



GROUPS, ORDERS, & CLASSIFICATION

JOINT WITH

D. MARKER (UIC), L. MOTTO ROS (TORINO), AND A. SHANI (CONCORDIA)

FILIPPO CALDERONI

DEPARTMENT OF MATHEMATICS
RUTGERS UNIVERSITY

BLAST 2025 - CU BOULDER
MAY 19, 2025

Partially supported by the National Science Foundation through the grant DMS – 2348819.



The Continuum Hypothesis (CH)

For any subset $A \subseteq \mathbb{R}$ either

1. A is countable; or
2. A is equipotent with $\mathbb{R}$, i.e. $|A| = 2^{\aleph_0}$.

The Continuum Hypothesis (CH)

For any subset $A \subseteq \mathbb{R}$ either

1. A is countable; or
2. A is equipotent with $\mathbb{R}$, i.e. $|A| = 2^{\aleph_0}$.

- Gödel's constructibility axiom (**V=L**) shows that **CH** cannot be refuted. (1940)

The Continuum Hypothesis (**CH**)

For any subset $A \subseteq \mathbb{R}$ either

1. A is countable; or
2. A is equipotent with $\mathbb{R}$, i.e. $|A| = 2^{\aleph_0}$.

- Gödel's constructibility axiom (**V=L**) shows that **CH** cannot be refuted. (1940)
- Cohen's method of forcing showed that $\neg$ **CH** cannot be refuted. (1963)

The Continuum Hypothesis (**CH**)

For any subset $A \subseteq \mathbb{R}$ either

1. A is countable; or
2. A is equipotent with $\mathbb{R}$, i.e. $|A| = 2^{\aleph_0}$.

- Gödel's constructibility axiom (**V=L**) shows that **CH** cannot be refuted. (1940)
- Cohen's method of forcing showed that $\neg$ **CH** cannot be refuted. (1963)

Whether **CH** is solved or not is an endless dispute. Settling **CH** is still one of the main motivation for current research on the foundation of mathematics. (E.g., Woodin's **V= Ultimate L** program.)

Vaught's CONJECTURE

Vaught's conjecture (1961)

Let T be a first order complete theory in some countable language. Then either

1. T has at most countably many models of size $\aleph_0$ up to isomorphism; or
2. continuum many.

VAUGHT'S CONJECTURE

Vaught's conjecture (1961)

Let T be a first order complete theory in some countable language. Then either

1. T has at most countably many models of size $\aleph_0$ up to isomorphism; or
2. continuum many.

Vaught's conjecture is one of the most long-standing and elusive open problems in mathematical logic.

VAUGHT'S CONJECTURE

Vaught's conjecture (1961)

Let T be a first order complete theory in some countable language. Then either

1. T has at most countably many models of size $\aleph_0$ up to isomorphism; or
2. continuum many.

Vaught's conjecture is one of the most long-standing and elusive open problems in mathematical logic.

Even though it was verified in some special cases, Vaught's conjecture remains open.

Mathematical objects are usually organized in equivalence classes under some notion of **equivalence** (or isomorphism).

Mathematical objects are usually organized in equivalence classes under some notion of **equivalence** (or isomorphism).

Classifying a given collection of mathematical object means fully understanding the relevant equivalence relation.

Mathematical objects are usually organized in equivalence classes under some notion of **equivalence** (or isomorphism).

Classifying a given collection of mathematical object means fully understanding the relevant equivalence relation.

As the formulation of Vaught's conjecture might suggest, this can be done simply by **counting equivalence classes**.

A BLAST... FROM THE PAST

Theorem (approx. 300 B.C.)

There are exactly five convex regular polyhedra (i.e. Platonic solids).



Figure: Tetrahedron, Cube, Octahedron, Dodecahedron, Icosahedron.

CLASSIFICATIONS BEYOND VAUGHT'S CONJECTURE

Let $\mathcal{M} = (M, <, \dots)$ be a first order structure, where M is the domain, $<$ is a binary relation symbol interpreted by a **total order** on M — there may be other symbols for relations, functions, or constants.

Let $\mathcal{M} = (M, <, \dots)$ be a first order structure, where M is the domain, $<$ is a binary relation symbol interpreted by a **total order** on M – there may be other symbols for relations, functions, or constants.

Definition

The above structure $\mathcal{M}$ is **o-minimal** if every definable subset of M is a **finite union of singletons and open intervals** (with endpoints in $M \cup \{\infty, -\infty\}$).

Let $\mathcal{M} = (M, <, \dots)$ be a first order structure, where M is the domain, $<$ is a binary relation symbol interpreted by a **total order** on M – there may be other symbols for relations, functions, or constants.

Definition

The above structure $\mathcal{M}$ is **o-minimal** if every definable subset of M is a **finite union of singletons and open intervals** (with endpoints in $M \cup \{\infty, -\infty\}$).

