

Local finiteness in varieties of MS4-algebras

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Joint work with Guram Bezhanishvili (NMSU)

Outline

1 S4-algebras

2 Duality

3 Monadic S4-algebras

4 M^+S4 -algebras

An S4-algebra is a tuple $\mathfrak{A} = (B, \diamond)$ such that

- B is a Boolean algebra (with $\wedge, \vee, -, 0, 1$)
- $\diamond : B \rightarrow B$ is a unary operation satisfying the identities of a closure operator:
 - ▶ $\diamond 0 = 0$
 - ▶ $\diamond(a \vee b) = \diamond a \vee \diamond b$
 - ▶ $a \leq \diamond a$
 - ▶ $\diamond \diamond a \leq a$
- We write $\Box := -\diamond-$, which is an interior operator.

S4-algebras provide algebraic semantics for the modal logic S4, which is the smallest normal modal logic in one modality \diamond containing

- $\diamond \perp \leftrightarrow \perp$
- $\diamond(p \vee q) \leftrightarrow \diamond p \vee \diamond q$
- $p \rightarrow \diamond p$
- $\diamond \diamond p \rightarrow p$

Note: We will conflate formulas with terms

- $\mathfrak{A} \models \varphi$ means $\mathfrak{A} \models (\varphi = 1)$
- $\mathfrak{A} \models \varphi \rightarrow \psi$ iff $\mathfrak{A} \models \varphi \leq \psi$

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If $\mathfrak{A} = (B, \diamond)$ is an S4-algebra, we write $H_{\Box} = \Box B$ for the set of \Box -fixpoints. In fact,

- H_{\Box} is a bounded sublattice of B that forms a Heyting algebra, where
- Heyting implication is given by $a \rightarrow b = \Box(-a \vee b)$.
- This is intimately related to the Gödel translation of intuitionistic propositional logic (IPC) into the modal logic S4.

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From any **set** X , we get a Boolean algebra $\mathcal{P}(X)$.

Conversely, Stone duality gives us a representation theorem:

- Every Boolean algebra is representable as a subalgebra of $\mathcal{P}(X)$ for some set X .
- We may identify the precise subalgebra as the **clopen** subsets in a certain **Stone topology** on X .

The duality is implemented as follows:

- From B , let $X = \text{Uf } B$; the map $a \mapsto \mathfrak{s}(a) = \{x : a \in x\}$ is an embedding $B \hookrightarrow \mathcal{P}(X)$.
- In fact, $\{\mathfrak{s}(a) : a \in B\}$ is a clopen basis for a zero-dimensional compact Hausdorff (Stone) topology on X , and B is isomorphic to the subalgebra of clopen subsets of X .

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From any relational structure (X, R) where R is a transitive reflexive relation (quasiorder) on X (an **S4-frame**), we get an S4-algebra $(\mathcal{P}(X), \diamond_R)$ where

$$\diamond_R(U) = R^{-1}(U) = \{x \in X : xRy \text{ for some } y \in U\}$$

Conversely, Jónsson–Tarski duality gives a representation theorem:

- Every S4-algebra is representable as a subalgebra of (X, R) for some S4-frame (X, R) .
- We may identify the precise subalgebra as the clopen sets of a Stone topology on X .

The implementation is

- From $\mathfrak{A} = (B, \diamond)$, let $X = \text{Uf } B$ and define xRy iff $x \cap H_{\square} \subseteq y$, the map $a \mapsto \mathfrak{s}(a) = \{x : a \in x\}$ is an embedding $\mathfrak{A} \hookrightarrow (\mathcal{P}(X), R_{\diamond})$.
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The relation R_\diamond will turn out to be ‘compatible’ with the topology on X , in that it satisfies:

- $R(x) = \{y \in X : xRy\}$ is closed
- $R^{-1}(U)$ is clopen whenever U is clopen.

A relation on a Stone space satisfying these properties is said to be *continuous*.

