Relation algebras of Sugihara, Belnap, Meyer, Church. Relevance logic of Tarski.

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Koelbel 210, 4:30 – 4:55 pm, May 21, University of Colorado, Boulder CO, BLAST 2019 (<u>B</u>oolean <u>A</u>lgebras, <u>L</u>attices, <u>A</u>lgebraic <u>L</u>ogic and Quantum <u>L</u>ogic, Universal <u>A</u>lgebra, <u>Set T</u>heory, <u>S</u>et-theoretic and Point-free <u>T</u>opology), May 20 – 24, 2019 (Monday – Friday). 1. Historical association: Names — Topic

Sugihara, Belnap, Meyer, Church — Relevance Logic Tarski — Relation Algebras

- 2. Relevance logic algebras of Sugihara, Belnap, Meyer, Church turn out to be definitional reducts of proper relation algebras
- 3. The provable-with-four-variables characterization of Tarski's equational theory of relation algebras yields a new relevance logic
- 4. By 2 and 3, the association can be inverted—

Belnap, Meyer, Sugihara, Church — Proper relation algebras Tarski — A relevance logic To get a relation algebra (RA)—

- 1. Start with a Boolean algebra $\langle B,+,\cdot,^{-}\!\!\!,0,1\rangle$
- 2. Add a binary operation
- 3. Add a unary operation
- 4. Add a constant 1'
- 5. Insist that Tarski's Ten Axioms [1987] hold

 $\langle B, +, \cdot, \bar{}, 0, 1, ;, \check{}, 1' \rangle \models (\mathsf{Ra} \mathsf{I}), \cdots, (\mathsf{Ra} \mathsf{X})$

6. The associative law (Ra IV) x;(y;z) = (x;y);z

To get a proper relation algebra (PRA)—

- 1. Start with a set U, or an equivalence relation E
- 2. Use all binary relations on U, or all subsets of E
- 3. Add union \cup and intersection \cap
- 4. Add complementation $\overline{}$ (with respect to $U \times U$ or E)
- 5. Add |, composition R|S of relations R, S
- 6. Add $^{-1}$, converse R^{-1} of R
- 7. Add constants $U \times U$ or E, \emptyset , the identity relation on U $\operatorname{Id}_U = \{ \langle u, u \rangle : u \in U \}$ or the identity part of E $\{ \langle u, u \rangle : \langle u, u \rangle \in E \}$
- 8. Take a subalgebra
- 9. Representable RA means isomorphic to a PRA

Characterizations of Equations True in PRA, RA [1987]

- 1. (Tarski 1940s) Every RA equation translates into a sentence of first-order logic (of binary relations) using at most 3 variables, and conversely
- 2. (Tarski 1940s) The associative law for | requires 4 variables to prove, but the other axioms only need 3
- 3. (Tarski 1940s) An equation is valid in every PRA iff its translation can be proved in first-order logic
- [1978] An equation is valid in every relation algebra iff it can be proved in first-order logic restricted to 4 variables — RA=provable-with-four-variables motivates a definition:
- [3/2019] A formula (using ∧, ∨, →, ~) is in Tarski's relevance logic if it can be proved in first-order logic restricted to 4 variables

To get a definitional subreduct —

- 1. Start with an algebra (a set with some operations)
- 2. Add some defined operations
- 3. Delete some operations
- 4. Take a subalgebra with respect to the remaining operations

To represent algebras of relevance logic with binary relations by the First Method

1. Start with a proper relation algebra of relations on U

$$\begin{array}{l} \left\langle \wp(U \times U), \cup, \cap, \bar{}, \emptyset, U \times U, |, \bar{}^{-1}, \operatorname{Id}_{U} \right\rangle \\ R|S = \left\{ \langle x, z \rangle : \exists y (\langle x, y \rangle \in R \land \langle y, z \rangle \in S) \right\} \\ R^{-1} = \left\{ \langle y, x \rangle : \langle x, y \rangle \in R \right\} \\ \operatorname{Id}_{U} = \left\{ \langle u, u \rangle : u \in U \right\} \end{array}$$

