

## Modular Exponentiation

Q: Given  $b, n$ , compute  $b^x \pmod{n}$  as efficiently as possible, as a function of  $x$ .

$$b^2 = b \times b$$

$$b^4 = b^2 \times b^2$$

$$b^8 = b^4 \times b^4 \quad (3 \text{ multiplications})$$

$$b^{16} = b^8 \times b^8$$

$$b^{32} = b^{16} \times b^{16}$$

$$b^3 = b^2 \times b$$

$$b^7 = b^4 \times b^3$$

$$b^{23} = b^{16} \times b^7$$

$$b^{55} = b^{32} \times b^{23}$$

Challenges.

$$b^{128}$$

$$b^2 = b \cdot b$$

$$b^4 = b^2 \cdot b^2$$

$$b^8 = b^4 \cdot b^4$$

$$b^{16} = b^8 \cdot b^8$$

$$b^{32} = b^{16} \cdot b^{16}$$

$$b^{64} = b^{32} \cdot b^{32}$$

$$b^{128} = b^{64} \cdot b^{64}$$

} 7 mult's.

$$b^{170}$$

$$128 + 32 = 160$$

$$160 + 8 = 168$$

$$168 + 2 = 170$$

10 mult's.

$$\underline{170 = 128 + 32 + 8 + 2}$$

1	0	1	0	1	0	1	0
128	64	32	16	8	4	2	1

$$n = \sum_{i=0}^{\infty} a_i 2^i, \quad a_i \in \{0, 1\}$$

↙ compute modular inverse

$$b^{127}$$

$$b^{128} \cdot b^{-1}$$

$$1 + 2 + 4 + 8 + 16 + 32 + 64 = 127$$

$$\dots \quad | \quad | \quad | \quad | \quad | \quad | \quad | \quad |$$

## Successive Squaring (alg. to compute $b^x \pmod n$ )

① Compute  $b^{2^i}$  for all  $i$  s.t.  $2^i \leq x$ , by successive squaring.

$$b \xrightarrow{sq} b^2 \xrightarrow{sq} b^4 \xrightarrow{sq} b^8 \rightarrow \dots$$

② Multiply together those which are needed to give  $x$ .

(i.e. the positions of the 1's in the binary expansion)

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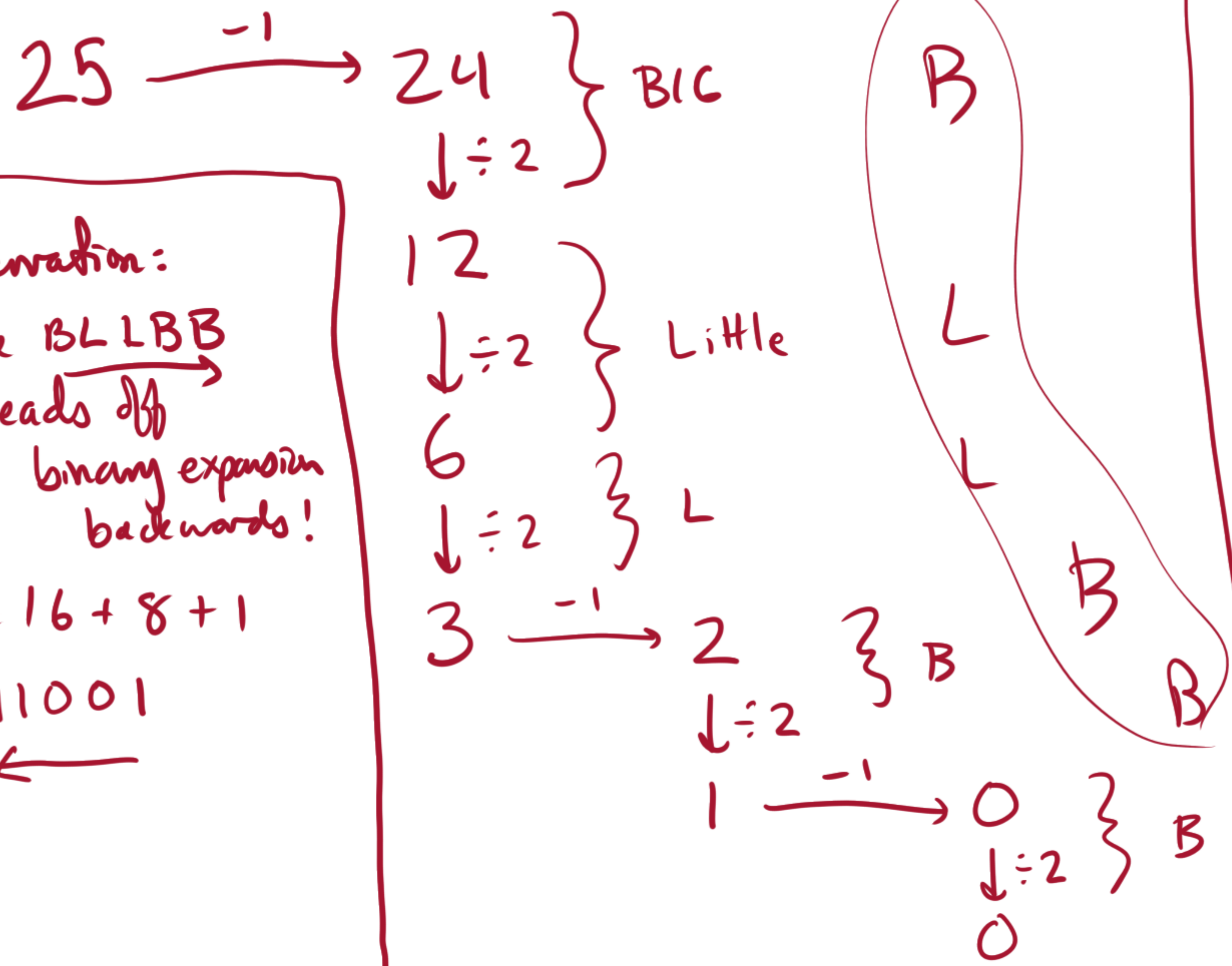
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# Square-and-multiply AKA Double-and-add.