Example

1. $(\mathbb{Q}, <)$

Let $\mathcal{M} = (M, <, \dots)$ be a first order structure, where M is the domain, $<$ is a binary relation symbol interpreted by a **total order** on M – there may be other symbols for relations, functions, or constants.

Definition

The above structure $\mathcal{M}$ is **o-minimal** if every definable subset of M is a **finite union of singletons and open intervals** (with endpoints in $M \cup \{\infty, -\infty\}$).

Example

1. $(\mathbb{Q}, <)$
2. $\mathcal{R} = (\mathbb{R}, <, +, -, \cdot, 0, 1)$ (Tarski QE)

Let $\mathcal{M} = (M, <, \dots)$ be a first order structure, where M is the domain, $<$ is a binary relation symbol interpreted by a **total order** on M – there may be other symbols for relations, functions, or constants.

Definition

The above structure $\mathcal{M}$ is **o-minimal** if every definable subset of M is a **finite union of singletons and open intervals** (with endpoints in $M \cup \{\infty, -\infty\}$).

Example

1. $(\mathbb{Q}, <)$
2. $\mathcal{R} = (\mathbb{R}, <, +, -, \cdot, 0, 1)$ (Tarski QE)
3. $(\mathcal{R}, \exp)$ (Wilkie 1996)

VAUGHT'S CONJECTURE: O-MINIMAL CASE

Definition

A theory T is **o-minimal** theory if every model of T is o-minimal.

Definition

A theory T is **o-minimal** theory if every model of T is o-minimal.

Theorem (Mayer 1988)

Let T be an o-minimal theory in a countable language. Either

- 1. T has $2^{\aleph_0}$ countable models; or*
- 2. T has exactly $6^a 3^b$ countable models, where $a, b \in \mathbb{N}$.*

Definition

A theory T is **o-minimal** theory if every model of T is o-minimal.

Theorem (Mayer 1988)

Let T be an o-minimal theory in a countable language. Either

- 1. T has $2^{\aleph_0}$ countable models; or*
- 2. T has exactly $6^a 3^b$ countable models, where $a, b \in \mathbb{N}$.*

Moreover, for all $a, b \in \mathbb{N}$ there exists an o-minimal theory T such that T has exactly $6^a 3^b$ countable models.

Definition

Let E be an equivalence relation on X . A **complete classification** for E is a map $c: X \rightarrow I$ such that for any $x, y \in X$,

$$x E y \iff c(x) = c(y).$$

Definition

Let E be an equivalence relation on X . A **complete classification** for E is a map $c: X \rightarrow I$ such that for any $x, y \in X$,

$$x E y \iff c(x) = c(y).$$

The elements of I are called **complete invariants** for E .

Definition

Let E be an equivalence relation on X . A **complete classification** for E is a map $c: X \rightarrow I$ such that for any $x, y \in X$,

$$x E y \iff c(x) = c(y).$$

The elements of I are called **complete invariants** for E .

Even in the presence of uncountably many E -classes, there might be a satisfactory complete classification theory for E .

Definition

Let E be an equivalence relation on X . A **complete classification** for E is a map $c: X \rightarrow I$ such that for any $x, y \in X$,

$$x E y \iff c(x) = c(y).$$

The elements of I are called **complete invariants** for E .

Even in the presence of uncountably many E -classes, there might be a satisfactory complete classification theory for E .

Example

Countable Abelian reduced p -groups are completely classified by their Ulm invariants.

THE DESCRIPTIVE SET THEORY VIEW

As a matter of fact, many mathematical objects (including countable first order structures) can be parametrized in a standard Borel space.

THE DESCRIPTIVE SET THEORY VIEW

As a matter of fact, many mathematical objects (including countable first order structures) can be parametrized in a standard Borel space.

Given a countable relational language $\mathcal{L} = \{R_i\}$ with R_i of arity a_i , an **enumerated $\mathcal{L}$ -structure** is an $\mathcal{L}$ -structure

$$\mathcal{M} = (\mathbb{N}, R_i^{\mathcal{M}}).$$

THE DESCRIPTIVE SET THEORY VIEW

As a matter of fact, many mathematical objects (including countable first order structures) can be parametrized in a standard Borel space.

Given a countable relational language $\mathcal{L} = \{R_i\}$ with R_i of arity a_i , an **enumerated $\mathcal{L}$ -structure** is an $\mathcal{L}$ -structure

$$\mathcal{M} = (\mathbb{N}, R_i^{\mathcal{M}}).$$

Note $R_i^{\mathcal{M}} \subseteq \mathbb{N}^{a_i}$.

THE DESCRIPTIVE SET THEORY VIEW

As a matter of fact, many mathematical objects (including countable first order structures) can be parametrized in a standard Borel space.