- 2. Add $R \rightarrow S = \overline{R^{-1}|\overline{S}} \qquad \sim R = \overline{R^{-1}}$ 3. Delete $\emptyset \quad U \times U \quad | \quad ^{-1} \quad \text{Id } U$
- 4. Take subalgebra wrt $\ \cup \ \cap \ o \ \sim$

To represent algebras of relevance logic with binary relations by the Second Method, relativize to diversity

1. Start with a proper relation algebra of relations on U

$$\left<\wp(U imes U),\,\cup,\,\cap,\,\bar{},\,\emptyset,\,U imes U,|,^{-1},\,\operatorname{\mathsf{Id}}_U
ight>$$

- 2. Add $R \to S = \overline{R^{-1}|\overline{S}} \cap \overline{\operatorname{Id}}_U \quad \sim R = \overline{R^{-1}} \cap \overline{\operatorname{Id}}_U$
- 3. Delete $\begin{array}{ccccc} & & \emptyset & U \times U & | & -1 & \operatorname{Id}_U \end{array}$
- 4. Take subalgebra wrt $\ \cup \ \cap \ o'$

Algebras of Relevance Logic are Definitional Subreducts of PRAs [2007][2010][1/2019]

Examples of algebras used in relevance logic that turn out to be definitional subreducts of proper relation algebras—

- Belnap [1960][1975][1982]: the M₀ matrices (8-element algebra)
- 2. Meyer [1982][2003]: the crystal matrices (6-element algebra)
- 3. Church [1982][2003]: the diamond matrices (4-element algebra)
- 4. Meyer [1975][1982]: the RM84 matrices (8-element algebra)
- 5. Sugihara [1955][1975]: the odd and even matrices (algebras of every finite size, two infinite ones) complete for **RM**

Belnap's Relation Algebra [2007][2010]

- 1. Start with the proper relation algebra of all relations on $U = \mathbb{Q}$, the set of rational numbers
- Belnap's relation algebra 𝔅_{{0}} is the subalgebra whose elements are these binary relations on ℚ:
 - \emptyset = < > \leq \geq \neq $\mathbb{Q} \times \mathbb{Q}$
- 3. Belnap's RA $\mathfrak{S}_{\{0\}}$ has been called The Point Algebra since the 1980s because its elements are all the ways two points on the rational line can be related to each other
- The algebra called Belnap's M₀ matrices is a definitional reduct of Belnap's RA 𝔅_{0} by the First Method

- 1. Start with the proper relation algebra of all relations on $U = \mathbb{Q}^2$, the set of pairs of rational numbers
- 2. Meyer's crystal relation algebra $\mathfrak{S}_{\{0,1\}}$ is the subalgebra generated by those relations whose images of $\langle 0, 0 \rangle$ are the positive and negative *x*-axes, and the upper and lower half planes (with *x*-axis excluded)
- 3. Meyer's crystal matrices occur as a 6-element definitional subreduct of Meyer's RA $\mathfrak{S}_{\{0,1\}}$ using the Second Method
- 4. Meyer's RA $\mathfrak{S}_{\{0,1\}}$ is isomorphic to RA 2_{83} in [2006]

Church's Diamond Relation Algebra [1/2019]

- 1. Start with the proper relation algebra of all relations on a set U with 6 elements
- 2. Church's diamond relation algebra 27 is the subalgebra generated by the equivalence relation with exactly two equivalence classes, three elements in each class
- The algebra called Church's diamond is a 4-element definitional subreduct of Church's RA 27 by the Second Method
- 4. Church's RA 2₇ has size 8, and is representable on sets of size 6 or more; see [2006]

Meyer's RM84 Relation Algebra [2010]

