The Amazing Power of PP Constructions

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Outline

- Primitive positive constructions
- Alternative title: clones on finite domains ordered by minor-preserving maps
- 2-element case, 3-element case.
- Digraphs

Primitive Positive Constructions

Three posets on finite structures:

Primitive positive (pp) definability: $\underline{A} \leq_{\mathsf{def}} \underline{B}$ if A = B and every relation in B has a primitive positive definition in \underline{A} .

$$\exists x_1,\ldots,x_n(\psi_1 \wedge \cdots \wedge \psi_m)$$

- Primitive positive interpretations: $\underline{A} \leq_{int} \underline{B}$ if there exists $d \in \mathbb{N}$ and partial $f \colon A^d \to B$ such that preimages of relations of \underline{B} are pp-definable in \underline{A} .
- Primitive positive constructions (Barto, Pinsker, Opršal): $\underline{A} \leq_{con} \underline{B}$ if \underline{B} is homomorphically equivalent to \underline{B}' and $\underline{A} \leq_{int} \underline{B}'$.

Motivation:

- lacksquare $\leq_{\text{def}}, \leq_{\text{int}}, \leq_{\text{con}}$ preserve the complexity of CSPs.
- Bulatov'2017, Zhuk'2017: $CSP(\underline{B})$ is in P if $\underline{B} \not\leq_{con} K_3$, and is NP-hard otherwise.
- Relevant for: which CSPs are in L? NL? NC?
- Relevant not only for CSPs

Posets on clones over finite sets

 $\leq_{\mathsf{def}}, \leq_{\mathsf{int}}, \leq_{\mathsf{con}}$: transitive.

 $Pol(\underline{A})$: the clone of polymorphisms of \underline{A} .

- **1** $\underline{A} \leq_{\mathsf{def}} \underline{B} \mathsf{iff} \mathsf{Pol}(\underline{A}) \subseteq \mathsf{Pol}(\underline{B}).$
- 2 $\underline{A} \leq_{int} \underline{B}$ iff there is a clone homomorphism $\xi \colon Pol(\underline{A}) \to Pol(\underline{B})$.

$$\xi(f(g_1,\ldots,g_n)) = \xi(f)(\xi(g_1),\ldots,\xi(g_n))$$
$$\xi(\pi_i^n) = \pi_i^n$$

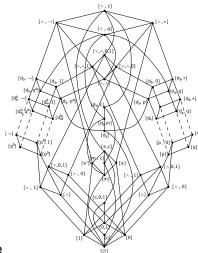
3 $\underline{A} \leq_{con} \underline{B}$ iff there is a minor-preserving map ξ : $Pol(\underline{A}) \rightarrow Pol(\underline{B})$.

$$\xi(f(\pi_{i_1}^k,\ldots,\pi_{i_k}^k)) = \xi(f)(\pi_{i_1}^k,\ldots,\pi_{i_k}^k)$$

(Every clone over a finite set equals $Pol(\underline{A})$ for some relational structure \underline{A} .)

Clones over two elements

 \leq_{def} on $\{0, 1\}$:



Post's lattice

Clones over three elements

 \leq_{def} on $\{0, 1, 2\}$:



Yanov-Muchnik: 2^ω



How about \leq_{int} ?

The interpretability poset on $\{0, 1, 2\}$



 \leq_{int} on $\{0, 1, 2\}$:

$$\begin{aligned} &C_3 := \big\{ (0,1), (1,2), (2,0) \big\} \\ &B_2 := \big\{ (1,0), (0,1), (1,1) \big\} \\ &R_3^- := \big\{ (x,y,z) \mid x \in \{0,1\} \land x = 0 \Rightarrow y = z \big\} \end{aligned}$$

Zhuk'15: 2^ω many clones between

$$\begin{array}{ccc} & \text{Pol}(\{0,1,2\}; C_3, R_3^=) \\ \text{and} & \text{Pol}(\{0,1,2\}; C_3, B_2) \end{array}$$



- Clones below $Pol(\{0, 1, 2\}; C_3)$: self-dual
- $Pol({0,1,2}; C_3, R_3^=)$ contains binary paper-scissor-stone operation

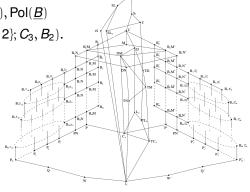
Non-collapse

Theorem. Let \underline{A} and \underline{B} be structures s.t.

$$\begin{split} (\{0,1,2\};\textit{C_3},\textit{$R_3^=$}) \leq_{\mathsf{def}} \mathsf{Pol}(\underline{\textit{A}}), \mathsf{Pol}(\underline{\textit{B}}) \\ \leq_{\mathsf{def}} (\{0,1,2\};\textit{C_3},\textit{B_2}). \end{split}$$

If $\underline{B} \leq_{int} \underline{A}$ then $\underline{B} \leq_{def} \underline{A}$.