We say that (X, R) is a *descriptive S4-frame* when X is a Stone space and R is a continuous quasiorder on X .

In the same way that Stone duality ultimately yields a dual categorical equivalence between

- The category **Bool** of Boolean algebras, with homomorphisms
- The category **Stone** of Stone spaces (zero-dimensional compact Hausdorff spaces), with continuous maps

We have a dual equivalence between

- The category **S4** of S4-algebras, with homomorphisms
- The category **DS4Fr** of descriptive S4-frames, with continuous *p-morphisms* or *bounded morphisms*.

A map $f : (X, R) \rightarrow (X', R')$ is *p-morphic* if it preserves and reflects the relation, e.g.

- xRy implies $f(x) R' f(y)$
- $f(x) R' y'$ implies xRy for some y with $f(y) = y'$.

$$\begin{array}{ccc}
 y & \overset{f}{\dashrightarrow} & y' \\
 \uparrow R & & \uparrow R' \\
 x & \xrightarrow{f} & x'
 \end{array}$$

Given $\mathfrak{A} = (B, \diamond)$, the *dual frame* or *dual space* of \mathfrak{A} is the descriptive frame $\mathfrak{F} = (X, R)$.

- Subalgebras of \mathfrak{A} correspond to *quotients* of \mathfrak{F} by *correct partitions*, which are ‘compatible’ equivalence relations = kernels of continuous p -morphisms.
- Quotients or homomorphic images of B correspond to closed *generated subframes* of \mathfrak{F} (subframes closed under R).

We also have dual characterizations of simplicity and subdirect irreducibility:

- $x \in X$ is a *root* if $R(x) = X$, \mathfrak{F} is strongly rooted if it has an open set of roots.
- \mathfrak{A} is subdirectly irreducible iff \mathfrak{F} is strongly rooted.
- \mathfrak{A} is simple iff \mathfrak{F} iff every point of X is a root (X is an R -cluster).

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- \mathfrak{A} is *subdirectly irreducible* iff \mathfrak{F} is *strongly rooted*.
- \mathfrak{A} is *simple* iff \mathfrak{F} iff every point of X is a root (X is an R -cluster).

For an S4-algebra \mathfrak{A} , define the *depth* of \mathfrak{A} to be the longest length of a proper R -chain in \mathfrak{F} . If there is no longest length, the depth is ω .

A variety $\mathbf{V} \subseteq \mathbf{S4}$ has *depth* $\leq n$ if the depth of any algebra $\mathfrak{A} \in \mathbf{V}$ is bounded by n . If \mathbf{V} contains algebra of arbitrary depth, then it is of depth ω .

Depth is controlled by the well-known identities:

$$P_1 = \diamond \Box q_1 \rightarrow \Box q_1 \quad P_n = \diamond(\Box q_n \rightarrow \neg P_{n-1}) \rightarrow \Box q_n$$

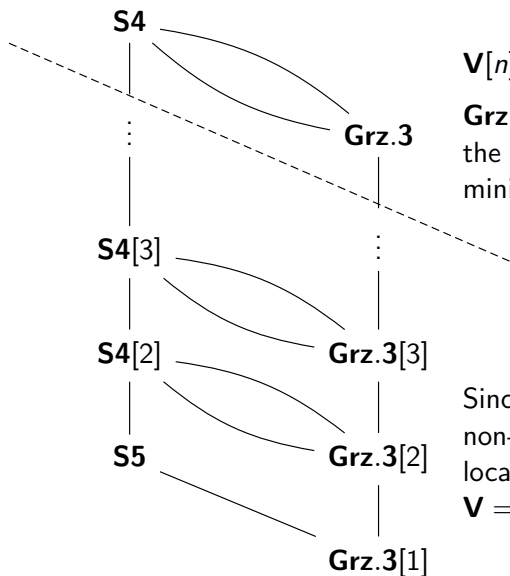
Theorem

$\mathbf{V} \models P_n$ iff the depth of \mathbf{V} is $\leq n$

Theorem (Seegerberg–Maksimova)

$\mathbf{V} \subseteq \mathbf{S4}$ is locally finite iff $\mathbf{V} \models P_n$ for some n .