Given a countable relational language $\mathcal{L} = \{R_i\}$ with R_i of arity a_i , an **enumerated $\mathcal{L}$ -structure** is an $\mathcal{L}$ -structure

$$\mathcal{M} = (\mathbb{N}, R_i^{\mathcal{M}}).$$

Note $R_i^{\mathcal{M}} \subseteq \mathbb{N}^{a_i}$.

Definition

By identifying $\mathcal{M}$ with $(R_i^{\mathcal{M}} \mid i \in I)$ we can define the **standard Borel space of (countably infinite) enumerated $\mathcal{L}$ -structures** as

$$X_{\mathcal{L}} = \{\mathcal{M} \mid \mathcal{M} \text{ is an enumerated } \mathcal{L}\text{-structure}\} = \prod_{i \in I} 2^{\mathbb{N}^{a_i}}.$$

THE DESCRIPTIVE SET THEORY VIEW (CONT'D)

This trick of coding structures by (products of) “reals” extends to

- Any countable language.

THE DESCRIPTIVE SET THEORY VIEW (CONT'D)

This trick of coding structures by (products of) “reals” extends to

- Any countable language.
- Models of countable first order theories.

$$X_T = \{\mathcal{M} \in X_{\mathcal{L}} \mid \mathcal{M} \models T\}$$

THE DESCRIPTIVE SET THEORY VIEW (CONT'D)

This trick of coding structures by (products of) “reals” extends to

- Any countable language.
- Models of countable first order theories.

$$X_T = \{\mathcal{M} \in X_{\mathcal{L}} \mid \mathcal{M} \models T\}$$

- Models of countable $\mathcal{L}_{\omega_1\omega}$ -theory.

THE DESCRIPTIVE SET THEORY VIEW (CONT'D)

This trick of coding structures by (products of) “reals” extends to

- Any countable language.
- Models of countable first order theories.

$$X_T = \{\mathcal{M} \in X_{\mathcal{L}} \mid \mathcal{M} \models T\}$$

- Models of countable $\mathcal{L}_{\omega_1\omega}$ -theory.

THE DESCRIPTIVE SET THEORY VIEW (CONT'D)

This trick of coding structures by (products of) “reals” extends to

- Any countable language.
- Models of countable first order theories.

$$X_T = \{\mathcal{M} \in X_{\mathcal{L}} \mid \mathcal{M} \models T\}$$

- Models of countable $\mathcal{L}_{\omega_1\omega}$ -theory.

Theorem (Lopez-Escobar)

A subset $X \subseteq X_{\mathcal{L}}$ is Borel (in which case it admits a standard Borel structure itself) if and only if $X = X_{\varphi}$ for some $\mathcal{L}_{\omega_1\omega}$ sentence φ .

THE DESCRIPTIVE SET THEORY VIEW (CONT'D)

This trick of coding structures by (products of) “reals” extends to

- Any countable language.
- Models of countable first order theories.

$$X_T = \{\mathcal{M} \in X_{\mathcal{L}} \mid \mathcal{M} \models T\}$$

- Models of countable $\mathcal{L}_{\omega_1\omega}$ -theory.

Theorem (Lopez-Escobar)

A subset $X \subseteq X_{\mathcal{L}}$ is Borel (in which case it admits a standard Borel structure itself) if and only if $X = X_{\varphi}$ for some $\mathcal{L}_{\omega_1\omega}$ sentence φ .

Church's thesis for real mathematics

Borel = explicit

BACK TO CLASSIFICATIONS (CONT'D)

To perform a finer analysis of the classification problem for $\cong_T$, descriptive set theorist has developed the theory of **Borel classification** (aka **Borel cardinalities**).

BACK TO CLASSIFICATIONS (CONT'D)

To perform a finer analysis of the classification problem for $\cong_T$, descriptive set theorist has developed the theory of **Borel classification** (aka **Borel cardinalities**).

Definition

Let E, F be Borel equivalence relations on the Polish spaces X, Y respectively.

BACK TO CLASSIFICATIONS (CONT'D)

To perform a finer analysis of the classification problem for $\cong_T$, descriptive set theorist has developed the theory of **Borel classification** (aka **Borel cardinalities**).

Definition

Let E, F be Borel equivalence relations on the Polish spaces X, Y respectively. We say that:

- E is **Borel reducible** to F (in symbols $E \leq_B F$) if and only if there exists a Borel map $f: X \rightarrow Y$ such that

$$x E y \iff f(x) F f(y).$$

In this case, f is called a **Borel reduction** from E to F .

BACK TO CLASSIFICATIONS (CONT'D)

To perform a finer analysis of the classification problem for $\cong_T$, descriptive set theorist has developed the theory of **Borel classification** (aka **Borel cardinalities**).

