

# Logarithmic geometry, curves, and moduli

## Lectures at the Wise course

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# Moduli of curves

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*Working with a non-complete moduli space is like keeping change in a pocket with holes*

Angelo Vistoli

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- allow only nodes as singularities
- What's so great about nodes?
- One answer: from the point of view of logarithmic geometry, these are the *logarithmically smooth* curves.



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## Definition

It is a *logarithmic structure* if  $\alpha : \alpha^{-1}\mathcal{O}_{\underline{X}}^* \rightarrow \mathcal{O}_{\underline{X}}^*$  is an isomorphism.

# Types of logarithmic structures

- We say that  $(\underline{X}, M_X)$  is *coherent* if *étale locally* at every point there is a *finitely generated* monoid  $P$  and a local chart  $P_X \rightarrow \mathcal{O}_{\underline{X}}$  for  $X$ .

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- We say that a logarithmic structure is *fine* if it is *coherent* with local charts  $P_X \rightarrow \mathcal{O}_X$  with  $P$  *integral*.
- We say that a logarithmic structure is *fine and saturated* (or *fs*) if it is coherent with local charts  $P_X \rightarrow \mathcal{O}_X$  with  $P$  *integral and saturated*.

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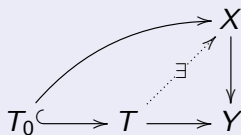
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- 2 For  $T_0$  fine and affine and  $T_0 \subset T$  strict square-0 embedding, given



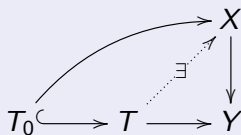
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The morphism is *logarithmically étale* if the lifting in (2) is unique.



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## Lemma

If  $X \rightarrow Y$  is *strict* and  $\underline{X} \rightarrow \underline{Y}$  *smooth* then  $X \rightarrow Y$  is *logarithmically smooth*.

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If  $X \rightarrow Y$  is strict and  $\underline{X} \rightarrow \underline{Y}$  smooth then  $X \rightarrow Y$  is logarithmically smooth.

## Proof.

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$$\begin{array}{ccccc} & & & & \underline{X} \\ & & & \nearrow & \downarrow \\ & & & \exists & \underline{Y} \\ \underline{T}_0 & \hookrightarrow & \underline{T} & \longrightarrow & \underline{Y} \end{array}$$

The diagram illustrates a commutative square. At the bottom left is  $\underline{T}_0$ , at the bottom middle is  $\underline{T}$ , and at the bottom right is  $\underline{Y}$ . A horizontal arrow points from  $\underline{T}_0$  to  $\underline{T}$ , and another from  $\underline{T}$  to  $\underline{Y}$ . A vertical arrow points from  $\underline{Y}$  up to  $\underline{X}$ . A curved arrow points from  $\underline{T}_0$  up to  $\underline{X}$ . A dotted arrow labeled with the Greek letter  $\exists$  points from  $\underline{T}$  up to  $\underline{X}$ .

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since  $\underline{X} \rightarrow \underline{Y}$  smooth, and the lifting of morphism of monoids comes by the universal property of pullback.



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## Proposition

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Then  $X \rightarrow Y$  is logarithmically smooth.

If also the cokernel is finite then  $X \rightarrow Y$  is logarithmically étale.

# Characterization of logarithmic smoothness

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Then  $X \rightarrow Y$  is *logarithmically smooth* iff there are extensions to local charts

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One direction:

$$\begin{array}{ccccc} \underline{X} & \longrightarrow & \underline{Y} \times_{\text{Spec } \mathbb{Z}[Q]} \text{Spec } \mathbb{Z}[P] & \longrightarrow & \text{Spec } \mathbb{Z}[P] \\ & & \downarrow & & \downarrow \\ & & \underline{Y} & \longrightarrow & \text{Spec } \mathbb{Z}[Q] \end{array}$$

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- the fibers are curves i.e. pure dimension 1 schemes.

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*where  $\mathbb{N} \rightarrow \mathbb{N}^2$  is the diagonal. Here the image of  $1 \in \mathbb{N}$  in  $\mathcal{O}_S$  is  $f$  and the generators of  $\mathbb{N}^2$  map to  $x$  and  $y$ .*

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# Moduli of stable log curves

We define a category  $\overline{\mathcal{M}}_{g,n}^{\log}$  of stable log curves: objects are log  $(g, n)$ -curves  $X \rightarrow S$  and arrows are fiber diagrams compatible with sections

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There is a forgetful functor

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So  $\overline{\mathcal{M}}_{g,n}^{\log}$  is a category fibered in groupoids over  $\mathcal{L}\text{og}\mathcal{S}\text{ch}^{\text{fs}}$ .

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Note that the Deligne–Knudsen–Mumford moduli stack  $\overline{\mathcal{M}}_{g,n}$  has a natural logarithmic smooth structure  $M_{\Delta_{g,n}}$  given by the boundary divisor.

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### Theorem (F. Kato)

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We also have a forgetful functor

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Note that the Deligne–Knudsen–Mumford moduli stack  $\overline{\mathcal{M}}_{g,n}$  has a natural logarithmic smooth structure  $M_{\Delta_{g,n}}$  given by the boundary divisor. As such it represents a category fibered in groupoids  $(\overline{\mathcal{M}}_{g,n}, M_{\Delta_{g,n}})$  over  $\mathcal{L}\text{ogS}\mathcal{c}\text{h}^{\text{fs}}$ .

### Theorem (F. Kato)

$$\overline{\mathcal{M}}_{g,n}^{\log} \simeq (\overline{\mathcal{M}}_{g,n}, M_{\Delta_{g,n}}).$$



# Minimality

Given a stable curve  $\underline{X} \rightarrow \underline{S}$  we define

$$X^{\min} = \underline{X} \times_{\overline{\mathcal{M}}_{g,n+1}} \overline{\mathcal{M}}_{g,n+1}^{\log} \quad \text{and} \quad S^{\min} = \underline{S} \times_{\overline{\mathcal{M}}_{g,n}} \overline{\mathcal{M}}_{g,n}^{\log}.$$

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We write

$$S^{\min} = (\underline{S}, M_{X/S}^S) \quad \text{and} \quad X^{\min} = (\underline{X}, M_{X/S}^X).$$

# Fundamental diagram

$$\begin{array}{ccccc} X & \longrightarrow & X^{\min} & \longrightarrow & \mathcal{X}_{g,n}^{\log} \\ \downarrow & & \downarrow & & \downarrow \\ S & \longrightarrow & S^{\min} & \longrightarrow & \overline{\mathcal{M}}_{g,n}^{\log} \\ & \searrow & \downarrow & & \downarrow \\ & & \underline{S} & \longrightarrow & \overline{\mathcal{M}}_{g,n} \end{array}$$

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- $\overline{\mathcal{M}}_{g,n}^{\log}$  parametrizes stable log curves over  $\mathcal{L}\text{og}\mathcal{S}\text{ch}^{\text{fs}}$
- $\overline{\mathcal{M}}_{g,n}$  parametrizes **minimal** stable log curves over  $\mathcal{S}\text{ch}$ .

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## Observation

The minimal object exists, with characteristic sheaf dual to the lattice in the corresponding space of tropical curves.