

Input: vector \vec{v} or function $f: \mathbb{Z}/m\mathbb{Z} \rightarrow \mathbb{C}$

Output: $F_m \vec{v}$ or $\hat{f}: \mathbb{Z}/m\mathbb{Z} \rightarrow \mathbb{C}$

Formula for new coefficients:

$$(F_m \vec{v})_n = \frac{1}{\sqrt{m}} \sum_{k=0}^{m-1} v_k \omega_m^{nk}$$

or

$$\hat{f}(n) = \frac{1}{\sqrt{m}} \sum_{k=0}^{m-1} f(k) \omega_m^{nk}$$

Computer
Demo

Runtime?

Classical: discrete fourier transform via matrix multiplication
 $O(m^2)$

Fast Fourier Transform = $O(m \log m)$

Quantum Fourier Transform: F_m is a unitary matrix

$$\sum_{x=0}^{m-1} f(x) |x\rangle \longmapsto \sum_{y=0}^{m-1} \hat{f}(y) |y\rangle$$
$$= \frac{1}{\sqrt{m}} \sum_{y=0}^{m-1} \left(\sum_{k=0}^{m-1} \omega_m^{yk} f(k) \right) |y\rangle$$

$m = 2^n$
n qubits
can store
 $0 = 0000 \dots 0$
up to
 $m-1 = 111 \dots 1$

Can be implemented with $O(\log_2^2(m))$ gates

(H and a controlled phase gates)

QFT of a periodic state

QFT reminder:

$$\sum_{x=0}^{m-1} \alpha_x |x\rangle \mapsto \frac{1}{\sqrt{m}} \sum_{y=0}^{m-1} \left(\sum_{k=0}^{m-1} \alpha_k \omega_m^{ky} \right) |y\rangle$$

Input:

$$m = 2^n$$

$$r | m$$

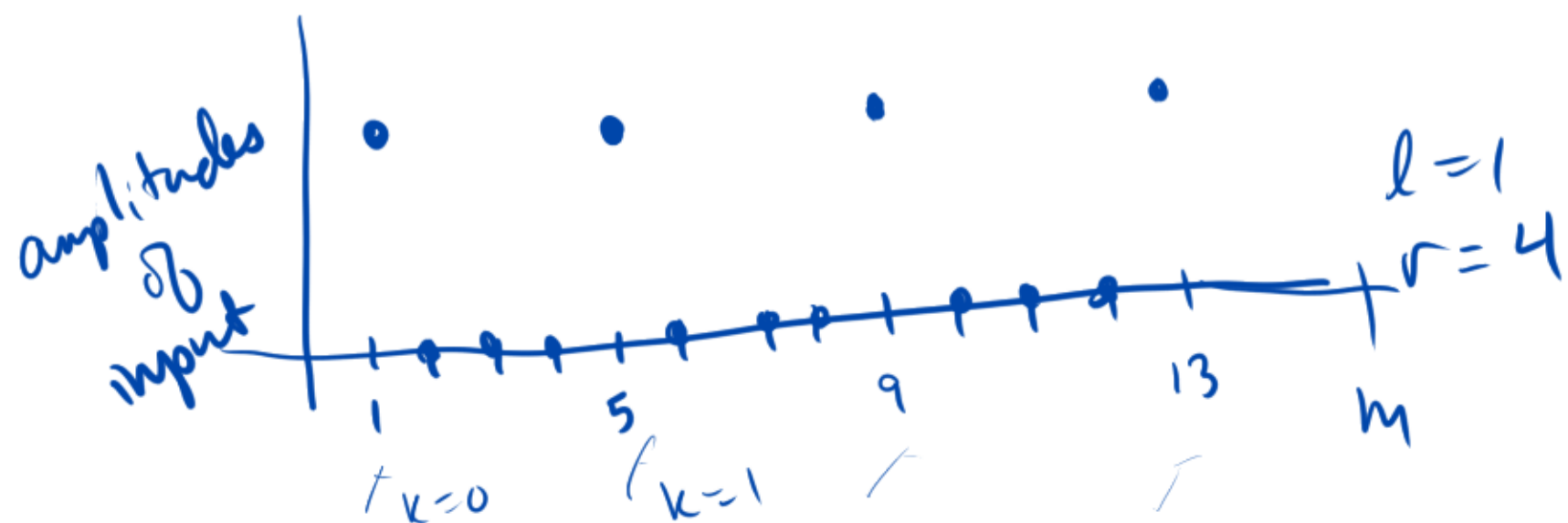
↑
period

$$\frac{1}{\sqrt{m/r}} \sum_{k=0}^{m/r-1} |l + kr\rangle$$

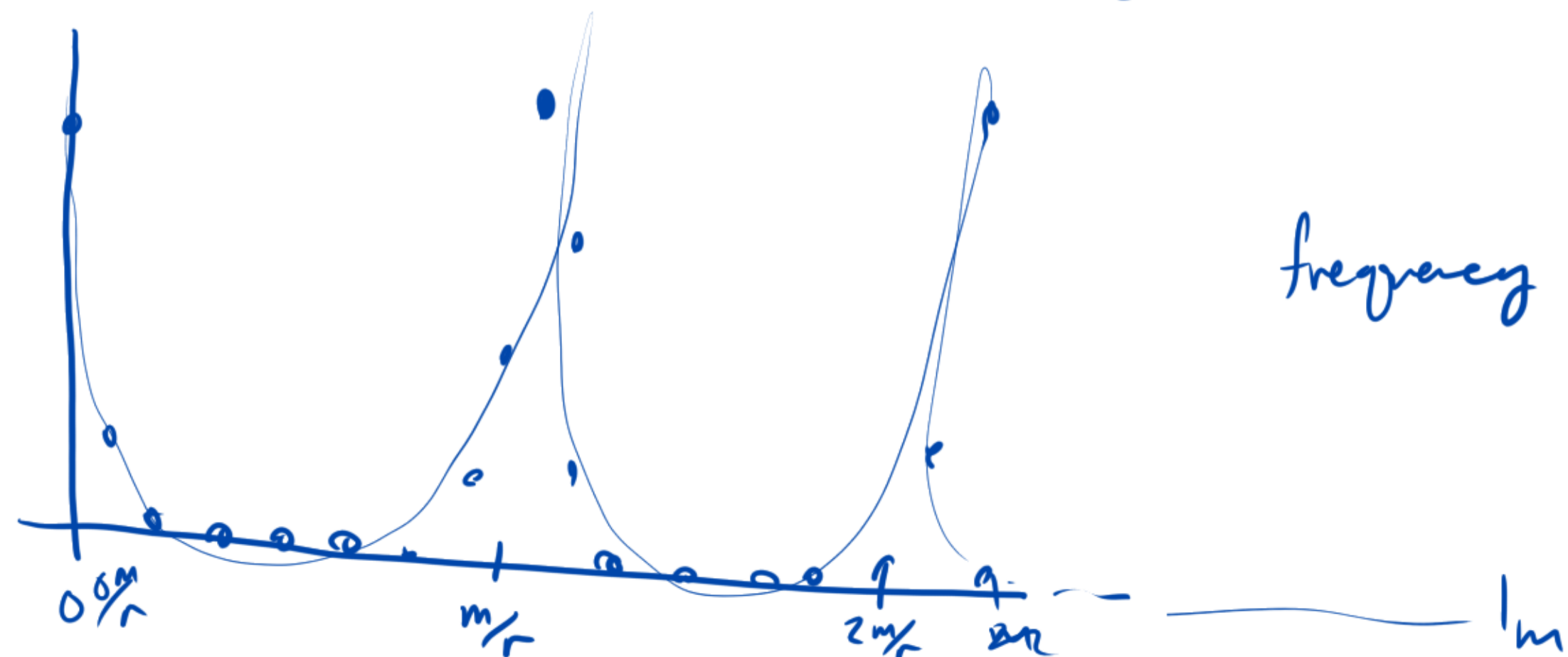
QFT_m →

$$\frac{1}{\sqrt{m} \sqrt{m/r}} \sum_{y=0}^{m-1} \left(\sum_{k=0}^{m/r-1} \omega_m^{(l+kr)y} \right) |y\rangle$$