The subvariety lattice of **S4**:



$$\mathbf{V}[n] = \mathbf{V} + P_n, \mathbf{S5} = \mathbf{S4}[1]$$

Grz.3 $[n]$ is the variety generated by the chain of length n , and is the minimal variety of depth n .

Since **Grz.3** is the unique minimal non-locally-finite subvariety of **S4**, local finiteness is decidable (given $\mathbf{V} = \mathbf{S4} + \varphi$, decide if $\mathbf{V} \supseteq \mathbf{Grz.3}$)

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An *monadic S4-algebra* or MS4-algebra is a tuple $\mathfrak{A} = (B, \diamond, \exists)$ such that

- B is a Boolean algebra (with $\wedge, \vee, -, 0, 1$)
- \diamond is an S4-operator (a closure operator)
- \exists is an S5-operator (an S4-operator satisfying $P_1 := \exists \forall a \leq \forall a$)
- $\exists \diamond a \leq \diamond \exists a$
- We write $\Box := -\diamond-$ and $\forall = -\exists-$

MS4-algebras provide semantics for the one-variable fragment of the standard predicate-logic extension of S4 propositional logic (QS4). The translation is, e.g.

$$\exists \diamond p \rightarrow \diamond \exists p \quad \xrightarrow{T} \quad \exists x \diamond p(x) \rightarrow \diamond \exists x p(x)$$

and φ is a theorem of MS4 iff $T(\varphi)$ is a theorem of QS4.

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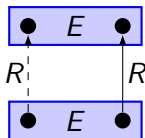
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A *descriptive MS4-frame* is a tuple (X, R, E) where

- X is Stone space
- R is a continuous quasi-order on X
- E is a continuous equivalence relation on X
- $RE \subseteq ER$



Jónsson–Tarski duality again yields a dual equivalence between

- The category **MS4** of MS4-algebras, with homomorphisms
- The category **DMS4Fr** of descriptive MS4-frames with continuous p -morphisms
 - ▶ Such morphisms are p -morphic with respect to both R and E .

Thus every MS4-algebra (B, \diamond, \exists) can be thought of as the algebra of clopen subsets of (X, R, E) , where

$$\diamond U = R^{-1}(U) \quad \exists U = E(U)$$

For any $\mathfrak{A} = (B, \diamond, \exists)$, let

- $B_0 = \exists B$ be the set of \exists -fixpoints
- $\diamond_0 = \diamond|_{B_0}$ the restriction of \diamond

Then $\mathfrak{A}_0 = (B_0, \diamond_0)$ an S4-subalgebra of (B, \diamond) . This follows from the fact that $\exists \diamond \leq \diamond \exists$.

Dually, for any (X, R, E) , let

- $X_0 = X/E$
- R_0 defined by $\alpha R_0 \beta$ iff $\forall x \in \alpha, \exists y \in \beta, xRy$.

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- In an algebra $\mathfrak{A} = (B, \diamond, \exists)$, define $\blacklozenge = \diamond\exists$
- In a descriptive frame $\mathfrak{F} = (X, R, E)$, define $Q = ER$

We may act as if \blacklozenge is an S4-operator/closure operator on B that corresponds dually to the quasi-order Q . We have the dual operator $\blacksquare = -\blacklozenge-$.

This is the appropriate notion to get an analogous characterization of simplicity and subdirect irreducibility:

- \mathfrak{A} is subdirectly irreducible iff \mathfrak{F} is strongly Q -rooted
- \mathfrak{A} is simple iff \mathfrak{F} every point of X is a Q -root (X is a Q -cluster).