- 1. Start with the proper relation algebra of all relations on $U = \{0, 1, 2, 3, 4, 5, 6\}$
- 2. Meyer's RM84 relation algebra 3_3 is the subalgebra generated by the relations $\{\langle n, n+i \mod 7 \rangle : n \in U, i = 1, 2, 4\}$ and $\{\langle n, n+i \mod 7 \rangle : n \in U, i = 3, 5, 6\}$
- 3. The algebra called Meyer's RM84 in [1982] (unnamed in [1975]) is a definitional reduct of Meyer's RA 3_3 by the First Method
- 4. Meyer's RA 3_3 has size 8 and is representable ...
 - ... on sets with 7 elements (uniquely),
 - ... on every set with more than 8 elements (not uniquely),
 - ... but not on 8 elements; see [2006]

Sugihara's Relation Algebras [1/2019]

 Start with the proper relation algebra of all relations on U_I, for any I ⊆ Z, the Z-indexed sequences of rationals that are zero outside I and eventually always zero:

$$U_I = \{q \in {}^{\mathbb{Z}}\mathbb{Q} : (orall i \in \mathbb{Z} \setminus I)(q_i = 0) \land (\exists n \in \mathbb{Z})(orall i > n)(q_i = 0)\}$$

2. Sugihara's relation algebra \mathfrak{S}_I is the subalgebra completely generated (use arbitrary unions) by the relations L_n^I , $n \in I$, "less-than" at n (the last place they differ)

$$L_n^I = \{ \langle q, r \rangle : q, r \in U_I, q_n < r_n \text{ and } q_i = r_i \text{ for all } i > n \}$$

- 3. Examples: Belnap's RA $\mathfrak{S}_{\{0\}}$ and Meyer's RA $\mathfrak{S}_{\{0,1\}}$
- 4. Defined for finite *I* in [2010]

- The 2- and 4-element Sugihara algebras are definitional subreducts of Belnap's RA 𝔅_{0} by the First Method [2007][2010]
- 2. Every Sugihara algebra of even cardinality is a definitional subreduct of Sugihara's RA $\mathfrak{S}_{\mathbb{Z}}$ by the First Method
- 3. The algebra originally defined in Sugihara [1955] is a definitional subreduct of Sugihara's RA $\mathfrak{S}_{\mathbb{Z}}$ by the First Method

- 1. The Sugihara matrices of odd cardinality have an element fixed by \sim (it is its own negation)
- The 1- and 3-element Sugihara algebras are definitional subreducts of Belnap's RA 𝔅_{0} by the Second Method
- 3. The 5-element Sugihara algebra is a definitional subreduct of Meyer's RA $\mathfrak{S}_{\{0,1\}}$ by the Second Method
- 4. Every Sugihara algebra of odd cardinality is a definitional subreduct of Sugihara's RA $\mathfrak{S}_{\mathbb{Z}}$ by the Second Method

Significance of Sugihara matrices/algebras —

- 1. Every Sugihara algebra is a linearly ordered chain, and its designated values are the top half
- When a Sugihara algebra is a subreduct of a PRA, the designated values are the relations that contain the identity relation [2010][1/2019]
- 3. (R.K.Meyer) the Sugihara matrices are characteristic for the Dunn-McCall system of logic called **RM** or **R**-mingle, i.e.,
- 4. (R.K.Meyer) A formula is a theorem of **RM** iff it is mapped to a designated value by every homomorphism from the formula algebra into a Sugihara algebra

A class of algebras of binary relations, characteristic for RM-

Definition

- 1. ${\it K}_{\rm RM}$ is the class of algebras $\mathfrak{K}=\langle {\it K},\cup,\cap,\rightarrow,\sim\rangle$, where
 - 1.1 K is a set of dense and transitive binary relations on a set U
 - 1.2 K is closed under \cup , \cap , \rightarrow , and \sim
 - 1.3 K is commutative under |
- 2. A formula is valid in $K_{\rm RM}$ if it is sent to a relation containing the identity relation by every homomorphism from the algebra of formulas into an algebra in $K_{\rm RM}$