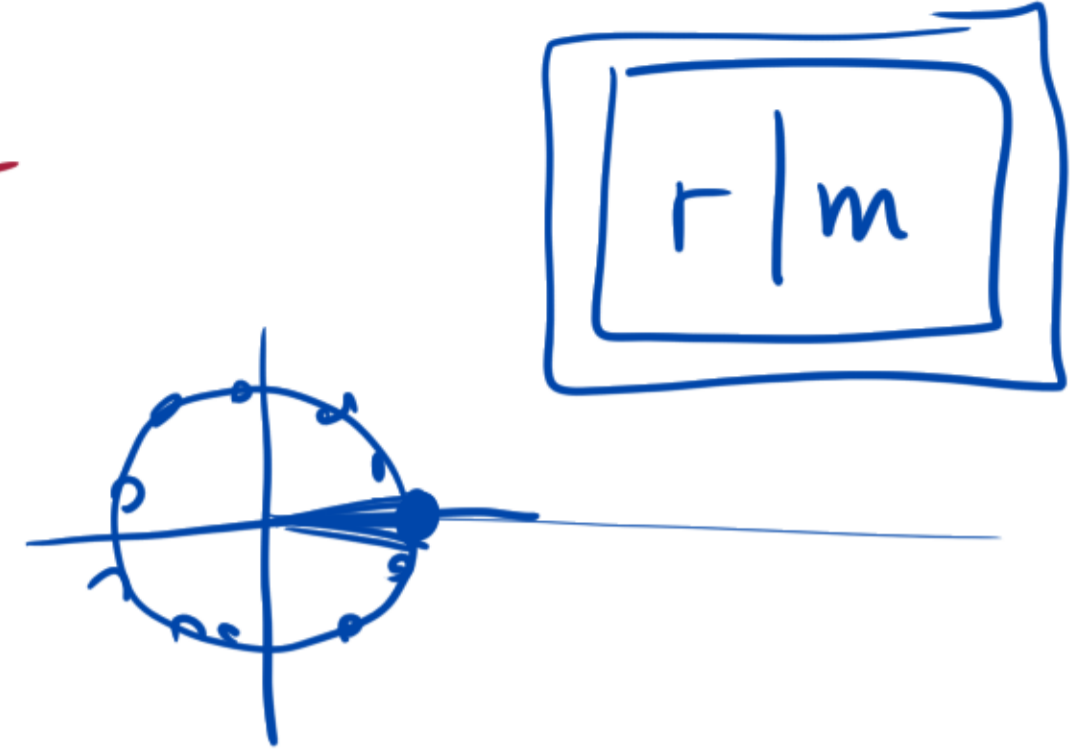
← |ampl.|²



↔



Amplitude @ $|y\rangle$ is $\frac{1}{\sqrt{m}\sqrt{m/r}} \sum_{k=0}^{m/r-1} \omega_m^{(2+kr)y}$



If $y = 0, \frac{m}{r}, 2\frac{m}{r}, 3\frac{m}{r}, \dots$ i.e. $y \in \frac{m}{r}\mathbb{Z}$.

$$\text{ampl} = \frac{1}{\sqrt{m}\sqrt{m/r}} \omega_m^{ly} \sum_{k=0}^{m/r-1} \omega_m^{kry}$$

$$\left\{ \begin{array}{l} kry \in km\mathbb{Z} \subseteq m\mathbb{Z} \\ \text{so} \\ \omega_m^{kry} = (\omega_m^m)^{\text{integer}} = 1^{\text{integer}} = 1 \end{array} \right.$$

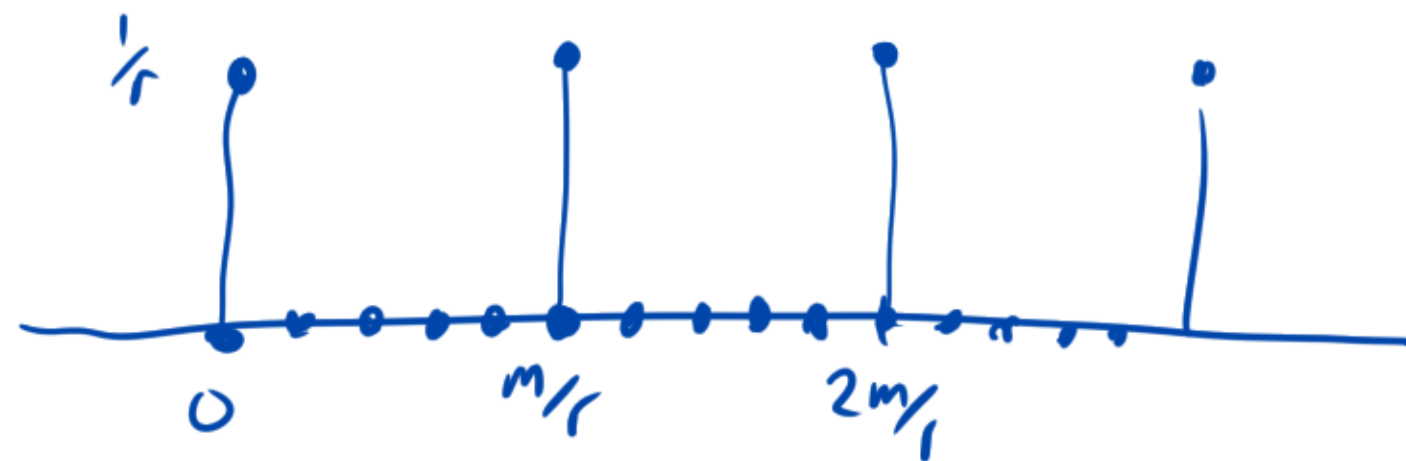
$$= \frac{1}{\sqrt{m}\sqrt{m/r}} \omega_m^{ly} \left(\frac{m}{r}\right)$$