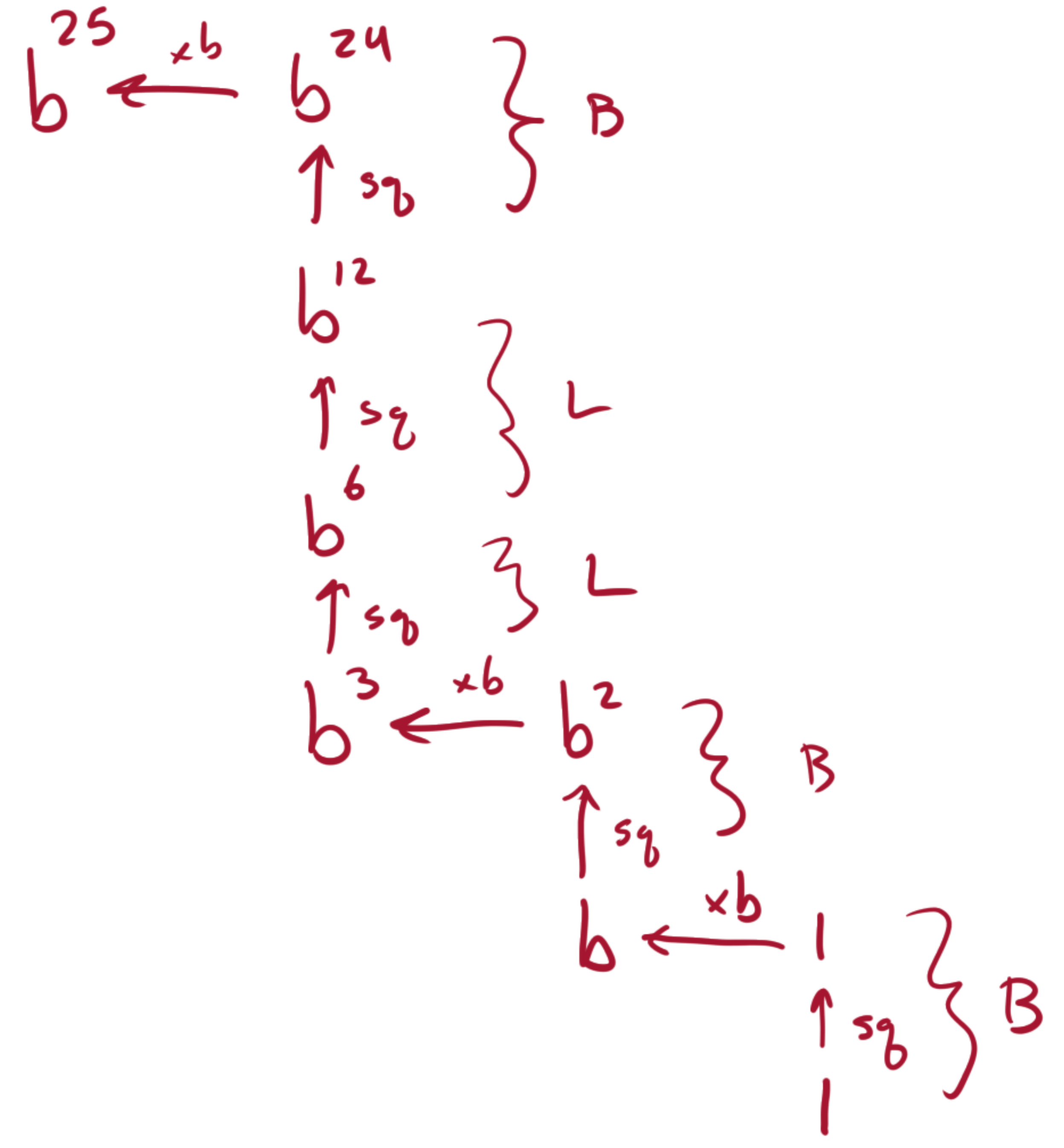
Exponent 25  
Analyse the exponent

Build  $b^x$



Observation:  
 the BLLBB  
 reads off  
 binary expansion  
 backwards!

$25 = 16 + 8 + 1$   
 11001  
 $\leftarrow$



How many multiplications?

Write  $x$  in binary  $a_{n-1} \dots a_0$  ( $n$  digits)

then  $n = \lceil \log_2(x) \rceil$  "bitlength of  $x$ "

$h =$  Hamming wt of binary expansion  $\leq n$   
 $=$  the # of 1's in " "

# multiplications is  $n-1 + h-1$   
for either algorithm

We get a runtime

of  $\sim \log(x)$  "linear"

Look ahead:

