

Relational coloring of varieties containing a cube term

Matthew Moore

Vanderbilt University

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Coloring varieties containing a cube term

- 1 Coloring
- 2 Clone \mathcal{L} -homomorphisms
- 3 Linear interpretability
- 4 Further directions

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Cube terms

An **n -dimensional cube term** for an algebra \mathbb{A} is a term $t(\dots)$ such that $\mathbb{A} \models t(u_1, \dots, u_m) \approx \bar{x}$ for some $u_i \in \{x, y\}^n \setminus \{x\}^n$.

Example

$$\mathbb{A} \models t \begin{pmatrix} y & x & x & y & y & x & y \\ x & y & x & y & x & y & y \\ x & x & y & x & y & y & y \end{pmatrix} \approx \begin{pmatrix} x \\ x \\ x \end{pmatrix}$$

Finite idempotent \mathbb{A} has a cube term...

$\Leftrightarrow \mathbb{A}$ has few subpowers: $\log_2 |\mathbf{S}(\mathbb{A}^n)| \in \mathcal{O}(n^k)$.

$\Leftrightarrow \mathbb{A}$ is congruence modular and finitely related.

$\Leftrightarrow \mathbb{A}$ has no cube term blockers: $\mathbb{D} < \mathbb{C} \leq \mathbb{A}$ with $C^n \setminus (C \setminus D)^n \leq \mathbb{A}^n$.

$\Rightarrow \text{CSP}(\mathbb{A})$ is tractable.

Let \mathcal{A} be a clone and $\mathbb{B} = (B; (R_j)_{j \in J})$ be a relational structure.

Let $F^{\mathcal{A}}(B)$ be the free algebra in $\mathcal{V}(\langle \mathcal{A}; \mathcal{A} \rangle)$ with generators B .

For $R_i \in (R_j)_{j \in J}$ let

$$R_i^{\mathcal{A}} = \text{closure of } \{(b_1, \dots, b_n) \in F^{\mathcal{A}}(B)^n \mid (b_1, \dots, b_n) \in R_i\} \text{ under } \mathcal{A}$$

Let $\mathbb{F}^{\mathcal{A}}(\mathbb{B}) = (F^{\mathcal{A}}(B); (R_j^{\mathcal{A}})_{j \in J})$.

A **coloring** of \mathcal{A} by \mathbb{B} is a relational homomorphism

$$c : \mathbb{F}^{\mathcal{A}}(\mathbb{B}) \rightarrow \mathbb{B} \quad \text{such that} \quad c(b) = b.$$

We say that \mathcal{A} is **\mathbb{B} -colorable**. We say that \mathcal{V} is **\mathbb{B} -colorable** if $\text{Clo}(\mathcal{V})$ is.

What is it good for?

Theorem (Sequeira)

\mathcal{V} is congruence k -permutable for some k iff \mathcal{V} is not $(\{0, 1\}; \leq)$ -colorable.

Theorem (Sequeira)

Let $D = \{1, 2, 3, 4\}$, $\alpha = 12|34$, $\beta = 13|24$, $\gamma = 12|3|4$, and $\mathbb{D} = (D; \alpha, \beta, \gamma)$. \mathcal{V} is congruence modular iff \mathcal{V} is not \mathbb{D} -colorable.

NB: Sequeira calls this “compatibility with projections”.

Theorem

Let $\emptyset \neq C \subsetneq B \subseteq A$ be sets, $R_n = B^n \setminus (B \setminus C)^n$, and $\mathbb{A} = (A; (R_n)_{n \in \omega})$.
Idempotent \mathcal{V} has a cube term iff \mathcal{V} is not \mathbb{A} -colorable.

Theorem

Let $R_n = (\omega^2)^n \setminus (\Delta_2)^n$ and $\mathbb{W} = (\omega; (R_n)_{n \in \omega})$.
 \mathcal{V} has a weak Taylor term iff \mathcal{V} is not \mathbb{W} -colorable.

Theorem

Let $\emptyset \neq C \subsetneq B \subseteq A$ be sets, $R_n = B^n \setminus (B \setminus C)^n$, and $\mathbb{A} = (A; (R_n)_{n \in \omega})$.
Idempotent \mathcal{V} has a cube term iff \mathcal{V} is not \mathbb{A} -colorable.