$$|\text{ampl}|^2 = \frac{1}{r}$$

So: $|\text{ampl}|^2 = \frac{1}{r}$ @ the r positions $y = 0, \frac{m}{r}, 2\frac{m}{r}, \dots$

Since $\sum |\text{ampl}|^2 = 1$, the rest must be 0.

output of QFT



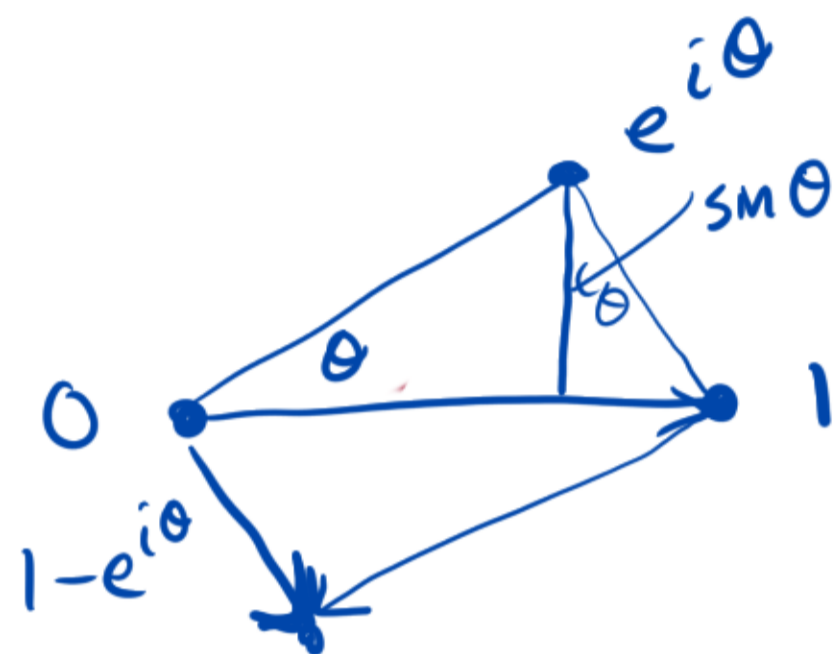
If $\frac{m}{r} \notin \mathbb{Z}$?

Amplitude @ $|y\rangle$ is $\frac{1}{\sqrt{m} \sqrt{s}} \sum_{k=0}^{s-1} w_m^{kry}$

$(l+kr)y$
 w_m

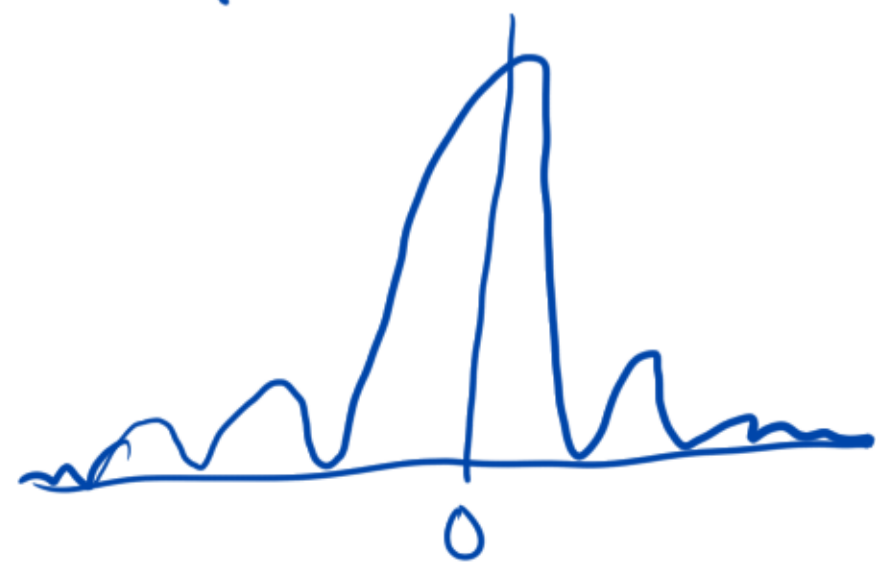
where $s = \lfloor \frac{m}{r} \rfloor$
 or $\lfloor \frac{m}{r} \rfloor + 1$

