Relative Maltsev definability (of some commutator properties)

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PALS September 27, 2022

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A Maltsev definable property

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This theorem has the structure $\mathscr{P} + \Gamma = \Sigma$, i.e. the class of varieties having property \mathscr{P} and satisfying the weak ground Maltsev condition Γ is definable by the Maltsev condition Σ .

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Answer for specialists. In 2013, K+Sz+W proved Park's Conjecture for varieties with a difference term. (*A finitely generated variety with a finite residual bound is finitely based.*) The proof uses properties of the TC-commutator to establish a version the Freese-McKenzie property (C1). To extend the result, it would help to understand TC-commutator arithmetic for varieties that have a Taylor term but not a difference term.

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$$\exists t \qquad \begin{array}{l} t(x,x,x,x,x,x) \approx x \quad \& \\ t(x,y,y,y,x,x) \approx t(y,x,y,x,y,x) \approx t(y,y,x,x,x,y). \end{array}$$

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An (ordinary) Maltsev condition is a sequence $\Sigma = (\sigma_n)_{n \in \omega}$ of successively weaker strong Maltsev conditions ($\forall n(\sigma_n \vdash \sigma_{n+1})$). A variety \mathcal{V} satisfies Σ if it satisfies σ_n for some n. Σ defines the class of those varieties which satisfy Σ .

Maltsev Definability, Part 2

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The clone of a variety \mathcal{V} satisfies the clone sentence σ from the previous slide iff there is a clone homomorphism $\operatorname{Clo}(\mathcal{O}) \to \operatorname{Clo}(\mathcal{V})$. Write $\mathcal{O} \leq \mathcal{V}$.

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The interpretability relation \leq is a lattice order on bi-interpretability classes of varieties.

Maltsev definability visually

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A strong Maltsev filter

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$$\begin{bmatrix} p & q \\ r & s \end{bmatrix} = \begin{bmatrix} t(\mathbf{a}, \mathbf{u}) & t(\mathbf{a}, \mathbf{v}) \\ t(\mathbf{b}, \mathbf{u}) & t(\mathbf{b}, \mathbf{v}) \end{bmatrix}$$

where $t(\mathbf{x}, \mathbf{y})$ is an (m + n)-ary term operation of \mathbf{A} , $\mathbf{a} \alpha \mathbf{b}$, and $\mathbf{u} \beta \mathbf{v}$.

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The (TC-)commutator, Part 2

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The variety of sets in the empty language has the property that for all $\mathbf{A} \in S\mathcal{ET}$ and all $\alpha, \beta, \delta \in Con(\mathbf{A})$ it is the case that $\mathbf{C}(\alpha, \beta; \delta)$ holds and $[\alpha, \beta] = 0$ holds.

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Results #1

Theorem.

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Results #2



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