Corollary: 2^{ω} clones over $\{0, 1, 2\}$ even when considered up to clone homomorphism equivalence!



Conclusion: Need stronger weapons.

The constructability poset

 $<_{con}$:



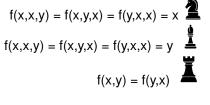
Clones \mathcal{C}_1 and \mathcal{C}_2 collapse if there is a minor-preserving map $\mathcal{C}_1 \to \mathcal{C}_2$ and a minor-preserving map $\mathcal{C}_2 \to \mathcal{C}_1$.

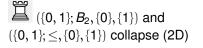
- lacktriangle all clones ${\mathcal C}$ with constant operation collapse.
- if \mathcal{C} has operation with image of size k, then \mathcal{C} collapses with a clone on k elements.
- if $\mathcal{C}^{(1)} \subseteq S_n$, then \mathcal{C} collapses with its idempotent reduct.
- consequence: to separate clones, can focus on idempotent strong linear Mal'cev conditions!

The constructability poset on {0, 1}

Pieces:

all clones without cyclic operation collapse







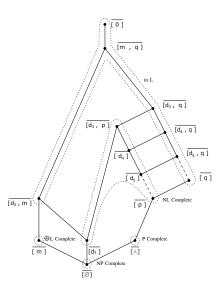


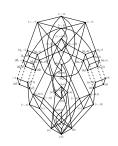




The constructability poset

 \leq_{con} on $\{0, 1\}$: outcome.





B., Vucaj 2020

The constructability poset on $\{0, 1, 2\}$



3-4 weak near unanimity

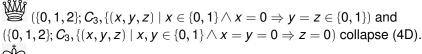
$$f(x, x, x, y) = f(x, x, y, x) = f(x, y, x, x) = f(y, x, x, x),$$

$$f(x, x, x, y) = g(x, x, y), g(x, x, y) = g(x, y, x) = g(y, x, x)$$



'guarded 3-cyclic':

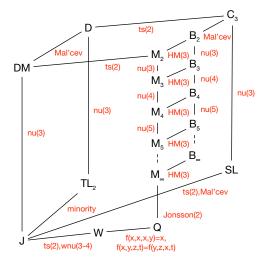
$$f(x, x, x, y) = x, f(x_1, x_2, x_3, y) = f(x_2, x_3, x_1, y)$$



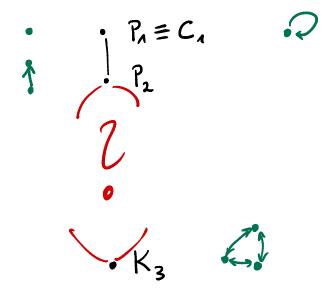
Further collapses ...

The constructability poset on $\{0, 1, 2\}$

 \leq_{con} for self-dual clones on $\{0, 1, 2\}$: outcome.



Digraphs



Digraphs: pieces

 $P_2 \leq_{con} D$ for every digraph D with a Mal'cev polymorphism and cyclic polymorphisms of all prime arities

 $D \leq_{\text{con}} T_3$ for every digraph D without a Mal'cev polymorphism $D \leq_{\text{con}} C_p$ for every digraph D without p-cyclic polymorphism

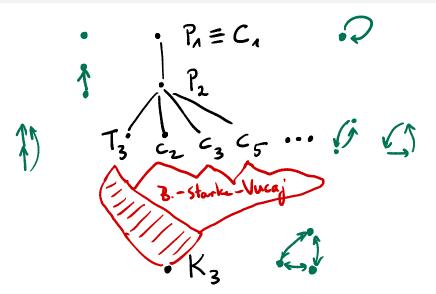
$$f(y, y, x) = f(x, y, y) = y$$

$$f(x_1, x_2) = f(x_2, x_1)$$

$$f(x_1, x_2, x_3) = f(x_2, x_3, x_1)$$

$$f(x_1, x_2, x_3, x_4) = f(x_2, x_3, x_4, x_1)$$

Digraphs: current state



Recruiting

\leq_{con} on finite structures:

- 1 What is the cardinality of \leq_{con} ?
- 2 Are there infinite ascending chains?
- 3 Is \leq_{con} a lattice?
- 4 What are the maximal elements below P_2 for general finite structures?
- 5 What are the maximal digraphs below T_3 ?