Proof.

(\Rightarrow): Suppose \mathcal{V} has an n -dimensional cube term and is \mathbb{A} -colorable.

There is some identity $\mathcal{V} \models t(u_1, \dots, u_n) \approx \bar{x}$ for some $u_i \in \{x, y\}^n \setminus \{x\}^n$.

Let $c \in C$ and $b \in B \setminus C$.

Substitute b for x and c for y to get $t(u_1, \dots, u_n) = \bar{b}$ and $u_i \in R_n^{\mathcal{V}}$.

Thus $\bar{b} \in R_n^{\mathcal{V}}$. Then $c(\bar{b}) = \bar{b} \in R_n \subseteq A^n$, a contradiction.

Proof (cont.).

(\Leftarrow): Suppose \mathcal{V} does not have a cube term.

Write $g \prec_t h$ if we have $u_i \in \{g, h\}^n \setminus \{g\}^n$ such that $t(u_1, \dots, u_n) = \bar{g}$.

Fix some $c_0 \in C$ and define $c : \mathbb{F}^{\text{Clo}(\mathcal{V})}(A) \rightarrow \mathbb{A}$ by

$$c(f) = \begin{cases} a & \text{if } f = a \in A \subseteq F^{\text{Clo}(\mathcal{V})}(A), \\ b & \text{else if } \exists b \in B, t(\dots) \text{ such that } b \prec_t f, \\ c_0 & \text{else.} \end{cases}$$

Clearly $c(a) = a$ for $a \in A$. Does c preserve relations?

Proof (cont.).

$$c(f) = \begin{cases} a & \text{if } f = a \in A \subseteq F^{\text{Clo}(\mathcal{V})}(A), \\ b & \text{else if } \exists b \in B, t(\dots) \text{ such that } b \prec_t f, \\ c_0 & \text{else.} \end{cases}$$

Suppose that c fails to preserve $R_n^{\text{Clo}(\mathcal{V})}$. Then there exists

- $(f_1, \dots, f_n) \in R_n^{\text{Clo}(\mathcal{V})}$,
- a term $s(\dots)$,
- and $u_1, \dots, u_k \in B^n \setminus (B \setminus C)^n \subseteq F^{\text{Clo}(\mathcal{V})}(A)$,

such that

- $s(u_1, \dots, u_k) = (f_1, \dots, f_n)$
- and $c(f_i) = b_i \in (B \setminus C)$.

Thus $b_i \prec_{t_i} f_i$. Let $t = t_1 * t_2 * \dots * t_n * s$.

Then $\exists v_1, \dots, v_m \in B^n \setminus (B \setminus C)^n$ such that $t(v_1, \dots, v_m) = (b_1, \dots, b_n)$.

This is a free algebra, so substitute x for $B \setminus C$ and y for C to obtain a cube identity for $t(\dots)$, a contradiction. □

A corollary

Theorem

Let $\emptyset \neq C \subsetneq B \subseteq A$ be sets, $R_n = B^n \setminus (B \setminus C)^n$, and $\mathbb{A} = (A; (R_n)_{n \in \omega})$.
Idempotent \mathcal{V} has a cube term iff \mathcal{V} is not \mathbb{A} -colorable.

Let $A = B = \{0, 1\}$ and $C = \{0\}$.

Then the polymorphism clone of \mathbb{A} is generated by $\hat{\mathcal{D}} = \langle \{0, 1\}; \nearrow \rangle$:
 $\text{Pol}(\mathbb{A}) = \text{Clo}(\hat{\mathcal{D}})$.

Corollary

Idempotent \mathcal{V} has a cube term iff \mathcal{V} is not $\text{Rel}(\hat{\mathcal{D}})$ -colorable.

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Clone homomorphisms

Definition

A homomorphism between clones \mathcal{A} and \mathcal{B} is an arity-preserving mapping $\varphi : \mathcal{A} \rightarrow \mathcal{B}$ such that

$$\varphi(\pi_k) = \pi_k \quad \text{and} \quad \varphi(f(g_1, \dots, g_n)) = \varphi(f)(\varphi(g_1), \dots, \varphi(g_n)).$$

An \mathcal{L} -homomorphism between clones \mathcal{A} and \mathcal{B} is an arity preserving mapping $\varphi : \mathcal{A} \rightarrow \mathcal{B}$ such that

$$\varphi(\pi_k) = \pi_k \quad \text{and} \quad \varphi(f(\pi_{i_1}, \dots, \pi_{i_n})) = \varphi(f)(\pi_{i_1}, \dots, \pi_{i_n}).$$

Let \mathcal{A} and \mathcal{B} be clones, and let $\varphi : \mathcal{A} \rightarrow \mathcal{B}$ be an arity preserving map.

Let $\mathbb{A} = (A; \mathcal{A})$, and $\mathbb{B} = (B; \mathcal{B})$.

φ is a clone homomorphism iff $f \approx g$ in \mathbb{A} implies $\varphi(f) \approx \varphi(g)$ in \mathbb{B} .

φ is a clone \mathcal{L} -homomorphism iff every identity $f \approx g$ in \mathbb{A} not involving composition implies $\varphi(f) \approx \varphi(g)$ in \mathbb{B} .

Example

Consider $\mathfrak{B} = \langle \{0, 1\}; \rightarrow \rangle$ and $\hat{\mathfrak{B}} = \langle \{0, 1\}; \nrightarrow \rangle$.

No clone homomorphisms between $\text{Clo}(\hat{\mathfrak{B}})$ and $\text{Clo}(\mathfrak{B})$ since \rightarrow cannot be defined as a term in \nrightarrow and \nrightarrow cannot be defined as a term in \rightarrow .

This is also witnessed by the identities

$$\mathfrak{B} \models y \rightarrow (x \rightarrow x) \approx (x \rightarrow x), \quad \hat{\mathfrak{B}} \models (x \nrightarrow x) \nrightarrow y \approx (x \nrightarrow x).$$

Given a function $f : \{0, 1\}^n \rightarrow \{0, 1\}$ define

$\delta(f(x_1, \dots, x_n)) = \neg f(\neg x_1, \dots, \neg x_n)$, where \neg is boolean negation.

$\delta : \text{Clo}(\mathfrak{B}) \rightarrow \text{Clo}(\hat{\mathfrak{B}})$ is a clone \mathcal{L} -homomorphism, but not a clone homomorphism.

Let $s(x, y, z) = x \rightarrow (y \rightarrow z)$ and $s'(x, y, z) = (x \nrightarrow y) \nrightarrow z$.

$$\delta(s) = s' \text{ and } \delta(s') = s.$$

δ translates the identities witnessing the non-existence of a clone homomorphism.

Theorem (Barto, Oprsal, Pinsker)

Let \mathcal{A} be a clone and \mathbb{B} a relational structure with polymorphism clone \mathcal{B} . \mathcal{A} is $\text{Rel}(\mathbb{B})$ -colorable iff there is a clone \mathcal{L} -homomorphism $\mathcal{A} \rightarrow \mathcal{B}$.

Corollary

Let \mathcal{V} be an idempotent variety. The following are equivalent.

- \mathcal{V} does not have a cube term,*
- there is a clone \mathcal{L} -homomorphism $\text{Clo}(\mathcal{V}) \rightarrow \text{Clo}(\hat{\mathbb{2}})$.*

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Interpretability

Let \mathcal{V} and \mathcal{W} be varieties and write $\mathcal{V} \leq \mathcal{W}$ if there is a clone homomorphism $\text{Clo}(\mathcal{V}) \rightarrow \text{Clo}(\mathcal{W})$. Say that \mathcal{V} is **interpretable into** \mathcal{W} .

Define an equivalence relation by $\mathcal{V} \equiv \mathcal{W}$ if $\mathcal{V} \leq \mathcal{W} \leq \mathcal{V}$.

Modulo this equivalence relation, under the \leq order the class of all varieties is a lattice. This is the **lattice of interpretability types**.

Write $\mathcal{V} \leq_{\mathcal{L}} \mathcal{W}$ if there is a clone \mathcal{L} -homomorphism $\text{Clo}(\mathcal{V}) \rightarrow \text{Clo}(\mathcal{W})$. Say that \mathcal{V} is **\mathcal{L} -interpretable into** \mathcal{W} .

