

## Review 2: Logic

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# Implications

Let  $P, Q$  be statements.

statement	equivalent meanings	negation
$P \Rightarrow Q$	If $P$ is true, then $Q$ is true. $\sim P \vee Q$ $\sim Q \Rightarrow \sim P$ (contrapositive)	$P \wedge \sim Q$
$P \Leftrightarrow Q$	$Q$ is true if and only if $P$ is true. $(P \Rightarrow Q) \wedge (P \Leftarrow Q)$	$P \Leftrightarrow \sim Q$ $\sim P \Leftrightarrow Q$

How to prove  $P \Rightarrow Q$ :

- ▶ **direct:** Assume  $P$ . Show  $Q$ .
- ▶ **contrapositive:** Assume  $\sim Q$ . Show  $\sim P$ .
- ▶ **by contradiction** (only as last resort): Assume  $P$  and  $\sim Q$ . Show some contradiction.

How to prove  $P \Leftrightarrow Q$ :

Show  $P \Rightarrow Q$  and show  $P \Leftarrow Q$ .

# Equivalent statements

## Theorem

Let  $f: A \rightarrow B$  be a function. Then the following are equivalent (TFAE):

1.  $f$  is bijective.
2.  $\forall y \in B \exists$  unique  $x \in A : f(x) = y$ .
3.  $f$  has an inverse function.
4.  $\exists \ell, r: B \rightarrow A : \ell \circ f = \text{id}_A$  and  $f \circ r = \text{id}_B$ .

This Theorem states that  $1. \Leftrightarrow 2. \Leftrightarrow 3. \Leftrightarrow 4.$

Instead of showing 3 times  $\Leftrightarrow$ , such statements can be proved more efficiently as a cycle of implications:

$$1. \Rightarrow 2. \Rightarrow 3. \Rightarrow 4. \Rightarrow 1.$$

## Quantified statements

Let  $A$  be a set,  $P(x)$  be a statement for  $x \in A$ .

statement	meaning	negation
$\forall x \in A: P(x)$	For all $x \in A$ , $P(x)$ is true.	$P(x)$ is not true for some $x \in A$ . $\exists x \in A: \sim P(x)$
$\exists x \in A: P(x)$	There exists $x \in A$ such that $P(x)$ is true.	$P(x)$ is not true for all $x \in A$ . $\forall x \in A: \sim P(x)$

How to prove  $\forall x \in A: P(x)$

Let  $x \in A$  (arbitrary but fixed). Show  $P(x)$ .

How to refute  $\forall x \in A: P(x)$

Give concrete explicit  $x \in A$  that does not satisfy  $P(x)$ .

How to prove  $\exists x \in A: P(x)$

Give a concrete explicit  $x \in A$  that satisfies  $P(x)$ .

How to refute  $\exists x \in A: P(x)$

Show  $\forall x \in A: \sim P(x)$ .

# Quantifiers and implications in Calculus

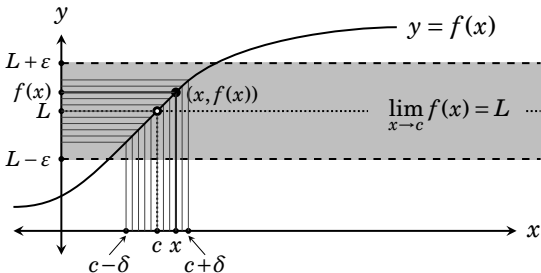
**Calculus:** Informally  $\lim_{x \rightarrow c} f(x) = L$  means that  $f(x)$  is arbitrarily close to  $L$  provided that  $x$  is sufficiently close to  $c$ .  
More precisely

### Definition (Limit of a function)

Let  $c \in \mathbb{R}$  and  $f: \mathbb{R} - \{c\} \rightarrow \mathbb{R}$ . Then  $\lim_{x \rightarrow c} f(x) = L$  if

$$\forall \varepsilon \in \mathbb{R}^+ \exists \delta \in \mathbb{R}^+ \forall x \in \mathbb{R} - \{c\}: |x - c| < \delta \Rightarrow |f(x) - L| < \varepsilon$$

Diagram taken from Hammack, Book of Proof, 2018.



How to prove  $\forall \varepsilon \in \mathbb{R}^+ \exists \delta \in \mathbb{R}^+ \forall x \in \mathbb{R} - \{c\}$ :

$$|x - c| < \delta \Rightarrow |f(x) - L| < \varepsilon$$

### Example

Prove  $\lim_{x \rightarrow 0} 2x \sin(\frac{1}{x}) = 0$ .