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A variety \mathbf{V} is *semisimple* if every subdirectly irreducible algebra is simple. The previous characterization yields that $\mathbf{V} \subseteq \mathbf{MS4}$ is semisimple iff any of the following equivalent conditions hold

- any Q -rooted dual frame from \mathbf{V} is a Q -cluster
- In any dual frame from \mathbf{V} , Q is an equivalence relation
- \mathbf{V} has Q -depth 1, e.g. $\mathbf{V} \models P_1^\blacklozenge := \blacksquare\blacklozenge p \leq \blacklozenge p$.

Theorem

$\mathbf{MS4}_S := \mathbf{MS4} + (\blacksquare\blacklozenge a \leq \blacklozenge a)$ is the largest semisimple subvariety of $\mathbf{MS4}$.

Theorem (Bezhanishvili, M.)

$\mathbf{MS4}_S$ is finitely approximable (generated by its finite members), and so the logic $\mathbf{MS4}_S$ has the finite model property and is decidable.

This generalizes the analogous situation for two well-known logics that are actually extensions of the logic $\mathbf{MS4}_S$.

- $S4_u$ – S4 augmented with a *universal modality*)
- $S5^2$ – the *product* of the modal logic S5 with itself.

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We will speak of the *depth* of an MS4-algebra (B, \diamond, \exists) or a variety of MS4-algebras always meaning the depth of the S4-reduct (B, \diamond) or, equivalently, the R -depth of the dual frame.

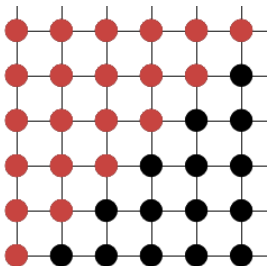
A direct analogue of the Segerberg–Maksimova theorem does not hold; already $\mathbf{MS4}[1] = \mathbf{S5}^2$ is not locally finite.

This was first noticed by Henkin, Monk and Tarski in their study of cylindric algebras.

- $\mathbf{S5}^2 =$ the variety of *two-dimensional diagonal-free cylindric algebras*.

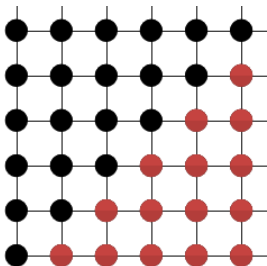
In $S5^2$ -frames, the relations are two commuting equivalence relations (X, E_1, E_2) .

Henkin–Monk–Tarski give a one-generated infinite subalgebra of the $(\mathcal{P}(\omega \times \omega), E_1, E_2)$ where the equivalence classes of E_1 and E_2 are rows and columns.



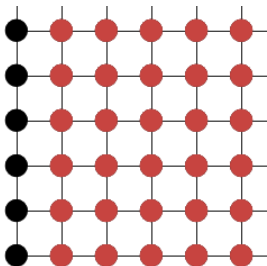
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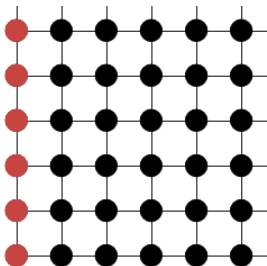
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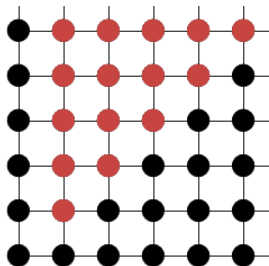
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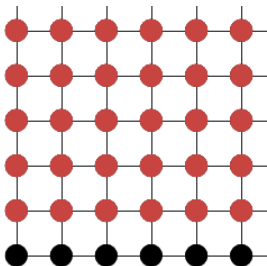
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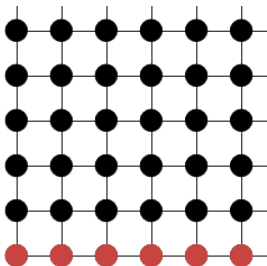
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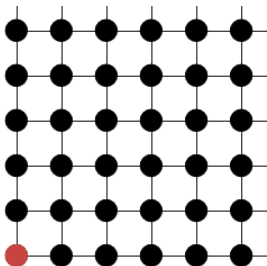
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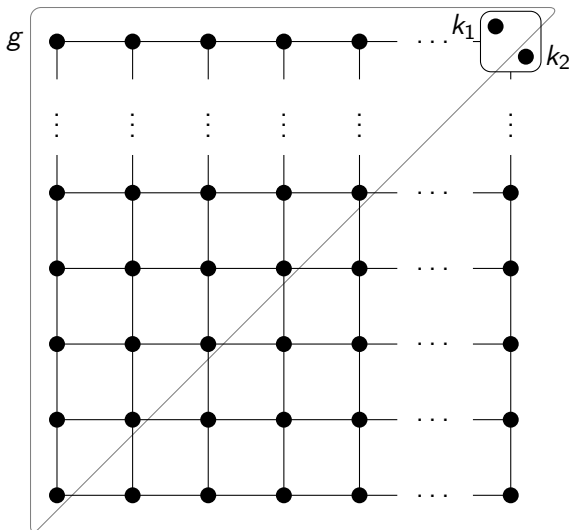