Define an equivalence relation by $\mathcal{V} \equiv_{\mathcal{L}} \mathcal{W}$ if $\mathcal{V} \leq_{\mathcal{L}} \mathcal{W} \leq_{\mathcal{L}} \mathcal{V}$.

Example

$$\mathcal{V}(\mathfrak{2}) \equiv_{\mathcal{L}} \mathcal{V}(\hat{\mathfrak{2}})$$

Let \mathcal{V} be a variety with signature σ .

Replace σ by $\text{Clo}(\mathcal{V})$, call the resulting variety \mathcal{V}_c so that $\mathcal{V}_c \equiv \mathcal{V}$.

Let $\Sigma \subseteq \text{Th}(\mathcal{V}_c)$ be the subset of all identities not involving composition.

Let $\mathcal{V}_{\mathcal{L}} = \text{Mod}(\Sigma)$.

Example

If \mathcal{V} has an identity $t(s(x, y), x, y) \approx s(x, y)$, then \mathcal{V}_c will have a new operation symbol $f(x, y, z, w)$ and $\text{Th}(\mathcal{V}_c)$ will have identities

- $t(s(x, y), x, y) \approx s(x, y)$,
- $f(x, y, z, w) \approx t(s(x, y), z, w)$,
- $f(x, y, x, y) \approx s(x, y)$.

Out of these, Σ will contain **only** the last one.

Example

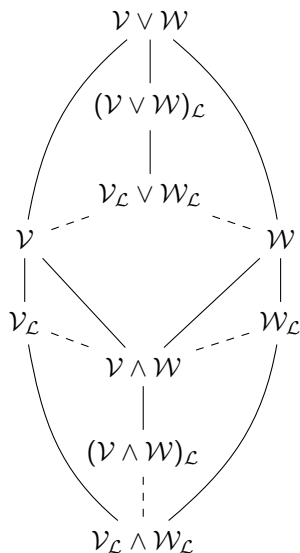
$(x \nrightarrow x) \nrightarrow y \approx (x \nrightarrow x)$ in $\hat{\mathcal{D}}$ becomes $s(x, x, y) \approx x \nrightarrow x$ in $\mathcal{V}(\hat{\mathcal{D}})_{\mathcal{L}}$.

Theorem

Let \mathcal{V} and \mathcal{W} be varieties. Then

- $\mathcal{V} \equiv_{\mathcal{L}} \mathcal{V}_{\mathcal{L}}$,
- $\mathcal{V}_{\mathcal{L}} \leq \mathcal{W}_{\mathcal{L}}$ iff $\mathcal{V} \leq_{\mathcal{L}} \mathcal{W}$,
- if $\mathcal{V} \leq \mathcal{W}$ then $\mathcal{V}_{\mathcal{L}} \leq \mathcal{W}_{\mathcal{L}}$.
- \mathcal{V} is defined by linear identities iff $\mathcal{V} \equiv \mathcal{V}_{\mathcal{L}}$.

Usual interpretability lattice:

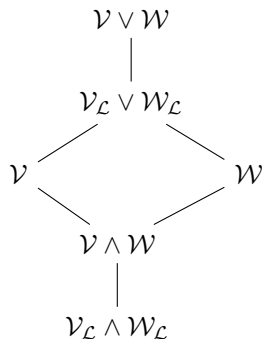


Theorem

Let \mathcal{V} and \mathcal{W} be varieties. Then

- $\mathcal{V} \equiv_{\mathcal{L}} \mathcal{V}_{\mathcal{L}}$,
- $\mathcal{V}_{\mathcal{L}} \leq \mathcal{W}_{\mathcal{L}}$ iff $\mathcal{V} \leq_{\mathcal{L}} \mathcal{W}$,
- if $\mathcal{V} \leq \mathcal{W}$ then $\mathcal{V}_{\mathcal{L}} \leq \mathcal{W}_{\mathcal{L}}$.
- \mathcal{V} is defined by linear identities iff $\mathcal{V} \equiv \mathcal{V}_{\mathcal{L}}$.

\mathcal{L} -interpretability lattice:



Corollary

Let \mathcal{V} be an idempotent variety. The following are equivalent.

