

## Classical Reversible Circuits:

Can construct any bijection  $f: \{0,1\}^n \rightarrow \{0,1\}^n$  as a reversible circuit s.t.



## Classical Reversible Circuits on a Quantum Computer:

**Fact:** Any reversible classical circuit  $f$  is a unitary transformation:  $U_f$

So we can perform it on a quantum computer:

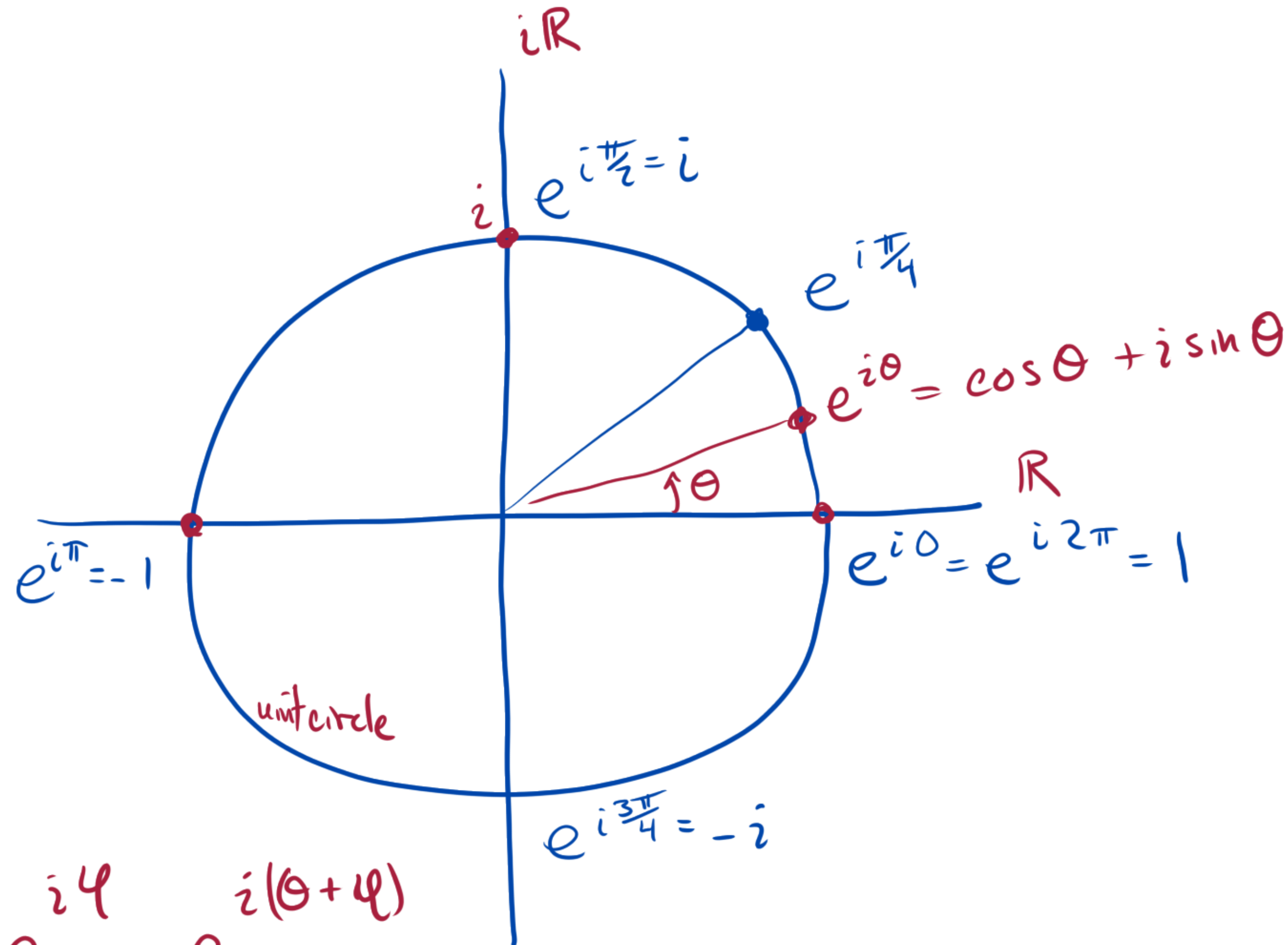


"massive parallelism"

$$U_f \left( \sum_{x \in \{0,1\}^n} \alpha_x |x\rangle |0\rangle \right) = \sum_{x \in \{0,1\}^n} \alpha_x |x\rangle |f(x)\rangle$$

"computes"  $f(x)$  for all  $2^n$  classical inputs simultaneously.

# Roots of Unity in $\mathbb{C}$



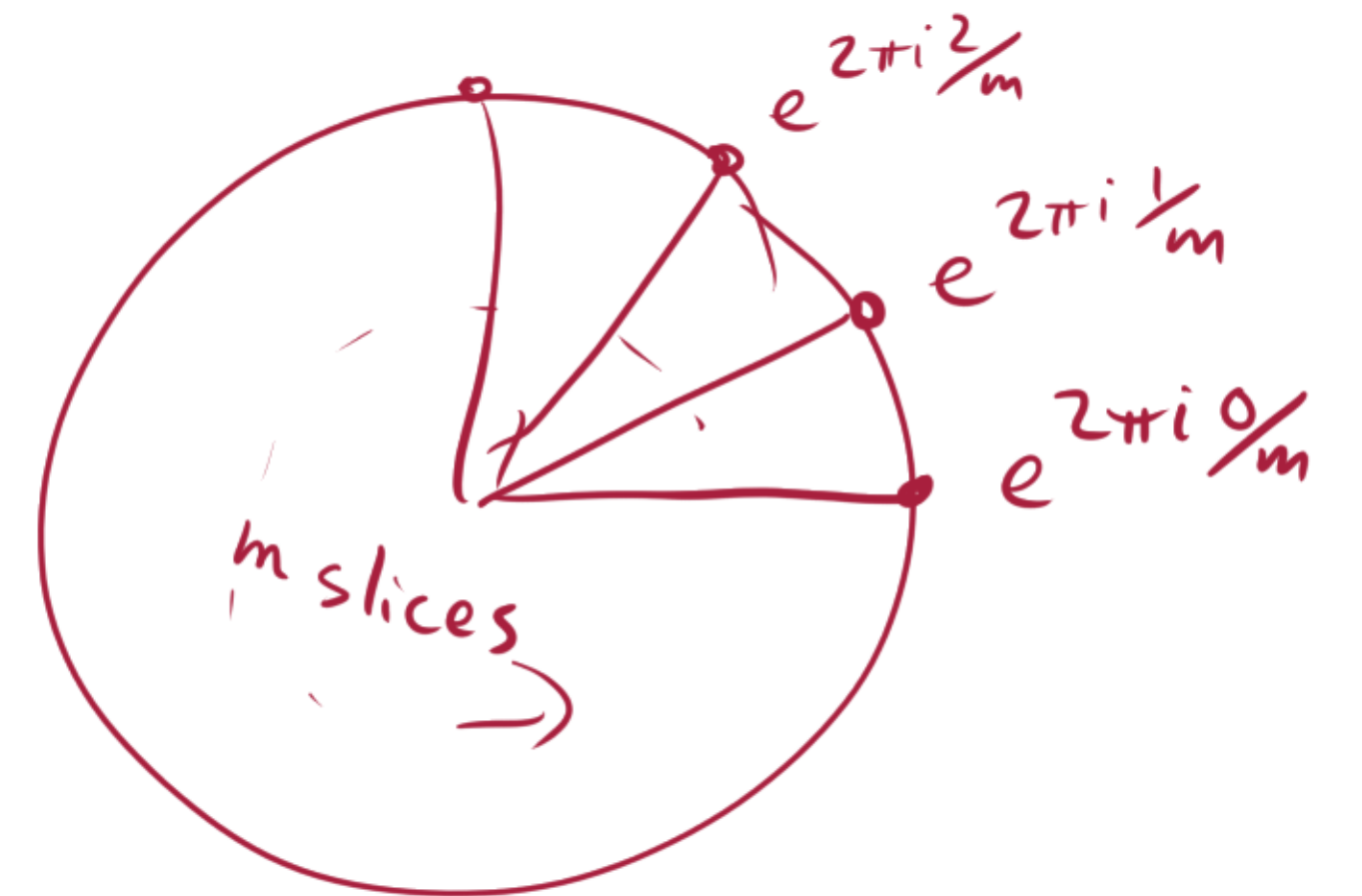
$$e^{i\theta} e^{i\varphi} = e^{i(\theta+\varphi)}$$