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We characterize the dual frame of this algebra, in order to give an explicit example of an MS4-algebra and descriptive MS4-frame that is one-generated and infinite.



The *layers* of a quasiorder can be defined as

$$D_1 = \max X \quad D_{n+1} = \max \left(X - \bigcup_{i=1}^n D_i \right)$$

For descriptive S4-frames,

- ① D_1 is always closed.
- ② If \mathfrak{F} is finitely generated, then every finite layer is clopen, and consists of isolated points.

This example shows that finitely generated MS4-algebras can have limit points of finite depth; we do not know if it remains true that the layers themselves are clopen in this case.

We would also like to know if there is an MS4-frame whose layers *only* consist of limit points.

Since every proper subvariety of $\mathbf{S5}^2$ is locally finite, and every layer of an MS4-algebra is an $S5^2$ -frame, this suggests the reasonable conjecture that a variety of MS4-algebras is locally finite if

- 1 It is finite depth (necessary)
- 2 Every layer is locally finite as an $S5^2$ -frame (“finite width”)

This conjecture fails drastically in general, but first some positive results.

We give the following opaque characterization of local finiteness:

Theorem (Bezhanishvili, M.)

$\mathbf{V} \subseteq \mathbf{MS4}$ is locally finite iff

- ① It is finite depth ($\mathbf{V} \models P_n$ for some n)
- ② There is a uniform bound $f : \mathbb{N} \rightarrow \mathbb{N}$ such that, for every n -generated s.i. $\mathfrak{A} \in \mathbf{V}$, the algebra \mathfrak{A}_0 is $f(n)$ -generated as an S4-algebra.

Using this criterion, we obtain

Theorem

A variety $\mathbf{V} \subseteq \mathbf{MS4} + \text{alt}_k^0$ is locally finite iff it is finite depth.

This generalizes the previously known results for $\mathbf{S4}_u = \mathbf{MS4} + \text{alt}_1^0$.

We give the following opaque characterization of local finiteness:

Theorem (Bezhanishvili, M.)

$\mathbf{V} \subseteq \mathbf{MS4}$ is locally finite iff

- ① It is finite depth ($\mathbf{V} \models P_n$ for some n)
- ② There is a uniform bound $f : \mathbb{N} \rightarrow \mathbb{N}$ such that, for every n -generated s.i. $\mathfrak{A} \in \mathbf{V}$, the algebra \mathfrak{A}_0 is $f(n)$ -generated as an S4-algebra.