Definition

Let E, F be Borel equivalence relations on the Polish spaces X, Y respectively. We say that:

- E is **Borel reducible** to F (in symbols $E \leq_B F$) if and only if there exists a Borel map $f: X \rightarrow Y$ such that

$$x E y \iff f(x) F f(y).$$

In this case, f is called a **Borel reduction** from E to F .

- $E \sim_B F$ if and only if $E \leq_B F$ and $F \leq_B E$.

THE FRIEDMAN-STANLEY JUMP HIERARCHY

Definition (Friedman–Stanley 1989)

Let E be an equivalence relation on a standard Borel space X .
For $x = (x_i : i \in \mathbb{N})$ and $y = (y_i : i \in \mathbb{N})$ in $X^{\mathbb{N}}$ let

$$x E^+ y \iff \{[x_i]_E : i \in \mathbb{N}\} = \{[y_i]_E : i \in \mathbb{N}\}.$$

THE FRIEDMAN-STANLEY JUMP HIERARCHY

Definition (Friedman–Stanley 1989)

Let E be an equivalence relation on a standard Borel space X .

For $x = (x_i : i \in \mathbb{N})$ and $y = (y_i : i \in \mathbb{N})$ in $X^{\mathbb{N}}$ let

$$x E^+ y \iff \{[x_i]_E : i \in \mathbb{N}\} = \{[y_i]_E : i \in \mathbb{N}\}.$$

Starting from equality $=_{\mathbb{R}}$ on the real numbers:

THE FRIEDMAN-STANLEY JUMP HIERARCHY

Definition (Friedman–Stanley 1989)

Let E be an equivalence relation on a standard Borel space X .
For $x = (x_i : i \in \mathbb{N})$ and $y = (y_i : i \in \mathbb{N})$ in $X^{\mathbb{N}}$ let

$$x E^+ y \iff \{[x_i]_E : i \in \mathbb{N}\} = \{[y_i]_E : i \in \mathbb{N}\}.$$

Starting from equality $=_{\mathbb{R}}$ on the real numbers:

1. The first jump $=_{\mathbb{R}}^+$ is defined on $\mathbb{R}^{\mathbb{N}}$ so that the map

$$(x_i : i \in \mathbb{N}) \mapsto \{x_i : i \in \mathbb{N}\} \in \mathcal{P}(\mathbb{R}).$$

is a complete classification of $=_{\mathbb{R}}^+$ by **countable sets** of reals.

THE FRIEDMAN-STANLEY JUMP HIERARCHY

Definition (Friedman–Stanley 1989)

Let E be an equivalence relation on a standard Borel space X .

For $x = (x_i : i \in \mathbb{N})$ and $y = (y_i : i \in \mathbb{N})$ in $X^{\mathbb{N}}$ let

$$x E^+ y \iff \{[x_i]_E : i \in \mathbb{N}\} = \{[y_i]_E : i \in \mathbb{N}\}.$$

Starting from equality $=_{\mathbb{R}}$ on the real numbers:

1. The first jump $=_{\mathbb{R}}^+$ is defined on $\mathbb{R}^{\mathbb{N}}$ so that the map

$$(x_i : i \in \mathbb{N}) \mapsto \{x_i : i \in \mathbb{N}\} \in \mathcal{P}(\mathbb{R}).$$

is a complete classification of $=_{\mathbb{R}}^+$ by **countable sets** of reals.

2. The second jump $=_{\mathbb{R}}^{++}$ is defined on $(\mathbb{R}^{\mathbb{N}})^{\mathbb{N}}$ and admits a complete classification by **hereditarily countable set** in $\mathcal{P}(\mathcal{P}(\mathbb{R}))$.

THE FRIEDMAN-STANLEY JUMP HIERARCHY

We can iterate the jump operator transfinitely and have a proper hierarchy.

THE FRIEDMAN-STANLEY JUMP HIERARCHY

We can iterate the jump operator transfinitely and have a proper hierarchy.

$$\begin{aligned} =_{\mathbb{R}}^{(\alpha+1)+} &:= (=_{\mathbb{R}}^{\alpha+})^+ \\ =_{\mathbb{R}}^{\lambda+} &:= \prod_{\alpha < \lambda} =_{\mathbb{R}}^{\alpha+} \text{ for } \lambda \text{ limit.} \end{aligned}$$

THE FRIEDMAN-STANLEY JUMP HIERARCHY

We can iterate the jump operator transfinitely and have a proper hierarchy.

$$\begin{aligned} =_{\mathbb{R}}^{(\alpha+1)+} &:= (=_\mathbb{R}^{\alpha+})^+ \\ =_{\mathbb{R}}^{\lambda+} &:= \prod_{\alpha < \lambda} =_{\mathbb{R}}^{\alpha+} \text{ for } \lambda \text{ limit.} \end{aligned}$$

Proposition

- For countable $\alpha < \beta$, we have $=_{\mathbb{R}}^{\alpha+} <_B =_{\mathbb{R}}^{\beta+}$