Theorem (K_{RM} characterizes **RM**) **RM** is the set of formulas valid in K_{RM}

Classification of Axioms and Rules of RM [1/2019]

Axioms of \mathbf{RM} that are valid for all binary relations A, B, and C

$$A \to A$$
 (A1)

$$A \wedge B \to A$$
 (A5)

$$A \wedge B \to B$$
 (A6)

$$(A \to B) \land (A \to C) \to (A \to B \land C) \tag{A7}$$

$$A \to A \lor B$$
 (A8)

$$B \to A \lor B$$
 (A9)

$$(A \to C) \land (B \to C) \to (A \lor B \to C)$$
 (A10)

$$A \land (B \lor C) \to (A \land B) \lor C \tag{A11}$$

$$\sim \sim A \to A$$
 (A13)

$$A, B \vdash A \land B \tag{R1}$$

$$A \to B, A \vdash B$$
 (R2)

Classification of Axioms of RM [1/2019]

Axioms that are valid...

1. ... if all relations commute under |:

$$(A \to B) \to ((B \to C) \to (A \to C))$$
 (A2)

$$A \to ((A \to B) \to B)$$
 (A3)

2. ... if and only if all relations are dense:

$$(A \to (A \to B)) \to (A \to B)$$
 (A4)

3. ... if and only if all relations commute under |:

$$(A \to \sim B) \to (B \to \sim A)$$
 (A12)

4. ... if and only if all relations are transitive:

$$A \to (A \to A)$$
 (A14)

Tarski's Relevance Logic ...

- ... includes axioms and rules of **RM** that are valid for all binary relations *A*, *B*, and *C*: A1, A5, A6, A7, A8, A9, A10, A11, A13, R1, R2.
- 2. ... excludes the axioms involving density, commutativity, and transitivity: A2, A3, A4, A12, A14
- 3. . . . includes (*), where $A \circ B = \sim (A \to \sim B)$,

$$(A \circ B) \land C \rightarrow ((A \land \sim D) \circ B) \lor (A \circ (B \land (D \circ C)))$$
 (*)

- 4. (*) is not in the Anderson-Belnap system R, nor in RM,
- 5. (*) is a counterexample to a theorem of Kowalski 2012, that the class of commutative dense RAs is characteristic for **R**

A Chronological Bibliography, Part 1 of 2

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Abstract. Algebras of Sugihara, Belnap, Meyer, Church

- Sugihara's relation algebra is a proper relation algebra containing chains of relations isomorphic to Sugihara's original matrices.
- Belnap's matrices form a definitional reduct of a proper relation algebra known as the Point Algebra.
- The 2-element and 4-element Sugihara matrices are definitional subreducts of the Point Algebra.
- Meyer's crystal matrices, Meyer's RM84 matrices, and Church's diamond matrices are also definitional subreducts of proper relation algebras.

• The representation of Sugihara matrices as algebras of binary relations, together with Meyer's characterization of R-mingle by Sugihara matrices, yields the characterization of R-mingle by sets of transitive dense relations that commute under composition.

• The axioms of the Dunn-McCall system R-mingle can therefore be understood as statements about binary relations.

Abstract. Relevance Logic of Tarski

- The definition of Tarski's relevance logic arises from his work on first-order logic restricted to finitely many variables.
- It is inspired by the characterization of equations true in all relation algebras as those provable in first-order logic restricted to 4 variables.
- It has Belnap's variable-sharing property and avoids the paradoxes of implication.
- It contains the Basic Logic of Routley, Plumwood, Meyer, Brady, along with several derived rules of inference.
- It does not include several formulas used as axioms in the Anderson Belnap system R, such as Contraposition.
- It contains a formula (outside both R and R-mingle) that is a counterexample to a completeness theorem of Kowalski (that the system R is complete with respect to the class of dense commutative relation algebras).