The  $m^{\text{th}}$  roots of unity are

$$\omega_m = e^{2\pi i/m}$$

$$\omega_m^k = e^{2\pi i k/m}$$

$$k=0, 1, \dots, m-1$$



$$\omega_m^m = 1$$

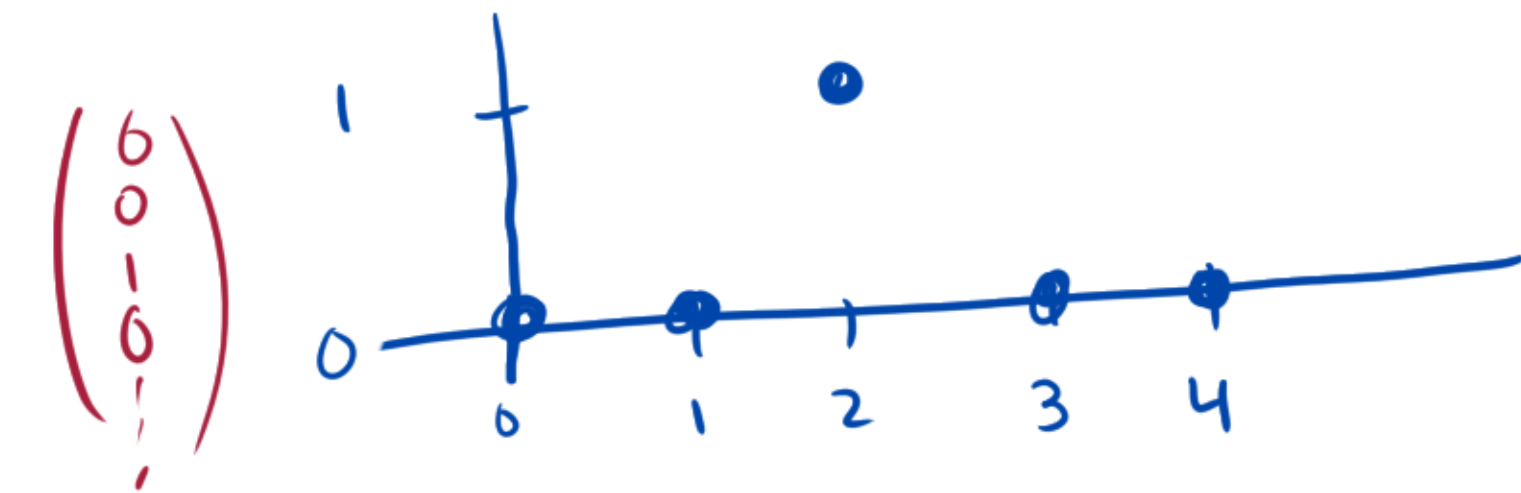
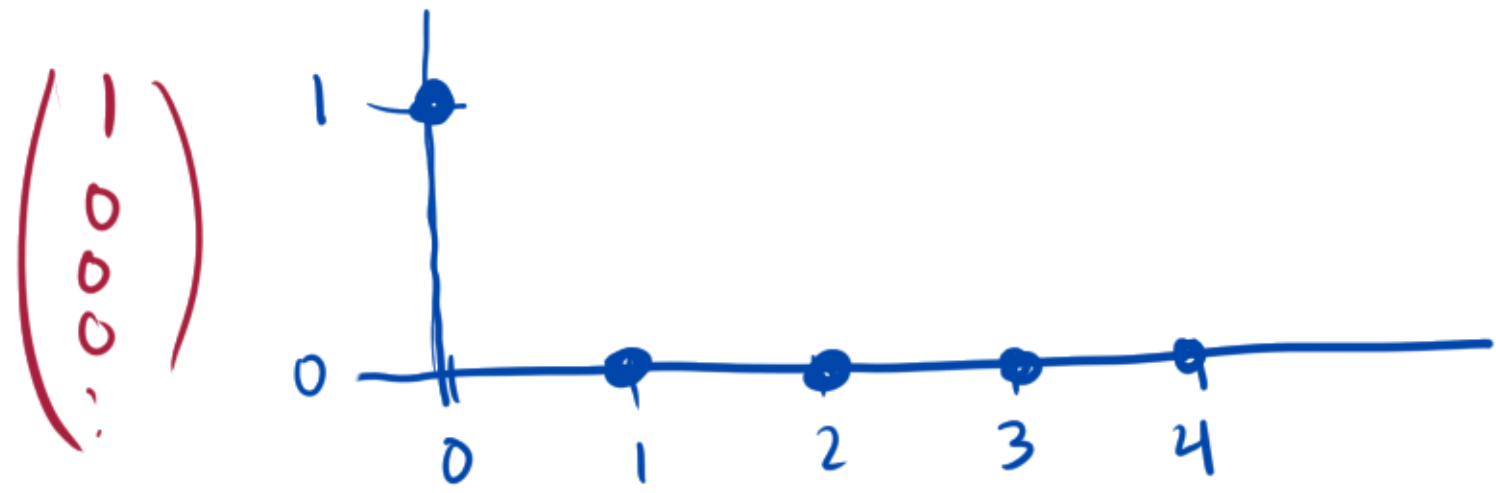
The Space of Functions  $\{ f: \mathbb{Z}/m\mathbb{Z} \rightarrow \mathbb{C} \}$  is a vector space.

Eg.  $f: \mathbb{Z}/3\mathbb{Z} \rightarrow \mathbb{C}$  given by  $f(0) = 1$   
 $f(1) = \pi$   
 $f(2) = i$

"is" the vector  $\begin{pmatrix} 1 \\ \pi \\ i \end{pmatrix}$

in "standard" basis  
 $\begin{pmatrix} 1 \\ 0 \\ 0 \end{pmatrix} \begin{pmatrix} 0 \\ 1 \\ 0 \end{pmatrix} \begin{pmatrix} 0 \\ 0 \\ 1 \end{pmatrix}$