THE FRIEDMAN-STANLEY JUMP HIERARCHY

We can iterate the jump operator transfinitely and have a proper hierarchy.

$$\begin{aligned} =_{\mathbb{R}}^{(\alpha+1)+} &:= (=_{\mathbb{R}}^{\alpha+})^+ \\ =_{\mathbb{R}}^{\lambda+} &:= \prod_{\alpha < \lambda} =_{\mathbb{R}}^{\alpha+} \text{ for } \lambda \text{ limit.} \end{aligned}$$

Proposition

- For countable $\alpha < \beta$, we have $=_{\mathbb{R}}^{\alpha+} <_B =_{\mathbb{R}}^{\beta+}$
- Every Borel isomorphism relation is Borel reducible to $=_{\mathbb{R}}^{\alpha+}$ for some $\alpha < \omega_1$.

Theorem (Rast–Sahota '17)

Let T be a complete o-minimal theory in a countable language.

1. $\cong_T \leq_B =_{\mathbb{R}}$.
2. $\cong_T \sim_B =_{\mathbb{R}}^+$.
3. $\cong_T$ is not Borel (in fact, $\cong_T$ is maximal among isomorphisms).

ORDERED DIVISIBLE ABELIAN GROUPS

One of theories of maximal complexity is ODAG — the theory of **ordered divisible abelian groups**.

ORDERED DIVISIBLE ABELIAN GROUPS

One of theories of maximal complexity is ODAG — the theory of **ordered divisible abelian groups**.

It suffices to show $\cong_{\text{LO}} \leq_B \cong_{\text{ODAG}}$.

ORDERED DIVISIBLE ABELIAN GROUPS

One of theories of maximal complexity is ODAG — the theory of **ordered divisible abelian groups**.

It suffices to show $\cong_{\text{LO}} \leq_B \cong_{\text{ODAG}}$.

Given a linear order $\mathbf{L} = (L, <_L)$ consider the group $\mathbf{G}_{\mathbf{L}} = (G_L, <_{lex})$ where

- $G_L = \{f : L \rightarrow \mathbb{Q} \mid \text{supp}(f) \text{ is finite}\}$
- $<_{lex}$ is the reverse lexicographic order

ORDERED DIVISIBLE ABELIAN GROUPS

One of theories of maximal complexity is ODAG — the theory of **ordered divisible abelian groups**.

It suffices to show $\cong_{\text{LO}} \leq_B \cong_{\text{ODAG}}$.

Given a linear order $\mathbf{L} = (L, <_L)$ consider the group $\mathbf{G}_{\mathbf{L}} = (G_L, <_{lex})$ where

- $G_L = \{f : L \rightarrow \mathbb{Q} \mid \text{supp}(f) \text{ is finite}\}$
- $<_{lex}$ is the reverse lexicographic order

For $x, y \in \mathbf{G}_{\mathbf{L}}$ we define

$$x \preceq y \iff \exists n \in \mathbb{Z} \setminus \{0\} (x \leq_{lex} ny)$$

$$x \approx y \iff x \preceq y \text{ and } y \preceq x$$

ORDERED DIVISIBLE ABELIAN GROUPS

One of theories of maximal complexity is ODAG — the theory of **ordered divisible abelian groups**.

It suffices to show $\cong_{\text{LO}} \leq_B \cong_{\text{ODAG}}$.

Given a linear order $\mathbf{L} = (L, <_L)$ consider the group $\mathbf{G}_{\mathbf{L}} = (G_L, <_{lex})$ where

- $G_L = \{f : L \rightarrow \mathbb{Q} \mid \text{supp}(f) \text{ is finite}\}$
- $<_{lex}$ is the reverse lexicographic order

For $x, y \in \mathbf{G}_{\mathbf{L}}$ we define

$$x \preceq y \iff \exists n \in \mathbb{Z} \setminus \{0\} (x \leq_{lex} ny)$$

$$x \approx y \iff x \preceq y \text{ and } y \preceq x$$

The $\approx$ -classes are maximal Archimedean subgroups. From the quotient set $\mathbf{G}_{\mathbf{L}} / \approx$ it is possible to recover the order $\mathbf{L}$, so that

$$\mathbf{G}_{\mathbf{L}} \cong_{\text{ODAG}} \mathbf{G}_{\mathbf{K}} \implies \mathbf{L} \cong_{\text{LO}} \mathbf{K}.$$

A (left-)order on a group $(G, \cdot)$ is a total order on G such that

$$g < h \implies fg < fh.$$

A (left-)order on a group $(G, \cdot)$ is a total order on G such that

$$g < h \implies fg < fh.$$

Definition

A (left-)order $<$ on G is **Archimedean** iff for all positive $g, h \in G$ there is $n \in \mathbb{N}$ such that $g < h^n$.