$$\begin{aligned}
 |\text{ampl}|^2 &= \frac{1}{sm} \left| \sum_{k=0}^{s-1} w_m^{kry} \right|^2 \\
 &= \frac{1}{sm} \left| \frac{1 - w_m^{rys}}{1 - w_m^{ry}} \right|^2 \\
 &= \frac{1}{sm} \left| \frac{\sin\left(\frac{2\pi r y s}{m}\right)}{\sin\left(\frac{2\pi r y}{m}\right)} \right|^2
 \end{aligned}$$



where $\theta = \frac{2\pi r y}{m}$ peaks @ $\theta \rightarrow 2\pi \mathbb{Z}$

$y \rightarrow \frac{m}{r} \mathbb{Z}$
 $\frac{y}{m} \rightarrow \frac{\mathbb{Z}}{r}$



Observation / Bottom Line:

output has magnitude peaks

@ $\frac{m}{r} \mathbb{Z}$

and the offset (l) is absorbed into the phase not magnitude.

Computer Demo

Computer Demo

Factoring N

Recall: If we can find $x \in \mathbb{Z}/N\mathbb{Z}$ s.t. $x^2 \equiv 1$, $x \not\equiv \pm 1 \pmod{N}$
 then we can factor N , by $\gcd(x-1, N)$ or $\gcd(x+1, N)$.

Ex. $N=15$ $\varphi(15)=8$

	1	α	α^2	α^3	α^4	α^5	α^6	α^7	$(\alpha^8=1)$	
$\alpha=2$	1	2	4	8	1					mult. order = 4 of 2
										$\alpha^4=1$ $\alpha^2=4$ $(\alpha^2)^2=1$

$\alpha=13$	1	13	11	7	1
$\alpha=11$	1	11	1		

Idea: Take a random α , find $\alpha^r = 1$

If r is even, compute $\alpha^{r/2}$

If $\alpha^{r/2} \not\equiv \pm 1$, we win

since $(\alpha^{r/2})^2 \equiv 1$

So take $\gcd(\alpha^{r/2} + 1, N)$

Finding mult. order of α
 \Downarrow
 factoring!

Shor's Algorithm

classical ① Take a random $\alpha \pmod N$, $\alpha \neq 0$.

quantum ② Find the order of α in $(\mathbb{Z}/N\mathbb{Z})^*$.

classical ③ Check if r is even, and $\alpha^{r/2} \not\equiv \pm 1 \pmod N$. } Luck
(If not, change α ,

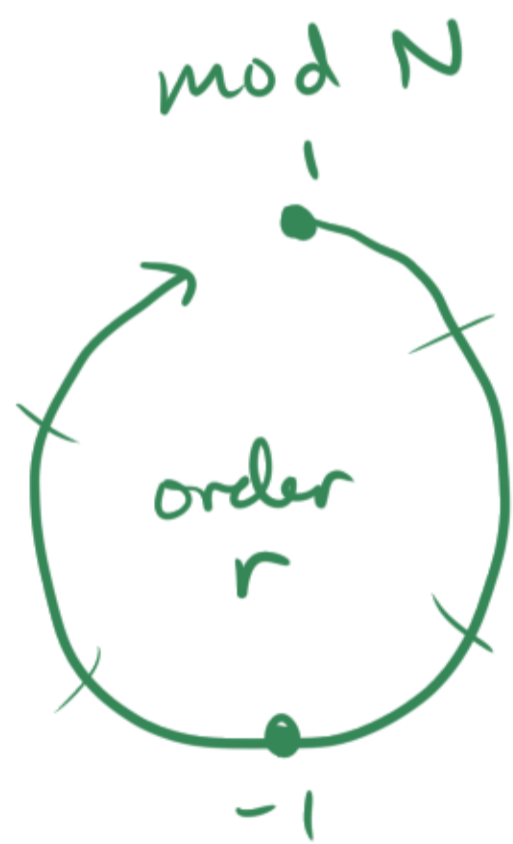
classical ④ If so, take $\gcd(\alpha^{r/2} - 1, N)$.
try again.)