Using this criterion, we obtain

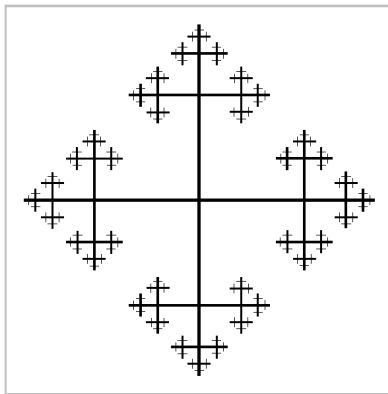
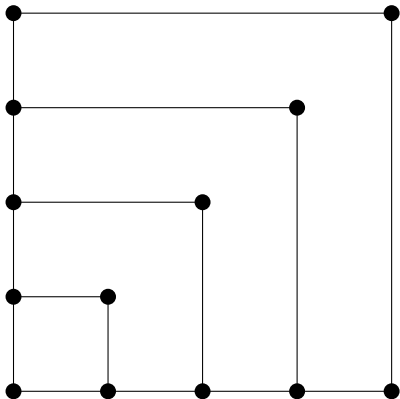
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$S5_2$ is the variety of $S5_2$ -algebras: tuples $(B, \exists_1, \exists_2)$ where the \exists_i are **unrelated** S5-operators.

- Dually, these correspond to frames (X, E_1, E_2) where the E_i are **unrelated** equivalence relations.
- $S5^2$ is the *product* of S5 with itself (= commutivity axiom)
- $S5_2$ is the *fusion* of S5 with itself (= no axioms relating the modalities)



A characterization of local finiteness in $\mathbf{S5}_2$ is wide-open.

- Fusions do not preserve local finiteness
- In general, only highly restrictive sufficient conditions are known to guarantee a fusion to be locally finite.

Let $\Lambda(\mathbf{V})$ denote the lattice of subvarieties of a variety \mathbf{V} .

Theorem (Bezhanishvili, M.)

There is a translation $T : \Lambda(\mathbf{S5}_2) \rightarrow \Lambda(\mathbf{MS4}_S[2])$ that preserves and reflects local finiteness (that is, \mathbf{V} is locally finite iff $T(\mathbf{V})$ is locally finite)

In particular, characterizing local finiteness even in $\mathbf{MS4}_S[2]$ is as hard as the corresponding problem for $\mathbf{S5}_2$.

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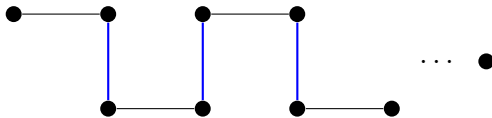
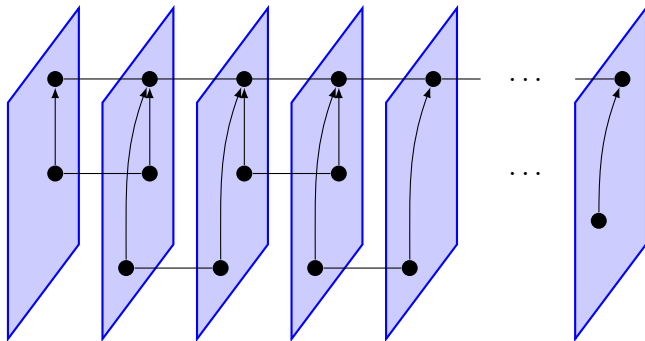
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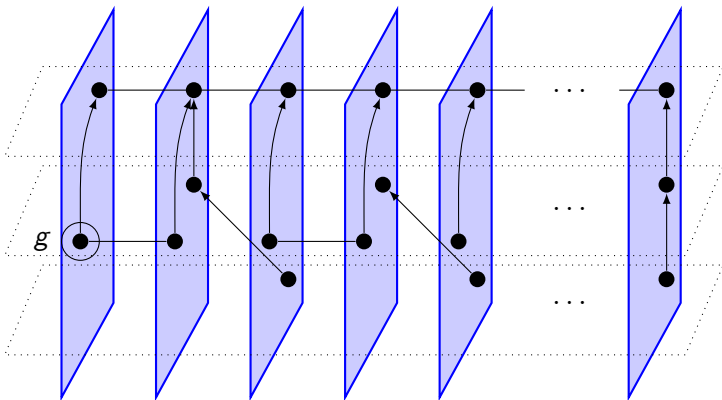
Proof in one slide

