0 && (\text{assumption}) \\ &= n && ((\mathrm{IC}), +) \end{aligned}$$

(Inductive Step: Assume that $m+k = n+k$ implies $m = n$. Prove that $m+S(k) = n + S(k)$ implies $m = n$.)

Assume that $m+S(k) = n+S(k)$. Then by $((\mathrm{RR}), +)$ we have $S(m+k) = S(n+k)$. But the successor function is injective, by Part (b) of the Laws of Successor. Thus, $m + k = n + k$. Now, by the inductive hypothesis, we derive that $m = n$.

Laws of multiplication (and addition).

(a) (Associative Law) $m \cdot (n \cdot k) = (m \cdot n) \cdot k$
(b) (Unit Law for 1) $m \cdot 1 = 1 \cdot m = m$
(c) (Commutative Law) $m \cdot n = n \cdot m$
(d) (0 is absorbing) $m \cdot 0 = 0 \cdot m = 0$
(e) ($\cdot$-Irreducibility of 1) $m \cdot n = 1$ implies $m = n = 1$
(f) (Distributive Law) $m \cdot (n + k) = (m \cdot n) + (m \cdot k)$

Laws of exponentiation (and multiplication and addition).

(a) $m^0 = 1$, $m^1 = m$, $0^m = 0$ (if $m > 0$), and $1^m = 1$.
(b) $m^{n+k} = m^n \cdot m^k$
(c) $(m \cdot n)^k = m^k \cdot n^k$
(d) $(m^n)^k = m^{n \cdot k}$