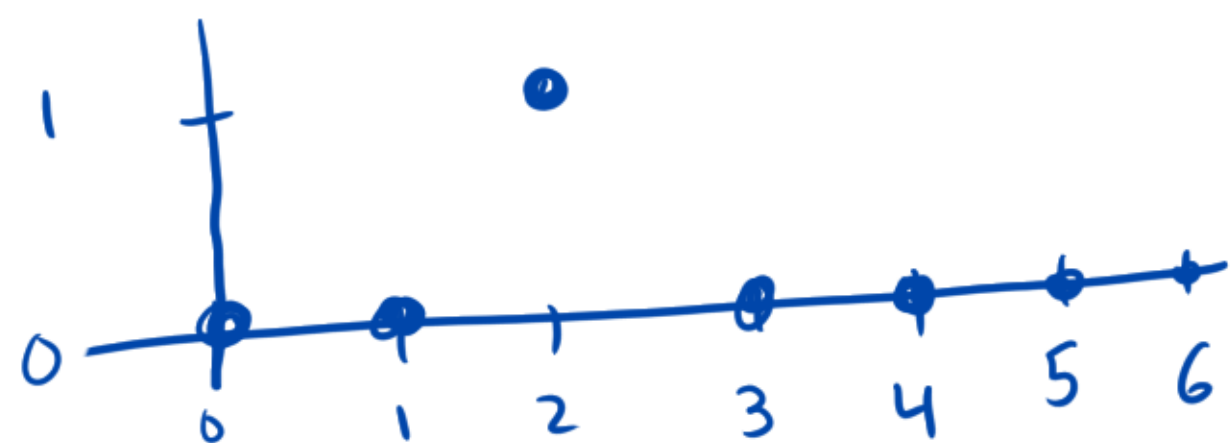
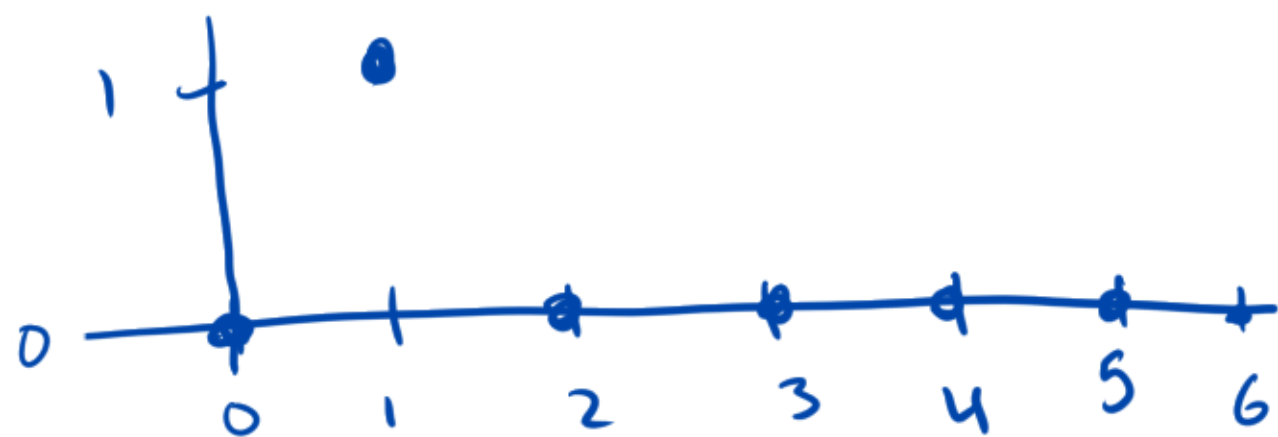
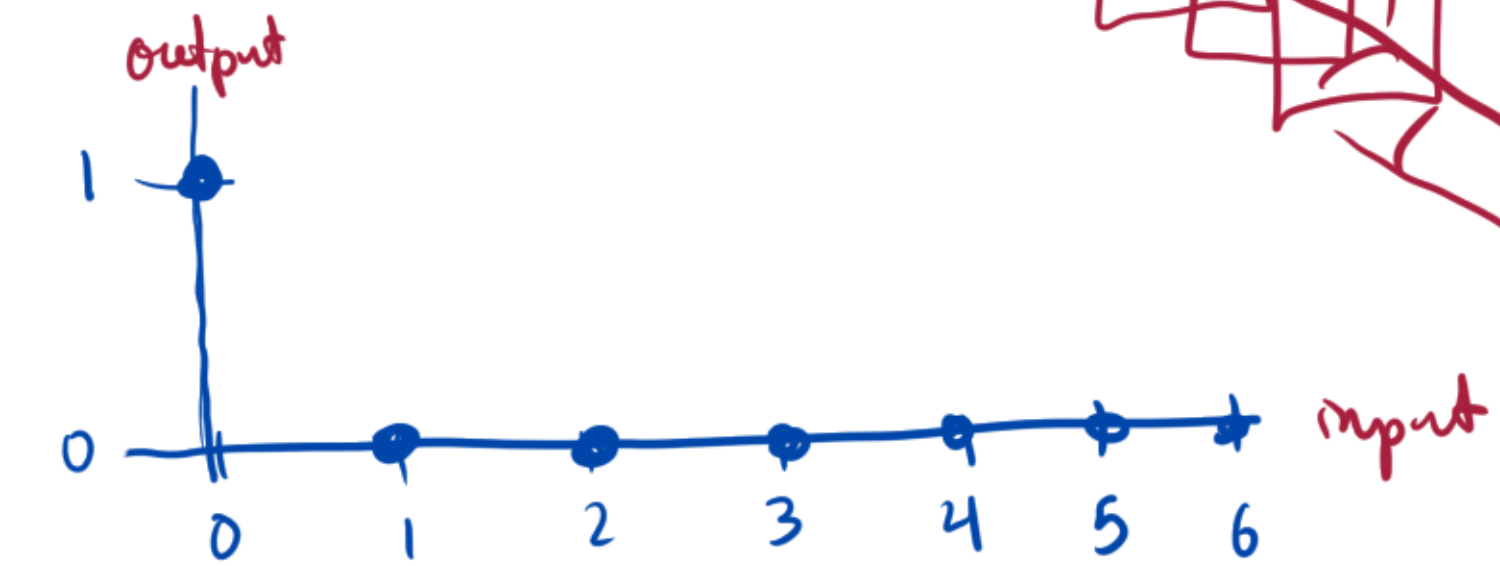
# Standard Basis