A (left-)order on a group $(G, \cdot)$ is a total order on G such that

$$g < h \implies fg < fh.$$

Definition

A (left-)order $<$ on G is **Archimedean** iff for all positive $g, h \in G$ there is $n \in \mathbb{N}$ such that $g < h^n$.

The argument showing $\cong_{\text{LO}} \leq_B \cong_{\text{ODAG}}$ uses non-Archimedean groups.

A (left-)order on a group $(G, \cdot)$ is a total order on G such that

$$g < h \implies fg < fh.$$

Definition

A (left-)order $<$ on G is **Archimedean** iff for all positive $g, h \in G$ there is $n \in \mathbb{N}$ such that $g < h^n$.

The argument showing $\cong_{\text{LO}} \leq_B \cong_{\text{ODAG}}$ uses non-Archimedean groups.

What is the Borel complexity of the isomorphism relation $\cong_{\text{ArGp}}$ for countable ordered Archimedean groups?

Definition

A **(left-)order** $<$ on G is **Archimedean** iff for all positive $g, h \in G$ there is $n \in \mathbb{N}$ such that $g < h^n$.

Definition

A **(left-)order** $<$ on G is **Archimedean** iff for all positive $g, h \in G$ there is $n \in \mathbb{N}$ such that $g < h^n$.

Theorem (Hölder 1901)

When G is countable, the following are equivalent:

Definition

A **(left-)order** $<$ on G is **Archimedean** iff for all positive $g, h \in G$ there is $n \in \mathbb{N}$ such that $g < h^n$.

Theorem (Hölder 1901)

When G is countable, the following are equivalent:

- *G has an Archimedean order.*

Definition

A **(left-)order** $<$ on G is **Archimedean** iff for all positive $g, h \in G$ there is $n \in \mathbb{N}$ such that $g < h^n$.

Theorem (Hölder 1901)

When G is countable, the following are equivalent:

- *G has an Archimedean order.*
- *G acts freely on $\mathbb{R}$ by orientation preserving.*

Definition

A **(left-)order** $<$ on G is **Archimedean** iff for all positive $g, h \in G$ there is $n \in \mathbb{N}$ such that $g < h^n$.

Theorem (Hölder 1901)

When G is countable, the following are equivalent:

- *G has an Archimedean order.*
- *G acts freely on $\mathbb{R}$ by orientation preserving.*
- *G is isomorphic to a subgroup of $(\mathbb{R}, +)$ equipped with the natural ordering on $\mathbb{R}$.*

Rather than working on the space of enumerated models X_{ArGp} , it is more convenient to work on the standard Borel space of countable subgroups of $\mathcal{A}$ and the isomorphism relation $\cong_{\mathcal{A}}$ defined there.

HION'S LEMMA

Rather than working on the space of enumerated models X_{ArGp} , it is more convenient to work on the standard Borel space of countable subgroups of $\mathcal{A}$ and the isomorphism relation $\cong_{\mathcal{A}}$ defined there.

Theorem (C.-Marker–Motto Ros–Shani 2023)

There is a continuous map $X_{\text{ArGp}} \rightarrow \mathcal{A}$ showing $\cong_{\text{ArGp}} \leq_B \cong_{\mathcal{A}}$.

HION'S LEMMA

Rather than working on the space of enumerated models X_{ArGp} , it is more convenient to work on the standard Borel space of countable subgroups of $\mathcal{A}$ and the isomorphism relation $\cong_{\mathcal{A}}$ defined there.

Theorem (C.-Marker–Motto Ros–Shani 2023)

There is a continuous map $X_{\text{ArGp}} \rightarrow \mathcal{A}$ showing $\cong_{\text{ArGp}} \leq_B \cong_{\mathcal{A}}$.

Lemma (Hion 54')

Suppose that A and B are two subgroups of $\mathbb{R}$ and $h: A \rightarrow B$ is an order preserving homomorphism. Then, there exists $\lambda \in \mathbb{R}^+$ such that $h(a) = \lambda a$, for every $a \in A$.

HION'S LEMMA

Rather than working on the space of enumerated models X_{ArGp} , it is more convenient to work on the standard Borel space of countable subgroups of $\mathcal{A}$ and the isomorphism relation $\cong_{\mathcal{A}}$ defined there.

Theorem (C.-Marker–Motto Ros–Shani 2023)

There is a continuous map $X_{\text{ArGp}} \rightarrow \mathcal{A}$ showing $\cong_{\text{ArGp}} \leq_B \cong_{\mathcal{A}}$.

Lemma (Hion 54')

Suppose that A and B are two subgroups of $\mathbb{R}$ and $h: A \rightarrow B$ is an order preserving homomorphism. Then, there exists $\lambda \in \mathbb{R}^+$ such that $h(a) = \lambda a$, for every $a \in A$.

In fact, such λ is computed as the ratio $\frac{h(a)}{a}$, for any nonzero $a \in A$.