- \mathcal{V} does not have a cube term,
- $\mathcal{V} \leq_{\mathcal{L}} \mathcal{V}(\hat{\mathfrak{D}})$ in the \mathcal{L} -interpretability lattice.
- $\mathcal{V}_{\mathcal{L}} \leq \mathcal{V}(\hat{\mathfrak{D}})_{\mathcal{L}}$ in the interpretability lattice.

Theorem (Sequeira; Barto, Oprsal, Pinsker)

If \mathcal{V} and \mathcal{W} are defined by linear identities and are \mathbb{A} -colorable, then $\mathcal{V} \vee \mathcal{W}$ is \mathbb{A} -colorable as well.

Corollary

If idempotent \mathcal{V} and \mathcal{W} are defined by linear identities and $\mathcal{V} \vee \mathcal{W}$ has a cube term, then one of \mathcal{V} or \mathcal{W} does.

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Question

Is having a cube term join-prime in the idempotent interpretability lattice?

Answer is tentatively “yes”.

Question

Is there a characterization of when $(\mathcal{V} \vee \mathcal{W})_{\mathcal{L}} \equiv \mathcal{V}_{\mathcal{L}} \vee \mathcal{W}_{\mathcal{L}}$?

In the class of idempotent varieties satisfying $(\mathcal{V} \vee \mathcal{W})_{\mathcal{L}} \equiv \mathcal{V}_{\mathcal{L}} \vee \mathcal{W}_{\mathcal{L}}$, both being congruence modular and having a cube term are join prime.

Question

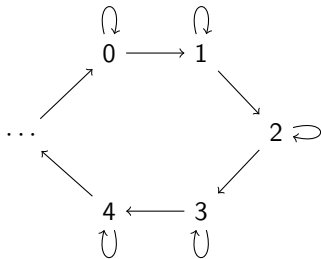
Does every \vee -prime filter in the \mathcal{L} -interpretability lattice come from a coloring characterization?

Question

Is there a coloring characterization for k -permutability for fixed k ?

We know that \mathcal{V} is not k -permutable for any k iff \mathcal{V} is $(\{0, 1\}; \leq)$ -colorable.

Let $\mathbb{P}_k = (\{0, \dots, k-1\}; \rightarrow)$ where \rightarrow is defined by



Then \mathcal{V} is congruence $(k-1)$ -permutable iff it is not \mathbb{P}_k -colorable.

Let $\hat{\mathbf{2}} = (\{0, 1\}; (\{0, 1\}^n \setminus \{1\}^n)_{n \in \omega})$ and $\hat{\mathbf{2}} = \langle \{0, 1\}; \nrightarrow \rangle$.

Theorem

For idempotent \mathcal{V} the following are equivalent.

- \mathcal{V} has no cube term,
- \mathcal{V} is $\hat{\mathbf{2}}$ -colorable,
- $\mathcal{V} \leq_{\mathcal{L}} \mathcal{V}(\hat{\mathbf{2}})$,
- $\mathcal{V}_{\mathcal{L}} \leq \mathcal{V}(\hat{\mathbf{2}})_{\mathcal{L}}$.

If $\mathcal{V} = \mathcal{V}(\mathbb{A})$ for finite \mathbb{A} , then these are equivalent to

- *there exists $\mathbb{C} < \mathbb{D} \leq \mathbb{A}$ such that $D^n \setminus (D \setminus C)^n \leq \mathbb{A}^n$ for all n (Markovic, Maroti, McKenzie).*

Question

Without repeating the proof of MMM, can we obtain the finite algebra result from the other results?

Conclusion

Let $\hat{\mathbf{2}} = (\{0, 1\}; (\{0, 1\}^n \setminus \{1\}^n)_{n \in \omega})$ and $\hat{\mathcal{2}} = \langle \{0, 1\}; \nrightarrow \rangle$.

Theorem

For idempotent \mathcal{V} the following are equivalent.

- \mathcal{V} has no cube term,
- \mathcal{V} is $\hat{\mathbf{2}}$ -colorable,
- $\mathcal{V} \leq_{\mathcal{L}} \mathcal{V}(\hat{\mathcal{2}})$,
- $\mathcal{V}_{\mathcal{L}} \leq \mathcal{V}(\hat{\mathcal{2}})_{\mathcal{L}}$.

Thank you.