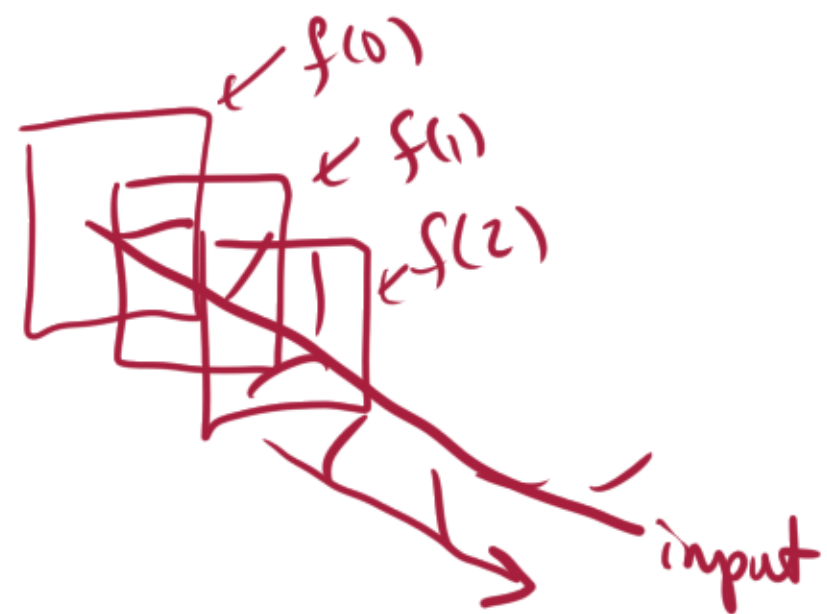
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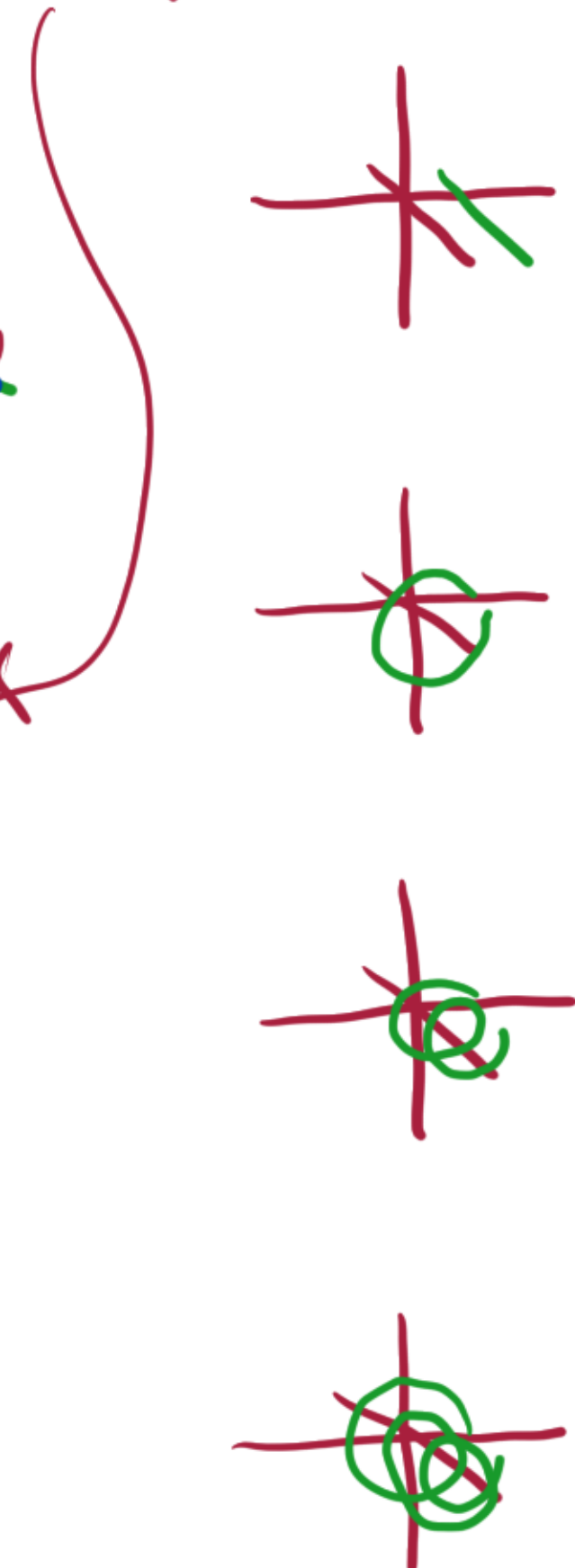