If G is a nontrivial subgroup of $\mathbb{R}$ let

$$A_G := \left\{ \underbrace{\left\{ \frac{g}{r} : g \in G \right\}}_{G/r} : r \in G \setminus \{0\} \right\}.$$

AN UPPER BOUND

If G is a nontrivial subgroup of $\mathbb{R}$ let

$$A_G := \left\{ \underbrace{\left\{ \frac{g}{r} : g \in G \right\}}_{G/r} : r \in G \setminus \{0\} \right\}.$$

Proposition

Let G and H be non-trivial subgroups of $\mathbb{R}$. Then

$$G \text{ and } H \text{ are order isomorphic} \iff A_G = A_H.$$

Therefore, $\cong_{\text{ArGp}} \leq_B =_{\mathbb{R}}^{++}$.

AN UPPER BOUND

If G is a nontrivial subgroup of $\mathbb{R}$ let

$$A_G := \left\{ \underbrace{\left\{ \frac{g}{r} : g \in G \right\}}_{G/r} : r \in G \setminus \{0\} \right\}.$$

Proposition

Let G and H be non-trivial subgroups of $\mathbb{R}$. Then

$$G \text{ and } H \text{ are order isomorphic} \iff A_G = A_H.$$

Therefore, $\cong_{\text{ArGp}} \leq_B =_{\mathbb{R}}^{++}$.

Note that for $r \neq s$, the sets G/r and G/s are not at odds with each other.

(Hjorth–Kechris–Louveau 1998) defined a refinement of the Friedman–Satanley hierarchy. In particular,

$$=^+ <_B \cong_{3,0}^* <_B \cong_{3,1}^* <_B =^{++}$$

(Hjorth–Kechris–Louveau 1998) defined a refinement of the Friedman–Satanley hierarchy. In particular,

$$=^+ <_B \cong_{3,0}^* <_B \cong_{3,1}^* <_B =^{++}$$

An invariant for $\cong_{3,1}^*$ is a hereditarily countable set $A \in \mathcal{P}_2(\mathbb{R})$ (i.e., a $=^{++}$ -invariant) together with

- a ternary relation $R \subseteq A \times A \times \mathbb{R}$, definable from A , such that given any $a \in A$, $R(a, -, -)$ is an injective function from A to $\mathbb{R}$.

Theorem (C.–Marker–Motto Ros–Shani 2023)

$$=^+ <_B \cong_{\text{ArGp}} <_B =^{++}$$

Theorem (C.–Marker–Motto Ros–Shani 2023)

$$=^+ <_B \cong_{\text{ArGp}} <_B =^{++}$$

In fact,

- $\cong_{\text{ArGp}} \leq_B \cong_{3,1}^*$;
- $\cong_{\text{ArGp}} \not\leq_B \cong_{3,0}^*$.

Theorem (C.–Marker–Motto Ros–Shani 2023)

$$=^+ <_B \cong_{\text{ArGp}} <_B =^{++}$$

In fact,

- $\cong_{\text{ArGp}} \leq_B \cong_{3,1}^*$;
- $\cong_{\text{ArGp}} \not\leq_B \cong_{3,0}^*$.

We cannot use countable sets of reals to classify $\cong_{\text{ArGp}}$.

A **circular order** on a set X is defined by a **cyclic orientation cocycle**, i.e., a function $c: X^3 \rightarrow \{\pm 1, 0\}$ satisfying:

1. $c^{-1}(0) = \Delta(X)$, where
$$\Delta(X) := \{(x_1, x_2, x_3) \in X^3 : x_i = x_j, \text{ for some } i \neq j\},$$
2. $c(x_2, x_3, x_4) - c(x_1, x_3, x_4) + c(x_1, x_2, x_4) - c(x_1, x_2, x_3) = 0$ for all $x_1, x_2, x_3, x_4 \in X$.

A **circular order** on a set X is defined by a **cyclic orientation cocycle**, i.e., a function $c: X^3 \rightarrow \{\pm 1, 0\}$ satisfying:

1. $c^{-1}(0) = \Delta(X)$, where
$$\Delta(X) := \{(x_1, x_2, x_3) \in X^3 : x_i = x_j, \text{ for some } i \neq j\},$$
2. $c(x_2, x_3, x_4) - c(x_1, x_3, x_4) + c(x_1, x_2, x_4) - c(x_1, x_2, x_3) = 0$ for all $x_1, x_2, x_3, x_4 \in X$.

Definition

A group G is **circularly orderable** if it admits a circular order c which is left-invariant in the sense that $c(g_1, g_2, g_3) = c(hg_1, hg_2, hg_3)$ for all $g_1, g_2, g_3, h \in G$.

Definition

A circularly ordered group (G, c) is said to be **Archimedean** if there are no elements $g, h \in G$ such that $c(1_G, g^n, h) = 1$ for all $n \geq 1$.