Proposition Let $N = pq$, p, q odd primes.

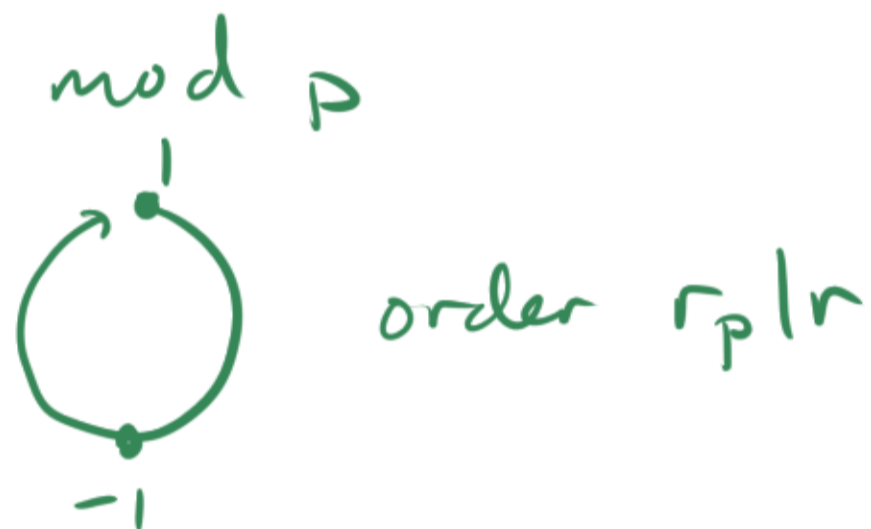
Then $\text{Prob}(\text{ord}(x) \text{ is even and } x^{\frac{\text{ord}(x)}{2}} \not\equiv -1 \pmod{N}) \geq \frac{1}{2}$ as x ranges over $(\mathbb{Z}/N\mathbb{Z})^*$.

Proof. Suppose $\text{ord}(x)$ is even and $x^{\frac{\text{ord}(x)}{2}} \equiv -1 \pmod{N}$.

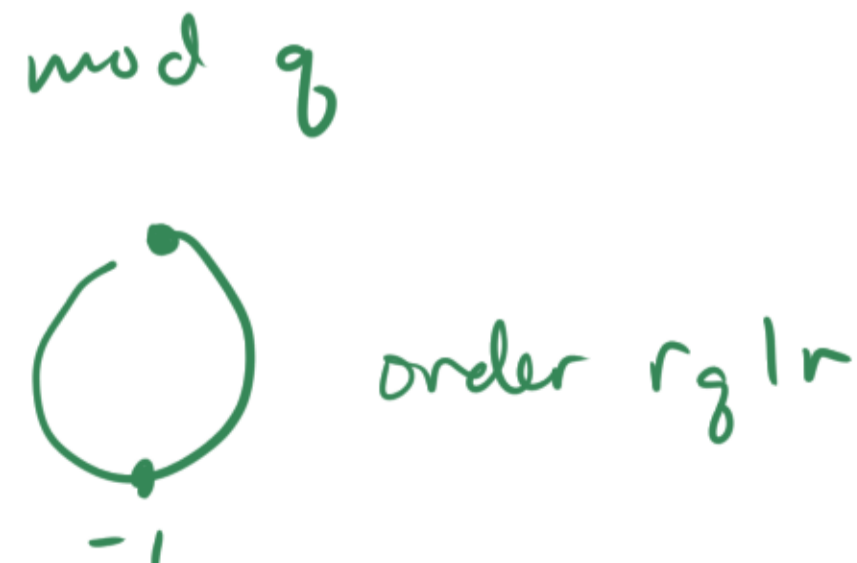
Mult. dyn. of x :



wrap
odd # of times



wrap
odd



-1 appears in a cycle mod p
 $\Leftrightarrow x^{r/2} \equiv -1 \equiv g^{\frac{p-1}{2}}$
 $\Leftrightarrow x = g^s$ where $s \mid \frac{p-1}{2}$.
 $p-1 = 2^{\sim}(\text{odd})$

This occurs at most $\frac{1}{2}$
of the time.

So at most $\frac{1}{4}$ of the residues of $(\mathbb{Z}/N\mathbb{Z})^*$ have this property mod p & mod q
i.e. land in this case.