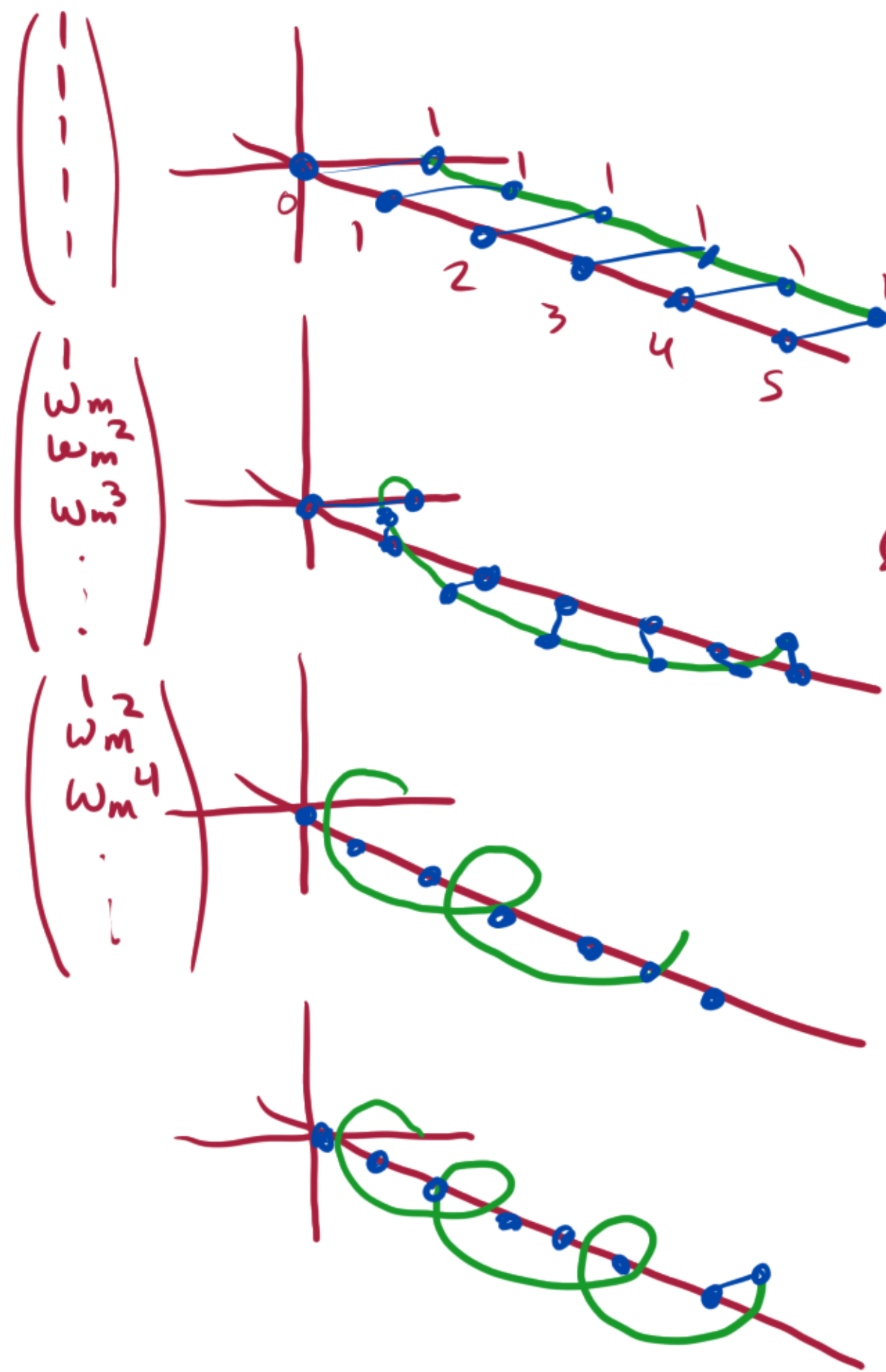
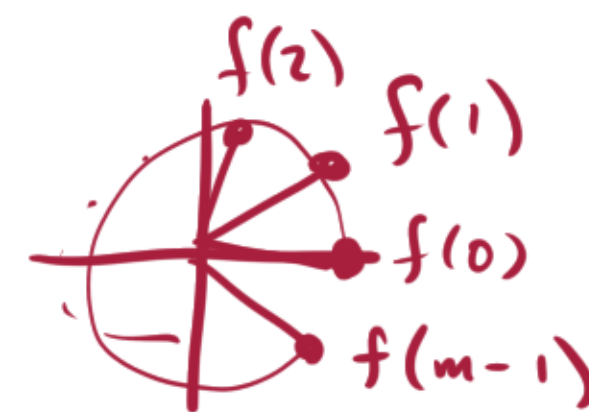
# Standard Basis