Definition

A circularly ordered group (G, c) is said to be **Archimedean** if there are no elements $g, h \in G$ such that $c(1_G, g^n, h) = 1$ for all $n \geq 1$.

A well-known example of Archimedean circularly ordered group is S^1 with the obvious circular order.

There is a classical functorial construction that starting from a circular order group (G, c) produces a corresponding ordered group $(\tilde{G}, <)$ called **central extension**. (Želeva)

Definition

A circularly ordered group (G, c) is said to be **Archimedean** if there are no elements $g, h \in G$ such that $c(1_G, g^n, h) = 1$ for all $n \geq 1$.

A well-known example of Archimedean circularly ordered group is S^1 with the obvious circular order.

There is a classical functorial construction that starting from a circular order group (G, c) produces a corresponding ordered group $(\tilde{G}, <)$ called **central extension**. (Želeva)

This can be used to show that

$$\cong_{\text{ArCO}} \leq_B \cong_{\text{ArGp}}.$$

We also showed the following:

Theorem (C.–Marker–Motto Ros–Shani 2023)

$$\cong_{\text{ArCO}} \leq_B =_{\mathbb{R}}^+.$$

We also showed the following:

Theorem (C.–Marker–Motto Ros–Shani 2023)

$$\cong_{\text{ArCO}} \leq_B =_{\mathbb{R}}^+.$$

There is NO Borel (explicit) converse to the central extension functor because $\cong_{\text{ArCO}} <_B \cong_{\text{ArGp}}$.

The proof of $\cong_{\text{ArGp}} \not\leq_B \cong_{3,0}^*$ is highly nonelementary and requires the analysis of the complete invariants of $\cong_{\text{ArGp}}$ and $\cong_{3,0}^*$ in choiceless models of set theory.

The proof of $\cong_{\text{ArGp}} \not\leq_B \cong_{3,0}^*$ is highly nonelementary and requires the analysis of the complete invariants of $\cong_{\text{ArGp}}$ and $\cong_{3,0}^*$ in choiceless models of set theory.

Definition

Suppose A is a set in some generic extension of V . Let $V(A)$ be the **minimal transitive model** of **ZF** containing V and A , denoted by $V(A)$.

The proof of $\cong_{\text{ArGp}} \not\leq_B \cong_{3,0}^*$ is highly nonelementary and requires the analysis of the complete invariants of $\cong_{\text{ArGp}}$ and $\cong_{3,0}^*$ in choiceless models of set theory.

Definition

Suppose A is a set in some generic extension of V . Let $V(A)$ be the **minimal transitive model** of **ZF** containing V and A , denoted by $V(A)$.

For any set X , there is some formula ψ , parameters $\bar{a} \in tc(A)$ and $v \in V$ such that X is the unique set satisfying $\psi(X, A, \bar{a}, v)$.

We use the following powerful technique:

Theorem (essentially Shani 2018)

Suppose E is a Borel equivalence relation on a standard Borel space X , and $x \mapsto A_x$ is an absolute classification of E by hereditarily countable sets.

INTERDEFINABLE INVARIANTS

We use the following powerful technique:

Theorem (essentially Shani 2018)

Suppose E is a Borel equivalence relation on a standard Borel space X , and $x \mapsto A_x$ is an absolute classification of E by hereditarily countable sets. Let x be an element of X in some generic extension of V .

INTERDEFINABLE INVARIANTS

We use the following powerful technique:

Theorem (essentially Shani 2018)

Suppose E is a Borel equivalence relation on a standard Borel space X , and $x \mapsto A_x$ is an absolute classification of E by hereditarily countable sets.

Let x be an element of X in some generic extension of V .

If $E \leq_B \cong_{3,0}^$, then there is a set of sets of reals $B \in V(A_x)$ so that:*

- *B is definable from A_x and parameters in V ,*
- *B is countable in $V(A_x)$,*
- *$V(A_x) = V(B)$.*

INTERDEFINABLE INVARIANTS

We use the following powerful technique:

Theorem (essentially Shani 2018)

Suppose E is a Borel equivalence relation on a standard Borel space X , and $x \mapsto A_x$ is an absolute classification of E by hereditarily countable sets.

Let x be an element of X in some generic extension of V .

If $E \leq_B \cong_{3,0}^$, then there is a set of sets of reals $B \in V(A_x)$ so that:*

- *B is definable from A_x and parameters in V ,*
- *B is countable in $V(A_x)$,*
- *$V(A_x) = V(B)$.*

Proof of the main theorem: Over the Cohen model, we force the existence of a generic subgroup G of $\mathbb{R}$ so that every set of reals $B \in V(A_G)$ which is definable from A_G and parameters in V alone, we have $V(B) \neq V(A_G)$.

THANK YOU!