# Relative Maltsev definability (of some commutator properties) 

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PALS
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This theorem has the structure $\mathscr{P}+\Gamma=\Sigma$, i.e. the class of varieties having property $\mathscr{P}$ and satisfying the weak ground Maltsev condition $\Gamma$ is definable by the Maltsev condition $\Sigma$.

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Answer for specialists. In 2013, K+Sz+W proved Park's Conjecture for varieties with a difference term. (A finitely generated variety with a finite residual bound is finitely based.) The proof uses properties of the TC-commutator to establish a version the Freese-McKenzie property (C1). To extend the result, it would help to understand TC-commutator arithmetic for varieties that have a Taylor term but not a difference term.

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\begin{array}{ll}
\exists t \quad & t(x, x, x, x, x, x) \approx x \quad \& \\
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An (ordinary) Maltsev condition is a sequence $\Sigma=\left(\sigma_{n}\right)_{n \in \omega}$ of successively weaker strong Maltsev conditions $\left(\forall n\left(\sigma_{n} \vdash \sigma_{n+1}\right)\right.$ ). A variety $\mathcal{V}$ satisfies $\Sigma$ if it satisfies $\sigma_{n}$ for some $n$. $\Sigma$ defines the class of those varieties which satisfy $\Sigma$.

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\langle t \mid t(x x x x x x) \approx x, t(x y y y x x) \approx t(y x y x y x) \approx t(y y x x x y)\rangle
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The interpretability relation $\leq$ is a lattice order on bi-interpretability classes of varieties.

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Definition. ( $\alpha, \beta$-matrices) If $\mathbf{A}$ is an algebra and $\alpha, \beta \in \operatorname{Con}(\mathbf{A})$, then an $\alpha, \beta$-matrix is a $2 \times 2$ matrix of elements of $\mathbf{A}$ of the form

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\left[\begin{array}{ll}
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r & s
\end{array}\right]=\left[\begin{array}{ll}
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where $t(\mathbf{x}, \mathbf{y})$ is an $(m+n)$-ary term operation of $\mathbf{A}, \mathbf{a} \alpha \mathbf{b}$, and $\mathbf{u} \beta \mathbf{v}$.

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Definition. (The centralizer) $\alpha$ centralizes $\beta$ modulo $\delta$ if $p \equiv q(\bmod \delta)$ implies $r \equiv s(\bmod \delta)$ whenever $\left[\begin{array}{ll}p & q \\ r & s\end{array}\right]$ is an $\alpha, \beta$-matrix. Write $\mathbf{C}(\alpha, \beta ; \delta)$ to denote this.

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## Remark.

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Definition. (The commutator) Let $[\alpha, \beta]$ be the least $\delta$ for which $\mathbf{C}(\alpha, \beta ; \delta)$ holds.

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follows from the definition of the commutator. For groups, the converse implication also holds. (The centralizer and the commutator carry the same information.)

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- But $\mathcal{S E} \mathcal{T}$ interprets into any variety, and
- each of the commutator properties on the preceding slide fails in some variety. (E.g. the variety of semigroups.)


## Results \#1

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