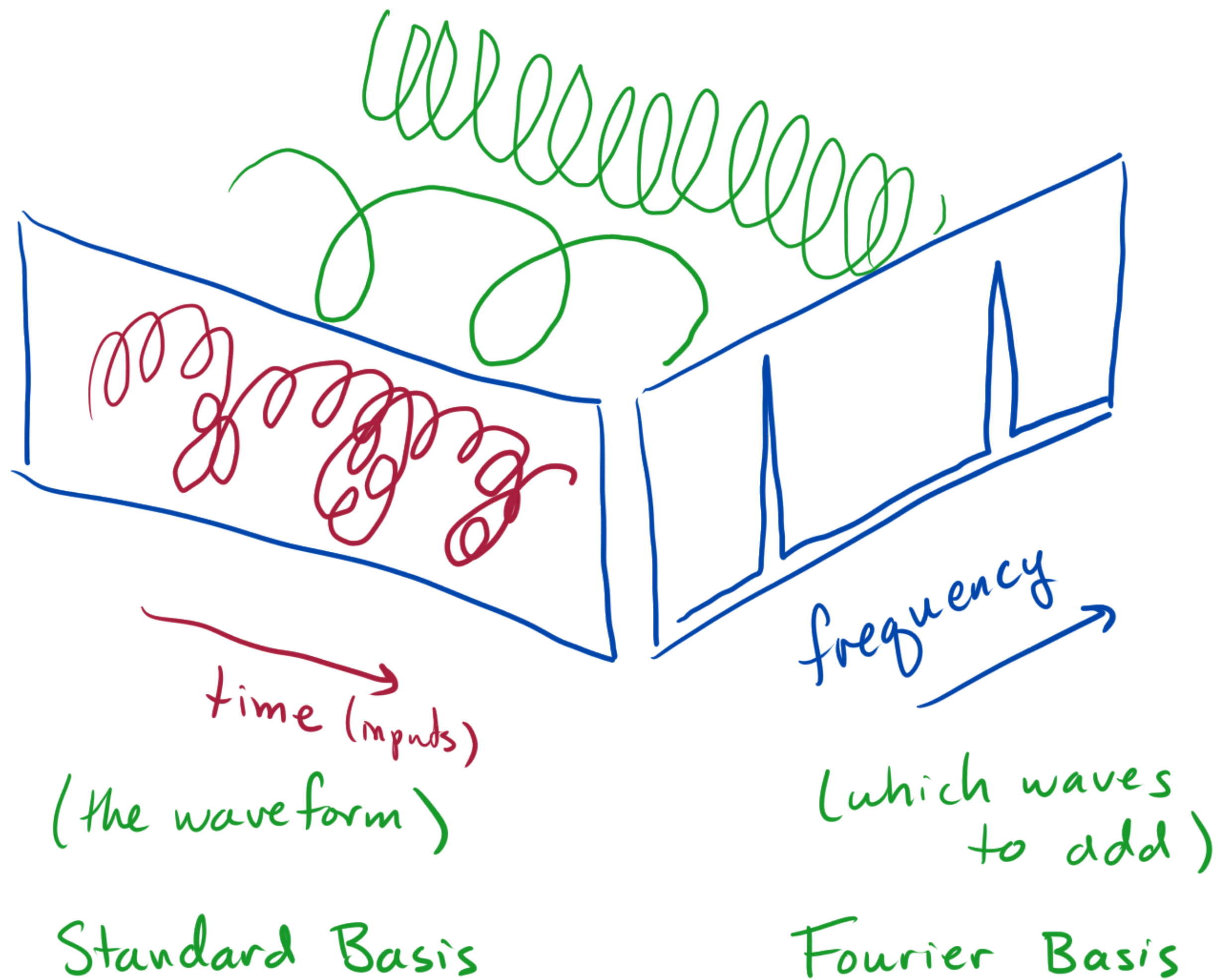


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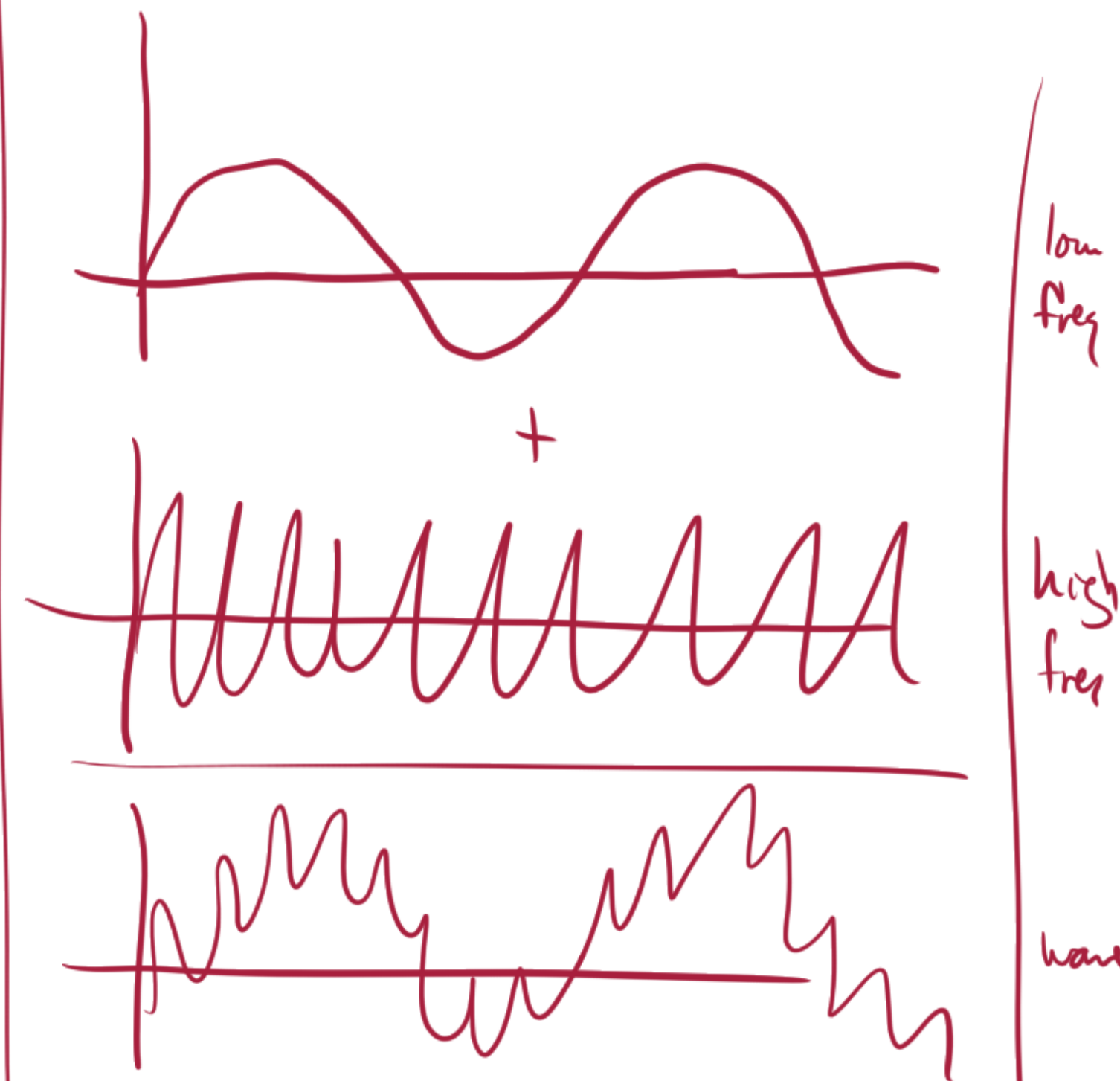