- ▶  $(\forall)$  Let  $\varepsilon > 0$  arbitrary, fixed for the remainder of the proof.
- ▶  $(\exists)$  Find  $\delta > 0$  ( $\delta$  may depend on  $\varepsilon$ ) such that  $\forall x \neq c$ :  
 $|x - c| < \delta \Rightarrow |f(x) - L| < \varepsilon. \quad (*)$
- ▶  $(\forall)$  Let  $x \in \mathbb{R} - \{0\}$  arbitrary, fixed. For  $(*)$  consider

$$|f(x) - L| = \left| 2x \sin\left(\frac{1}{x}\right) - 0 \right| = 2|x| \underbrace{\left| \sin\left(\frac{1}{x}\right) \right|}_{\leq 1} \leq 2|x|$$

- ▶ The latter is  $< \varepsilon$  if  $|x| < \varepsilon/2$ .
- ▶ We found  $\delta := \varepsilon/2$  such that  $|x - 0| < \delta \Rightarrow |f(x) - 0| < \varepsilon$ .
- ▶ Thus  $\lim_{x \rightarrow 0} 2x \sin(\frac{1}{x}) = 0. \quad \square$

## How to prove a limit does not exist?

We need to show that **for all**  $L \in \mathbb{R}$ ,

$$\sim (\forall \varepsilon \in \mathbb{R}^+ \exists \delta \in \mathbb{R}^+ \forall x \in \mathbb{R} - \{c\}: |x - c| < \delta \Rightarrow |f(x) - L| < \varepsilon),$$

equivalently,

$$\forall L \in \mathbb{R} \exists \varepsilon \in \mathbb{R}^+ \forall \delta \in \mathbb{R}^+ \exists x \in \mathbb{R} - \{c\}: |x - c| < \delta \text{ and } |f(x) - L| \geq \varepsilon.$$

### Example

Prove  $\lim_{x \rightarrow 0} \frac{1}{x}$  does not exist.

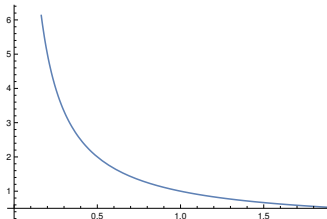
▶ **( $\forall$ )** Let  $L \in \mathbb{R}^+$ , **( $\exists$ )** guess  $\varepsilon > 0$ ,  
**( $\forall$ )** let  $\delta > 0$  arbitrary.

▶ **( $\exists$ )** Find  $x \in \mathbb{R} - \{0\}$ :  
 $|x - 0| < \delta$  and  $|\frac{1}{x} - L| \geq \varepsilon$ .

▶ Say  $0 < x < \delta$  and  $\frac{1}{x} - L \geq \varepsilon$ .

Hence  $\frac{1}{\varepsilon + L} \geq x$  and e.g.  $x := \min(\frac{\delta}{2}, \frac{1}{\varepsilon + L})$  works!

▶ Thus  $\lim_{x \rightarrow 0} \frac{1}{x}$  cannot be any  $L \in \mathbb{R}^+$ .





## Sum rule for limits

$\lim_{x \rightarrow c} f(x) = L$  if  $\forall \varepsilon > 0 \exists \delta > 0 \forall x \neq c: |x - c| < \delta \Rightarrow |f(x) - L| < \varepsilon$

### Theorem

If  $\lim_{x \rightarrow c} f(x)$  and  $\lim_{x \rightarrow c} g(x)$  both exist, then

$$\lim_{x \rightarrow c} [f(x) + g(x)] = \lim_{x \rightarrow c} f(x) + \lim_{x \rightarrow c} g(x).$$

### Proof.

- ▶ Let  $\lim_{x \rightarrow c} f(x) =: L$  and  $\lim_{x \rightarrow c} g(x) =: M$ . Show  $\lim_{x \rightarrow c} [f(x) + g(x)] = L + M$ .
- ▶ Let  $\varepsilon > 0$  arbitrary. Find  $\delta > 0$  such that  $|x - c| < \delta \Rightarrow |f(x) + g(x) - L - M| < \varepsilon$ .
- ▶  $|f(x) + g(x) - L - M| \leq |f(x) - L| + |g(x) - M|$  by the triangle inequality.
- ▶ We have  $\delta_1 > 0: |x - c| < \delta_1 \Rightarrow |f(x) - L| < \frac{\varepsilon}{2}$ ,  
 $\delta_2 > 0: |x - c| < \delta_2 \Rightarrow |g(x) - M| < \frac{\varepsilon}{2}$ .
- ▶ Then  $|x - c| < \underbrace{\min(\delta_1, \delta_2)}_{=: \delta} \Rightarrow |f(x) + g(x) - L - M| < \underbrace{\frac{\varepsilon}{2} + \frac{\varepsilon}{2}}_{=: \varepsilon}$ .

▶ Do you want to know more? Take Math 3001 – Analysis