 \mathfrak{F}  $T(\mathfrak{F})$

Some other consequences:

- MS4 is the *expanding relativized product* (or *semiproduct*) of S4 and S5 (impose left commutivity axiom $\exists\Diamond \rightarrow \Diamond\exists$).
- Our construction actually extends to the full *product*, where one imposes the full commutivity condition $\exists\Diamond \leftrightarrow \Diamond\exists$.
- By adding a “bottom rail” to the translation, one can ensure that the relations R and E actually commute in the resulting frame.

We also provide an example demonstrating that local finiteness of layers (as S_5_2 -algebras) is not even sufficient for local finiteness already in depth-3:



Outline

1 S4-algebras

2 Duality

3 Monadic S4-algebras

4 M⁺S4-algebras

The following formulas are well-known in the study of intuitionistic predicate logic

$$\text{Cas} := \forall x((P(x) \rightarrow \forall yP(y)) \rightarrow \forall yP(y)) \rightarrow \forall xP(x)$$

$$\text{K} := \forall x\neg\neg P(x) \rightarrow \neg\neg\forall xP(x)$$

Both give rise to distinct intermediate predicate logics, and Cas is strictly weaker than K.

Cas plays a role in extending the translations

$$\text{IPC} \xrightarrow{\text{Gödel translation}} \text{Grz} \xrightarrow{\text{splitting translation}} \text{GL}$$

where

- IPC = Intuitionistic propositional logic
- Grz = S4 + grz is the Grzegorzczuk logic
- GL = K4 + Löb's axiom = the logic of the Peano Arithmetic provability predicate.

Let

- $\text{MCas} = \forall((p \rightarrow \forall p) \rightarrow \forall p) \rightarrow \forall p$
 - ▶ the monadic version of the Casari formula
- $\text{MCas}^t = \blacksquare(\Box(\Box p \rightarrow \blacksquare p) \rightarrow \blacksquare p) \rightarrow \blacksquare p$
 - ▶ the Gödel translation of MCas
 - ▶ ($\blacksquare = \Box \forall$)

Bezhanishvili, Brantley and Ilin showed that these formulas are exactly what is required to lift the previous translations to the monadic setting:

$$\underbrace{\text{MIPC} + \text{MCas}}_{\text{M}^+\text{IPC}} \xrightarrow{\text{Gödel translation}} \underbrace{\text{MGrz} + \text{MCas}^t}_{\text{M}^+\text{Grz}} \xrightarrow{\text{splitting translation}} \text{MGL}$$

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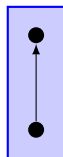
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It is natural then to examine the variety $M^+S4 = MS4 + MCas^t$.

Roughly, M⁺S4-frames do not have *dirty clusters*,
an *E*-class containing properly *R*-related points.



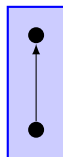
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Theorem (M.)

- In an M⁺S4-frame \mathfrak{F} , every layer of finite depth is *E*-saturated, and hence every layer is a (refined) S5²-frame.
- There is a unique minimal subvariety of MS4 of depth ω .

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Finally, the local finiteness situation already seems more under control:

Theorem (M.)

$\mathbf{V} \subseteq \mathbf{M}^+\mathbf{S4}$ is locally finite iff

- ① \mathbf{V} is of finite depth (*decidable*)
- ② \mathbf{V}_T is locally finite – equivalently, $\mathbf{V}_T \subset \mathbf{S5}^2$.

Here \mathbf{V}_T is a subvariety of $\mathbf{S5}^2$ defined semantically from the layers of algebras in \mathbf{V} ; work is ongoing to investigate the decidability of (2).

Thanks!