# Fourier Basis





$$(f_1 + f_2)(x) = f_1(x) + f_2(x)$$



# The Quantum Fourier Transform $m \times m$ , $\omega = \omega_m = e^{\frac{2\pi i}{m}}$

$$F_m = \frac{1}{\sqrt{m}} \begin{pmatrix} 1 & 1 & 1 & 1 & 1 & \dots & 1 \\ 1 & \omega & \omega^2 & \omega^3 & \omega^4 & \dots & \omega^{m-1} \\ 1 & \omega^2 & \omega^4 & \omega^6 & \omega^8 & \dots & \omega^{2m-2} \\ 1 & \omega^3 & \omega^6 & \omega^9 & \omega^{12} & \dots & \omega^{3m-3} \\ 1 & \omega^4 & \omega^8 & \omega^{12} & \dots & \dots & \omega^{4m-4} \\ \vdots & \vdots & \vdots & \vdots & \vdots & \dots & \vdots \\ \vdots & \vdots & \vdots & \vdots & \vdots & \dots & \vdots \\ 1 & \omega^{m-1} & \omega^{2m-2} & \omega^{3m-3} & \dots & \dots & \omega^{(m-1)m-1} \end{pmatrix}$$

$(ij)^{\text{th}}$  entry =  $\omega_m^{ij}$



# Examples.

$$\begin{pmatrix} 1 & 1 \\ 1 & w \end{pmatrix}$$

$$m=2$$

$$w_2 = -1$$

$$(-1)^2 = 1$$

$$F_2 = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix}$$

Hadamard

$$m=4$$

$$w_4 = i$$

$$i^2 = -1$$

$$i^3 = -i$$

$$(-i)^2 = -1$$

$$(-i)^3 = i$$

$$F_4 = \frac{1}{\sqrt{4}} \begin{pmatrix} 1 & 1 & 1 & 1 \\ 1 & i & -1 & -i \\ 1 & -1 & 1 & -1 \\ 1 & -i & -1 & i \end{pmatrix}$$

Input: vector  $\vec{v}$  or function  $f: \mathbb{Z}/m\mathbb{Z} \rightarrow \mathbb{C}$

Output:  $F_m \vec{v}$  or  $\hat{f}: \mathbb{Z}/m\mathbb{Z} \rightarrow \mathbb{C}$

Formula for new coefficients:

$$(F_m \vec{v})_n = \frac{1}{\sqrt{m}} \sum_{k=0}^{m-1} v_k \omega_m^{nk}$$

or

$$\hat{f}(n) = \frac{1}{\sqrt{m}} \sum_{k=0}^{m-1} f(k) \omega_m^{nk}$$

Computer  
Demo

Runtime?

Classical: discrete fourier transform via matrix multiplication

$$O(m^2)$$

Fast Fourier Transform =  $O(m \log m)$



Quantum Fourier Transform:  $F_m$  is a unitary matrix

$$\sum_{x=0}^{m-1} f(x) |x\rangle \longmapsto \sum_{x=0}^{m-1} \hat{f}(x) |x\rangle$$
$$= \frac{1}{\sqrt{m}} \sum_{x=0}^{m-1} \sum_{k=0}^{m-1} \omega_m^{xk} f(k) |x\rangle$$

Quantum Fourier Transform:  $F_m$  is a unitary matrix

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Can be implemented with  $O(\log_2^2(m))$  gates

( $H$ ) and a controlled phase gates)